

SQL

Class notes

By
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**SRI RAGHAVENDRA
XEROX**

Software languages Material Available

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Introduced in
2003. ← Oracle 11g → greedy Technology

DBMS

— oracle also RDBMS s/w files. RS = BC

— Services all are works on Port no. 1521
(for alter)

→ data store → Permanently in Secondary storage devices. (Hard disk).

— These called dbms s/w.

← To store data

Permanently is Secondary Storage

device (Hard disk)

Oracle d.b.
SQL (language) PLSQL. Dynamic SQL

↓
Structured Query language (Procedure language extension for sql)

→ DBMS S/W → Relation DB management System

— Data :- Collection of facts.

Student marks → 90 80 90 20 0.

Collection of nos.

names - Shekhar

Collection strings.

— Information:-

— Data store:- Data store is DBMS s/w,

→

SOL
↓
isql
ASCII language

→ Every org'n deals with lot of data

Data: Collection of raw facts.

Ex:- (1) Student marks

(2) Customer names

Sequel

data

is high

Information: When we are processing data we achieve

meaningful results.

Ex:- (1) marksheet.

(2) Invoice of a customer.

Data Store: It is a place where we can store

data.

(1) books & papers.

D.E

(2) flat files.

E

(3) Data base.

→ w

flat files: This is a traditional mechanism which is used to store data or information in files

- Only

Interface

with

Programs

appn

Disadvantages:-

(1) data security

(1) data retrieval

(2) data indexing

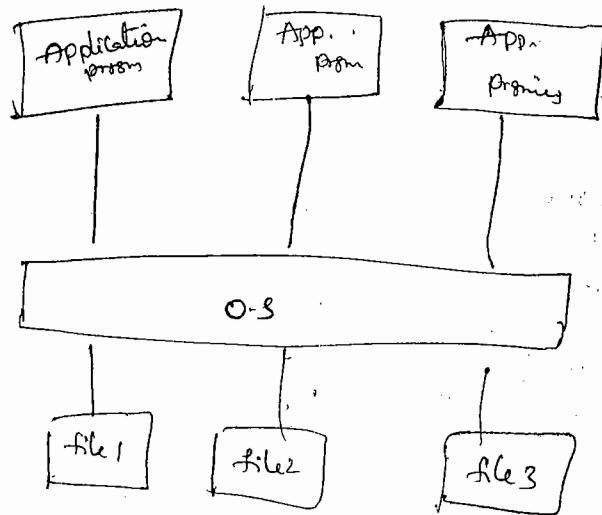
(2) data redundancy

(3)

(3) data integrity

⑪ Date Retrieval :-

To retrieve data from date bases we are using Sequel language , whereas as if we want to retrieve data from flat file we must develop Application program in high level languages.



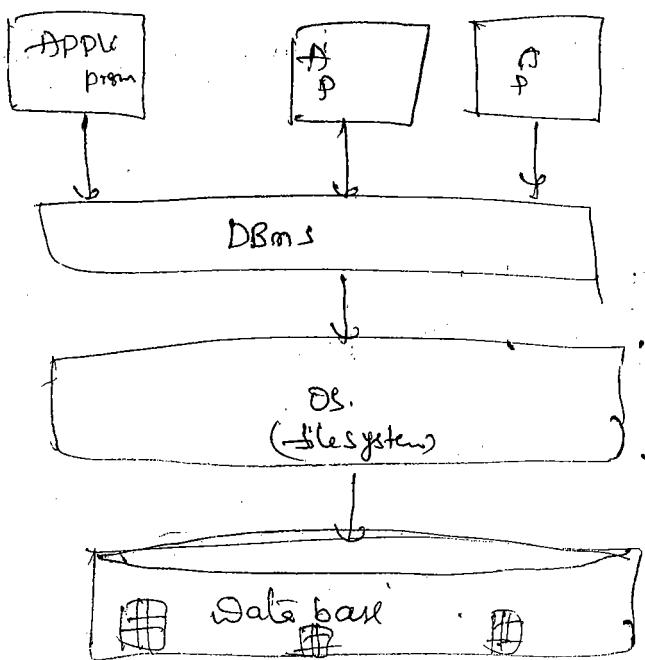
flat file approach.

D.B.M.S S/hds - oracle, SQL.

D.B - Collection of files.

→ When ever we are loading d.b.m.s shd's automatically d.b. is created & also automatically interface is created that's why end user directly interacts with date base using interface otherwise application programs indirectly interacting with a database using appn programs.

DBMS approach



Transactions designs all are in emp no. only.

↳ primary key.

→ In Dept duplicate not available.

→ unique values contains Dept table.

→ Contains primary key → master table (Dept table).

→ foreign key → child table.

(Emp table)

→ we can't try to delete master table record if not possible.

| Dept no loc | |
|-------------|--|
| 10 | |
| 10 | |
| 30 | |

To delete not possible.

bcz we first delete the records in related emp (child) table record

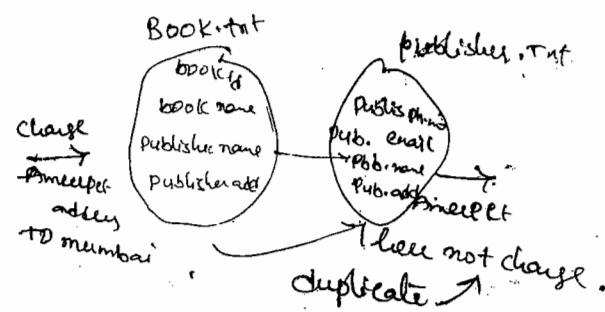
→ On - delete - cascade

when ever delete master table, all deleted in child table

Data redundancy: (duplicate data) 11-6

Inconsistency is a biggest problem in flat file.

Consistency data
in DBMS



→ Sometimes we are maintaining multiple copies of the same data in diff. locations.

→ This data is also called as "dependent data".

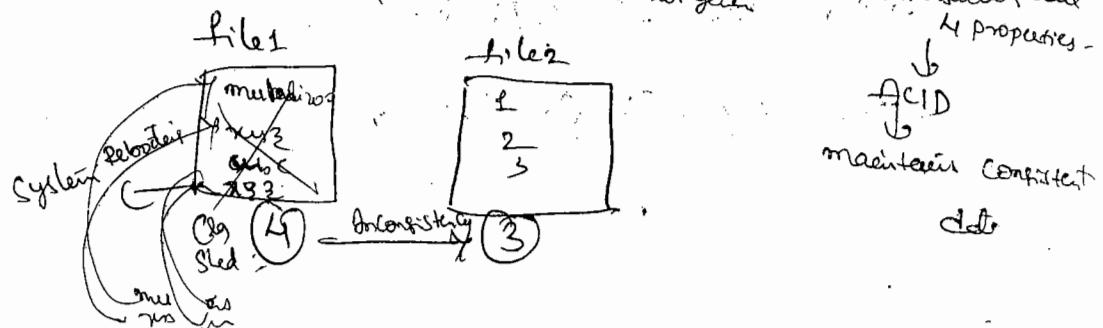
→ In flat files when we are modifying data in 1 location it is not effected in another location (files) this is called inconsistency.

→ where as DBMS like (SQL servers) automatically

maintains consistent data through transactions

Ex:- Inconsistency ATM : Transaction 3000, but money not given

↓
Every transaction has
4 properties -



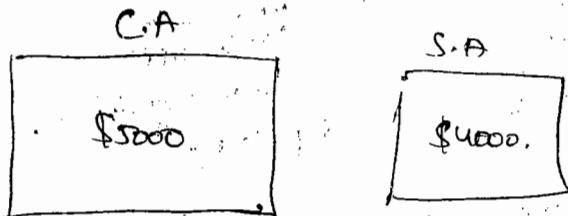
Transaction Algorithm.

before Trans + After Trans = Total data is same

Consistency property.

Atomacity \rightarrow Either all or none.

Roll back:



ACID
↓
Only consistency
data

Task: To transfer \$1000 from C.A. \rightarrow S.A.

rollback.

Step 1: Read amount C.A

2. deduct \$1000 from C.A

3. Write new amount into C.A.

4. Read S.A

5. add \$100

6. write new amt to S.A

is this
atomicity
broken
live arm

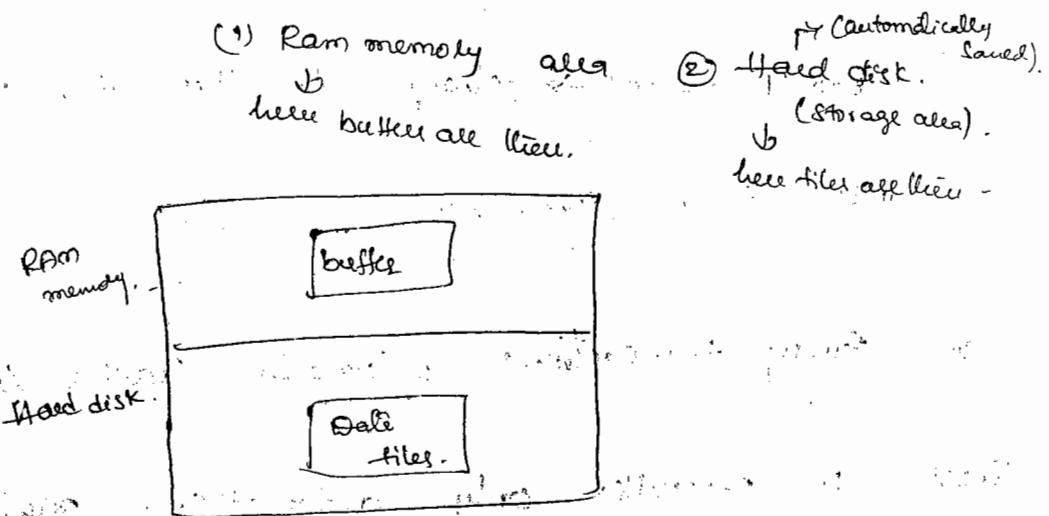
\rightarrow Permanently store all in Hard disk is \rightarrow durability.

(Commit)

\rightarrow Automacity \rightarrow Rollback (cancel)

\rightarrow Durability \rightarrow Commit (ok)

→ Every d.b. session having 2 parts.



→ Create 'I' table. It is stored in Hard disk only.
with out saving also bcz ddl commands
automatically saved.

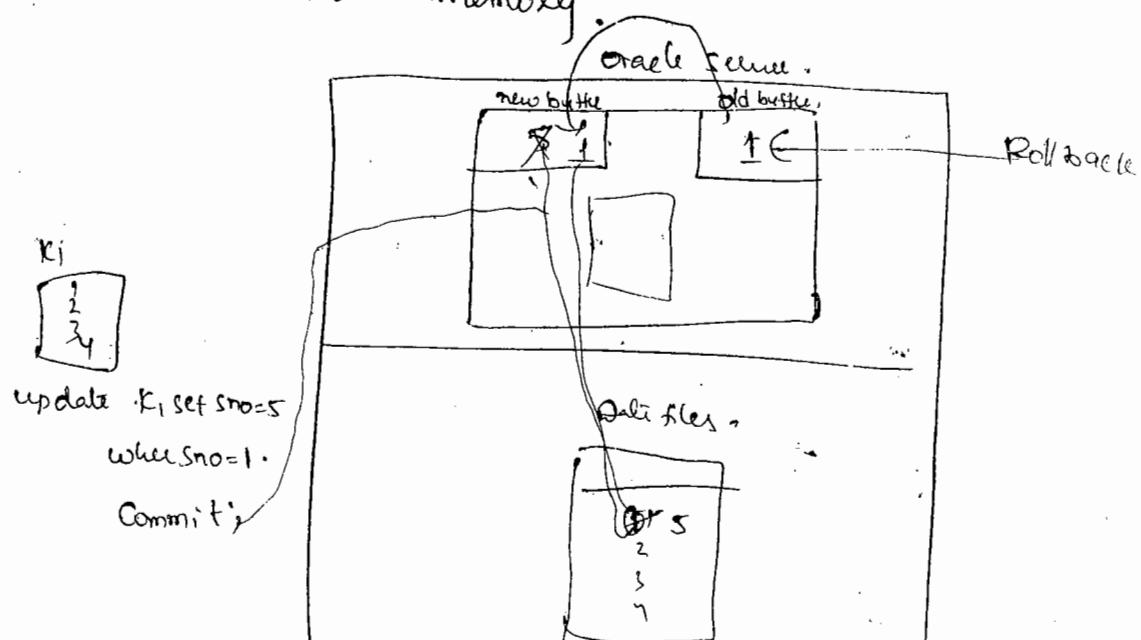
→ Commit

Is used to save permanently in Hard disk.

→ Save transactions into disk.

→ Rollback (Cancel the transaction).

** undo from memory.



→ After Commit, Roll back is not possible.

Commit means → changes in hard disk permanently.

insert, delete → These possible.

→ Every transaction internally having 4 properties

These 4 properties only automatically maintains consistent data in database. These are also called as "ACID" properties.

A → Automacity (either all or none)

C → Consistency (before, after total data remains same)

I → Isolation (running transactions individually;
Trns cannot be overlapped).

D → Durability (Same like transactions permanently
in disk).

→ All db. products introduced Commit Command

based on durability property, and also all db.

products introduce roll back Command based on
Automacity property.

→ Reduce the redundancy we having Normalization

→ In d.b.s, if you want to reduce Redundancy

d.b. designer uses normalization process.

12-6

(3) Data Integrity - ^(constraint)

Set of rules are there in projects

then set of rules called Business Rules

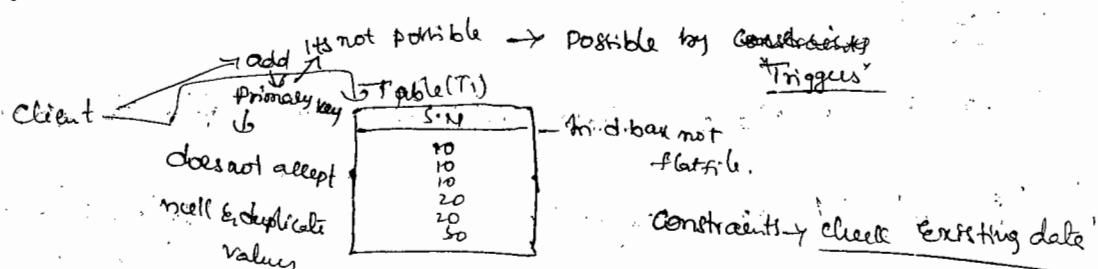
(To maintain proper data) → 2 constraints are there triggers.

→ Integrity means to maintain proper data

→ To maintain proper data developer defines set of rules
based on client requirement

→ These set of rules are also called as Business Rules

→ In d.b.s we are defining business rules through ^(trigger, by) Constraints,
triggers.



→ whereas in flat-files we are defining business rules
using application programs (say) in high level languages.

(4) Data Security: It is problem in flat-files (cause no security)

d.b. provides security internally - role-based security

→ data stored in flat files cannot be secured b.cz files
does not provide security mechanism.

→ whereas d.b.s (role) provides role based security.

data indexing (Btree Index, Bitwise indexing)

↓
very fast retrieval data

Airtel (very fast indexing)
B&NL (slow)

→ In databases to retrieve data very fastly then we are using indexing mechanism.

→ whereas flat files does not provide indexing mechanism

→ Orgns. Suffering from flat file approach to store data

→ In 1960 onwards orgns uses dbms's to store data permanently in hard disk.

(Interrelated data)

DBMS It is a

Collection of programs (^(cls)) internally written to manage

database

→ It is a place where we can store data in hard disk.

→ i.e. It is an organized collection of interrelated data (Emp → Dept).

→ Once data stored in database it can be integrated and also it can be shared by no. of users.

working

On 1970 E.F. Codd written a paper "relational

model of data for large shared data banks'

→ In this paper E.F.Codd introduced relational data model

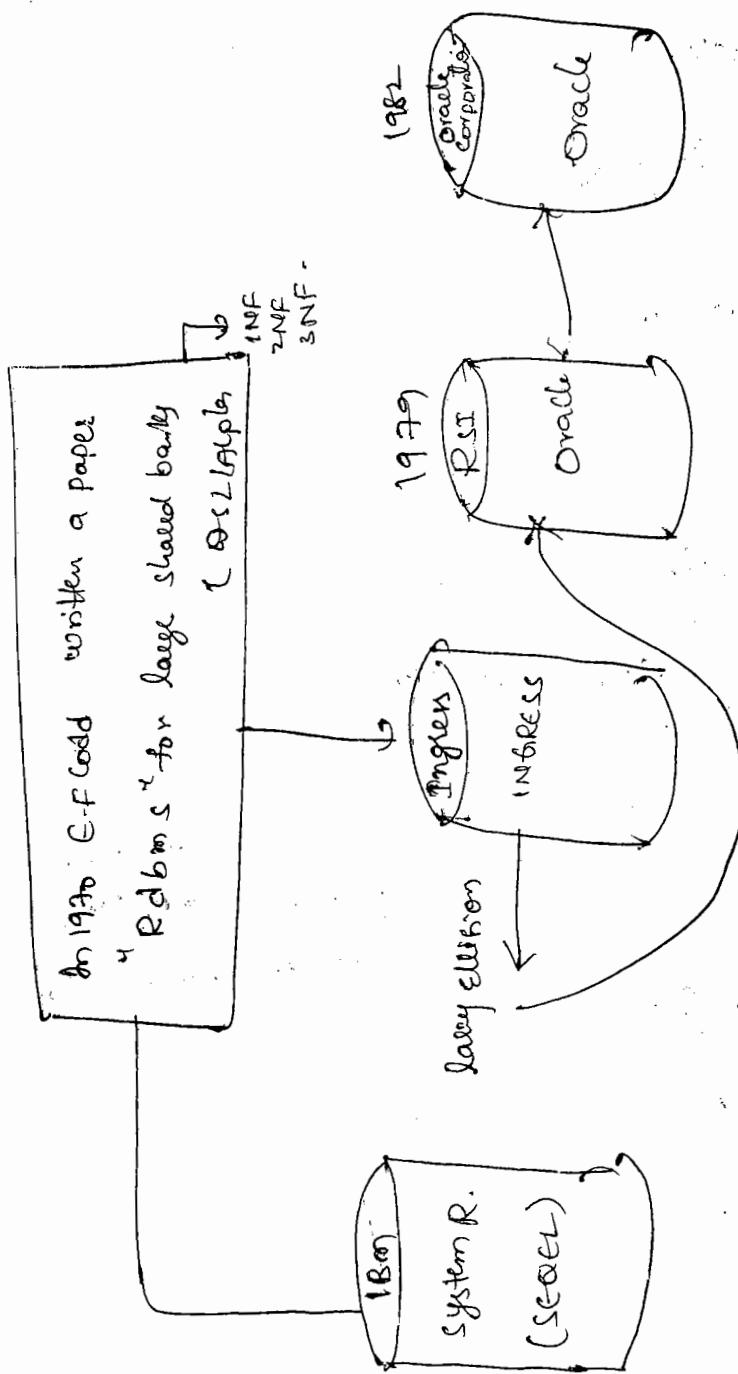
i.e. In this model we all store data in the format

format of table.

- And also in this paper E.F. Codd introduced 1st 3 normal forms.
- and also introduced DDL/ALPHA language to operate relational databases.

Now → Oracle → SQL.

IBM [1970's] System R → DDL/Alpha (SQL)



In 1979 : Oracle Versions: (1) SQL

(2) PLSQL.

- ① 1979 1st version → 2.0.
 - (1) first public release
 - (2) basic sql functionality, joins.
- ② 1983 → oracle 3.0.
 - (1) Commit, rollback.
 - (2) rewritten in C language.
- ③ 1984 → oracle 4.0.
 - (1) read consistency
- ④ 1985 → oracle 5.0.
 - introduced Client - Server architecture.
- ⑤ 5.1
 - distributed d.b.
- ⑥ 1986 → oracle 6.0.
 - introduced PLSQL.
 - low level locks.

there is no
up to 6.0 ↪ PLSQL

PLSQL there from 6.0.

⑧

1992 → Oracle 7.0.

- Integrity Constraints (foreign key) introduced.
- Stored procedures, stored functions.
upto 7.0.
d.types
↓
Varchar only
- Triggers are introduced.
- d.t. Varchar change into Varchar2.
available.
In 7.0 it
changes
Varchar2
- Truncate table.
- View Compilation (snapshot) introduced

⑨

1992 → Oracle 7.1.

- Introduced dynamic SQL.
- ANSI/ISO SQL92 Standard.

⑩

Oracle 7.2.

- Inline Views.
- PL/SQL Introduced Ref Cursors.

⑪

Oracle 7.3.

- bitmap indexes.
- (load file, write file) (UTL_file package).

⑫

Oracle 8.0 → 1996,

- Object technology (class, C++)
- nested table.
- Varray
- Instead of triggers. Oracle 8.0

(13) Oracle 8i (Internet) → 1997

- rollup, cube
- columns increased per a table upto 1000.
- analytical functions

(14) Oracle 9i → 2000 ^{two colls.}

- renaming a column
- anti joins
- multi-table inserts.
- flashback queries.
- merge stmt.

(15) 9.2, 9.3.

(16) Oracle 10g (grid technology) → 2005.

- introduces administration side
- Recycle bin concept
- flashback table.
- rename tablespa
- wri_concat()
- indices of char (P1/P2)
- regular expressions.

Oracle 11g (2007).

- Introduced Continue stmt. in PL/SQL loops
- Loops read only tables.
- Simple integer datatype introduced in PL/SQL.
- Virtual Columns
- follow clauses introduced in triggers.
- Using sequences without using dual table in PL/SQL.
- Compound triggers.
- enable, disable clauses used in trigger specification.
- named, mixed notations used in a subprogram executed using select stmt.
- regexp_count().

13-6

purpose
 → load date
 modified date
 deleted date
 retrieve date

SQL is directly not interact with db.

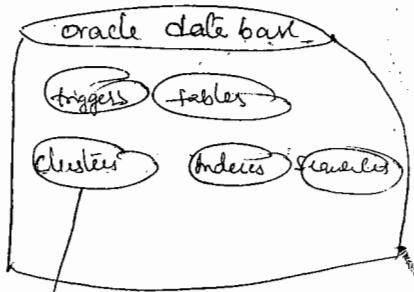
→ Only d.bane

Through the "query tools" SQL select

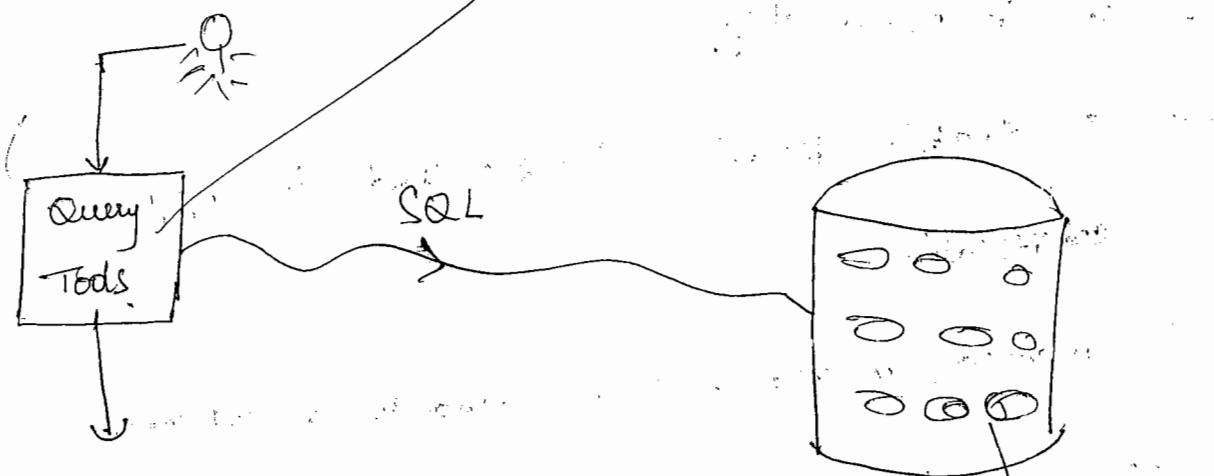
Interact with d.b.

Tods

- (1) SQL plus
- (2) isqlplus
- (3) Toad
- (4) PL/SQL developer (run in projects)
- (5) SQL navigator.



Collection of objects



Purpose -

- (1) load date
- (2) modify date
- (3) delete date
- (4) Retrieve date

→ don't load diff. stufs

Oracle
 — Express Edition.
 — Standard u.

good deb. is

↓

Enterprise Edition \$/o.

→ log, triggers & p.editions

Oracle

→ Each version having so many editions

date
double
decimal
float

→ start prompts → any Oracle SQL we know the "SQLplus Tel by default" application
↓
SQL

→ Express Edition we use user name: System-developer.

but here no tables.

P.W.: Edition.

→ So we use Enterprise edition to work in production

previous days.

→ 10g, 11g → SCOTT, TIGER → not work in Enterprise edition
but ^{first time} 10g onwards security mechanism is there. (locking mechanism)

Environment (real time)

U.N.: [sys AS SYSDBA] → unlock user.

Oracle log.

U.N.: [scott]

P.W.: [Tiger]

Error: Account locked.

unlock user → U.N.: [sys AS SYSDBA]

automatically asked like password.

Enter password: [sys]

SQL) Alter user scott account unlock;

SQL) Conn scott /Tiger

connecting keyword

pwd.

Pwd: Tiger

Confirm Pwd: Tiger

→ Clear Screen:-

(2)

Shift + delete or cl spc scrs;

→ To view all tables

→ Select * from tab;

Predefined ^{tables} we must R-table
and 4 tables → DB-tables.

(3)

- (1) Dept ^{m.t}
(2) Emp ^{r.u.k}
(3) Bonus
(4) Salgrade.
- { most powerful tables, then 2 are
normalized table.

(4)

Dept → master table. Primarily by ↓
↓
Emp → child table. master key.

Foreign key
child table

→ To view particular table;

(5)

→ Select * from Dept;

→ SQL is divided into sublanguages

(1) DDL (Data Definition language): - (All DDL Commands are

(1) Create (automatically saved)

(2) Alter

In universal standard SQL

(3) Drop

base comment base table

(4) Truncate

(5) Rename (oracle 9i)

② DML (Data manipulation language):-

(4) merge. oracle giving. (Oracle 9i)

(1) Insert

(2) update } Common commands

(3) delete

③ Data retrieval language / Data query language :- (DQL)

(1) Select.

(2)

④ Transaction Control language:- (TCL):-

(1) Commit. (Save). Permanently store in disk

(2) rollback. (undo from the memory)
cancel (ctrl+z)

(3) Savepoint.

⑤ Data Control language:-

(1) Grant (Giving permissions to another user)

(2) Revoke. } Powerful

GRANT

REVOKE

dsale

)

1 DQL

and the
table

SQL

DQL, DML

TCL

PL/SQL

TCL, DML.

DCL

Data types

PL/SQL ~~Variable~~ ~~a number(10)~~

SQL Column number(10) | ~~int a;~~

S&L

→ Dtypes identifies type of data in a column

Data type:-

- ① number(p,s)
- ② char
 - ↳ varchar2(maxsize)
- ③ date.

① Number (P,S) :-

Precision $\leftarrow P \rightarrow$ Total no. of digits.

Scale $\leftarrow S \rightarrow$ floating point nos.

It is used to store fixed, floating pt nos.

Syntax:-

Column name number(p,s).

| T1 |
|--------|
| SNO |
| 567.67 |

sno. number(5,2)

sno. number(50)

max. limit of 'P' is '38',
 \downarrow
 precision '1 to 38'.

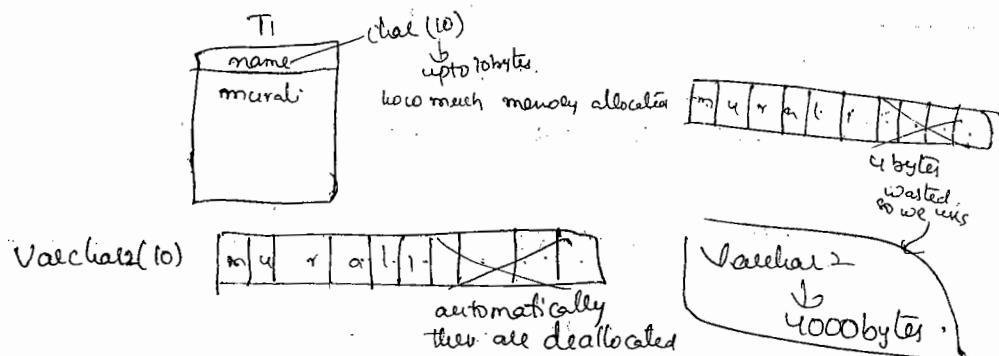
② char :- Alpha numeric data.

It is used to store fixed length alpha numeric data in bytes.

→ max. limit upto '2000' bytes.

→ By default character data leaving 1 byte.

Syntax:- Columnname char(size).



>Create table $t_1(\text{name} \text{char}(10))$;

↓ table create ↓ column name

Insert into t_1 Values ('murali');

↓ row creation

↓ insertion

Select * from t_1 ;

Name
murali.

→ memory storage is displayed by "dump"

Select dump(name) from t_1 ;

dump(name)

Type=96 Len=10; 20 40 30 40 50 30

waster
↑
32 32 32 32

32 ASCII of 'Space'

Variable

↓ dump(name)

Type=1 Len=2; 20 30 40 50 60

↓

→ Varchar2(maxsize)

→ It is used to store variable length, Alpha-numeric.

data in bytes.

→ max. limit is upto '4000' bytes.

Syntax: Columnname Varchar2(maxsize)

→ DATE:

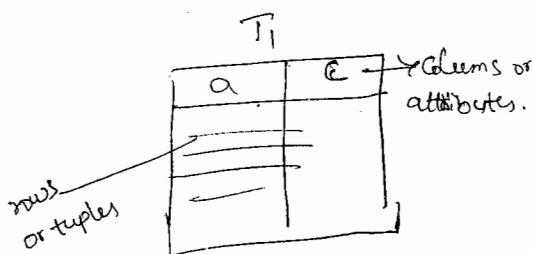
It is used to store dates in oracle date format

→ DD-MON-YY. → By default date format in Oracle
13-JUN-13.

→ Select SYSDATE from Dual;

Syntax: Columnname Date
 14-6-13.

→ Data Definition Language:-



These commands are used to define structure of the table.

(1) ~~Drop~~

- Create:- not only for table any thing like triggers, database.
- It is used to create database objects like tables, views, sequences.
- Creating a table:-

Create table tablename (column1 datatype(size), column2 datatype(size), ...);

Ex:-
 Create table d1 (sno number(10), name varchar2(10));
 4 table created

To View Structure of the table:-

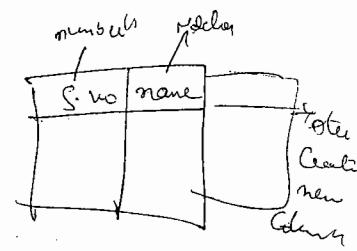
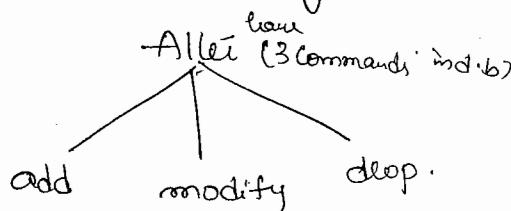
Syntax:- desc tablename;

Ex:-

Sql> desc d1;
 SNO Type
 name Number(10)
 Value2(10).

- Alter:- Add or remove Column..

It is used to change existing table structure.



→ Add: It is used to add columns into the Existing table.

Syntax:- alter table tablename (Col1 datatype, ...);

Ex:- Add sal column.

→ alter table d1 add sal number(10);

→ If Column is altered

Modify: It is used to modify Column datatype.

Or Column datatype size only.

| sno | name | sal |
|-----|------|-----|
| | | |

Syntax:- alter table tablename modify (Col1 datatype);

Empty table

s.no. is changed to date.

Ex:- alter table d1 modify sno Varchar(10);

desc d1;

| | |
|------|-------------|
| sno | Varchar(10) |
| name | Varchar(20) |

drop? removing the Column.

existing

2 methods.

① using () →

② using single column → use column keyword

→ It is used to remove columns from the table
(column set)

method 1:- If u want to drop a single column

at a time with out using parenthesis then we
must use column key in front of the column name.

Syntax:- alter table tablename drop Column Column name;
K.W.

Ex:- alter table d1 drop Column sal ;
(no need to write datatype)

method 2:- If you want to drop a single or multiple
columns with using parenthesis then we are using
following syntax.

Syntax:- alter table tablename drop (col1, col2..);

Ex:- alter table d1 drop(sal, name);

Note:- We can not drop all columns in a table.

→ If we have only '1' col. It is not possible.

→ If we have 100 we can drop only bgs, but '1' is not dropped error will be occurred.

drop:- drop a tables, triggers, sequences, clusters, objects etc.

→ It is used to remove database objects from the database.

Syntax:- (Objects)

Ex:- drop object-type objname;

- drop table tablename

- drop view viewname

- drop procedure procedurename

Ex:- drop table d1;

up to 109 ^{tables} → permanently dropped

but 109, 119 → not permanently dropped bcz here

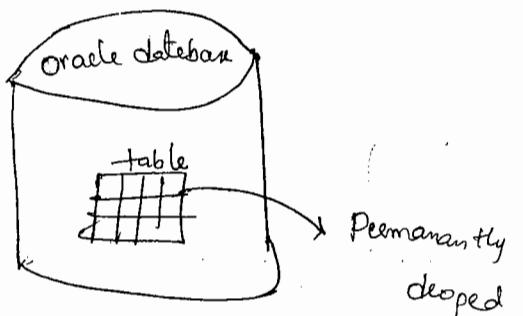
we have C:\Recycle bin

→ we remove permanently we use ~~trunc~~ in 109, 119 Enterprise Edition

dropping a table.

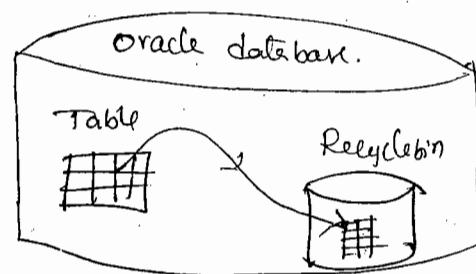
before oracle log.

Syntax: drop table tablename;



Oracle log (Enterprise edition)

Syntax: drop table tablename;



get it back from Recyclebin:-

Syntax:

Flashback table tablename to before
drop;

To drop permanently:-

Syntax: drop table tablename purge;
or drop table d1 purge;

→ drop table d1;

Table dropped.

→ get it back:-

➤ flashback table d1 to before drop;

➤ desc d1;

To drop permanently:-

➤ drop table d1 purge;

➤ again apply flashback comment.

➤ Error:- object not in Recycle bin.

→ Truncate— (7.0 onwards) (⁸² late command)

don't apply.

Original tables in projects. not columns

→ d.bas sever always delete data permanently.

[
→ Delete Command → is not freely.
→ rollback → it comes back data]

→ It is used to delete data permanently

from table (not ~~columns~~ ^{do delete} columns)

Syntax:

truncate table tablename;

Copy Table:- ^{Copy} date one table to d.table

➤ Create table d1 as select * from emp;

➤ Select * ^{all columns} from d1;

➤ truncate table d1;

Table truncated.

Testing:

➤ Select * from d1;

No rows selected.

➤ desc d1;

Only columns are displayed

Syn

→

XY
N

C

1)

Rename- It is used to rename a table & renaming a column also (nowadays)

Syntax:-

(Rename Parameter 'Renamed')

rename oldtablename to newtablename;

Ex:-

rename d1 to z1;

→ at a time only 1 table not a no. of tables.

Renaming a Column (Oracle q1):-

Column related to → Alter.

Only 1 column at a time rename is possible.

Syntax:-

alter table tablename rename Column _{IC.W.} oldcolname to newcolname;

oldcolname to newcolname;

→ 2 column is same is not possible in oracle.

→ alter table emp rename Column empno to sno;
Ex:-

Note:- By default all DDL Commands are automatically committed (Saved).

DML (Data manipulation language)

These Comments are used to manipulate

data in a table these are

① Insert

② Update

③ Delete

④ Merge (oracle 9i)

① Insert :-

It is used to insert data into table

3 methods are there.

method 1:-

Syntax:- Insert into tablename values (value1, value2, ...);

Eg:- Create table test (sno number (10), name

 varchar2(10));

Insert into test values (1, 'meali');

1 row created

Insert into test (2, 'Sachin');

1 row created

Select * from test;

| sno | Name |
|-----|--------|
| 1 | Meali |
| 2 | Sachin |

Eg:-

method 2:-

Using Substitutional operator (&));

Syntax:-

Insert into tablename Values (&col1, &col2, ...)

Eg:- > Insert into test values (&sno, 'ename');

Enter value for sno: 4

Enter value for name: XYZ

> Select * from test;

method 3:-

Skipping Columns

Syntax:-

Insert into tablename (Col1, Col2, ...) Value,
(Value1, Value2, ...);

Eg:- insert into test (name) values ('bbb');

> Select * from test;

| Sno | name |
|-----|--------|
| 1 | Murali |
| 2 | Sachin |

Update:-

It is used to modify data within a table.

Syntax:- Update tablename,

Set columnname = newvalue where column = oldvalue;

Eg:- update emp set sal = 1000 where ename = 'Smith';

> Select * from emp;

If you want to insert a value into particular cell then we also using update statement.

Eg:- > alter table test add address Value2(10);

> update test set address = 'mumbai' where name = 'Sahein';

> Select * from test;

| Sno | name | address |
|-----|--------|---------|
| 1 | murali | |
| 2 | Sahein | mumbai |
| 3 | abc | |

> update test set address = null where address = 'mumbai';
1 row updated.

> Select * from test;

Delete:- It is used to delete all rows or

Particulars rows from a table.

Syntax:- delete from tablename; (to delete all rows)

Syntax:- Delete from tablename where condition;
↳ (to delete particular rows)

Eg:- > Delete from test (temporarily deleted)
4 rows deleted

> rollback; → (get it back)

Rollback completed

➤ [truncate - permanently delete]

➤ Select * from test

➤ Commit;

Commit;

➤ roll back; rollback completed.

➤ Select * from test;

no rows selected.

Difference between delete and truncate.

→ whenever we are using delete from table name
the total data deleted is a table and also this
data temporarily stored in a buffer. We can get it
back this data using rollback (without using commit)

→ whenever we are using truncate-table-name
total data permanently deleted because by default
all ddl commands are automatically committed
that's why we can get it back deleted data.

Data Query Language (or) data retrieval language

Select: This command is used to retrieve data from table.

Syntax: Select Col1, Col2 from tablename where Condition group by Column name having condition
Order by columnname [asc/desc];

- ① Select all ~~columns~~ and all rows
- ② Select all columns and particular rows ~~where~~
- ③ Select ~~particular~~ ^{Col 1 to 6} particular Columns and all rows
- ④ Select particular Columns in particular rows

Creating a new table from another table.

17-6

Syntax:-

`Create table tablename as select statement;`

Ex:- `>Create table X1 as select * from emp;`

`* Select * from X1;`

`display as like emp table.`

Columns — primary key, foreign key,
 $\underbrace{\text{constraints.}}$

when we are copying table constraints are
not copied. Except (NOT NULL)

①

→ Only Copy Columns

→ data ($\begin{matrix} \nearrow \text{set of rows} \\ \text{where condition} \end{matrix}$)
will false only columns will selected

Ques

Creating a new table. from another table without

Copying data:-

Syntax:- `Create table newtablename as select`

`* from tablename where false condition;`

Exn > Create table X₃ as Select * from emp

where I=2; no rows selected > select * from X₃;

like emp%0000000; & desc - X₃; → only columns all opted.

I=9; } false condition.

I=1 } true It displays data.

→ when ever Condition is false data will not copy.

→ Operators used in Select Statement.

→ after select some columns we'll use (Select col1 col2)

(1) Arithmetic operators (*, /, +, -) (with mod (min)) instead of the

(2) Comparison op's (=, <, <=, >, >=, << or !=)

(3) logical op's (AND, OR, NOT)

(4) Special op's.

} where

used in where clause

(1) Arithmetic op's:- all used in number, date datatype columns.

Select display ename, sal, annsal from Emp table.

> Select ename, sal, sal*2 "annsal" from emp;

ename . sal annsal

Column alias.

➤ select sysdate from dual;

SysDate

17-FUN-13

➤ select sysdate+1 from dual;

Tomorrow date.

➤ select sysdate-1 from dual;

yesterdays date.

➤ Se

Query: display the employees except JOB as clerks
from emp table.

➤ Select * from emp where job != 'clerk';

(After where compulsory write column)

Q> display the employees who are getting more

than 2000 sal from emp table.

➤ Select * from emp where SAL > 2000;

using AND opr.

➤ Select * from emp where JOB = 'CLERK' AND
SAL > 2000;

➤ Select * from emp where job = 'CLERK' OR
sal > 2500.
OR → Clerk and
all sal above 2500. display;
OR is nothing but 'g' (comma)

→ display the employee who all belongs to
deptno 10, 30, 50, 70, 90. from emp

➤ Select * from emp where deptno = 10 OR
deptno = 30 OR deptno = 50 OR deptno = 70 OR deptno = 90;
Instead of we can another.

→ Special opr:

| <u>opr:</u> | <u>OPP:</u> |
|-------------|-------------|
| ① in | not in |
| ② between | not between |
| ③ is null | is not null |
| ④ like | not like .. |

Used in
where clause

In: This opr is used to ~~take~~ pick the values

one by one from list of values.

→ generally we are using In opr instead of 'OR' operator.

→ Bcz 'In' opr performance is very high compare to 'OR' opr

→ In multiple row subqueries we must use In operator.

System [select * from tablename where

Columnname in (list of values);

↓
separated with ,.

➤ Select * from emp where deptno in(10,20,30,40);

➤ Select * from emp where ename in('King', 'Smith'); →

➤ NOT in \rightarrow NULL
does not work with

↓
No rows selected

Y Select * from emp where Deptno

Not in (20,30);

- It displays deptno 10.

Y Select * from emp where Deptno

NOT IN (20,30, NULL);

- no rows selected.

Note: NOT IN opr does not work with null values.

Null is diff.
0 is diff.

NULL-

NULL is an undefined, unavailable,

Unknown value

- It is not same as 'Zero'

Null ≠ 0

Any arithmetic opr with NULL = NULL

→ Any arithmetic operations performed on NULL again it will become NULL

NULL + 50 = NULL

Q) display ename, sal, comm, sal+comm,
of the employee Smith from emp table.

> select ename, sal, comm, sal+comm
 $2400 + \text{NULL} = \text{NULL}$.

from emp where ename='SMITH';

O/P:- ename sal comm sal+comm
Smith 2400 (NULL) (NULL)

→ Sal + Comm

$2400 + \text{NULL} = \text{NULL}$

— If ^{like} the above ex. any problem is there to overcome
problem
NULL).

→ To overcome this problem oracle provided

NULL) function

Ex:- $\text{NULL} + 50$

→ NULL

In place of NULL we replace '0'

NULL) $\rightarrow 0 + 50$
 $\rightarrow 50$.

fun automatically substitute 'Zero'
user defined value

NVL: NVL is an predefined fun. which is used to replace or substitute user defined value in place of 'NULL'

Syntax:

$\boxed{\text{NVL(expr1, expr2)}}.$

must be same date type only.

→ These 2 expressions must belongs to same date type.

Ex:

$\text{NVL}(\text{Null}, 0) = 0$, $\text{NVL}(\text{Null}, 20) \rightarrow 20$

1st exp - Null, 2nd exp value = 2nd exp.

- If Expr1 is Null It returns expression 2.
- Otherwise It returns expression 1,

Ex:- $\text{NVL}(30, 20) = 30$.

→ select NVL(Null, 20) from Dual;

$\boxed{\text{NVL(Null)}}$
20

→ Select ename, sal, comm, sal+NVL(comm, 0)

from emp where ename = 'SMITH';

| ename | sal | comm | Sal+Comm |
|-------|------|--------|----------|
| Smith | 2000 | (null) | 2000 |

→ Sal + NVL(Comm, 0)

$2000 + NVL(\text{NULL}, 0)$

$2000 + 0$

= 2000

18 - 6

BETWEEN: This opr is used to retrieve range of values. [always 1st write low value, not work 1st high value]

Syntax:

Select * from tablename where

- Colname between lowvalue AND highvalue;

→ use AND opr only

Ex: Select * from emp where sal

between 2000 AND 5000;

↓

If given 5000 AND 2000;

↓

It is no rows selected.

→ BETWEEN - 2000 AND 5000:

↓

Including 2000, 5000 also.

~~Be~~ NULLS NOT NULL; → Other
not indexes not work.
Index not used.

These oprs are used in "where clause only".

→ These oprs are used to Test whether a Column having null values or not.

Dept No

10 — Display Employees 10th dept.

20 → Select * from emp where

50

70

90

DeptNo=10

— who are working in 10th dept.

→ Select * from emp where

~~DeptNo~~ Null. IS NULL;

Syntax:-

① Select * from tablename where ColumnName is null;

② Select * from tablename where ColumnName is not null;

Employees ^{whose}

not getting Commission

→ Select * from emp where {Commission}
Comm is null; {no rows selected}

→ Display the employees who are getting Comm. from emp table. using Special operator.

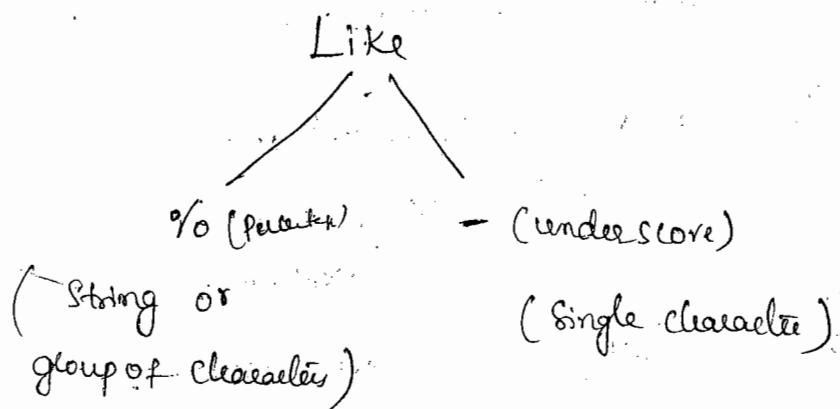
→ Select * from emp where

Comm is NOT null;

{Comm values}
display)

LIKE: This opr is used to retrieve data very fastly based on character pattern.

→ along with like opr we are using 2 special characters.



Syntax: Select * from table name

where columnname Like 'character pattern'.

dates, no's, times, names all are in single quotes ' '.

Ex:-

ALLEN MARTIN
%o. o%
 ends with 'N'

Display the employees whose ename start with 'M' from emp table using like opr.

→ Select * from emp where

ename like 'M%';

— Martin
— Miller

Revy
Recal
editing
M

Y Display the employees whose 4th character of ename second letter would be 'L'. from emp using like opr.

Y Select * from emp where
ename like '_L%';

- ALLEN
- BLAKE
- CLARK

- Concatenation Opr.: || (2 pipes).

It is used to Concatenate diff. data-type columns
and also used to Concatenate column values with
Strings.

Y Select 'my employee name is '||ename from emp.

my employee name is Smith
" " " Allen

Functions

- Functions are used to solve particular task and also fun's return a value.
- All dban systems having 2-types of fun's

- (1) predefined fun's (in SQL)
- (2) user defined (not possible in SQL).

(1) Predefined functions

→ ① Number fun's.

→ ② Character fun's

→ ③ Date fun's.

→ ④ group fun's / aggregate fun's

(1) Number functions:- These fun's operate over number data.

(1) abs():- Convert -ve values to +ve

Convert -ve values into +ve side we using "absolute fun"

Ex:- Select abs(-50) from dual;

| O.P | abs(-50) |
|-----|----------|
| | 50 |

*→ Test Predefined fun we use 'dual' table;
Virtual table → Predefined Y rows
generate sequel values.

Y Select abs(-50) from emp; Select abs(-50) from dept;

① 50
50
|

② 50 - 4 rows selected.

50
50
50
50 — 4 rows selected

→ Dual: Dual is an predefined Virtual Table

which Contains only 1 row and 1 column.

→ By default Dual Table column datatype is
VARCHAR2.

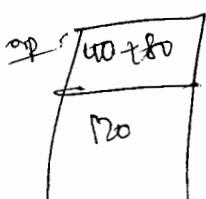
→ generally this table is used to test Predefined,
User defined fun's functionalities. and also

→ This table is used to generates Sequence Values
through Sequence object. & also

→ This table is used to perform mathematical
operations (in projects like calculator). Y select 40+80 from dual;

| |
|-------|
| 40+80 |
| 120 |

Y Select 40+80 from dual;



*
➤ select * from dual;

column.

Ø (dummy)

-

X

row

➤ desc dual;

Dummy

datatype

Namechar(10)

→ Select predefined or user defined

1st check on dual bcz it is proper table or not.

➤ select mul(null,50) from dual;

OP:- 50

➤ select mul(30,50) from dual;

OP:- 30.

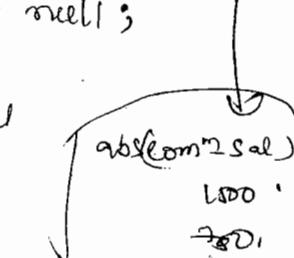
➤ select ename, comm, sal, com-sal emp

abs(comm-sal)

where comm is not null;

OP:- ename comm sal com-sal

| | | | |
|---|---|---|------|
| 1 | 1 | 1 | -100 |
| 1 | 1 | 1 | -50 |



* $\text{mod}(m,n)$:- It returns a remainder after m divided by n . 

Y Select $\text{mod}(10,2)$ from deal;

| | |
|-------|----------------------------|
| O.P:- | $\boxed{\text{mod}(10,2)}$ |
| | 0 |

Y greatest(expr, expr, ..., expr), least(expr, expr, ..., expr)

→ greatest returns "maximum value" within given expressions

→ where as least returns min value among given expressions.

Y Select greatest(3,6,9) from deal;

| | |
|-------|---|
| O.P:- |  |
|-------|---|

— Display greatest sal from emp;

Y Select greatest(sal) from emp;

O.P:-

1800
1200 } Same values but not greatest value,

→ Single column don't apply

greatest → expr, expr ..

max → Single column.

ename com sal

| | | |
|-------|-----|---------------------------|
| Smith | 300 | 1800 |
| | | → now greatest (sal, com) |

→ If u want to compare no. of explosions
we can greatest.

(E.C)
113V

→ Select ename, comm, sal,

greatest (comm, sal) from emp where
comm is not null;

19-6

when ever necessary flng pt nos then round.
round (m,n) → It rounds given floated value number 'm'
based on 'n'.

→ Select round(1.8) from dual; (1.4) (1.5)

OP '2'

1. 2

→ Select round(1.23456, 3) from dual;

after decimal 3 places.

4.23410,3
↓
1.234 only

OP+ 1.234 | 56.
↓
1.234 . | above 50%.
|
--- 1.235
↓ add 1

Note:- round always checks remaining no. If
remaining no. above 50% then '1' added to the
rounded member.

→ Select ename, sal from emp;

➤ Select ename, sal, round(sal/2), round(sal/2,1)
from emp;

O.P. Sal 1132.5
192.7

Trunc(m,n) :-

It ^{removal} Truncates gives floaled member m,
based on 'n'.

➤ Select trunc(1.7) from deal;
^{→ High place}
_{Truncates}
O.P :- '1'.

➤ Select trunc(1.23456, 3) from deal;
_↓
^{Truncated}
O.P :- 1.234.

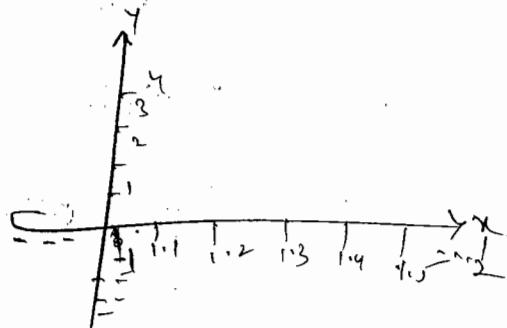
→ Ceil(), floor()
_↓
returns a
nearest greatest
integer value

→ Select ceil(1.3) from deal;
O.P :- 2. (nearest greatest integer)

→ select floor(1.9) from deal;
'1' (nearest lowest integer)

→ Ceil() returns nearest greatest integer.

→ Whereas floor() returns nearest lowest integer



Character functions

These fun's operate over character data.

① upper() :-

upper('abc')

upper ('ename')

↳ col name do not use ''.

It is used to return column values or string
into upper case.

>Select * from emp;

all enames are upper case

→ select upper ('ename') from emp;
↳ col values.

>Select upper ('abc') from dual;

② lower() :-

Select lower('ename') from emp;

↳ displays lower case.

Select * from emp;

↳ displays original case (Caps)

> update emp set ename = lower(ename);

14 rows updated

> select * from emp;

displays all enames in small letters

initcap() :-

It returns 1st letter would be capital & all remaining should be small

> select initcap(ename) from emp;

o/p Allen

James

> select initcap('ab cd ef') from dual;

In every d.b. when ever space is there next letter will be new letter.

o/p Ab Cd Ef.

length() :-

always returns number data type

- It returns number of characters and also includes spaces also eni- lsh. ru. thi- g/ length

> select length('AB CD') from dual;

o/p 5 emp when length(ename)=5;
select
from
diles

*** used in every d.b.

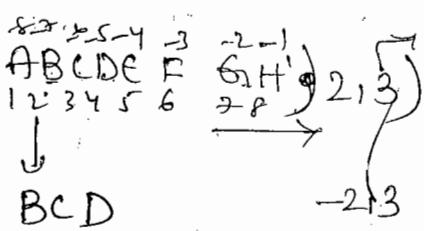
Syntax

Substr () :-

min - - 

It will extract position of the string with
given string based on last '2' parameters

> Substr ('ABCDEF', 2, 3)


Counting:
2, 3
→ 2, 3

~~ABC~~. DEF

→ S(4) (-5)

DEF6. DEF6, H

> Select Substr ('ABCDEF', 2, 3) from dedb;

OP: BCD

(-614)

CDEF.

Syntax:-



Substr (Colname (or) 'String name', Starting position, no.of characters, from position)

must be
nts

Ne

?

➤ Select * from emp;

→ Using Substring display the employees whose
their
names
starts
with
the
second
letter
would
be 'LA' (substring)
from emp table.

[BLARK]
[CLARK]

Powered

➤ ~~Substr(ename, 2)~~ → wrong.

➤ Select * from emp where

➤ Select * from emp where

Substr(ename, 2, 2) = 'LA';

O.P:- BLARK, CLARK.

→ After Select in most fun & character fun's
where clause is possible.

after selection
select max

→ Select * from emp where

group fun is not accepted

Note:- We are not allowed to use group fun's

In where clause. But we ^{all} allow to use
no. fun's, character fun's, date fun's in
where clause.

→ InSTR() → Most powerful. But dead easy

Always this fun returns no. data type

- It returns position of the DELIMITER (special symbols ^{20,14})
position of the character, position of the string within given string.

Y Select Instr ('ABCD', '*') From deal;

O.P:- 0. 
0
0

Instr ('AB*CD', '*')

O.P:- 3. 'A'
 1 'B'
 2 'BC'
 2 'BD' → sequence is not true.
 0C }

Instr ('AB*CD*EFG*KJH*', '*', 13); molecule no. of occures

O.P:- 3. (3, *, 2) (1, *, 3)

O.P:- 9. O.P:- 13

Instr() is always contains '4' parameters

→ goes to 2nd star.

(1, *, 5)

O.P:- 0.

(1, *, 5, 2)

O.P:- 13

"easy"

-type

ITER(special
symbols
\$,,*)
f (line)

AB*
Count
 $(11, 'x'), -5, 2)$.

O.P.L 3.

AB*
Counts
 $(11, 'x'), -4, 2)$

O.P.L 3.

$(11, 'x'), -3, 2)$

O.P.L 1?

-3-t is also count.

Syntax:-

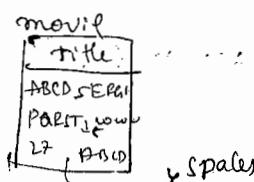
Instr (columnname (or) 'String name', 'Str',

Searching , no. of occurrences
Position from Position

(must be nos.)

Note:- Always Instr() returns position of
the String 'band' on last '2' parameters

But Oracle Sees Counts no. of characters
left side 1st position onwards (AB*)



> Select substr (Title, 1, *), instr>Title, 1') from movie;

→ lpad:- (and is hierarchical query) 20.6
 adding some thing.

lpad ('ABCD', ^{total length} 10, 'x')


→ It will fills remaining spaces with specified characters on left side.

→ Always here second parameter refers total length of the string.

>Select lpad ('ABCD', 10, '#') from dual;

O.P:- ##### ABCD.

(ENAG, 10, 'x') from emp;

O.P:- ***** Smith.

→ RPad:-

Select RPad ('ABCD', 10, '*') from dual;

O.P:- ABCD*****.

> Select Rpad ('Ename', 40, ('-')) from emp;

O.P:- Smith -----

> Select Rpad ('Ename', 40, ('-')) || Sal from emp.

O.P:- Smith ----- 2000

Allen ----- 4000

→ LTRIM():-

→ removed.

It removes left side specified characters.

> Select LTrim ('MISSTHSS', 'S') from dual;

O.P:- MISSTHSS.

> Select JOB || ltrim (JOB, ('C')) from emp;

O.P:- CLERK → LERK.

> Select ltrim (JOB, ('CS')) from emp;

O.P:- Salesman - aleman.

Clerk - leek.

→ (JOB, ('CSM')) from emp:

O.P:- Clerk - leek.

Sales man - aleman.

Manager - anager.

> (JOB, ('CL')) from emp.

O.P:- Clerk - leek.

→ RTRIM:

>Select RTRIM ('ssmissThss'), from dual;

O.P.H.
= SEMI-SSTH

TRIM(): Oracle 8i introduced trim function.

➤ Select Trim('mississipi', 's') from deal;

MITH (wrong).
D.P. - ~~MISS TH~~ wrong system

→ Trim remove only left & right side characters only.

→ middle characters are not dimmed

→ And also need to remove 1st, & last plates leading up.

Syntax :-

Trim (character @
string , From
Class) given string @
Column name)

Y Select Trim ('s' from 'semisness') from deal.

exp' miss tit.

➤ Select Trim ('Smith') from deal;

o.p. SMITH for clearing o.p. bcz there is spacing
↓

➤ Select length (Trim ('Smith')) from deal

o.p. 5 (SMITH).

→ Translate(), Replace();

➤ Select translate ('JOHN', 'H', 'N') from deal;
+ Replace with

o.p. 'TONN'.

➤ Select Replace ('JOHN', 'H', 'N') from deal;

o.p. TONN'

➤ Select Replace ('JOHN', 'H', 'NHFDFN') from deal;
+ Single character with string.

o.p. TONHFDFN.

- ➤ Select Translate('JOHN', 'H', 'NHFDFN') from deal;
o.p. TONN'

➤ Select Translate ('JOHN', 'HNOJ', 'MQPQ') from deal;
o.p. QP*19.

→ Translate is used to replaces character by character where as

→ Replace is used to replaces character by string or string by string

➤ Select translate ('ABCD', 'DCBA', '1234') from dual

O.P. → 4321

➤ Select Replace('A B C', ' ', 'ORACLE');

O.P. → AoracleBoracleC

→ String with string also Replace.

➤ Select JOB, Replace (JOB, 'salesman', 'marketeis') from emp;

O.P. → Salesman, Converting into marketeis.

~~Replaces~~ Select Replace ('SS mississ', 'S', 'M') from dual;

O.P. → MITH.

Deleted 'S' from total string
not only left, right, middle
also.

DATE Functions: [Sysdate - 1 → yesterday]
[Sysdate + 1 → tomorrow]

→ In oracle by default date format is

dd-mon-yy

① SYSDATE → no(.) (reduce column)

⑤ months-BETWEEN()

② ADD_MONTHS()

③ LAST_DAY()

④ NEXT_DAY()

① SYSDATE— It retrieves Current System Date.

> Select Sysdate From Dual;

Op: 20-Jun-13.

② ADD-MONTHS()— It is used to add or subtract no. of months to the specified date based on 2nd parameter.

> Select Add-months(sysdate, 1) from dual;

Op: 20-JUL-13.

> (sysdate+15) > (sysdate, -1)

Op: 20-SEP-14.

Op: 20-MAY-13.

> Select add-months('15-aug-05', 1) from dual

15-Sep-05

↳ rounddate

③ LAST_DAY()— It retrieves last date of the Specified month.

↳ passing only 1 parameter

> Select last-day(sysdate) from dual

Op: 30-JUN-13.

(4) NEXT_DAY():—

→ It retrieves next occurrence day from the specified date based on 2nd parameter.

> Select Next-day(sysdate, 'Sunday') from dual
next sunday

23-JUN-13.

(5) MONTHS_BETWEEN():

[~~retrieves number date type~~
~~no. only~~]
months-between(Date1, Date2)

Always this fun. retrieves no. date-type. i.e., It retrieves no. of "months" b/w 2 dates. Always Date1 more than Date2. [Date1 > Date2] otherwise worse It retrieves -ve sign.

> Select months-between('16-SEP-04', '16-JUL-04')

from dual

> Select months-between('16-FEB-04', '16-JUL-04');

O.P. -5

> Select months-between(abs('16-FEB-04'), '16-JUL-04');

O.P. 5

→ Always returns floating pt values bcz date is '31'. So we use round().

↳ Select months - between (sysdate, hiredate) for emp;

→ OPT - 390.1600
bcz 1st, lowest date, 2nd highest date. } floating pts.

↳ Select months - between (sysdate, hiredate). → 1/12

OPT: 390
36 → 1/12
years.

→ This fun always returns floating pt values bcz internally this fun uses every month having 31 days.

↳ Select ename, round (months - between (sysdate, hiredate)/12) from emp;

OPT: Smith - 22
Allen - 33.

Date Arithmetic
in

① Date + Number ✓ (possible)

② Date - Number ✓

③ Date1 + Date2 ✗

④ Date1 - Date2 ✓ - 8

① Select sysdate + 1 from deal;

Tomorrow.

② Select sysdate - 1 from deal;

Yesterday.

③ Select sysdate - sysdate from deal;

Zero.

→ 1st day of month

→ Way to display 1st date of the current month

using predefined fun's.

→ Previous month last day + 1.

→ Select ADD_MONTHS(LAST_DAY(sysdate), -1) + 1
from deal;

01-JUN-03.

Date Conversion — To date.

To date.

21-6

Simple, but more powerful

Date Conversion functions:-

① to_char();

② to_date();

① to_char(); — It is used to convert date type into character type. (date, string).

i.e., it is used to convert date type into date string.

> Select TO_CHAR(SYSDATE, 'Day month yyyy') from dual

O/P Friday JUNE 2013

→ Convert any date into client requirement format

> Select TO_CHAR(SYSDATE, 'DD/MM/YYYY') from dual.

O/P 21/06/2013.

Day
dy
Month
mon
Year

Return Value
Character

dates
nos. } { D
DD
DDD
MON
YY

→ Select TO_CHAR (SYSDATE, 'DAY') from dual;

O.P.: → Friday FRIDAY
'day'
(DY) (DD) (D) (DDO)
FR) 21 6 172.
SAT

Note: Basically TO_CHAR is case sensitive func.

DAY - FRIDAY DAY - FRI.

day - friday

→ Select TO_CHAR (SYSDATE, 'D') from dual;

O.P.: → '6' (21 → date)

Sunday - 1

6-friday

mon - 2 ----- Friday - 6'

→ DDD → day of the year (completed days)

→ DDth → 20th, 23rd

→ Ddsp → Date Spell out. O.P.: Twenty-first

→ mon → JUN.

→ month - June

→ MM → 06.

→ YEAR → Twenty thirteen. → YYYY - 2013

→ HHT: m T:SS → 12:04:46.
hours → minutes → seconds

→ YY - 13

FM - Full mode

- > If any space is occur in To-char O.P. we use 'fm'.
- > Select to-char (sysdate, 'DD Month YY') from dual;

O.P.: 21 JUN 1 2013

- > Select to-char (sysdate, 'DD Month YY') from dual.

O.P.: 21/JUNE/2013

→ 15 - Aug - 05 convert into 15 / August / 05.

- > Select to-char ('15-Aug-05', 'DD/month/yy') from dual;

Error O.P.: Invalid number.

- > In TO-char() we must & should follow rules:

→ 1st Parameter → date-type.

→ whenever we are giving a single quote ' ' it is taken as string.

→ String is convert into date using TO-date().

- > Select to-char (to-date ('15-Aug-05', 'DD Month YY')) from dual;

O.P.: ?

15/August/05.

DAY

Monday/August/05.

→ If 1st parameter is date, no need to-date.

Error Note:- when ever we are using `TO-Char()` always 1st parameter must be date type. otherwise oracle server returns an error. Ora-1722: Invalid number. (ⁱⁿ PL/SQL side) to over come this problem we are converting date string into date type using `TO-Date()` conversion function.

TO-Date():- It is used to convert date string into date type. (Oracle date format) (Any string is date.)

> `Select TO-Date('15/JUNE/07'), from dual;`

O.P:- 15-Jun-07 (Oracle format).
dd-mon-yy.

> `Select TO-Date('15/06/07') from dual;`

O.P:- Error. Not a valid month

→ Follow rules for TO-Date():- (most powerful concept)

`TO-Date('15/JUNE/07')`

When we write TO-Date it returns types must match with default

DD - mon - yy
↓ ↓ ↓
no - char no - no - no
no - no - no - no - no
↓ ↓ ↓ ↓ ↓
date format
15/06/07
↓
invalid month

no date no
DD-mon-yy
15 + 06/07
me. no invalid
error :- month.

- when ever we get an error in To-date in SQL
we use a 2nd Parameter Same as 1st parameter.
technically

>Select to_date ('15/06/07', 'DD/MM/YY') from dual

o/p : 15-JUN-07

Note: whenever we are using To-date
between values match with the default date format
otherwise use a 2nd parameter as same as 1st
parameter format.

Then only Oracle Server internally converts
date string into date type.

→ any date + no. is not working so it convert

in to date type so use to-date

>Select to_date ('09-feb-04') + 5 from dual

o/p: 14-feb-04

Select '09-Feb-04' + 5 from dual

↓
error so use to-date

> Select

year month and palandat

" " " "

to_date ('09-02-04', 'DD-mm-YY') + 5 days

O.P 14 Feb-03

Q) Display the employees who are joining in the month December from emp table using To_date()

(*)

> Select To_date(*.Hiredate, 'DD-Dec-YY') from emp;

Error

>

Select To_date(

This is conditionally this query. So use where clause

> Select * from emp where To_date(Hiredate, mon)=Dec;

O.P = Dec
Mon

(Hiredate, mon)=Dec;

→ display the employees who are joining in the

year 81 from emp table using To_date()

> Select * from emp where To_date(Hiredate, YY)='81';

→ we want to know year is 1987 or 1981 or 1881, 1781.

> Select * from emp where To_date(Hiredate, YYYY) =

Y Select min date & To-Date (min date, '4444)

From Emp;

O.P.: - 17 Dec - 80 To Date
 1980
 1981

27-06

memps;

17-06-80

1980

1981

date

date (min)=15

is the

class

1) = '81';
~~then~~

, 17-1.

(4444) ~~1000~~

Date Conversion Functions :-

13-6-2014.

(1) to-char () :- In where clause.

(2) to-date () :-

Select to-char (to-date ('24-JUN-05'), 'DD/MM/YYYY')

O/P:- 24/06/2005. from dual

PLA query to display the employees who are joining in the month December from emp table using to-char function.

X Select ename to-char (hiredate, 'dd-DEC-yy')
from emp; X

Select * from emp where to-char (hiredate, 'mon') =
'Dec';
→ O/P

PLA query to display the employees who are joining in the year '81' from emp table.

Select * from emp where to-char (hiredate,
'yy') = '81';

a character is to-clear so gaps is there in o.p.

Fill Mode:— (FM).

when ever to-clear returns gaps then we must use fill mode this mode is represented by 'FM'.

Ex:
Select to-clear (sysdate, 'DAY/MON/YY') from dual

O.P.— FRIDAY---[JUN] 2014

Select to-clear (sysdate, 'FM DAY/MON/YY')

O.P.— FRIDAY/JUN/2014

Full length character to-clear is a problem. (Internal concept)

> Select *From emp where to-clear (H1ledate,

'month' = December;

O.P.— NO rows Selected.

> — — — 'Fmonth' = December;

O.P will occurs.

to-clear helps in another situation

does Year is 1980 or 2080 @ uro.

> Select H1ledate, to-clear (H1ledate, 'YYYY')

from emp;

O.P.— H1ledate
1980 1981 1982

Note :- In Oracle whenever we are passing date string '05-Jun-06' into date function then oracle server automatically converts date string into date type.

But here passed parameter format must be by default date format ('dd-mon-yy') otherwise oracle server returns an error.

For this case we must use to_date (It error will come)

>Select last-day ('15-JUN-05') from dual;

O.P.F 30-JUN-05.

Select last-day ('15-06-05') from dual

O.P.F not a valid month.

Solutions :-

Select last-day (to_date ('15-06-05'), 'DD-mm-yy')) from dual;

O.P.F 30-JUN-05.

⇒ round(), trunc() functions used in dates :-

* Every d.b. date data type ^{contains} is 2 parts

① date part ② time part

(visible)

(not visible)

⇒ In oracle when ever we are using round(), trunc() fun's of dates ^{then} automatically date part is changed based on timeportion.

⇒ whenever we are using a round() fun then oracle Server automatically add 1 day to the date if time portion is 12 Noon. And also time portion automatically set to 'Zero' ('00:00:00')

Testing at 7:40 PM :-

➤ Select to_char (sysdate , ' DD-MON-YY HH24:MI:SS ') from dual;

O.P:- 13-JUN-14 19:10:45.

Using Round() :-

➤ Select to_char (round (sysdate) , ' DD-MON-YY HH24:MI:SS ') from dual;

O.P:- 14-JUN-14 00:00:00.

trunc() :- whenever we are using trunc() oracle server returns same date and also time portion automatically set to 'Zero' (00:00:00). If time portion \geq 12 Noon also.

➤ Select to char (trunc(sysdate), ('DD-MON-YY

HH24:MI:SS') from dual;

O.P:- 13-JUN-14 00:00:00,

WAP query to display the employees who are joining today from emp table

➤ Exe Insert into emp(empno, ename, hiredate)

values (1, 'murali', sysdate);

➤ Select * from emp where hiredate = sysdate;

\Rightarrow no rows selected.

Solution :-

➤ Select * from emp where trunc(hiredate) = trunc(sysdate)

1 murali 13-JUN-14.

| Group functions (or) aggregate fun's: | | 14-6-14. |
|---------------------------------------|-------------------------|----------------------------|
| | | compulsory part |
| (1) max() | (4) sum() | Column name not our own |
| (2) min() | (5) count(*) | |
| (3) avg() | (6) count (column name) | |

All group fun's operate on group of data and returns a single value.

(1) max:- It returns maximum value within a column.

➤ Select max(sal) from emp;
O.P:- 5500.

➤ Select max(Hiredate) from emp; (freshie)
O.P:- 12-Jan-83 ..

➤ Select max(ename) from emp;
WARD (alphabetical order)

(2) min:-

➤ Select min(sal) from emp;
O.P:- 950.

➤ Select min(Hiredate) from emp; (Senior)

Note:- We are not allowed to use group fun's in where clause.

(3) avg():

➤ Select avg(sal) from emp;

O.P:- 2433.92857.

Note:-

➤ Select avg(comm) from emp;

O.P:- 550.

Note:- In oracle by default all group fun's ignores null values except "Count(*)".

➤ With null value we are using nvl().

➤ Select avg(nvl(comm,0)) from emp;

O.P:- 157.142857

Sum(): It returns total in a no. data type column.

➤ Select sum(sal) from emp;

34075

Count (*) :- It Returns Count no. of rows within a table.

Y Select Count(*) from emp;

O.P:- 14 (rows).

Count (Column name) :-

It Counts no. of not null values within a Column

Y Select Count (Comm) from emp;

O.P:- 4.

Y Select Count (distinct (deptno)) from emp;

O.P:- 3
/ (10,20,30)

Group by :- group by clause is used to arrange similar data items into set of logical groups

this clause used in select stmt

Syntax :- Select Column name from table name

Group by Column name;

Y Ques to display no. of employees in each department from emp table using group by.

Y Select deptno, Count(*) from emp group by deptno;

| Deptno | Count(*) |
|--------|----------|
| 30 | 6 |
| 20 | 5 |
| 10 | 3 |

W.A.Q to display no. of employees jobwise from emp table using group by.

➤ Select job, count(*) from emp group by job;

| | | | | |
|-------|----------|---|-----------|-----------|
| O.P:- | clerk | 4 | president | 1 |
| | salesman | 4 | manager | 3 |
| | | | | Analyst 2 |

➤ Select deptno, min(Sal), max(Sal) from emp group by deptno

| O.P:- | <u>Deptno</u> | <u>min(Sal)</u> | <u>max(Sal)</u> |
|-------|---------------|-----------------|-----------------|
| | 30 | 950 | 3150 |
| | 20 | 1500 | 3600 |
| | 10 | 2900 | 5500. |

Note:- In all d.b. Systems we can also use group by clause without using group functions.

➤ Select deptno from emp group by deptno.

| <u>Deptno</u> , |
|-----------------|
| 30 |
| 20 |
| 10 |

Rule:- Other than group-fun columns specified after Select those all columns must be used after group by otherwise Oracle Server returns an error "not a group by expression".

Ex: Select deptno, max(sal), ename from emp
group by deptno.

Error: "not a group by expression"

Y Select deptno, max(sal), ename from emp

group by deptno, ename,
↓

O.P. will come. here we do not use group-fun

Y Select deptno, job from emp

group by deptno, job.

O.P. ✓

Y Select deptno from emp
→ here less cols is workers

group by deptno, job;

→ bt here after select how many cols here
also that no. of cols

| <u>Deptno</u> |
|---------------|
| 20 |
| 30 |
| 20 |
| 10 |
| 20 |
| 10 |

Note: when we are try to display group fun columns
with another column then database servers return
errors to overcome this problem we must use
group by clause

Ex: Step 1
Select sum(sal) from emp;

➤ Step 2

➤ Select deptno, sum(sal) from emp;

Error: ~~sum~~ not a single-group function.

Solution:- using group by

Select deptno, sum(sal) from emp

group by deptno;

| O.P: | Deptno | Sum(sal) |
|---------|--------|----------|
| 16-6-14 | 30 | 9100 |
| | 20 | 13625 |
| | 10 | 11350 |

➤ W.A.Q to display those departments having more than 4 employees from emp table using group by clause.

➤ Select deptno, count(*) from emp
group by deptno

where count(*) > 4;
~~deptno = 30;~~

Error: ~~sum~~ after group by we can't use where clause

Sol:-

Select deptno, count(*) from emp group by
deptno having count(*) > 4;

| O.P:- | Deptno | Count(*) |
|---------|--------|----------|
| 16-6-14 | 30 | 6 |
| | 20 | 5 |

Having:- After group by clause we are not allowed to use "where" clause.

In place of this one ANSI/ISO SQL provided another clause "having".

⇒ Generally if you want to restrict rows in a table then we are using "where clause". (where is in table)

Whereas if u want to restrict groups after group by then we are using "Having clause".

⇒ Generally we are not allowed to use group func in "where clause".

Whereas we can also use group func's in "Having clause".

Order by:- not allowed to use in a subquery

⇒ This clause is used to arrange data in sorting Order along with order by clause we are using 2 key words. They are Asc, Desc.
① Asc
② Desc

→ by default Order by clause having 'Asc' order

Syntax Select * from table name

Order by column name [asc/desc];

↳ Select Sal from emp Order by Sal desc

OP:
5600
3500
950.

↳ Note: In Order by clause we can also use Column Position (no.) in place of column name for improve performance of the query.

Ex: ↳ Select * from emp order by 2 desc
OP:
eno ename - must use only when use the column position.
1
X
X
:

↳ Select deptno, Count(*) from emp

where Sal > 1000

group by deptno

group having Count(*) > 2

Order by deptno asc.

| deptno | COUNT(*) |
|--------|----------|
| 30 | 3 |
| 20 | 5 |
| 10 | 4 |

Roll up, Cube: along with group by. (Oracle 8i)

→ Oracle 8i introduced roll up, cube clauses. These clauses are optional clauses use along with group by clause only.

→ There clauses are used to calculate Sub-total, Grand total automatically

Syntax:-

Select Col1, Col2 ... , ...

From table name group by rollup (Col1, Col2);

② Select Col1, Col2 ... , ...

From table name group by cube (Col1, Col2);

→ "Rollup" is used to calculate Sub-total, grand-total band on a single column.

Whereas if you want to calculate Subtotal, Grand total band on no. of columns then we are using 'Cube'.

→ Select deptno, job, sum(sal) from emp group by
rollup (deptno, job);
only this column total

Y Select deptno, job, sum(sal), count(*) from
emp group by cube (deptno, job);

⇒ WAP to sumof sal and with ename with grand total
Select ename, sum(sal) from emp

group by ~~ename~~ rollup (ename);

Joins.

Joins are used to retrieve data from multiple tables.

Rule- If you are joining 'n' table then we are using (n-1) join conditions.

- Oracle Server having following types of joins

- 8th joins.
in Oracle d.b.
- (1) Equi join (or) inner join.
 - (2) Non Equi join
 - (3) Self join
 - (4) Outer join.
- left outer
right outer (9th)

These joins are called "8th joins"

9th joins (or) ANSI joins.

- (1) Inner join.
- (2) left outer join. (left join)
- (3) right outer join. (right join)
- (4) full outer join.
- (5) Natural join.

These joins are working any d.b. that is why only in oracle.

Internally Oracle Server uses Cross join by default.

Note:- We can also retrieve data from multiple tables without using joining Condition. In this case Oracle Server internally uses "Cross join" (by default).

Cross join between more Duplicate data because Cross join works based on Cartesian product.

Ex:- > Select ename, sal, dname, loc
from emp, dept.

⑥ Equi join (or) INNER JOIN :- common column & same data type
Banded on equality Condition we are retrieving data

from multiple tables. Here joining Conditional Columns
must belongs to same data type.

— When tables having Common Column names then
only using this join.

Syntax :- Select col1, col2, ... ,

from tablename1, tablename2

where tablename1. Common column name = tablename2.

Common column name
Joining condition

→ Select ename, sal, deptno, dname, loc from emp, dept

where emp.deptno = dept.deptno;

Error:- Column ambiguously defined.

→ always execution goes to first from only.

Note:- In databases if we want to avoid ambiguity then we must specify table name along with common column name using '.' operator.

Ex:-

Select ename, sal, emp.deptno, dept.deptno, dname, loc
from emp, dept

where emp.deptno = dept.deptno.

Note:- Generally to avoid ambiguity feature we can specify every column name along with table using '.' operator.

Using alias names:- We can also use alias names for the table in from clause. Then alias names also called as reference names created in from clause and also their reference names must be different names.

Syntax from tablename1 aliasname1,
tablename2 aliasname2.

Ex: Select ename, sal, d.deptno, dname, loc
from emp e, dept d
where e.deptno = d.deptno;

Note: using equi join - we are retrieving "matching rows" only.

> Select ename, sal, d.deptno, dname, loc
from emp e, dept d
where e.deptno = d.deptno

(Here deptno 40 does not display if we are using
d.deptno also)

- Way to display the employees who are working
from emp, dept table.
In the location Chicago using equi join

Select e.ename, d.lo e.~~location~~

from emp e, dept d

where e.deptno = d.deptno AND ~~loc = 'Chicago'~~
AND operator

Note:- If we want to use filter Condition (restriction) after joining Condition then we must use 'AND' operator in '8i' joins whereas in '9i' joins we are using either 'and' (or) 'OR' 'where' clause.

Ques - to display dname, sum(sal) from emp, dept table using equi join

Select d.dname, sum(sal)

from ~~group by~~ ~~deptno~~ dname, emp ~~e~~

where d.deptno = e.deptno;

> Select d.dname, sum(sal), d.deptno

from emp e, dept d

where d.deptno = e.deptno

group by d.dname, d.deptno.

| <u>DeptNo</u> | <u>Dname</u> | <u>Sum(Sal)</u> | <u>Outer join</u> |
|---------------|--------------|-----------------|-------------------|
| 10 | Accounting | 18247.5 | |

Using clause :-

Ques - to display dname, deptno, sum(sal) from emp, dept table of those employees, whose salary is greater than 15000

> Select d.deptno, dname, sum(sal) from emp e, dept d
where e.deptno = d.deptno group by d.deptno, dname

where = on

having sum(sal) > 15000

| Deptno | Dname | Sum(sal) |
|--------|------------|----------|
| 10 | Accounting | 18247.5 |

Q1 to display deptname & sum of salary of those department in which no. of employees in dept > 3.

Select dname, sum(sal) from emp e, dept d
where e.deptno = d.deptno group by dname having
Count(*) > 3

→ DNAME SUM(SAL)

RESEARCH

SALES.

Q2 to display deptname & sum of salary of department where no. of employees > 3 & also calculate total of salary automatically

X Select dname, sum(sal) from emp e, dept d
where e.deptno = d.deptno group by dname having
Count(*) ≥ 3 rollup (dname);

Select dname, sum(sal) from emp e, dept d
where e.deptno = d.deptno group by rollup (dname)
having Count(*) > 3

Note :- equi join :-

- Based on other than equality condition ($<$, $<=$, $>$, $>=$, in , between) we are retrieving data from multiple tables.
- Generally when tables does not have common columns then only we are using this join.

Eg:-
x) Select ename, sal, losal, hisal from emp, salgrade
where sal between losal and hisal;

(or)

y) Select ename, sal, losal, hisal from emp, salgrade
where sal \geq losal and sal \leq hisal;

Self Join :-

- Joining a table to itself is called self join.
- Here joining conditional columns must belong to same datatype.
- Generally if you want to compare 2 different column values in a single table then we must use self join.
In this case these 2 columns must belong to same datatype.
- If you want to compare multiple values within a single column also then we use self-join

— whenever we are using self join then we must create table aliases in "from" clause

from tablename aliasname1, tablename aliasname2 ...

Table Alias Names are Reference Name:-

→ In all database, when we representing a same table, twice (or) if we are referring a table using diff table name then we must use table as aliasname using clause → All Joins

These are also called referenced for a table.

These aliasnames must be diff. name.

These aliasnames behaves as table where query executed.

Name in oracle. table : aliasname are used self-join
Correlated Subqueries, merge start,

Q) write to display the employees who are getting same salary in emp table. using self join (Employee Smith Salary)

→

| e1 | |
|-------|-----|
| ename | sal |
| Smith | x |

e2

| ename | sal |
|-------|-----|
| ? | x |

Smith
Salary

→ Select e2.ename, e1.sal from emp e1, emp e2 where e1.sal = e2.sal And e1.ename = 'SMITH';

ename sal

SMITH 1600

ALLAN 1600

WAP to display employee names & their manager names from emp table using selfjoin
As all employees are not manager but all managers are employees so select searches each manager member in a column empno

emp e₂

empno emp mgr
7369 SMITH 2902
7499 ALLEN 3693

Y Select e₁.ename "employees", e₂.ename "managers" from emp e₁, emp e₂ where e₁.mgr = e₂.empno;
(or)

Y Select e₁.ename "employees", e₁.mgr, e₂.empno, e₂.ename "managers" from emp e₁, emp e₂ where e₁.mgr = e₂.empno;

Outer join:-

This join is used to retrieve all rows from 1-table & matching rows from another table.

Generally equi join retrieves matching rows only if we want to retrieve non matching rows also then we are using join operator (+) in joining condition of the equijoin this is called 8) outer join

Note: This join operator can be used only 1 side at a time in joining condition.

Y select ename, sal, d.deptno, d.dname, loc from emp e, dept d where e.deptno (+) = d.deptno

matching rows → All rows

O.P.: 40 operations
BOSTON

9i joins (or) ANSI Joins

- Inner join
- left outer join
- right outer join
- full outer join
- Natural join

Inner join:

This join also returns matching rows only. This join performance is very high compared to 8i equijoin.

When tables having common columns then only we are using this join. Here also joining conditionnal Columns must belongs to Same data type

> Select ename, sal, d.deptno, dname, loc
from emp e join dept d on
e.deptno = d.deptno

Q) Write to display the employees who are working in the location 'chicago' from emp, dept table using '9i' inner join

> Select ename, loc from emp e join dept d on
e.deptno = d.deptno where loc = 'CHICAGO';
(or)

> Select ename, loc from emp join dept on

emp.deptno = dept.deptno

where loc = 'CHICAGO';

| <u>ENAME</u> | <u>LOC</u> |
|--------------|------------|
| WARD | CHICAGO |
| TURNER | CHICAGO |
| ALLEN | CHICAGO |
| JAMES | CHICAGO. |

22-7-14

Using clause :-

→ In qⁱ joins, we can also use using clause in place of 'on' clause using clause performance is very high compared to on clause. Always using clauses returns common columns one time only

Syntax:-

Select * from tablename1 join tablename2
using (CommonColumnName);

Ex:- > Select * from t1;

| A | B | C |
|---|---|---|
| x | y | z |

> Select * from t2;

| A | B |
|---|---|
| x | y |

> Select * from t1 join t2 on t1.a = t2.a and
t1.b = t2.b;

| A | B | C | A | B |
|---|---|---|---|---|
| x | y | z | x | y |

> Select * from t₁ join t₂ using (a**b**);

| A | B | C |
|---|---|---|
| x | y | z |

Note:- whenever we are using "using" clauses, then we are not allowed to use alias name on joining conditional column.

Eg:- Select ename, sal, deptno, dname, loc from emp e join dept d using (deptno);

11. Then

joining

m

left outer join :- (left side table).
This join always retains all rows from left side table  & matching rows from right side table & also returns null values in place of non-matching rows in another table.

Select * from t_1 ;

| | | |
|---|---|---|
| A | B | C |
| X | Y | Z |
| P | q | r |

> Select * from t_2 ;

| | |
|---|---|
| A | B |
| X | Y |
| S | T |

> Select * from t_1 . left outer join t_2 on
 $t_1.a = t_2.a$ and $t_1.b = t_2.b$;

| | | | | |
|---|---|---|--------|--------|
| A | B | C | A | B |
| X | Y | Z | X | Y |
| P | q | r | (null) | (null) |

23-6-14.

right outer join :- (right side table)

This join returns all rows from right side table and matching rows from left side table and also returns 'NULL' values instead of non matching rows in another table.

Ex:-

>Select * from t₁ right outer join t₂ on t₁.a=t₂.a and t₁.b=t₂.b;

| | | | | |
|-------------------------|--------|--------|--|--------|
| A X | B Y | C Z | A X | B Y |
| <u>matched row.</u> | | | right side table | |
| <u>not matched data</u> | | | all rows from right side matching from left side | |

→ Full outer join:- It is the combination of left, right outer join i.e., it returns all rows from all the tables and also returns 'NULL' values, in place of non-matching rows, in another table.

| | | | | |
|--------|--------|--------|--------|--------|
| A X | B Y | C Z | A X | B Y |
| P Q | R S | T U | V W | X Y |
| W X | - - | - - | S T | - - |
| Z U | - - | - - | - - | - - |

Select * from t₁ full outer join t₂ on t₁.a=t₂.a and t₁.b=t₂.b;

on clause requires in sql, q1.

Natural join

where not required

this join also returns matching

rows only. ~~the~~

① In this join we are not allowed to use joining condition.

② * when tables having common columns then only we are using this join.

③ This join performance is very high compare to inner join.

④ Natural join internally uses using clause that's why their joins also returns common columns 1 time only.

Syntax: Select * from

t_1 natural join t_2 ;

| A | B | C |
|---|---|---|
| X | Y | Z |

AB one time only display.

Note:- When we are using natural join also we are not allowed to use alias name on joining conditional column. Bcz natural join internally uses using clause.

Ex:- > Select

dDept X

ename, sal, deptno, dname, loc

from emp e natural join dept d;

Cross Join :- (default join)

> Select ename, sal, dname, loc
 from emp cross join dept;

O.P.T. To rows (56 rows) 14x4

Syntax:
 (2 conditions)

Joining more than 2 tables.

Ex:- joining 3 tables. \rightarrow 2 join conditions.
 \rightarrow 8^o join.

Syntax:-

Select col₁, col₂

from table name₁, table name₂, table₃.

where

table name₁. Common col = table name₂. Common col_m,

AND.

table 2. Comm col = table 3. Comm col;

9^o join.
 (3 tables) (ANSI joins)

Syntax:-

Select col₁, col₂

from table₁ join table₂

on

table₁. Common col = table₂. Common col

Join table₃.

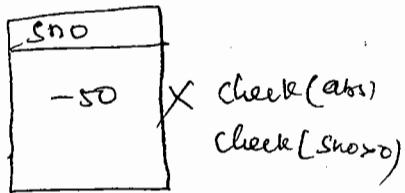
on

table₂. Common col = table₃. Common col.

2. Constraints

23-7-14

t1



← Constraints are used to prevent invalid data entry into our tables.

- Constraints are created on table columns.
- Oracle Server having following types of constraint

- (1) not null (3) primary key. (5) check.
- (2) unique (4) foreign key.

> For all the above constraints are created in 2 ways.

- (1) Column level.
- (2) Table level.

(1) Column level: In this method we are defining ^(Creating) constraints on individual columns.

For this method

e.g., When ever we are defining the column then immediately we are defining Constraint type specifying.

Syntax:-

Create table tablename (Col1 datatype(size),

constraint-type, Col2 datatype(size), constraint-type, ...);

*
2)

Table level :-

Create Bank(Accno varchar(20) primary key,

Accname varchar(20) primary key; XX

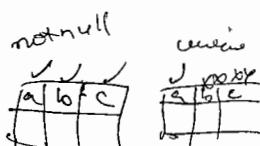
In this method we are defining constraints on group of columns. In this method i.e., 1st we are defining all columns & last only we can specify constraint type along with group of columns (primary key (Accno, name))

Syntax:-

Create table tablename (Col1 datatype(size),

Col2 datatype(size), ..., constraint-type (Col1, Col2, ...));

| Constraint Type | ① Not null. | Column | Table | Create table bank (Acc number(10), name varchar(20), balance number(10)) |
|-----------------|-------------|--------|-------|---|
| or | ② Unique. | ✓ | ✓ | |
| Constraints | ③ P.K. | ✓ | ✗ | Primary Key (Accno, name) |
| | ④ F.K. | ✓ | ✓ | |
| | ⑤ Check | ✓ | ✓ | |



→ Only 1 P.K. in 1 table.

... p.k.

→ Uniqueness of record in a table

Acc no name balance
1001 mukesh 00000

→ If 1 P.K. column not possible program assumption

1002 abc 20000

Combination of columns will take as P.Key

① NOT NULL :- Not null constraint does not support table level.

It does not accept "null" values but it will accept duplicate values.

Column level :-

Ex:- > create table d₁ (sno number(10) not null,
name varchar2(10));

| col | Name | Null? | Type |
|-----|------|----------|--------------|
| | sno | not-null | number(10) |
| | name | | varchar2(10) |

testing:-

> insert into d₁ values (null,'abc');

Error Ol400:- Cannot insert null int Sno.

> insert into d₁ values (20,'abc');

> insert into d₁ values (20,'xyz'); } duplicate allows.

Table level :-

Create table d₂ (sno number(10), name varchar2(10),
not null (sno));

Error invalid identifier

bz not null not supported Table level.

② Unique :- This Constraint is Created on column-level, table level.

It does not accept duplicate values but it will accept null values (oppo. not null). unique in P.K. column using where clause bcz (internally db server creates indexes) bcz improving performance.

Note :- when ever we are creating unique constraint oracle server internally automatically creates 'btree' indexes on those columns.

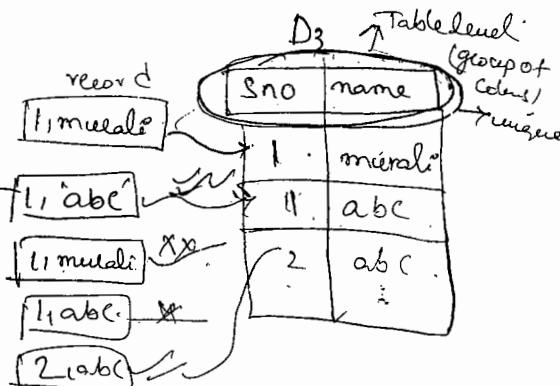
Column level :-

> Create table d2 (sno number(10) unique, name varchar(10));

Table level :- record

> Create table d3 (sno number(10), name varchar(10), unique (sno, name));

Table level so checking record



> Select * from d3;

| Sno | Name |
|-----|--------|
| 1 | merali |
| 1 | abc |

> Insert into d3 values (1, 'abc');

Error: ORA-00001: unique constraint violated.

(3) Primary key :- P.K. uniquely identifies ~~one~~ record in a table

- In all d.b. systems there can be only 1 primary key in a table and also primary key does not accept
- Duplicate, null values. and also
 - When we are creating p.k. Oracle ~~it~~ automatically creates btree index on p.k. columns.

Column level:-

> Create table d4 (sno number(10) primary key,
name varchar2(10));

Table level:-

> Create table d5 (sno number(10), name varchar2(10),
primary key (sno, name));

(4) Foreign key

24-7-14

- If u want to establish relationship b/w tables
then we must use referential integrity constraint

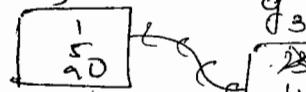
"Foreign key".

→ Here 1 table foreign key must belongs to another table primary key.

→ Here then 2 columns must belongs to same datatype

→ Always F.K values are based on P.K. values only.

d₃ P.K.

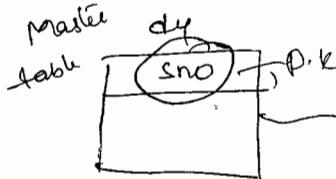


g₃ F.K.



— duplicate & null also accepted

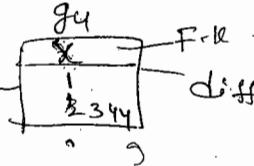
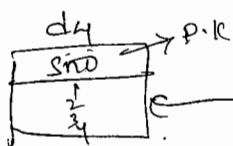
→ Generally P.K. does not accept duplicate, null values, whereas F.K. accepts duplicate, null values.



Column level:-

create table g₄

(sno number(10) references d₄(sno));



diff col name ✓
Pdt type is same only it works. ✓
accepted -

create table g₄

(x number(10) references d₄(sno));

→ Column level :- (referenes) :-

Syntax :-

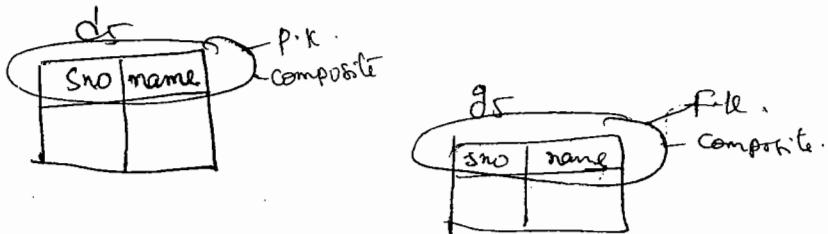
Create table tablename (Col1 datatype (size))

referenes mastertablename (Primary Key Colname),

Col2 datatype (size), ...);

- > Create table g_4 (sno number(10) references d_4);
- > Create table g_4 (x number(10) references $d_4(sno)$);
- (Table level: (F.K, references))
- Syntax:

Create table tablename (Col1 datatype(size), Col2
 datatype(size),,
 foreignkey (col1, col2, ...)
 references
 mastertablename (primary colnames));



> Create table g_5 (sno number(10), name varchar2(10),
 foreign key(sno, name) references d_5);

→ when ever we are establishing relationship
 b/w tables then oracle seems automatically violates
 2 rules.

(1) Deletion in master table.

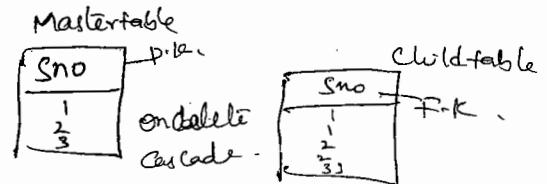
(2) Insertion in child table.

① Deletion in Master table:- In oracle if we all try to delete a master table record in master table then oracle server returns an error Ora-2292 if child table record exist.

→ To overcome this problem 1st we are deleting child table record in child table then only we are deleting appropriate master table record in master table.

→ Otherwise we use an "on delete Cascade clause"

→ On delete Cascade clause:-



when ever we are using "on delete Cascade" clause in child table then if we are deleting a master table record in master table. Then automatically master, child table records are deleted in both the tables. (not all records)

Syntax:-

```
Create table tablename (Col1 datatype(size) references  
master tablename (P.K. Colname) on delete cascade);
```

Ex:- > Create table mas (Sno number(10) primary key);

> Create Insert into mas Values (.....);

> Select * from mas;

| |
|-----|
| Sno |
| 1 |
| 2 |
| 3 |

- > Create table ~~table~~ child (sno number(10) references mas on delete cascade);
 - > insert into tablename child values (---);
 - > Select * from child; Sno

Testing:-

deletion in master table :-

- > delete from mas where sno = 1;

1 row selected.

- >Select * from mas;

- 7 Select * from Child;

On delete set null :-

Oracle also supports another clause on delete set null in child table.

- whenever we are using this clause if we are deleting P.K. values in master table then automatically appropriate foreign key values are set to null in child table.

Syntax Create table tablename (Col1 datatype(size)

References master table (Primary Key Column) on delete set null;

ales

Insertion in child table :- In oracle if we are trying to insert other than P.K. values into foreign key then oracle server returns an error ORA-2291.

→ BECAUSE in all d.b. system F.K values are based on P.K. values only.

⑤ Check :- This constraint is used to define logical conditions according to our business rules 25-7-14.

Column level :-

Syntax:-

```
Create table tablename (col1 datatype(size)
check(logical condition), ...);
```

➤ Create table test (sal number(10) check(sal > 5000));

➤ Insert into test values (2000);

error: Check constraint violated (SCOTT.SYS-00533)

➤ Insert into test values (8000); // \$ is now created.

➤ Select * from test;

SAL
8000.

Note:- Check constraint does not work with "Sedate".

It allows user defined date.

> Insert into

7 Create table test1(name varchar2(10))

```
check( name = upper (name) ) );
```

Y insert into test1 values ('murali');

ORA-02290: check constraint (SCOTT.SYS-0005336)

> insert into test1 values ('MURALI'); Violated

— ~~Row~~ Created

Table level :- (AND) operator

```
> create table test2 (name varchar(10), salt float(10));
```

Check (name = upper(name) and sal > 5000);

Assigns undeclare names to Constraints :-

- Generally when ever we are creating constraint

Oracle Server automatically generates an unique

identification member for identifying a constraint

uniquely in the format of Sys → 00056789 (or) Sys - cn sh

this is called "Predefined Constraint name".

In place of this one we can also create one

Own name using Constraint keyword. This is called

use defined Constraint name.

Syntax :-

| Constraint | user defined name | Constraint type |
|------------|-------------------|--------------------|
| Create | Pre defined | Constraint name :- |

- > Create table test3 (sno number(10) primary key);
error:- unique Constraint.
(SCOTT . SYS_C005338) violated

user defined Constraint name:-

- > Create table test4 (sno number(10) constraint P_fan primary key);
> Insert into test4 values(1);
> Insert into test4 values(1);
error:- unique Constraint

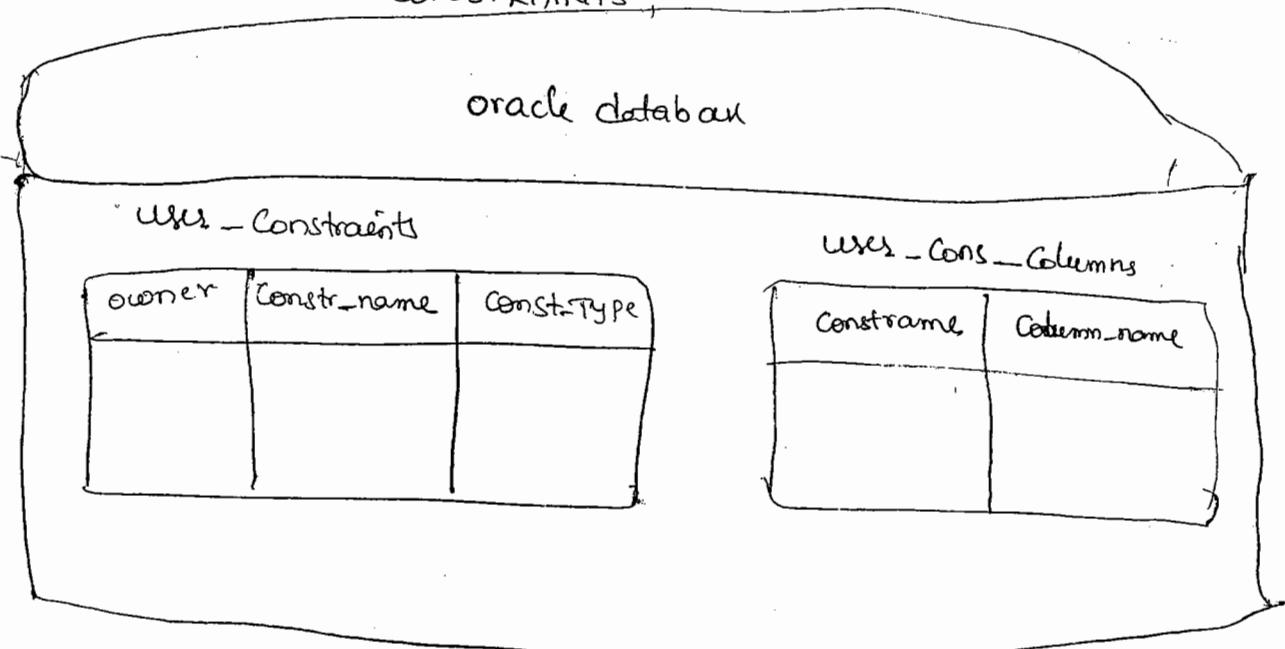
(SCOTT . P_fan) (SCOTT . P_FAN) violated.

Data Dictionary :-

when ever we are installing oracle Seem like a lot of tables are created
These read only tables are also called as Data Dictionaries

all constraints are stored in

Y Yes ! USER_CONSTRAINTS ;



Note 1: In oracle if u want to view constraint names
constraint types then we are using "USER_CONSTRAINTS"
Table

Data dictionary

SQL> desc USER_CONSTRAINTS;

Y Select constraint_name, constraint_type from

user_constraints where table_name = 'EMP';
after = Joins only

Constraint-Name Constraint-type

PK-EMP. P.

FK-EMP_DEPT. R. 26-7-14.

Note 2: If u want to view column names along
with constraint names then we are using

USER_CONS_COLUMNS user_cons_columns dictionary

Note: In Oracle if u want to view logical condition of the check constraint then we are using SEARCH - CONDITION property from user_constraint data dictionary.

> desc user_constraints;

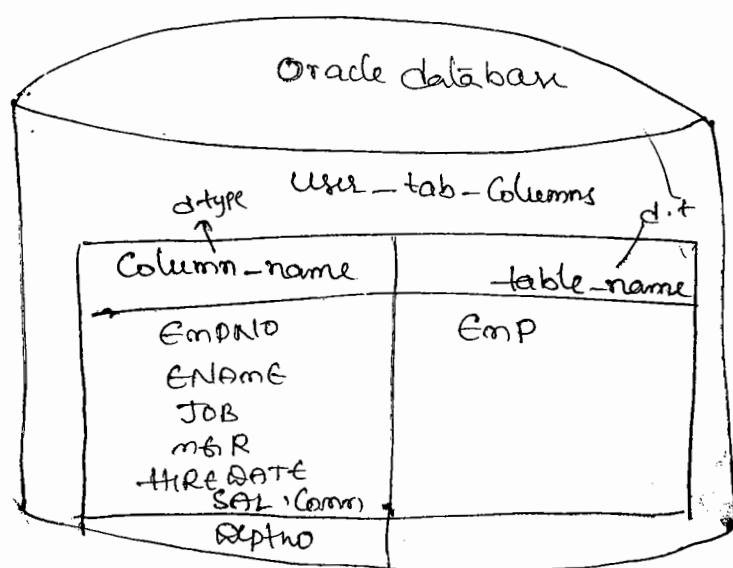
> Select SEARCH_CONDITION from user_constraints
where TABLE_NAME = 'TEST1';

SEARCH_CONDITION

name = upper(name)

→ Select count * emp; no. of rows → Count(*)
14
→ no. of columns in table.

Note: In Oracle all columns information stored under "user_tab_columns" data dictionary.



Ex: > desc user_tab_columns;

> Select Column_Name from user_tab_columns where table_name = 'EMP';

Column - Name

EMPNO

ENAME

JOB

MGR

HIRE DATE

SAL

Comm

Deptno

no. of columns no.

* If interview question

SQL to display no. of columns no. from emp table.

> Select Count(*) from user_tab_columns

where table_name = 'EMP';

O.P

| Count(*) |
|----------|
| 8 |

Default Constraint :-

In oracle we can also create default values for a column using default constraint.

Syntax :-

Columnname datatype (size)

default actualvalue

O.P:-

➤ Create table test (name Varchar2(10), sal number(10)
default 2000);

➤ insert into test(name) values ('abc');

➤ Select * from test;

name sal

abc 2000 - default

➤ Create table test (n

➤ insert into test(name) values ('abc'), sal=5000;

| name | sal |
|------|------|
| abc | 5000 |

— If u want to view default value for a column
then we are using DATA_DEFAULT property from
usr-tab-columns date dictionary.

➤ desc usr-tab-columns;

➤ Select Column-name, DATA_DEFAULT from usr-tab-
columns where table-name = 'TEST';

| Column-Name | DATA-DEFAULT |
|-------------|--------------|
| Sal | 2000 |

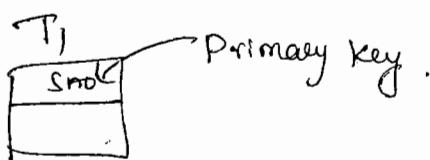
adding or dropping constraint on existing table
 we can also add and dropping constraint on existing table using 'alter' command.

Note :- new constraint :- Primary key

(1) If u want to add constraint on existing table existing column then we are using table level

Syntax method. (1st constraint then column ex primary (sno))

>

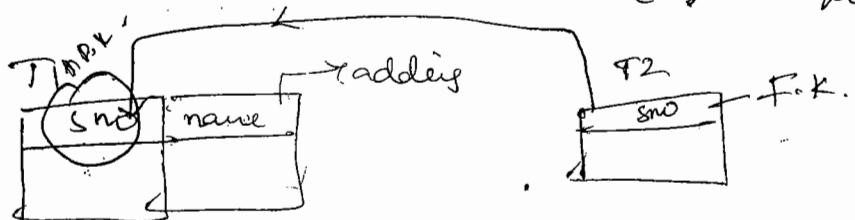


> Alter Table T1 add primary key (sno);

new column unique

(2) If u want to add a new column along with constraint then we are using "Column level Syntax" method

> alter table t1 add name varchar(10) unique;

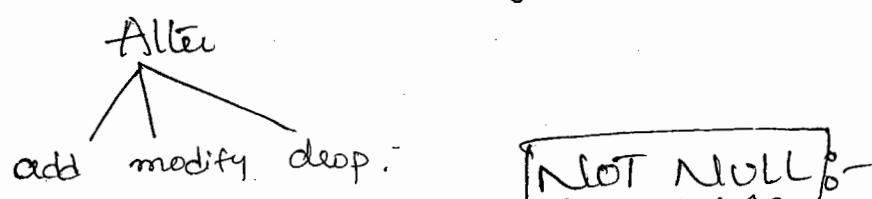


> adding foreign key } add 2 references

> Create table t2 (sno) number(10);

Table Created

> alter table2 add foreign key (sno) references t1;



Note: If u want to add not null constraint to
the existing table & existing column then we
are using alter with modify.

↓ problem is we are
modifying.

Syntax :-

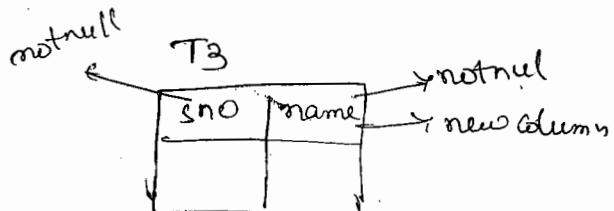
alter table tablename modify Columnname
not null;

> Create table t3(sno number(10));

> alter table t3 modify sno not null;

→ new column:-

> alter table t3 add name varchar(10) not null;



→ Copy 1 table to another table. Constraints are not copied.

> Create newemp as ~~copy~~ Select * from emp;

No6

Note:- Generally when we are copying 1 table from another table constraints are not copied.

In this case we must use alter command which is used to add constraint on existing table to existing column.

- > Create table newemp as Select * from emp;
- > Create table newdept as select * from dept;
- > alter table newdept add constraint D-hh
Primary key (deptno);
- > alter table newemp add constraint P-ab
Primary key (empno);
- > F.key
- > alter table newemp add foreign key (deptno)
references newdept (deptno) on delete cascade;
dropping Constraints Abs 28-9-14 monday.

Method 1:-

Syntax:- Alter table tablename drop Constraint Constraintname

Method 2:-

Syntax:- ① Alter table tablename drop primary key;

② Alter table tablename drop unique (col1, col2 ...);

from

Eg:- > Create table test (sno number(10) Constraint
P-ji primary key);

(or)

> alter table test drop primary key;

Note:- If you want to drop primary key along with referenced foreign key then we are using Cascade clause with alter-drop.

Syntax:- Alter table tablename drop primarykey cascade;

Eg:- > Create table t1 (sno number(10) primary key);

> Create table t2 (sno number(10) reference t1);

> alter table t1 drop primary key;

— error:- this unique/primary key is referenced by some foreign keys.

Solution:-

> alter table t1 drop primary key cascade;

→ Table altered

Eg!:- > Create table Y1 (sno number(10) Constraint ab notnull);

> alter table Y1 drop Constraint ab;

-- Search constraint name from user-constraint table

M. Imp Sub-queries

- Query within another query is called "Subquery" (or) "nested query"
- Sub queries are used to retrieve data from single or multiple tables based on more than 1 step process.
- All database Systems have 2 types of Subqueries
 - (1) Non Correlated Subqueries, 1st → C.Q
 - (2) Correlated Subqueries, 1st → P.Q
- In non correlated Sub queries, child query is executed 1st then only parent query is executed. whereas in for correlated Subqueries, parent query is executed 1st, then only child query is executed.

Non Correlated Subqueries :-

- Child Query : A query which provides value to another query is called child query.
- Parent Query : A query which receives value from another query is called parent query.

Non-Correlated Subqueries :-

- (1) Single row Subquery
- (2) Multiple row Subquery
- (3) Multiple Column Subquery
- (4) Inline views (or) Subqueries and in from clause

WAP to display the employees who are getting more than the average salary from emp table.

> Select * from emp where sal > (Select avg(sal) from emp);
→ Here, child query returns single value. That's why, these type of subqueries are also called as single row subqueries. In single row subqueries we are using =, >, >=, <, <=, <>, (!=) operators.

Execution process:-

Step 1:- > Select avg(sal) from emp;
→ 2101.78571.

Step 2:- > Select * from emp where sal > 2101.78571.

WAP to display the employees who are working in 'SALES' department using emp, dept table;

> Select * from emp where deptno = (Select deptno from dept where dname = 'SALES');

- { 1. Same columns
- 2. Group functions
- 3. Expression
- 4. diff. columns, but of same data type

WAP to display senior most employee from emp table

Select * from emp where HIREDATE = (Select min(HIREDATE from emp)),

(Hiredate = 14-2-1987 DATE)

30-7-14.

WAP to display the employees who are working in Smith dept from emp table.

> Select * from emp where

deptno = (Select deptno from emp
where ename = 'SMITH');

- WAP to display the employees who are working under Blake has manager from emp table using employee no, mgr columns.

- 1st write (find out) child query. result in parent query.

- Select * from emp where

mgr = (Select empno from emp where ename = 'BLAKE')

WAP to display lowest avg salary job from emp table.

Select job, avg(sal) from emp group by job

where ~~avg(sal)~~. avg(sal) = (Select low(avg(sal)) from emp);

- where clause is a parent clause
- Parent clause in having } { either where or having } it parent query

➤ Select job, avg(sal) from group by job having
 $\underline{\text{avg(sal)}} = (\text{Select } \underline{\min(\text{avg(sal)})} \text{ from emp})$
 nested group fun.

⇒ Oracle ORA - nested group fun without GROUP BY

Note:- In all database systems when child query contains nested group fun's then we must use group by clause in child query.

➤ Select job, avg(sal) from emp group by job having
 $\underline{\text{avg(sal)}} = (\text{Select } \underline{\min(\text{avg(sal)})} \text{ from emp group by job})$

JOB Avg(SAL).

CLERK 1137.5

SQL to display the jobs whose avg salary more than the managers avg salary from emp table.

⇒ Select job, avg(sal) from emp group by job having
 $\underline{\text{avg(sal)}} > (\text{Select } \underline{\text{avg(sal)}} \text{ from emp where job='MANAGER'})$

JOB Avg(SAL)

PRESIDENT 5000.

ANALYST 3000.

Select deptno, min(sal) from emp group by deptno

having min(sal) > (select min(sal) from emp where deptno=20).

| Deptno | Min(sal) |
|--------|----------|
| 30 | 600 |
| 10 | 1400. |

Each dept wise max salary,

6. Select deptno, max(sal) from emp
group by deptno.

| Dept no | max(sal) |
|---------|----------|
| 30 | 2850 |
| 20 | 3000. |
| 10 | 5000. |

Ques to display the employee details who are getting max-salary in each department from emp table

— select * from emp where sal = (select max(sal) from emp group by deptno).

— after where 1 column so we write subquery 1 column only

Error: Single-row subquery returns more than 1 row.
child query.

Allen 1000
Smith 900

900 \neq (2850
3000
5000)

900 = 2850
3000
5000

900 in (2850
3000
5000) ✓ 900 \neq n()

This is a multiple row Subquery bcz here child query returns multiple values,

For multiple row sub queries we are using "in", "all", "any" operators

Note:- We can also use "in" operator in single row sub queries.

> Select * from emp

where sal in (Select max(sal) from emp group by deptno).

> Select deptno, sal, ename from emp

where sal in (Select max(sal) from emp group by deptno).

| Deptno | Sal | ename |
|--------|------|-------|
| 30 | 2850 | BLAKE |
| 20 | 3000 | SMITH |
| 20 | 3000 | FORD |
| 10 | 5000 | KING |

31-07-14.

→ WAP to display the employees
who are work in either Sales or Research dept. from
emp, dept table.

> Select * from emp where deptno IN (select deptno from
dept where dname = 'SALES' OR dname = 'RESEARCH');
Here = is not allowed, bcz. child query returns multiple values
{ So, we cannot assign multi values to single variable }

TOP-N-Anaylsis:-

(1) Inline View.

(2) rownum.

① Inline View:

→ Oracle 7.2 introduced Inline Views.

→ Generally, we are not allowed to use "order-by"
clause in child query.

→ To overcome this problem, Oracle 7.2 introduced
Subqueries in from clause of the parent query this type
of queries are also called as inline views.

Syntax: Select * from (Select Statement);

Imp

② rownum :-

→ rownum is pseudo-column. It behaves like a table column.

→

→ rownum automatically assigns numbers (1, 2, 3...) to each row in a table at the time of selection.

Eg:- > Select rownum, ename from emp;

| RowNum | ENAME | |
|--------|-------|---------------|
| 1 | SMITH | Smith-AAA BCD |
| 2 | ALLEN | |
| 3 | WARD | |
| 4 | JONES | |

→ rownum having temporary values.

> Select rownum, ename from emp where deptno=10;

| RowNum | ENAME |
|--------|--------|
| 1 | CLERK |
| 2 | KING |
| 3 | MILLER |

Q:- Write a query to display 1st record from emp table using rownum,

→ Select * from emp where rownum = 1;

→ Q:- Write to display 2nd record from emp table using rownum.

& Select * from emp where rownum = 2;
Q:- no rows selected.

rownum does not work with more than 1 positive integer
but it work with <, <= operator.

Q:- write a query to display 1st 5 rows from emp table using rownum.

Select * from emp where rownum <= 5;

→ WAP to display first 5 highest salary employees from emp table using rownum.

Select * from (Select * from emp order by sal desc) where rownum <= 5;

→ WAP to display 5th highest salary of employee from emp table using rownum.

➤ Select * from (Select * from emp order by sal desc) where rownum <= 5

minus

Select * from (Select * from emp order by desc where rownum <= 4);

⇒ LTAQ to display 2nd row from emp table using rownum,

➤ Select * from emp where rownum >= 2
minus

Select * from emp where rownum <= 1;

Eg: ➤ Select * from emp where rownum between 1 and 5.

{ between 1 and 5 }
rownum <= 1 and >= 5 } o/p works

players

- > Select * from emp where rownum >= 1;
 - rows are selected.
- > Select * from emp where rownum > 1
 - no rows selected.

imp

Q) What to display last 2 records from emp table

using rownum?

> Select * from emp

minus

Select * from emp where rownum <= (Select
Count(*) - 2 from emp);

Note:- Whenever we are creating a respect
a reference name from Rownum in Line view that
alias names works with all SQL operators.

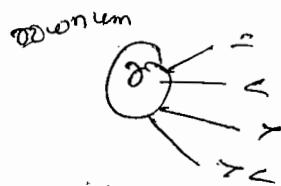
to 30

11.11

others

WAP to display 5th row from emp table using rownum alias name. (1-8-14)

Select * from (Select rownum δ , ename,



Sal from emp) where $\delta=5$;

| | | |
|---|--------------|------------|
| R | <u>ename</u> | <u>sal</u> |
| 5 | MARTIN | 1250. |

(or)

Select * from (Select rownum δ , e.* from emp e)
where $\delta=5$

WAP to display rows b/w 2 to 9 from emp table using rownum alias name.

Select * from (Select rownum δ , ename, sal from emp) where δ between 2 and 9;

| | | | |
|-----|----------------------------|--------------|------------|
| O.P | <u>δ</u> | <u>ename</u> | <u>sal</u> |
| | 2 | | |
| | 3 | | |
| | 4 | | |
| | 5 | | |
| | 6 | | |
| | 7 | | |
| | 8 | | |
| | 9 | | |

WAP to display 2nd, 5th, 7th, 9th rows from emp table using rownum alias name.

Select * from (Select rownum δ , ename, sal from emp)
where δ in(2,5,7,9);

or
rownum in(2,5,7,9)

using

WAP to display 1st, last row from emp table
using rownum alias name. $\text{Count}(\star) = 13^{\text{th}}$ record

Select * from (Select rownum σ , ename, sal from emp)

where $\sigma=1$ or $\sigma = (\text{select Count}(\star) \text{ from emp})$.

(or) $\sigma \in (1, 14)$.

R ENAME SAL.
1 SMITH 800
14 MILLER 1300

— WAP to display even no. of records from emp table using rownum alias name.

— Select * from (Select rownum σ , ename, sal from emp)

where $\sigma \mod 2 = 0$ $\sigma \mod 2 = 0$

where $\sigma \mod 2 = 0$. even

$\mod(\sigma, 2) = 1$. odd 3

2
4
6
8
10
12
14

(%)
mod
 $\sigma \mod 2 = 0$

WAP to display 5th highest salary employee from emp table using rownum alias name.

— Select * from (Select rownum σ , ename, sal from

(Select * from emp order by sal desc)) where $\sigma=5$

where R ENAME SAL.
5 BLAKE 2850.

σ en sal

Analytical Functions used in SQL views :- (A.6)

Oracle '8i' introduced analytical fun's there are

(1) row-number()

(2) rank()

(3) dense_rank()

These 3 analytical fun's automatically assigns no's to the rows either group wise or row wise in a table.

Syntax:-

Analytical function name() over

(partition by Columnname order by Columnname [asc/dsc])

| | denserank | rank |
|------|-----------|----------|
| 2000 | 1 | 1 - 2000 |
| 2000 | 1 | 2 - 2000 |
| 2000 | 2 | 3 - 2000 |
| 1000 | 3 | 4 - 1000 |

row-number():- This analytical fun automatically

assigns diff. rank no's when values are same

where as rank, denserank analytical fun's automatically

assigns same rank numbers when values

are same. ($\frac{2000}{2000}$)

- All(a,b) → Rank skips next consecutive rank no's
 i.e. → where as dense rank does not skip next consecutive rank numbers.

| <u>emp</u> | <u>deptno</u> | <u>ename</u> | <u>sal</u> | <u>r</u> |
|------------|---------------|--------------|------------|----------|
| | 20 | SCOTT | 3000 | 1 |
| rows | 20 | FORD | 3000 | 2 |
| rows | 20 | JONES | 2975 | 3 |

Rank() :-

| <u>deptno</u> | <u>ename</u> | <u>sal</u> | <u>r</u> |
|---------------|--------------|------------|----------|
| 20 | SCOTT | 3000 | 1 |
| 20 | FORD | 3000 | 1 |
| 20 | JONES | 2975 | 3 |

} 2 skips i.e. rank().

dense_rank() :-

| <u>deptno</u> | <u>ename</u> | <u>sal</u> | <u>r</u> |
|---------------|--------------|------------|----------|
| 20 | Scott | 3000 | 1 |
| 20 | Ford | 3000 | 1 |
| 20 | Jones | 2975 | 2 |

Select * from (Select deptno, ename, sal, row_number()
 over (partition by deptno order by sal desc))

My
 ne
 emp) where r <= 10;

interview → to display 2nd highest salary employee in each department from emp table using analytical fun.

> Select * from (Select deptno, ename, sal, dense_rank() over (partition by deptno order by sal desc) & from emp)
 where r = 2

O.P :-

| Deptno | ename | sal | r |
|--------|-------|------|----|
| 10 | clark | 2450 | 2 |
| 20 | Jones | 2975 | 2 |
| 30 | Allen | 1600 | 2. |

WAP to display 5th highest sal employee from emp table using analytical fun.

- Select * from (Select deptno, ename, sal,

dense_rank() over (Partition by ~~Deptno~~ Order by ~~Deptno~~
sal asc) r from emp) where r=5;

O.P
10 — 2
20 — 5
30 — 3

but we want only 5th highest overall
table.

> Select * from (Select deptno, ename, sal,

dense_rank() over (Order by sal desc) r from
emp) where r=5

deptno ename sal r
10 clark 2450 5

Note:-

In analytical fun's partition by clause is an optional
clause.

5-8-We.

AAB MXAAA AB. now add

AAB
AAB

P's are

rowid :- row address. (deleting duplicate rows).

row id is a pseudo column it behaves like a table

Column - row id uniquely identifying a record in a table

→ Generally when we are inserting data into table then Oracle Server automatically creates an unique identification no. for identifying a record uniquely within the block.

→ Using row id we are retrieving data very fastly from the database.

→ Generally "rownum" having temporary values where as "rowid" having fixed values.

→ Select rownum, rowid from emp;

→ Select rownum, rowid, ename from emp where deptno=10;

→ We can also use max, min fun's in rowid

Ex: ① Select min(rowid) from emp;

② Select max(rowid) from emp;

→ Generally row id's also used to delete duplicate rows in a table.

W.A.Q to delete duplicate rows except 1 row in each group from a table.

> Select * from test

of 20
Interview
@

> delete from test where
from test group by Sno;

> Select * from test;

SNO
10
20
30
40
50.

~~min(rowid)~~
~~(00)~~

~~rowid not in (Select max(rowid))~~

~~not correlated~~

Multiple Column Subqueries

In all d.b. Systems we can also Compare multiple

Column Values of the child query with the multiple

Column Values of the parent query.

These type of queries are also called as

"M.C. Subqueries"

Parent where clause compulsory exists '()'

multiple,

Parent query \rightarrow where '()'

When we are using multiple Column Subqueries when we must specify Parent query where Conditional

each
Glimpses within "Panther's U"

> Select * from tablename
where (col1, col2...) in (Select col1, col2 from tablename
where Condition);

WAP to display the employees whose job, mgr match
with the job, mgr of the employee 'Scott' from emp-table

> Select * from emp where (job, mgr) in (Select job, mgr
from emp where ename = 'Scott');

| ename | job | mgr |
|-------|---------|-----|
| ford | Analyst | |
| scott | Analyst | |

WAP to display the employees who are getting max.
Salary in each department from emp table using
multiple row Subquery.

> Select deptno, sal, ename from emp where
sal in (Select max(sal) from emp group by deptno);

| Deptno | Sal | ename |
|--------|------|-------|
| 30 | 2850 | blake |
| 20 | 3000 | scott |
| 20 | 2850 | FORD |
| 10 | 5000 | King |

- > Was to display the employees who are getting max. salary in each department from emp table using multiple column Subquery.
- > Select deptno, sal, ename from emp where(deptno, sal) in (Select deptno, max(sal) from emp group by deptno);

| Deptno | Sal | ename |
|--------|------|-------|
| 30 | 2800 | Blake |
| 20 | 3000 | Scott |
| 10 | 5000 | King |

Was to display Senior most employees in each job from emp table using multiple Column Subquery.

- > Select job, hiredate, ename from emp ~~where(job, hiredate)~~ in (Select job, min(hiredate) from emp group by job).

| job | hire date | ename. |
|-----|-----------|--------|
| | | |

WAP to display ename, dname, sal of the employees whose salary, comm match with the sal, comm of the employees working in the location 'DALLAS'.

→ Select ename, dname, sal from emp e, dept d where e.deptno = d.deptno.
and (sal, NVL(comm, 0)) in (Select sal, NVL(comm, 0))
from empe, dept d
where e.deptno = d.deptno and loc = 'DALLAS'.

6-8-W.

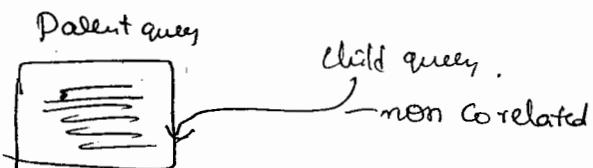
Correlated Subquery:- Generally in non-correlated subqueries child query is executed first. Then only parent query is executed. Whereas in Correlated Subqueries parent query is executed 1st then only child query is executed.
→ When ever we are submitting Correlated Subquery then database servers get a Candidate row from parent query table and then control passed into child query where clause and based on the evaluation value of the where clause it compare values in parent query.

→ In Correlated Subqueries we must Create alias name for the parent query table in parent query and use this alias name in child query 'where' Condition.

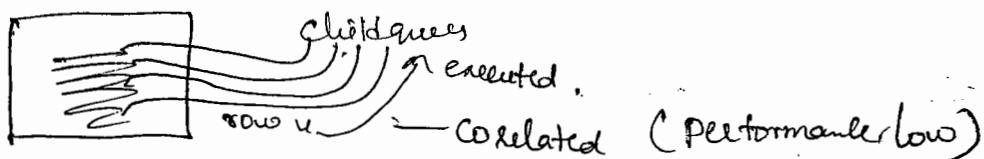
Syntax :-

Select * from tablename aliasname1

Where Columnname operator (Select * from tablename
where columnname = aliasname1. Columnname)



In non correlated Subqueries child query is executed only 1e per parent query table.



→ whereas in correlated Subqueries child query is executed for each and every row for parent query table.

→ Generally Correlated Subqueries are used in de-normalization process. (combined)

In this process they are using Correlated updates.

→ Generally If we want to modify 1 table Column values based on another table Columns then we are using Correlated update.

date tablename1 aliasname1 set

Column name = (select Columnname from tablename2

aliasname1 where aliasname1. Common Colname =
aliasname2. Common Colname);

Ex:-
in

SQL> alter table emp add dname VARCHAR2(10);

SQL> update emp e set dname = (select dname from
dept d where e.deptno = d.deptno);

SQL> Select * from emp;

| ename | eno | sal | mg | | dname |
|-------|-----|-----|----|-------|-------|
| | | | | | |

Ques - To display the employees who are getting more
than the avg salaries of their jobs from emp table
using Correlated Subquery.

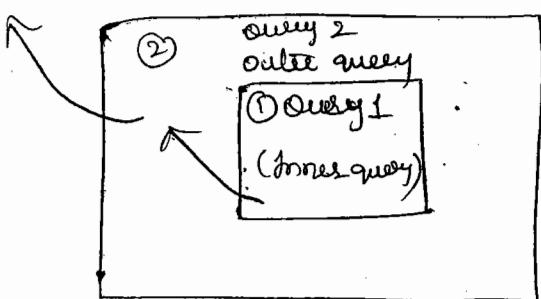
SQL> Select * from emp e where sal > (Select avg(sal)
from emp where job = e.job)

empno ename job mbr hiredate sal comm deptno.

- In non Correlated Subqueries always inner query is independent from the outer query, whereas in
- In Correlated Subquery always inner query depends on outer query. i.e., inner query, outer queries are correlated.

non Correlated Subquery

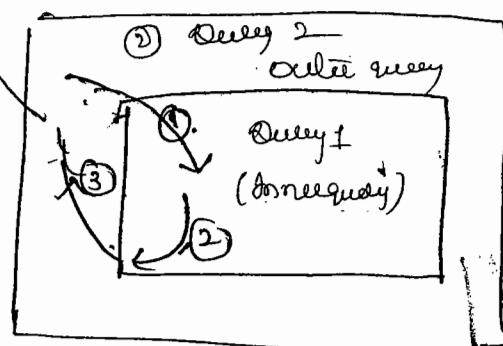
final output



uni direction

Correlated Subquery

final P.



Bi-direction

Ques to display 1st highest salary employee from using Correlated Sub query.

> Select * from emp e₁ where 1 = (Select count(*) from emp e₂ where e₂.sal >= e₁.sal).

O.P:

| |
|------------|
| KING \$100 |
|------------|

How to display 2nd highest Salary employee from the following table using Correlated Subquery.

➤ Select * from test;

| ENAME | SAL |
|-------|-----|
| abc | 100 |
| xyz | 150 |
| bbb | 200 |
| ccc | 300 |

Select * from test e₁ where 2 = (Select Count(*) from test e₂ where e₂.sal >= e₁.sal);

| ENAME | SAL |
|-------|-----|
| bbb | 200 |

Execution process.

Phase 1 :-

Step 1 :- get a Candidate row (1st row) → abc 100

Step 2 :- Chq → Select Count(*) from test e₂ with 100
where e₂.sal >= e₁.sal

Step 3 :- P.Q → Select * from test e₁ where 2=4; (False).

Phase 2 :-

Step 1 :- get a Candidate row (xyz 150)

Step 2 :- Select Count(*) from test e₂ with 150.
where e₂.sal >= e₁.sal

Step 3 :- Select * from test e₁ where 3=4; (False).

Phase 3 :-

Step 1 :- (bbb, 200)

Step 2 :- $e_2.\text{sal} >= 200$.

[2]

Step 3 :- $2 = 2$ (true)

Retrieves O.P is

| | |
|-----|-----|
| 200 | bbb |
|-----|-----|

7-8-14.

Note :- when our resource table contains duplicate data then above query does not return values to overcome this problem we must use "distinct clause".

Error :- $\times \text{insert into test values ('zzz', 200);}$

$\times \text{Select * from test e1 where } 2 = (\text{select count(*)}$
from test e2 where $e_2.\text{sal} >= e_1.\text{sal})$;

— no rows selected. (bcz of duplicate values)

Solution :-

$\text{Select * from test e1 where from test}$

$2 = (\text{select count}(\text{distinct}(\text{sal})) \text{ from test e2 when}$
 $e_2.\text{sal} >= e_1.\text{sal})$ \rightarrow duplicate values will display.

| ename | sal |
|-------|-----|
| bbb | 200 |
| zzz | 200 |

WAP to display 2nd highest salary employee from the above table using Correlated Subquery ($n-1$) method

- > Select * from test e₁ where (2-1) = (Select count (distinct (Sal)) from test e₂ where $e_2.\text{Sal} \geq e_1.\text{Sal}$)

O.P

bob 200.
bob 200.

④

don't use '='
in ($n-1$)

WAP to display n^{th} highest salary employee from emp table using Correlated Subquery.

- > Select * from emp e₁ where &_n = (Select count (distinct (Sal)) from emp e₂

where $e_2.\text{Sal} \geq e_1.\text{Sal}$.

= is there in n

> Enter value for n: 1

KING 5000.

EXISTS operator \exists (In Correlated only)

- Exists operator used in Correlated Subqueries.

- Exists operator always returns boolean value either true or false.

- Exists operator performance is very high compare to in operator.

- Exist operator used in "where" clause only.

* When we are using exists operator then we are not allowed to use Column

~~We are not allowed to use Colname.~~

Syntax:

Select * From tablename aliasname

where exists (Select * From tablename

where Colname = aliasname • Colname);

→ exists {} \Rightarrow False. {}₃-set means
↳ empty

Child query gives values.

→ exists {1,2,3} \Rightarrow True.
↳ nonempty

⇒ Exist operator is used to test whether a given set is empty or non empty

⇒ Exist operator on empty set returns False

whereas Exist operator on non-empty set

returns "True". (child query test true or false)

Ex:-

exists {1,2,3} \Rightarrow True

exists {} = False.

→ Generally if we want to test 1 table Col Values are available in another table then we are using Exists operator.

WAP to display those departments from dept table
those departments ^(employees) are available in emp table using

Correlated Subquery exists

Parent query alias name mandatory, child query alias name \rightarrow

> Select * from dept d where exists (Select * from
emp where deptno = d.deptno);

O.P.

10 Accounting N.Y
20 Research Dallas
30 Sales Chicago.

values
given

ex

\rightarrow exists always Test whether a given set is empty or
non empty in child query. (*, const, for..., char also works)

WAP to display those departments from dept table
does not have employees in emp table using Correlated
Subquery. (not exists)

> Select * from dept d where not exists (Select * from
emp where deptno = d.deptno);

| Deptno | Dname | Loc |
|--------|------------|--------|
| 110 | Operations | Boston |

WAP to display those departments from dept table
does not have employees in emp table using not-
exists

Correlated Subquery.

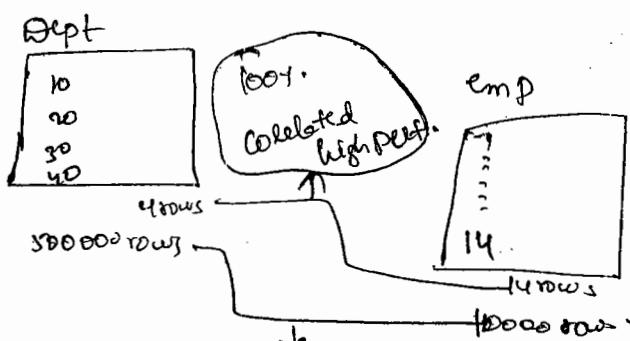
> Select * from emp;

> Select * from dept where deptno not in (select deptno from emp);

| | | | |
|--------|----|------------|--------|
| O.P :- | 40 | Operations | Boston |
|--------|----|------------|--------|

→

co
or 40 → none



This query is correlated also

O.P :- 40. bt Correlated
is very high perf.

not in operator does not work with Null values.

Note :- Generally not in operator does not work with null values to overcome this problem we must use not exists operator using correlated subqueries.

Ex :- > Insert into

emp (empno, ename, deptno) values

(1, 'murali', null);

> Select * from dept where deptno not in

Select deptno from emp);

→ no rows selected.

> Select * from dept d where not exists

(Select * from emp where deptno = d.deptno);

| eno | ename | loc |
|-----|------------|--------|
| 40 | operations | Boston |

Select

In where condition Is null, Is not null.

in correlated

related sub

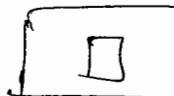
Index does not Search scenes working

8-8-14.

bt Correlated
very high perf.

Correlated Sub Query:-

non-correlated



all.
any.

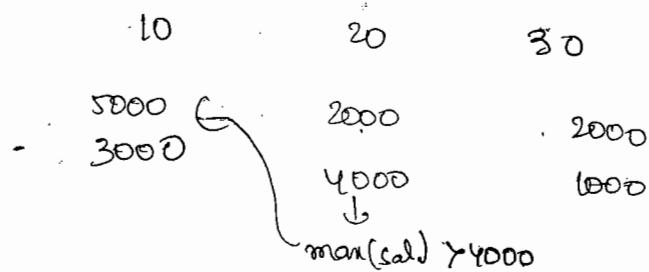
Correlated



exists.

mane.
Subquery:-

work with
most case
to display the employees who are getting more
than the highest paid employee in 20th dept from
emp table.



>Select * from emp where sal > (Select max(sal) from
emp where deptno = 20) O.P.

Who to display the employees who are getting more
than the lowest paid employee in 10th dept. from emp table

>Select * from emp where sal > (Select min(sal) from
emp where deptno = 10) O.P. \equiv grows

— When ever resource table having large amount of data, & also child query contains max or min functions & also we are comparing values using relational operators then those types of queries may slowly retrieve data from db.

This process automatically degrades performance of the application.

To overcome this problem for improve performance of a query, ansi/iso sql introduced Subquery special operators. There are "all, any".

These operators are used along with relational opr's ($<$, $>$, $=$, ...)

➤ Select * from emp where

min/max = all.
max/min = any

Sal > all (Select Sal from emp where
deptno=20);

➤ Select * from emp where

Sal > any (Select Sal from emp where
deptno=10);

amount

2 or min

sing

ees very

performance

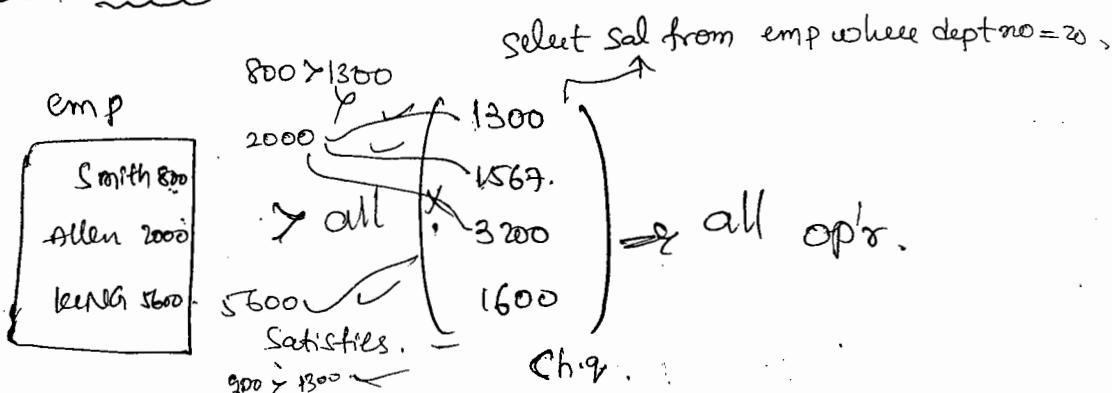
le prence

ey special

with

all, any operators and in multiple row

Sub queries :-



In \Rightarrow It returns same values in the list.

All (\geq) :- It satisfies all values in the list

Any ($>$) :- It satisfies any value in the list

more performe
than 'In'
 $=$ Any is nothing bt internally In op's

Q1A Q to display the employees who are getting more than the all salaries of the 'clerk's' from emp table using Subquery Special operator (All)

Order by clause \Rightarrow not in child query It won't work
that's w/ we have views.

after child query possible
closing

\Rightarrow Select * from emp where sal \geq all (Select sal
from emp where job = ('clerk')) Order by sal desc -

O.P :-
800
2000
500

without
O.P
3200
5000
3200

Note:- when we are using Subquery Special opr

"All" then optimizer internally uses logical opr "AND"

→ when ever we are using Subquery special opr

"Any" then internally optimizer uses logical ops "OR"

$$\begin{aligned} \text{All} &= \text{AND} \\ \text{Any} &= \text{OR} \end{aligned}$$

Ex:-

Select * from emp where

deptno > all (10,20);

O.P:- 30.

(10) \uparrow
deptno > 10 and
deptno > 20.

Tak

>Select * from emp where

deptno > any (10,20);

O.P:-
20.
30.

(10) \uparrow (20)
deptno > 10 or
deptno > 20.
 \downarrow
(30)

View

and

→

→

Note:- " \geq any" is also same as " \in " operator.

call

Bt " \geq any" not same as not in operator.

$$\begin{aligned} \geq \text{any} &= \in \\ \geq \text{any} &\neq \text{not in} \end{aligned}$$

Pa

not

↓ o

real opr
! opr "AND"
 \neg

Y Select * from emp where deptno not in (10,20); \neg any
O.P. :- 30.

real opr
! opr "OR"
 \exists

Y Select * from emp where deptno
 \neg any (10,20);
O.P. :- 10,20,30

\neg any
not equal or
 $\neq 10$ means - 20,30
 $\neq 20$ & - 10,30
SD 10,20,30

Views.

and
Table, cluster, d.b. objects.
synonyms

View is a database object which is used to provides authority level of security.

- Generally views are created from "base tables"
- View does not store data. View is also called as "Virtual table" & also "Window of table"



Select col. only give to another user not all cols. Then restrict cols.

Give to another user we use Views

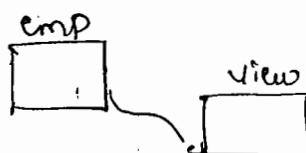
→ Generally if we want restrict table columns from one to another we are using 'Views'.

→ Based on the base tables views are categorised into 2 views.

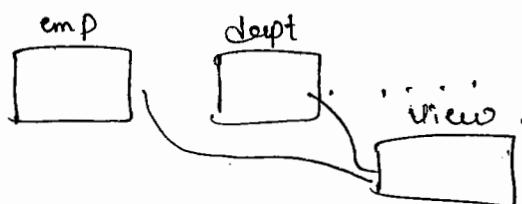
(1) Simple View

(2) Complex View. (or) Join View.

→ Simple View: Is a view which is created from only 1 base table.



→ Complex View: Is a view which is created from no of base tables.



Simple View:-

Syntax:-

Create or replace view Viewname as

Select Stmt;

e Colms

are using

Categorised

(Select stmts are internally stored)
~~view~~-views

9-8-14

- In all d.b. Systems when we are creating views then automatically view definitions are permanently stored in database.
- In oracle if you want to view, view definitions then we are using "user_views" data dictionary.

Ex) > desc user_views;

> Select TEXT from user_views where
VIEW_NAME = "V1";

TEXT

Select "SNO", "Name" from bat.

- Create a view.
- Create or replace view V1

as

Select * from emp where
deptno=10;

- select * from V1;

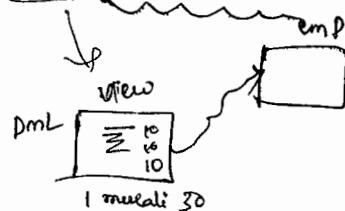
DML Operations on Simple View (to base-table):

Rule 1: If a simple view contains group fun's, group by clause, rownum, distinct, set operators, Joins then we cannot perform DML ops through simple view to base-table.

Rule 2: We must include base table "not null" column into the view then only we can perform "insertion" operation through View to base table.

Note
or e
an c

Ex-1 Create or replace view V₁ as



Select * from emp where deptno=10;

> Select * from V₁;

> Insert into V₁ (empno, ename, deptno)

effect in
base table
Here for V₁

Values (1, 'murali', 30);

1 row created.

> Select * from emp;

We can not see
this record bcz

in where condn deptno
is '10' so

We can see in base
table (emp) only

This 1 murali 30 is available.

Example 2:

> Create or replace view V₂

as

Select * from ename, sal, deptno

from where deptno=10;

> Select * from V₂;

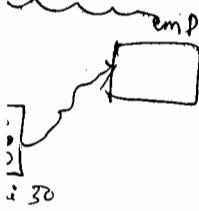
there is no
emp no.

> Insert into V₂ (ename, sal, deptno)

Values ('abc', 2000, 20);

Error: Cannot insert Null into empno.

if null
in perform
x-table.



Effect is
tabletable
ee for V,
as

- Can not see record bcz
2. Cond'n deptno
' 10
- See in ban
emp only

Note:- when a simple view contains group fun's or expressions or rownum then we must create an alias name for these expressions.

Ex:- Create or replace View V₃
as.

Select deptno, max(sal) from
emp

group by deptno;

Ques: must name this expression with a column alias

➤ Create or replace View V₃ ①
as

Select deptno, max(sal) "maximum"
alias name

from emp group by deptno;

➤ Select * from V₃;

Deptno maximum

Or

| | |
|----|------|
| 10 | 5600 |
| 20 | 3200 |
| 30 | 1525 |

→ alias name creation is 2 types

① with u v

② after select @ View create alias name

like V₃(deptno, maximum)

→ Create or replace view V₃ (deptno, max(sal))

as

Select deptno, max(sal) from emp

group by deptno; O.P. same

| |
|------------------------|
| >desc V ₃ ; |
| Name |
| Deptno |
| max(sal) |

→ Rec
in

→ Create or replace view V₄

as

Select rownum num, ename from emp;
alias

then

→ Select * from V₄ where num > 5;

| O.P. | rownum | ename |
|------|--------|----------|
| 6 | 1 | |
| 7 | 1 | |
| 8 | 1 | |
| 9 | 1 | |
| 10 | 1 | |
| 11 | 1 | |
| 12 | 1 | |
| 13 | 1 | |
| 14 | 1 | |
| 15 | 1 | parimali |

W.A.Q to create a view from Emp table to display each deptmt wise total salary (sum) from Emp, Dept tables & also display alias names as dname, total salary.

Ex:

→ Create or replace view V₅ (dname, total(sal))

→ test:

as

Select dname, sum(sal) from emp, dept d

error

Where e.deptno = d.deptno

group by dname;

➤ Select * from V4;

| DName | TotalSalary |
|------------|-------------|
| Accounting | 9325 |
| Research | 10887.5 |
| Sales | 7775. |

➤ Read only Views:-

→ When ever we are using with read only clause then we can not perform DML operations through View to base-table.

Syntax :-

Create or replace View Viewname

as

Select * from tablename where condition

With read only;

Ex:-

➤ Create or replace View V4

as

Select * from emp where deptno=10

With read only;

testing:-

➤ delete from V4 where deptno=10;

Error :-

Cannot delete from View.

With check option :- If we want to provide constraint type mechanism on views then we are using "With check option" clause.

When a view contains with check option clause then we can insert other than where condition values through view to base table.

Syntax :-

Create or replace view Viewname
as
Select * from tablename
where condition with check option;

Ex:- Create or replace view V5
as
Select from emp
where deptno=10
with check option;

> Insert into V5 (empno, ename, deptno)
Values (2, 'abc', 30);

Error:- View WITH CHECK OPTION where - clause
Violated

> Insert into V5 (empno, ename, deptno)
Values (2, 'abc', 10);
1 row inserted.

forced
we C
those
Syntax:-

Ex:-
G

Error:-

Sol:-
G

a
S

Warning

> creat

> alter

> des

to provide
so then we
check option
'Condition'

forced View or force View (Interview Purpose)
we can also create view without base table
those views are also called as forced views

Syntax:-

Create or replace force view Viewname
as
Select * from Anyname;

Ex:-

Create or replace View V₁
as
Select * from Sunday;

Error:-

table or view does not exist.

Sol:-

Create or replace^{force} View V₁
as
Select * from Sunday;

Warning:- View created with compilation errors.

- > Create table Sunday (Sno number(10));
- > Alter View V₁ Compile;
- > desc V₁;

Learn

11-8-14.

To

Complex View (or) join views-

Complex View is a view which is created from multiple tables.

Ex:- Create or replace View V₅

as

Select ename, sal, dname, loc
from emp, dept

where emp.deptno=dept.deptno;

— View Created

> Select * from V₅

> update V₅ Set ename='abc' where ename='SMITH';

I now updated.

> update V₅ Set dname='xyz' where dname='SALES';

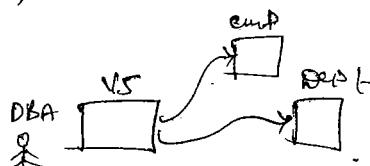
Error:-

In all d.b. Systems generally we can not perform

DML operations through Complex view to Base tables

i.e., Here Some table cols are effected some

Other table cols are not effected.



trigger

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Ex:-

> S

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TR

Trig

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To overcome this probm oracle introduced

"Instead of triggers".

By default these triggers are low level triggers.

In oracle if u want to view effectable or non effectable cols then we are using

"User-updatable-Columns" data dictionary.

Ex:- DESC user-updatable-columns;

Select COLUMN-NAME, UPDATABLE from user-updatable-columns where TABLE-Name = 'W6';

| COLUMN-NAME | UPDATABLE |
|-------------|-----------|
| ENAME | Yes |
| SAL | Yes |
| DNAME | No |
| LOC | No. |

TRIGGERS (PL/SQL) :-

Triggers is also same as "Stored Procedure" and also it will automatically invoked when ever DML oprns performed on Table.

All D.B. Systems having 2 types of Triggers

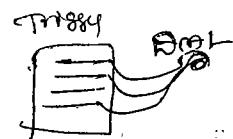
(1) Statement level Triggers.

(2) row level Triggers.

① In Statement level Triggers :- Trigger body is executed only once per DML Starts.



② row level Triggers :- Trigger body is executed for each row for per DML Starts.



Syntax:-

Create or replace Trigger Triggername

before/after insert/update/delete on tablename.

[for each row] → stmt level
→ row level trigger.

Trigger Specification

Trigger body
begin
end; / compulsory

Difference b/w Stmt level, Row level triggers :-

Ex:- row level Triggers :-

> Create table test (cols data);

-triggers

Start level trigger :-

Create or replace trigger TK1

after update on emp

begin

insert into test values (sysdate);

end;

/

testing :-

> update emp set Sal = Sal + 100 where Deptno=10.

3 rows updated.

> Select * from test;

Col

11-AUG-14. - ①

> drop trigger TK1;

> delete from test;

Row level trigger :-

Create or replace trigger TK2

after update on emp for each row

begin insert into test values (sysdate);

end;

/

say.

triggers :-

testing :-

update emp set sal = sal + 100

where deptno = 10;

3 rows updated.

> Select * from test;

COL1

11-AUG-14

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11-AUG-14.

Row level triggers :- In oracle in row level trigger

Trigger body is executed for each row for

DML Stmt - That why we are using for each row

class in trigger specification.

And also DML transaction values are automatically stored in two rollback start qualities there are "old,new".

These buffers are used in either in triggers specification or in trigger body.

When we are using this qualities in trigger body then we must use ';' in front of the

we

Sy

:ol

Sy

:n

:

WIA

when

auto,

Ex:-

>

>

>

>

>

a

f

We must use `:` in front of the qualified name

Syntax:-

:old . column name

Syntax:-

:new . column name

| | insert | update | delete |
|------|--------|--------|--------|
| :new | ✓ | ✓ | ✗ |
| :old | ✗ | ✓ | ✓ |

WFA PL/SQL row level trigger on employee table

whenever user deleting data on emp table then automatically deleted data stored in another table.

Ex:-

Create or replace trigger

> Create table backup as
Select * from emp where 1=2;

> Select * from backup;
no rows Selected.

> desc backup;
> Create or replace trigger th1
after delete on emp

for each row begin

| | merge

Sys
Create or replace trigger th,

after delete on emp

for each row

begin

insert into backup

values (:old.empno, :old.ename, ?old.job, ?old.sal,
----- :old.deptno);

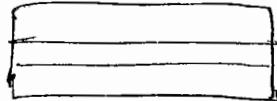
end;

/

testing:-

- Y delete from emp where sal > 2000;
- Y select * from backup;

rows



Instead of triggers

generally we can not perform DML operations
through Complex views to base table.

To overcome this problem Oracle 8.0
introduced Instead of triggers.

Instead of triggers are created on views.

→ By default instead of triggers are low level triggers.

Syntax :-

create or replace trigger trigger name

instead of insert | update | delete on view name

for each row

begin

=====

end ;

12-8-14.

Sol:-

create or replace trigger tm1

instead of update on v5

for each row

begin

update dept set

dname = :new.dname where

dname = :old.dname;

update dept set loc = :new.loc where loc = :old.loc;

end ;

Testing :-

> update v5 set dname = 'XYZ'

where dname = 'SALES';

rows updated -

on views,

triggers.

> desc user_updatable_Columns;

wh

> Select COLUMN_NAME, UPDATABLE from

date

user_updatable_Columns

where TABLE_NAME = 'V5';

COLUMN_NAME

UPDATABLE

ENAME

YES.

SAL

YES.

DNAME

YES.

LOC

YES.

Materialized Views :-

Oracle 8i introduced materialized views these views
are handle by db. administrator.

materialized views are used in d.b.t.h.
applications.

generally views does not store data where
materialized views stores data.

→ generally materialized views are used to
improve performance of the join or aggregatable
queries.

materialized view stores result of the query,
" " " replication (copy) of the remote
d.b. into local node.

when we are refreshing mat'zed views or Synchronize
data based on "base table".
mld view only Adminstrator

Syntax:-

```
Create materialized View Viewname  
as  
Select Statement ;
```

* Before we are creating materialized view DBA's
must give create any materialized view privilege to
(Permission)

Syntax:-

```
grant create any materialized View to Username;
```

then views

Ex:- Create materialized View mv1 ,
as
Select * from emp;

Error:-

InSufficient Privileges.

note view creates
only Adminstrator
So when we create
error occurs
So we will use
priv to grant

- > Conn Sys as sysdba;
- Enter password: sys
- > Conn scott / tiger;
- > Create materialized View mv1 ,
 as
 Select * from emp;

```
Materialized View Created.
```

1.6.11.

state where

are used to

negotiable

key,

the remote

Note: In oracle sometimes views also returns an error insufficient privileges.

To overcome this problem we must give

Create any view privilege to any user

Syntax:-

grant create any view to username;

Y Create or replace view v1

as

select * from emp;

Error: Insufficient privileges.

Y Conn sys as sysdba;

Enter Pwd : Sys

Y grant create any view to scott;

Y Conn scott/tiger;

Y Create or replace view v1,

as

Select * from emp;

View created

retrives an Note: In oracle 1 of the materialized base table must contain primary key otherwise oracle server returns an error.

Ex: Create table test (sno number(10));

>Create materialized view mn,
as

Select * from test;

Error: table 'test' does not contain a primary key

Constraint

Diff b/w View, materialized view :-

>Create table base (sno number(10) primary key,
name Vaishnav (10));

Insert into ^{base} values (.);

Select * from base;

Create view

| SNO | NAME |
|-----|------|
| 1 | a |
| 2 | b |
| 3 | c |
| 4 | d |

Create or replace view v1

as

Select * from base;

View created.

Create mat'd view.

➤ Create materialized view mK1
as

Select * from book;

⇒ In this case materialized view is also same as
view.

— But when ever we are creating mat'd view also
mat'd view definitions automatically stored in dbase
Same like a view definitions.

— For oracle if u want to view their materialized view
definitions then we are using user_mviews catalog

➤ desc user_mviews;

➤ Select query from

user_mviews where

MVIEW_NAME = "mK1";

➤ Select rowid, sno, name from book.

rowid sno name

➤ Select rowid, sno, name from V1;

In this case View rowid's all same as book-table
rowid's

— That's why view does not store data

— Through the view we can alter book-table data

That's why Views are also called as
"Virtual tables".

➤ Select rowid, Sno, Sname from mki;

→ Same as
Here mat'd view rowid's are diff from base table
rowid's.

→ View also stores data.

➤ update base set

name = upper(name);

➤ select * from V1;

| SNo | Name |
|-----|------|
| 1 | A |
| 2 | B |
| 3 | C |
| 4 | D |

➤ Select * from mki;

| SNO | Name |
|-----|------|
| 1 | a |
| 2 | b |
| 3 | c |
| 4 | d |

If it is not effected ↪

→ Mat'd View also stores data same like a-table
but when we are refreshing material'd view it
synchronizes data based on base-tables.

→ In oracle If u want to refresh mat'd view then
we all using refresh method from

dbms_mview package.

Syntax:

Ex:-

```
dbms_mview.refresh('materialized Viewname');
```

Package name.

→ ^{R.W} exec dbms_mview.refresh('mri');

| SNO | Param |
|-----|-------|
| 1 | A |
| 2 | B |
| 3 | C |
| 4 | D |

② -

13-8-14.

The

Oracle having 2 types of Mat'd Views:-

Oracle

(1) Complete refresh mat'd Views.

(2) Fast Refresh mat'd Views.

high

(1) Complete Refresh materialized Views:-

- In oracle by default mat'd views are complete refresh mat'd views

- When we are refreshing then mat'd views then refresh

internally rowid's are deleted if u are not modified
base table date also.

Syst

Syntax:-

Create materialized View Viewname

Refresh complete

as

Select Statement;

13-8-14.

① > Select rowid, sno, name from mkt;
> exec

dbms_mview.refresh ('mkt');
> Select rowid, sno, name from mkt;

(Other rowids are changed)

② Fast refresh material View :-

These type of mat'd Views are also called as

Incremental refresh mat'd views.

These mat'd views performance is very high compare to default mat'd views BCZ. For like

mat'd views rowid's does not change when we are

s. When refreshing mat'd views no. of times also

modified

Syntax:-

```
create materialized view Viewname  
Refresh fast  
as  
Select Statement;
```

— Before we are creating 'Fast refresh' mat'd view On
we must create mat'd view log on base tables. →

Syntax:-

Create materialized view log on basetablename;

(1) M

- Example :-
 - Select * from base;
- Create materialized view log on base;
- Create materialized view mk2
refresh fast
as
Select * from base;
- Select rowid, sno, name from mk2;
- update base set name = 'xy' where sno = 1;
- Select * from base;

| SNO | Name |
|-----|------|
| 1 | xy |
| 2 | B |
| 3 | C |
| 4 | D |
- exec dbms_mview.refresh('mk2');
- Select rowid, sno, name from mk2;
(here rowids are not changed),

Select * from base;

| 1 | A |
|---|---|
| 2 | B |
| 3 | C |
| 4 | D |

(2) Au

using

on con

Sy

mat'd view
bar tables

On demand, on Commit or (with out PL/SQL)

→ generally we are refreshing mat'd view in 2 days

(1) Manually.

(2) Automatically.

(1) Manually: In this method we are refreshing mat'd view using dbms_mview package this method is also called as ondemand method.
By default method is ondemand.

(2) Automatically: We can also refresh mat'd views without using dbms_mview package this method is called oncommit method.

Syntax:-

Create materialized view View name

Refresh Complete / refresh fast ondemand / oncommit
as by default without writing is also by default.

Select Statement;

→ Create materialized view mks

refresh fast onCommit

as

Select * from bar;

Y Select * from bar;

Ans

| | |
|---|----|
| 1 | Ay |
| 2 | B |
| 3 | C |
| 4 | D |

➤ update base set name = 'zz' where sno = 2;

1 row updated.

➤ Select * from base;

| SNO | Name |
|-----|------|
| 1 | Ay |
| 2 | zz |
| 3 | C |
| 4 | D |

➤ Select * from mk5;

| SNO | NAME |
|-----|------|
| 1 | Ay |
| 2 | B |
| 3 | C |
| 4 | D |

→ no changes

bcz we explicitly gives

a Commit then only View will
be changed.

➤ Commit;

➤ Select * from mk5;

| SNO | NAME |
|-----|------|
| 1 | Ay |
| 2 | zz |
| 3 | C |
| 4 | D |

Cre

Ent

Synt

U

Synt

D

Synt

X

Y

Y

Y

Y

err

DCL (Data Control Language) :-

(1) grant (giving permission).

(2) revoke (cancel permission).

Creating a User :-

> Conn Sys as sysdba ;  > Conn system / manager
Enter psw: sys;  

Syntax:-

(1) Create user username identified by password;

Syntax:-

(2) Grant connect, Resource to username;
or
DBA

Syntax:-

> SQL> Conn username / password;

> SQL> Conn Sys as sysdba;

Enter psw: sys.

> SQL> Create user india identified by india;
 username password

> Grant connect, resource to india;

> Conn india / india
 username password

> Select * from emp;
 In india view there is no data.

Error: table or view does not exists.

→ from scott@scott we transfer emp and all data to india@india

➢ Conn scott/tiger;

➢ grant all on emp to india;

➢ Conn india@india;

➢ Select * from emp;

Error: table or view does not exists.

➢ Select * from scott.emp; (username.object);

➢ Create synonym xy for scott.emp;

➢ Select * from xy;

⇒

Data Security point of view all d.b. Systems having

2 types of privileges (Permission)

(1) System privileges.

(2) Object privileges.

14-8-14.

(1) System privileges :- System privileges are given by database administrator these privileges are used to create or alter database objects.

India user

Oracle having more than 80 System privileges,

These are :- Create Table, Create Session, Create procedure, Create any materialized View, Create any Index, Create any Index, Create public Synonym, --- -

Syntax:-

grant System privileges to username1, username2, ...;

object).

oms having

> Conn sys as sysdba;

Enter password : sys.

> grant Create procedure Create any materialized View, Create any Index to Scott, murali, India;

Note:-

In Oracle If we want to Connect any user to the Server then we are using "Create Session privileges"

> ~~Conn~~ Create Session is used to Connect to Server

Role ^{db.} - allows an only admin user

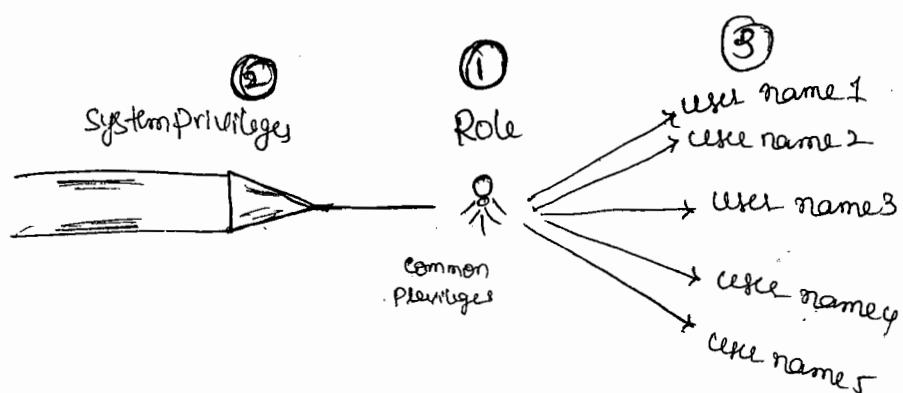
Role is nothing but collection of either System or Collection of Object Privileges.

User +

- user defined roles are created by DBA's.
- In multiuser environment no. of users works on same project.

In this case some users requires Common (Same) Set of privileges.

In this case only DBA's creating a user defined Role & assigns common Set of privileges into Role & then only that role gives to the no. of users..



Creating user defined roles:-

Step 1:- Creating a role :-

Syntax:- Create role any name;

Step 2:- Assign a System privileges to rolename :-

Syntax:- grant System privileges to rolename;

Step 3:- Assign role name to no. of users:-

Syntax:- grant rolename to

(username1, username2, ...)

> Conn sys as sysdba;

Entee password — sys.

> Create role ab;
 ^{filename}

> Grant create procedure Create any materialized view,
Creating
of
users to
Create trigger to ab;

> Grant ab to

Scott, minali, Pndia;

⇒ In oracle if we want to view system privileges
related to role then we are using role-sys-privs
data dictionary.

Example:-

> desc role-sys-privs;

> Select ^{ROLE} role, privilege from role-sys-privs
where ROLE = 'AB';

Predefined roles

When ever we are installing oracle Server then
automatically 3 predefined roles are created.

These are (1) Connect designed for end user purpose

(2) Resource → developer

(3) DBA → database administrator

purpose.

- > desc role-sys-privs;
- > Select ROLE, PRIVILEGE from role-sys-privs
where role ROLE
in ('CONNECT', 'RESOURCE');
- > Select ROLE, PRIVILEGE from role-sys-privs.
where ROLE
in ('DBA');

DBA
 ↴
 create material
 ↴
 generate
 ↴
 insert
 ↴
 d.d.
 create table

Object privileges:-

Object privileges are used to perform some operations of the object.

These privileges are given by either db developer or DBA.

Oracle having following object privileges
these are {Insert, Delete, Update, Select}

These 4 are also called as all in oracle side
execute, read, write, (PLSQL side)

Syntax:-

```
grant object privileges on Objectname
to usernames / role name / public;
```

> Conn Scott/tiger;

Connected.

- > grant all on emp to murali, India;
- > grant all on emp to ab;
- > grant all on emp to public;

With grant option:- (XX not in roles)

who receives with grant option clause those users giving same object privileges to another user.

Syntax:-

```
Grant object privileges on object name  
to username/public with grant option;
```

> grant all on emp to india, murali with grant option;

Privileges
Grant Succeeded.

> grant all on emp to ab with grant option;

Error:- Cannot GRANT to a role with grant option

Note:-

In all d.b. system roles does not work with grant option.

18-8-14.

In oracle, if we want to view object privileges related to user then we are using user-tab-prive d.d.
SQL> desc user-tab-prives;

Revoke:-

It is used to cancel either system or object privileges from the user

Syntax:-

Revoke System privileges From username, username....;

Syntax:-

Revoke Objectprivileges on objectname
From usernames /rolename /public;

In all d.b. systems If u want to restrict table columns from 1 user to another user then only we are using views.

i.e, d.b. administrators creating a view with required columns then the view given to the no. of users.

Syntax:-

grant all on Viewname to
username, username....;

> Create or replace view V₁

as

Select empno, ename, sal from emp

where deptno=10;

> Select * from V₁;

> Grant all on V₁ to india;

> Conn india/india;

> Select * from scott.V₁;

→ Can also drop view using drop View View name;

→ When ever we are dropping a base table we can't access a view.

But view definitions are still exists in d.b.

Ex- > Create or replace view V₁;

as

Select * from base;

> Select * from V₁;

> drop table base;

> desc user-views;

> Select Text from user-views where

VIEW-NAME = 'V₁';

Text

- Select "SNO", "Name" from bat.

MERGE Stmt (AML) 9i

oracle 9i introduced "Merge Stmt".

- This Stmt is used to transfer data from source table into target table.
- But here these 2 tables structures must be same.
- merge Stmt used in data warehousing applications.
- In merge Stmt we are using update, insert Stmt that's why this Stmt is also called as :

"UPSERT" Stmt.

- In merge Stmt we must create different name for the table.

Syntax:-

```
merge into target table name  
using Source tablename  
on (Join Condition)  
when matched then  
update set target table col1 = Source table col1, ...  
when not matched then  
Insert (target table column names) Values  
(Source table column names))
```

- Select * from dept; (target table);
 - Select * from dept;

| |
|----|
| 10 |
| 20 |
| 30 |
| 40 |
 - Create table depts as Select * from dept;
 - Insert into depts Values (1, 'a', 'b');
 - Select * from depts; (source table);
- 10
20
30
40
1 a b

Merge into dept d

using depts s

on (d.deptno = s.deptno)

when matched then

update set

d.dname = s.dname , d.loc = s.loc

when not matched then

insert (d.deptno, d.dname, d.loc) values

(s.deptno, s.dname, s.loc);

- Select * from dept;

| Deptno | Dname | Loc |
|--------|-------|-----|
| 10 | | |
| 20 | | |
| 30 | | |
| 40 | | |
| 1 | a | b |

If
then

Note:- In merge Stmt we can not update

"On clause Columns":

d.deptno = s.deptno — error.

SEQUENCE (Independent Obj)

Syntax

→ Sequence is a d.b. object which is used to generate sequence no's automatically.

Syntax:

→ generally sequences are used to generate primary key values automatically.

These

→ Sequence is an independent d.b. object once a sequence has been created no. of users simultaneously access that sequence.

Select

using

Syntax

Create Sequence Sequence name

Start with n

increment by n

min value n

max value n

cycle / nocycle

cache / no cache;

E

Syntax

SQL



update

we want to generate Sequence Values
then we are using 2 pseudo columns.

(1) Currval

(2) nextval.

Syntax:-

SequenceName . currval;

Syntax:-

SequenceName . nextval;

These pseudo cols are used in insert, update, delete,
select statements (Ans)

Final-

→ In Oracle if we want to generate Sequence values
using Select stat then we must use dual-table.

Syntax:-

Select SequenceName . currval from dual;

Syntax:-

Select SequenceName . nextval from dual;



Create Sequence S1

Start with 5

increment by 2

maxValue 100;

➤ Select s_1 .currval from dual;

Error:

Sequence s_1 .currval is not yet defined for this session.

➤ Select s_1 .nextval from dual;

5

➤ Select s_1 .nextval from dual;

7

→ In all d.b.sys' if u want to generate 1st sequence no. then we must use next val pseudo column.

→ Bcz currval pseudo column returns (Current value of the Sequence if Sequence Session already having a value) (or)

➤ Create Sequence s_1

start with 5

increment by 2

max value 100;

➤ Select s_1 .nextval, s_1 .currval from dual;

| Nextval | currval |
|---------|---------|
| 5 | 5 |

Note:- If u want to view default values for a sequence then we are using user_sequences d.d.

Ex

➤

➤

➤

Note

only

➤

➤

➤

= S

➤

Ex:- Create Sequence S₁

Start with 1;

➤ Select S₁. nextval from deal;

1.

➤ desc user_sequences;

➤ Delete 1 from user_sequences;

Note:- we can alter all sequence parameter values except "start with"

1st sequence

columns:

1 value

by leaving

Syntax:-

alter sequence S₁

Start with 7

Increment by 1

min value 5

max value 20;

➤ Select S₁. nextval from deal;

7.

➤ Select S₁. nextval from deal;

8.

➤ alter sequence S₁

Incremented by -1;

= Sequence altered.

➤ Select S₁. nextval from deal;

7

→ alter sequence S_1

Start with 8;

error: Cannot alter starting sequence number.

20-8-14

Note: Start with cannot be less than min value.

→ Create sequence S_1

Start with 4

incremented by 1

minValue 5;

error: START WITH cannot be less than minValue

Cycle / no cycle:-

No cycle:

→ Create sequence S_1

Start with 5

incremented by 1

minValue 4

maxValue 8

cycle;

no cache;

→ Create sequence S_1

Start with 5

incremented by 1

minValue 4

maxValue 8,

no cycle ~~no cache~~,

→ Select S_1 nextval from deal;

O.P:- 5.

O.P:- 5

6

7

8

| | | | |
|---|---|---|---|
| 1 | 6 | 4 | 4 |
| | | 5 | 5 |
| | | 6 | 6 |
| | 8 | 7 | 7 |

error

→ Error Sequence S₁.NEXTVAL exceeds MAXVALUE
and cannot be instantiated.

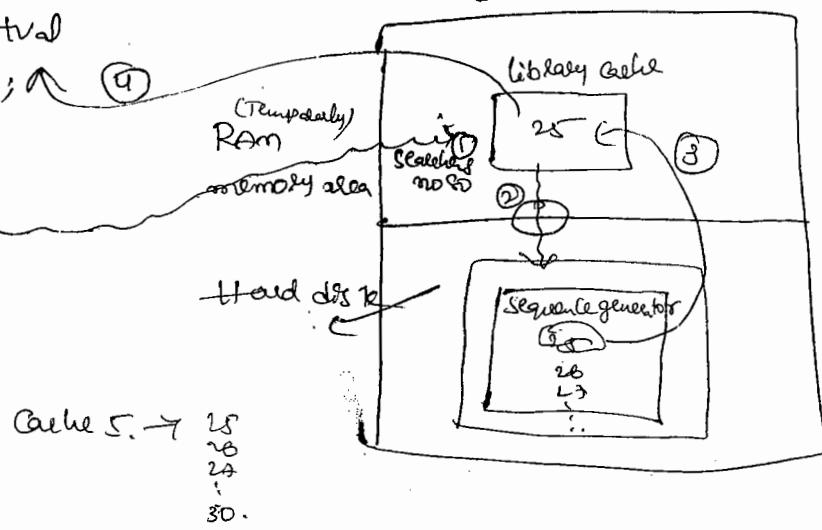
> Select S₁.nextval

from dual;

25.

Value

oracle server .



Caches Cache is a memory area which is used to access Sequence values very fastly.

→ generally if we want to improves performance of the application then only DBA define Cache values for the Sequence through cache option.

→ Cache is a memory area which is used to store set of Sequence no's

→ generally when ever we are creating Sequence the d.b. servers generates Sequence values using Sequence generator within Hard Disk.

→ when ever we are requesting Sequence value then Server process checks requested Sequence no's are available or not available in Cache memory Area

→ If 1st time Sequence values are not available in cache memory area then only Server process fetches cache values from disk & then only those values are transferred into cache memory area.

→ This process automatically (reduces) degrades performance of the apps.

→ To overcome this problem we are defining set of sequence no's using cache option.

In oracle by default Cache values 20 and also cache min value is '2'

Note :- When ever Server is shut down or system crashes then automatically Cache values are "lost"

> Create sequence sr

- Start with 1
- increment by 1

Cache 1;

Error: the no. of values to cache must be greater than 1

→ Generally if u want to generate primary key values automatically then only we are using sequences

available

Ex:

Create table test (sno number (10) primary key,
name varchar(10));

- process

only those

area,

for performance

Set of

→ Enter value for name : abc
: xyz.
: klm.

→ Select * from test;

| sno | name |
|-----|------|
| 1 | abc. |
| 2 | xyz. |
| 3 | klm. |

System

"losted"

→ All sequences information stored under
"User-Sequences" d.d.

desc User-Sequences;

→ we can also drop sequence by using drop

drop sequence Sequence name;

values

entry

LOCKS

- Locking is a mechanism which prevents unauthorized access for our resources (tables, d.d....)
- All db. Systems having 2 types of locks.
 - (1) Row Level Locks.
 - (2) Table Level Locks.

(1) Row Level Locks:-

In this method we are locking set of rows by using "for update" clause.

→ This clause is used in Select Statement only.

→ In Oracle Row Level Locks are introduced in Oracle 6.0.

Syntax:-

Select * from tablename

Where Condition For update [nowait];
↳ optional

→ rows are locked but it does not display any msgs but internally rows are locked. But DML operations working bcz authorised person allowed but unauthorized users can not perform DML ops. Select only working.

⇒ when ever we are performing locks another user query (select stmt) the data but they can not perform DML operations. & also

In all db. Systems when ever we are using Commit then only automatically locks are released.

Scott/Tiger

- Unauthorised)
- > Select * from emp where deptno = 10 for update;
 - > Commit; (or) > rollback;
[For releasing locks].

minali/minali

- > update Scott.emp set sal = sal + 100 where deptno = 10;
- [We cannot perform DMLs]

21-8-14.

by using

NoWait:-

- Oracle 6.0.
- "nowait" is an optional clause used along with "for update" clause.
- When ever we are using this clause Oracle Server automatically get the control into current session if another user not releasing locks also.
- In this case Oracle Server returns an error ORA-54

resource busy.

Scott/Tiger

- the case
- > Select * from emp where deptno = 10 for update no wait;

ORA = 00054; resource busy.

minali/minali

- > Select * from Scott.emp where deptno = 10 for update;

Table level Locks:-

In this method we are locking a table, table level locks are handled by DBA's only.

In oracle db. developers uses 2 types of table level locks.

(1) Share lock.

(2) Exclusive lock.

(1) Share lock :- (more than 1 user)

when we are using this lock another user query the data but we can not perform DML operations

& also no. of users locks the resource(table)

Syntax:- Still waiting all the users release the locks.

```
lock table tablename in share mode;
```

> Show user;

user is scott.

> lock table Scott.emp in share mode;

scott | tiger

> lock table emp in share mode;

murali | murali

> select * from Scott.emp; (✓)

> lock table Scott.emp ^{another lock} in share mode; (✓)

> update Scott.emp set sal=sal+100;

— (we cannot perform DMLs)

(2)

wh

que

ope

Syn

[

lock

:

Not

then

Excl

> l

er

wt

② Exclusive locks (only 1 user like row level lock)

when we are using this lock another user query the data but they can not perform DML operations.

& also at a time only 1 user lock the resource

Syntax:-

lock table tablename in exclusive mode;

ie query operations
e)

lock table emp

in exclusive mode;

Scott/Tiger
~~~~~

meali/meali  
~~~~~

> select * from Scott.emp(✓)

> lock table Scott.emp

In Share mode : (X)

Exclusive mode : (X) any locks are working
Stuckies to commit

Note: In Oracle when ever we are using DML statements then automatically Oracle Server Internal uses

Exclusive locks [some data performs then only exclusive locks applied]

Scott/Tiger
~~~~~

(In same user)

Scott/Tiger  
~~~~~

> update emp set

ename = 'abc'

where ename = 'SMITH';

> update emp set ename='abc'

where ename = 'SMITH';

[Can not perform DMLS bcz internally
Server uses exclusive locks]

NORMALIZATION.

- "Normalization" is a "Scientific process". (designers only use)
 - which is used to decomposing a table into no. of tables.
 - This process automatically reduce redundancy (dupl.)
and also automatically avoids insertion, updation, deletion problems.
- In Design phase of the SDLC db. designer uses (1) E "normalization" process i.e.,
 - In Design phase of the SDLC D.b. designer designs logical model of the database.
 - In this logical model only designer uses normalization process through "Normal Forms".
- In 1970 "E.F.Codd" written a paper "relational model of data for large shared data banks".
 - In this paper only E.F.Codd introduced 1st 2nd Normal Forms.

Candidate
No

Normal Forms :-

- (1) first normal form
- (2) Second normal form
- (3) third normal form.
- (4) BCNF
- (5) Fourth normal form.
- (6) Fifth normal form.

(1) First Normal Form :- (Domain Key NF)

Item table. → multiple values

| Item name | Color | price | tax. |
|-----------|-------------|-------|------|
| marker | black, red | 20 | 0.02 |
| Pen | blue, green | 30 | 0.03 |

Not in 1st NF

↓ 1NF

Item table.

↑ duplicates to

Candidate key

| Item name | color | price | tax |
|-----------|-------|-------|------|
| marker | black | 20 | 0.02 |
| marker | red | 20 | 0.02 |
| Pen | blue | 30 | 0.03 |
| Pen | green | 30 | 0.03 |

→ A table is in 1NF in that table every cell having an atomic value and also.

② Identifying a record uniquely using a key

(newly married couple) (couple for electrical shop)
 (B. straw) (cool down)
 Order Form 1 F 2
 orderno

orderdate

Customer

ordno

Or.d

custno

add

Ph.no

I1

A1

I2

A2

I3

A3

Item name

Amount

uniquely using a key

①

child-table

| orderno | D: | Cust no | add | Phy | I1 | A1 |
|---------|----|---------|-----|-----|----|----|
| | | | | | | |

②

| orderno | Or.d | Cust no | add | Ph.no | I1 | I2 | A1 | I3 | A2 | I4 | A3 |
|---------|------|---------|-----|-------|----|----|----|----|----|----|----|
| | | | | | | | | | | | |

not possible in
d.b.

Submit

Order Form 3

Order no

Order date

Customer

Address

Ph.no

Item name

Refrigerator
LED TV

P.K.

Order master Table

| orderno | orderdate | customer | add | ph.no |
|---------|-----------|----------|-----|-------|
| 1 | | | | |

Amount

Submit

Order detailed Table

F.K.

1NF Table

| orderno | Item name | Amount |
|---------|-----------|--------|
| 1 | | |
| 1 | | |
| 1 | | |

(child Table)

bcz duplicates allowed

Every Cell

Process:-

Identifying repeating groups & putting into separate table in more atomic form.

2 Key

| data-table | | |
|------------|-----|-----|
| add | phy | I A |
| | | |
| | | |
| | | |

→ By default 1NF process table is a child table. Bcz in this table 1 col having duplicate data
must

22-8-16-

2nd Normal Form :-

→ A Table is in 1NF & also all non-key attributes are "Fully Functionally" dependents on total Candidate Key.

→ Always 1st normal form deals with atomicity whereas 2nd NF deals with relationship b/w key, non key attributes.

→ If a table is in 1NF and also the table contains partial non-key attributes then that table not in 2NF.

Process :-

→ Identify partial Non-key attributes which depends on Partial key attributes. Those all attributes put into separate Table.

366

Amount

see this

→ The Separate Table is called 2NF Table.

→ By default this is a master table. In this table only - all Nonkey attributes are Fully Functionally dependent on total Candidate Key.

| Candidate key | | Item name | Color | Price | Tax | Non key attributes |
|----------------|--|-----------|-------|-------|-----|--------------------|
| key attributes | | marker | black | 20 | 0.2 | P |
| | | marker | Red | 20 | 0.2 | P |
| | | Pen | blue | 30 | 0.3 | P |
| | | Pen | green | 30 | 0.3 | P |

Partial bc+ Price only
dep's Itemname only
on not color to
Total.

not P in 2NF

| Item name | | Color |
|-----------|--------|-------|
| F.K. | marker | black |
| | marker | Red |
| | Pen | blue |
| | Pen | green |

| Item name | | Price | Tax |
|-----------|--------|-------|-----|
| | Marker | 20 | 0.2 |
| | Pen | 30 | 0.3 |

Partial
to
partial

(master-table) 2NF.

| Stno | Stname |
|------|--------|
| 101 | Shreya |
| 102 | madhu |
| 103 | poorva |
| 104 | Jyo |

| Stno | Activity |
|------|----------|
| 101 | Cricket |
| 101 | Football |
| 102 | Cricket |
| 102 | Hockey |
| 103 | Football |
| 103 | Swimming |
| 104 | Swimming |

| Activity | Cost |
|----------|-------|
| Cricket | \$ 10 |
| Football | \$ 20 |
| Hockey | \$ 30 |
| Swimming | \$ 40 |

Non key
attr.

Partial key attr.

2N

2NF (masterTable)

Student Sports Activity Table.

ble
this table

smallly

butler

Price only

Name only

Color to

Partial

R

Partial
to
partial

1NF

N. key
attr.

Primary key:

ble

| Sno | Sname | Acti 1 | Cost 1 | Act 2 | Cost |
|-----|----------|----------|--------|----------|------|
| 101 | Shreethi | Cricket | \$10 | Football | \$20 |
| 102 | madhu | Cricket | \$10 | Hockey | \$30 |
| 103 | Pooja | Football | \$20 | Swimming | \$40 |
| 104 | Jyothsna | Swimming | \$40 | | |



table

| Sno | Sname |
|-----|----------|
| 101 | Shreethi |
| 102 | madhu |
| 103 | Pooja |
| 104 | Jyo |

| Sno | Acti 1 | Cost 1 | Act 2 | Cost 2 |
|-----|----------|--------|----------|--------|
| 101 | Cricket | \$10 | Football | \$20 |
| 102 | Cricket | \$10 | Hockey | \$30 |
| 103 | Football | \$20 | Swimming | \$40 |
| 104 | Swimming | \$40 | | |

| Sno | Sname |
|-----|----------|
| 101 | Shreethi |
| 102 | madhu |
| 103 | Pooja |
| 104 | Jyo |

Key by attr

| Sno | Activity | Cost |
|-----|----------|------|
| 101 | Cricket | \$10 |
| 101 | Football | \$20 |
| 102 | Cricket | \$10 |
| 102 | Hockey | \$30 |
| 103 | Football | \$20 |
| 103 | Swimming | \$40 |
| 104 | Swimming | \$40 |

↓ 1NF

candidate key

P

Nonkey attribute

not in

2NF

depends on
activity

← 2NF

2NF

Convert to 2NF

(2) F - F 10

| Ecode | ProjCode | Dept only code | Depthead only code. | Hours eCode + ProjCode | F |
|-------|----------|----------------|---------------------|------------------------|----------------|
| E101 | P1 | Systems | D1 | 3 | not P2 2NF. |
| E102 | P1 | SALES | D2 | 5 | |
| E101 | P2 | Systems | D1 | 8 | |
| E103 | P2 | Research | D3 | 9 | |
| E102 | P3 | SALES | D2 | 10 | |
| E101 | P3 | Systems | D1 | 7 | |

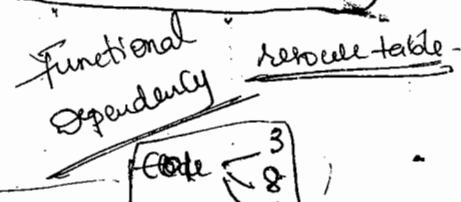
insertion error

XP

Phase 1:-

FD

① Ecode \rightarrow hours (X)



② ProjCode \rightarrow Hours (X)

hours

③ (Ecode + projCode) \rightarrow Hours (V)

Insert

for

Insert

meest

This

if

dept

part

Phase 2:-

① Ecode \rightarrow Dept (V)

E101 \rightarrow systems

② ProjCode \rightarrow Dept (X)

③ F - F 10

Phase 3:-

① Ecode \rightarrow Depthead (✓)
FD

② Proj code \rightarrow Dept head (✗)

not in
2NF.

P.K.
2NF

| Ecode | Dept | Depthead |
|-------|----------|----------|
| E101 | Systems | D1 |
| E102 | Sales | D2 |
| E103 | Research | D3 |

2NF master table

F.K.
2NF

| Ecode | Proj code | Hours |
|-------|-----------|-------|
| E101 | P1 | 3 |
| E101 | P2 | 8 |
| E101 | P3 | 7 |
| E102 | P1 | 5 |
| E102 | P3 | 10 |
| E103 | P2 | 9 |

* (child table)

23-8-11.

→ Before Normalization process above result table having insertion, update, deletion problems

Insertion problem :-

For the above result table when we are try to insert particular department employee details then we meet assign project code. (1st day ^{unday} we does not assign project) (so)

This is called insertion problem.

② update problem → In the above result table Ecode, dept, depthead attribute values are repeated whenever particular deptmt employee transferred 1 dept in another deptnt

then we must modify all their 3 attribute values correctly. (update sales dept to eng-~~dept~~).

→ otherwise inconsistency probm occurred. This is called "update probm".

(3) Deletion probm: In the above resource table when we are try to delete a particular employee then automatically deptmt details also deleted. This is called "deletion problem".

→ when ever we are using normalized process then automatically insertion, update, deletion problems are avoided (void)

3NF.

→ If a table is in 2NF & also all Non-key attributes are only dependents on Candidate (P.k) key

⇒ If a table is in 2NF & also Non-key attributes which ~~also~~ depends on another Non-key attributes then that table "not in 3NF".

attribute

This is

table
employee then
is called

its items

elements

non-key

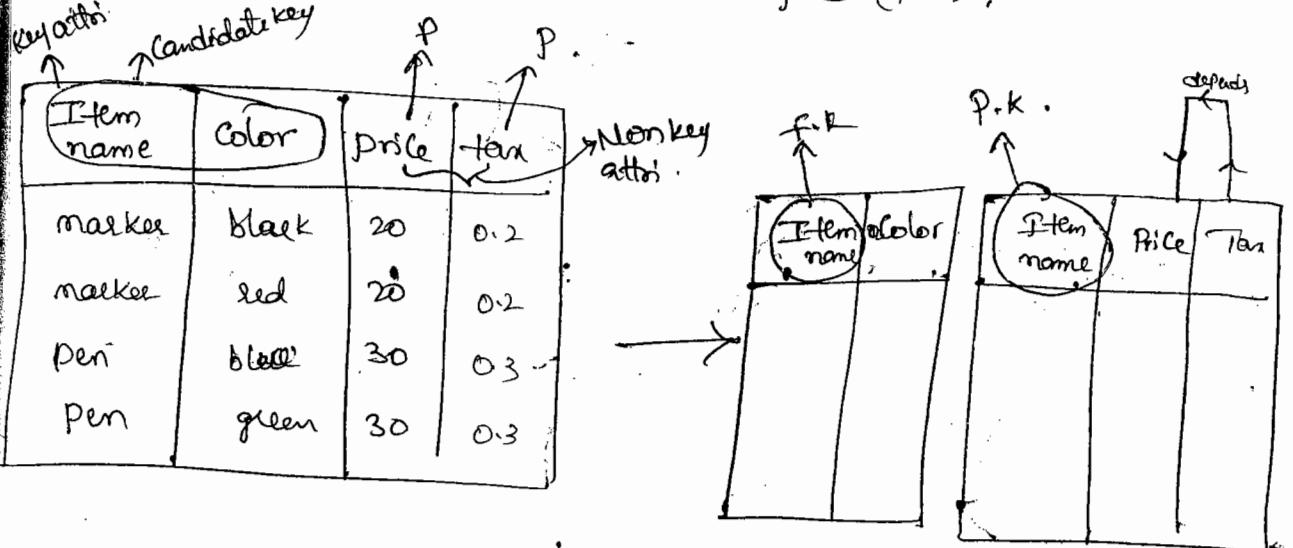
(P.K) key

attribute

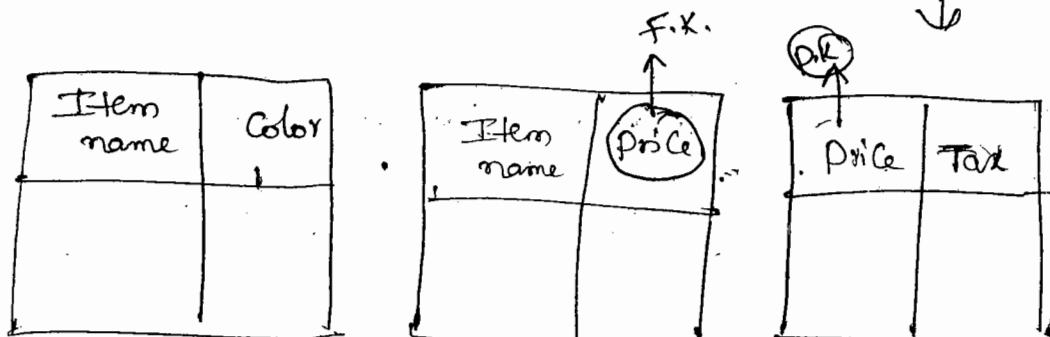
that

Process :-

- Identify Non key attributes, which depends on another Non-key attributes put into separate table.
- This table is called "3NF table". By default this table is master table.
- In this table only all non-key attributes are only dependents on Candidate key \oplus (P.K).



in 2NF master table
but not in 3NF



master table (3NF table).

Order Form.

Order no

Order date

Cust no

Item name Item no

Amount add another item

Qty Discount

Customer Enquiry Form

Cust.no
Cust.name
Address
Ph.no

Key α¹

S

P.K customer master table

(3NF)

| Custno | Cust' name | Address | Ph.no |
|--------|------------|---------|-------|
| | | | |
| | | | |
| | | | |
| | | | |

P.K order master table

(1C)

| Orderno | Orderdate | Custno |
|---------|-----------|--------|
| | | |
| | | |
| | | |
| | | |

Item master

| Items no | Item name | Amount |
|----------|-----------|--------|
| | | |
| | | |
| | | |
| | | |

order detailed master table

| Order no | Item no. | Qty | Discount |
|----------|----------|-----|----------|
| | | | |
| | | | |
| | | | |
| | | | |

Form

key attributes → Candidate key

25-8-14.

2NF

(not in 3NF) ↗
N-key FD on N. key

Non-key attributes ↗

| Stname | Courseid | Grade | GradeValue |
|--------|----------|-------|------------|
| abc | CS111 | A | 4.00 |
| xyz | CS111 | B | 3.00 |
| zzz | CS222 | C | 2.00 |
| aaa | CS222 | A | 4.00 |

Student table.

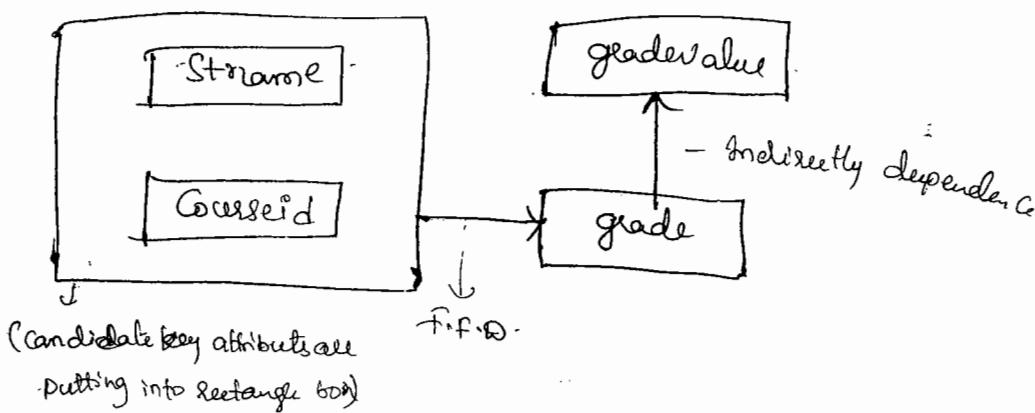
Child table

Master table

3NF

| Stname | Courseid | Grade | GradeValue |
|--------|----------|-------|------------|
| abc | CS111 | A | 4.00 |
| xyz | CS111 | B | 3.00 |
| zzz | CS222 | C | 2.00 |
| aaa | CS222 | A | 4.00 |

logical diagram.



Functional Dependency :- (Uniqueness) :- (FD) :-

If any given 2 tuples in a relation 'x' then,
 x (an attribute set) agrees then y (another attribute set)
cannot disagree ^(agree) then $x \rightarrow y$ is called "F.D."

$x \rightarrow y$.

here y is functionally dependent on x ($x \rightarrow y$)

(or)

* functionally determines ' y ' ($x \rightarrow y$)

| attribute | | |
|----------------------------|-------|------|
| Stno | Sname | Hall |
| t_1 $\rightarrow S_1$ | abc | H1 |
| t_1 $\rightarrow S_2$ | xyz | H1 |
| S_3 | aaa | H2 |

Relation 'x'.

1NF :- Remove repeating groups

2NF :- Remove partial attributes

3NF :- Remove attributes which are not dependent
on candidate key (P.K).

(or)

Remove non key attributes which are
dependent on another non-key attributes.

en
UNF
- Cowl
- Goo
- tutu
- tur
- Stn
- Stn
- @.o
- Gien
- lastre

1st & 2nd & 3rd NF.

Course id []

Course name []

Tutor id []

Tutor name []

| Stno | Sname | D.O.B | Gender | Last attendance |
|------|-------|-------|--------|-----------------|
| [] | [] | [] | [] | [] |

en

UNF

- CourseId
- Course name
- tutor id
- tutor name
- Stno
- Sname
- D.O.B
- Gender
- last.attende

1NF

Course id
Course name

Tutor name

1NF
not in
2NF

- Course id
- Stno
- Sname+P
- D.O.B.+P
- Gender+P
- last.attende+P

Some not in 3NF

- Courseid
- Course name
- Tutor id (M.R)
- Tutor name (M.A)

2NF

Stno
Sname
D.O.B
Gender

Courseid
Stno
last.attende

ident

Course id
Course name
Tutor id

3NF

- Tutor id
- Tutor name

Stno
Sname
D.O.B.
Gender
last.attende

try -

Cluster :-

Cluster is a d.b. object which contains group of tables together & also cluster shares same data.

Block :

→ clusters are used to improve performance of the joins.

→ clusters are created by DBA's

→ clusters are created based on a common column name.

→ This common col also called as "cluster key".

→ Generally in all d.b. systems when ever we are submitting inner joins or outer joins then d.b. servers checks "From" clause tables are available in clusters or not.

→ If those tables are available in cluster then d.b. servers very fastly derive data from cluster tables.

→ This process automatically improves performance of the application.

Cre

Step

Step

Sy

Step 2

cyo

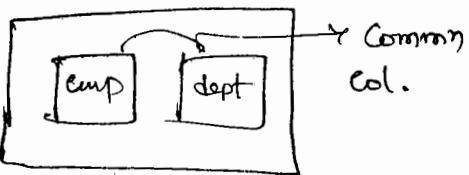
C

Synt

Step 3

C

cluster



group of

2 data

- Select dname, ename, sal, d.deptno, thndate, loc
- 1 From emp e, dept d
 - 2 where e.deptno = d.deptno; → These
common cols

Creating cluster:-

26-8-14.

Steps to creating clusters:-

Step 1) Create a cluster using clusterkey (Common Col)

Syntax:

Create cluster clustername (Common Colname d.ty (size));

Step 2) Create an index on cluster:-

Syntax:

Create index indexname on cluster clustername;

Syntax: Create actual table:-

Step 3:-

Create table tablename (col1 d.ty (size), col2 d.ty (size) ...).
cluster

clustername (Common Colname);

➤ Create cluster

emp-dept (deptno number (10));

➤ Create index emp_idx on cluster emp-dept;

➤ Create table emp1 (empno number (10), ename

varchar2 (10), sal number (10), deptno number (10))

cluster emp-dept (deptno);

➤ desc emp1;

➤ Create table dept1 (deptno number(10), dname

varchar2 (10), loc varchar2 (10)) cluster emp-dept(deptno);

➤ desc dept1;

➤ Insert into emp1 values (1, 'abc', 2000, 10);

➤ Insert into dept1 values (10, 'S/w', 'americpet');

➤ Select * from emp1;

➤ Select * from dept1;

➤ Select rowid from emp1;

➤ Select rowid from dept1;

* * * in cluster
Same rowid's only.

Note:- Generally two rowid's are not same in both

-Bt cluster tables having same rowid's.

No

cl

inclus

table

Synt

Ex

er

89 or

Ex

dr-dict

Ex

su

Sup

whic

Super

Note:- Generally we cannot drop cluster if cluster having tables.

→ To overcome this problem oracle 8i introduced "including tables" clause to drop clusters along with tables.

Syntax:-

```
drop cluster clustername including tables;
```

Ex:-

```
drop cluster emp_dept ;
```

Error:- cluster not empty;

8i onwards.

Ex:-

```
drop cluster emp_dept including tables;
```

→ All clusters information stored under "user-clusters" dictionary.

Ex:- desc user-clusters;

Super key, Candidate key :-

Super key :- A column or combination of columns which uniquely identifying a record \downarrow is called "super key".
in a table

Candidate Key

A minimal Super key which uniquely identifying a record in a table is called "Candidate Key".

Note: A Super key which subset of another S.Key (Combination) is not a Candidate key.

| empno | ename | skillid | SKPL | voterid |
|-------|--------|---------|-----------|---------|
| 1 | murali | 1 | Oracle | V1 |
| 1 | murali | 2 | Sybase | V2 |
| 1 | murali | 3 | Ingles | V3 |
| 2 | abc | 4 | DB2 | V4 |
| 2 | abc | 1 | oracle | V5 |
| 3 | xyz | 5 | Informix | V6 |
| 4 | zzz | . | | V7 |
| 5 | aaa | 6 | SQLserver | V8 |
| 5 | aaa | 4 | DB2 | V9 |

Superkey

- ① Empno + Skill
- ② Empno + ename + skill
- ③ Empno + skill + voterid
- ④ ename + skill
- ⑤ ename + skill + voterid
- ⑥ voterid + skill
- ⑦ empno + ename + voterid + skill

not a S.Key

- ① empno + ename
- ② Empno + voterid

this

Can

Candidate Key

- (1) empno + Skills.
- (2) ename + Skills.
- (3) Voterid + Skills.

BCNF

→ If a table is in BCNF then Every determinant is a "Candidate key".

→ If a table contains multiple Candidate keys & also those Candidate keys are overlapped and also if Candidate key non key attribute which depends on another candidate key nonkey attribute then only we are using "BCNF" process.

Process:

→ Identify 1 Candidate key n.key.attribute which depends on another candikey nonkey attribute put into separate Table this table is called "BCNF Table".

In this table only Every determinant is a candidate key.

BCNF table does not have nonkey attributes

→ Generally if a table contains only 1 Candidate key then BCNF is also same as 3NF.

→ Generally when a table contains multiple Candidate keys then deletion problem occurs.

→ To avoid deletion problem db. designer uses BCNF process.

→ Generally in BCNF we are specifying relation ship b/w candidate keys itself.

| Prof Code | Dept | H.O.D | Hours | Non key attributes. |
|----------------|-----------|----------------|-------|---------------------|
| P ₁ | Physics | H ₁ | 3 | |
| P ₁ | Computer | H ₂ | 5 | |
| P ₂ | Chemistry | H ₃ | 7 | |
| P ₂ | Botany | H ₄ | 9 | |

① Prof code + Dept → Hours

② Prof code + H.O.D → Hours

1 Candidate

1 plb Candidate

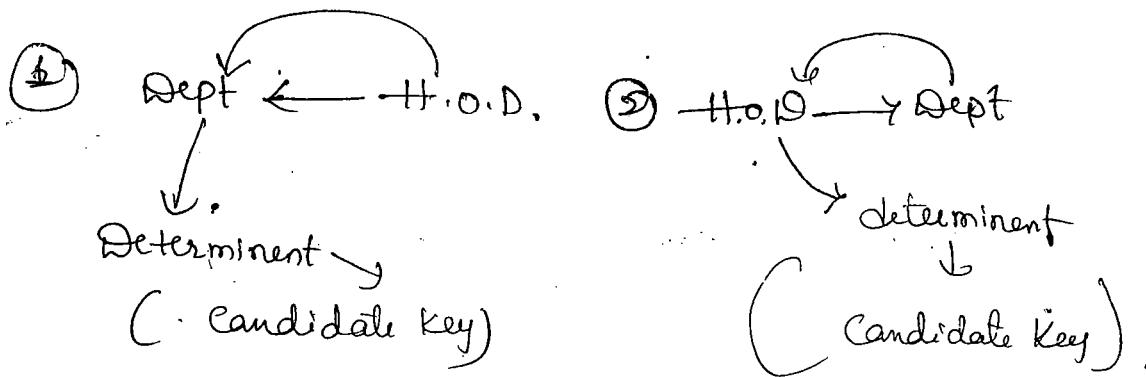
is BCNF

| Dept | H.O.D. |
|------|--------|
| phy | H1 |
| com | H2 |
| che | H3 |
| Bot | H4 |

| ProfCode | Dept | Hours |
|----------|------|-------|
| P1 | phy | 3 |
| P1 | com | 5 |
| P2 | che | 7 |
| P2 | Bot | 9 |

relation

non key
Attributes.



27-8-11.

4NF

4NF:-

If a table is in 4NF, then that table does not contain more than one independent multi-valued attribute.

→ If a table contains more than two attributes & also identifying a record uniquely using combination of all these attributes.

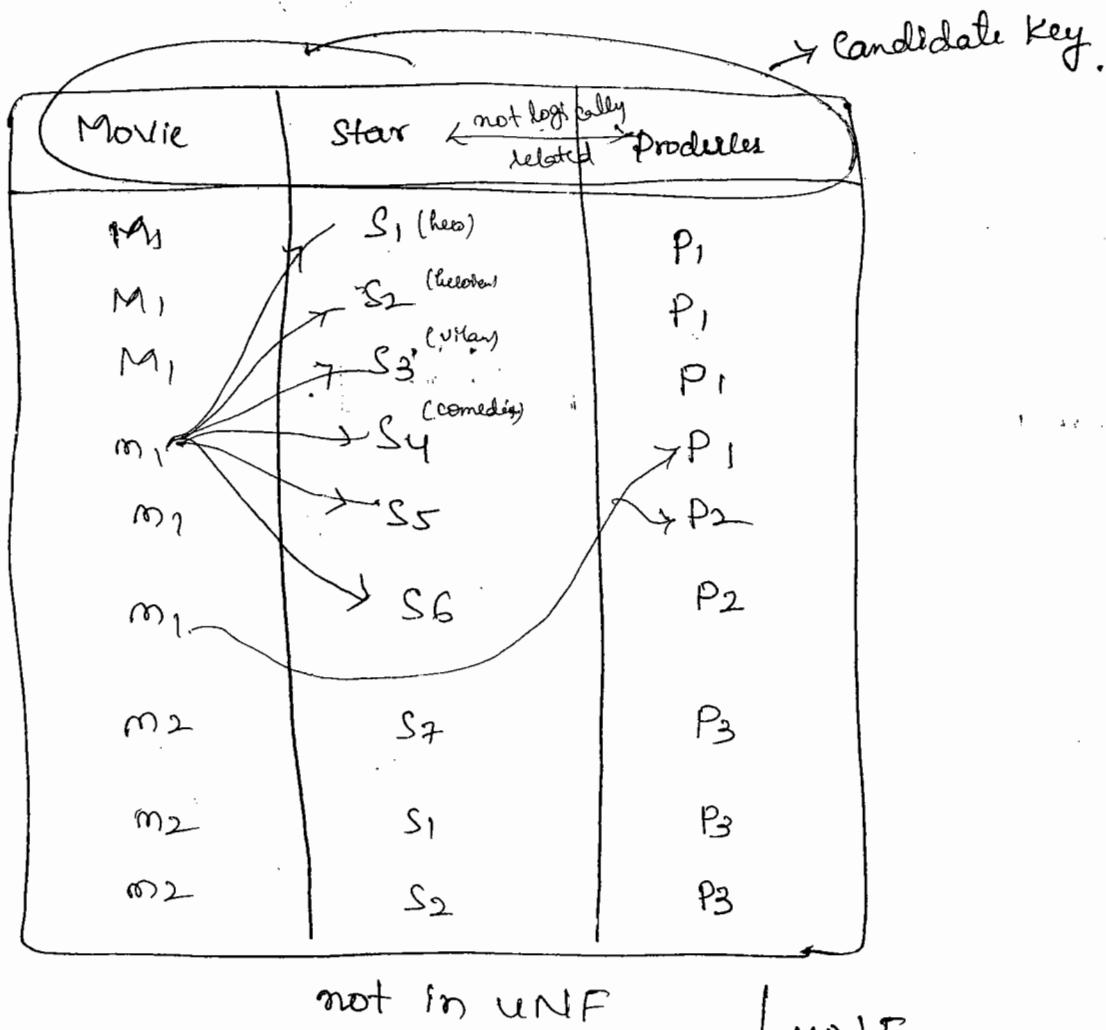
→ & also if attribute set of values which depends on another attribute set of values & also same attribute set of values which depends on another set of values & also some attributes are not logically related then only we are using 4NF.

In 4th NF result table also there is no non-key attributes.

A
~
proj

Process:-

- Identify a column which depends on another independent multi-valued attribute column. Put onto separate table.
- These tables are called 4NF tables



| Movie | Star |
|----------------|--------------------------|
| M ₁ | S ₁ (hero) |
| M ₁ | S ₂ (heroine) |
| M ₁ | S ₃ (comedy) |
| m ₁ | S ₄ (comedy) |
| m ₁ | S ₅ |
| m ₁ | S ₆ |
| M ₂ | S ₇ |
| m ₂ | S ₁ |
| m ₂ | S ₂ |

4NF

| movie | producer |
|----------------|----------------|
| M ₁ | P ₁ |
| m ₁ | P ₂ |
| M ₂ | P ₃ |
| m ₂ | P ₃ |
| m ₂ | P ₃ |

4NF

for
→ g
then
in

on-key

Assumption:- Each employee having multiple projects & each employee having multiple skills.

| Empid | Project id | Skills | Key |
|-------|---------------------|------------|-----|
| 1111 | Railway reservation | JSP | |
| 1111 | Railway reservation | Asp.net | |
| 1111 | Railway reservation | oracle | |
| 1111 | Library management | php | |
| 1111 | Library management | Salesforce | |
| 1111 | STO | Null | |
| 1111 | Null | SAP | |

(not in 4NF.)

if we insert null value problem will occur,

↓ 4NF

| Empid. | projectid |
|--------|-----------|
| 1111 | Rail.Ren |
| 1111 | Librarian |
| 1111 | STO |

problem

| Empid | Skills |
|-------|--------|
| 1111 | JSP |
| 1111 | .net |
| 1111 | oracle |
| 1111 | PHP |
| 1111 | SQL |
| 1111 | SAP |

→ Before 4NF process whenever an employee got a new project then we need to supply null values for the skills in the above resource table.

→ & also when ever an employee got a new skills then we need to supply null values for the project in the above resource table.

→ These problems are automatically avoided when we are using 4NF tables.

Projection Joined NF (5NF):

→ If a table is in 5NF there are no ① Cyclic dependencies

→ When a table contains more than 2 attributes and also identifying a record uniquely using combination of all attributes (or also all fields sets of values related to one another).

→ Then only we are using 5NF process.

→ This NF is also called as "projection-joined NF" bcz when ever we are joining process a-table then resultant record must be available in respective table.

Note:-

→ In 4NF Some field are not logically related whereas in 5NF all fields are logically related

| Buyer | Itemname | Company |
|-------|----------|---------|
| B1 | Shirt | Pepe |
| B1 | Tshirt | levis |
| B1 | Jeans | denim |

Buyer table

not in 5NF

when

| Buyer | Itemname |
|-----------|----------|
| Item name | Company |
| buyer | Company |

①

SET OPERATORS

→ SET operators are used to retrieve data from single or multiple tables.

→ These operators are also called "Vertical Joins".

These are:

(1) union

(2) union all

(3) intersect

(4) minus.

① union: It returns values 1 time only, (unique)

② union all: unique + duplicate (all values).

③ intersect: Common values.

④ minus: It returns values in 1st query those values are not in 2nd query.

→ union:

Select job from emp where deptno=10

Union (small) (in) (large)

Select job from emp where deptno=20;

Note: whenever we are using set operators
always corresponding expressions must belongs to same
data type

& also set operator between 1st query column names
as col headings.

→ we can also retrieve data from set operators if
corresponding expressions not belongs to same data type
also, In this case we must use appropriate type
conversion func.

Ex: Select deptno "deptno", to-char(mvll) "deptname"
from emp , union

Select to-number (mvll), dname from dept.

| deptno | dname |
|--------|-------|
| 10 | |
| 20 | |
| 30 | |

Account
Fleck
Allen

Conversions: Converting 1 dty. into another

dty. is called "Conversion".

28-8-14.

relations

to same Oracle having 2 types of Conversions.

(1) Implicit Conversions.

(2) Explicit con .

1) Implicit Conversion &
more common

These Conversions are also called as "automatic"
Conversions".

In this Case Oracle Server only automatically
Converts 1 d.t.y. into another d.t.y.

Note 1:

In Oracle when ever string representing pure number
then Oracle Server automatically converts string type
to no-type.

Ex:- Select Sal + '100' from emp;

Sal+100 ↓
convert into pure no.

2) Note 2: In Oracle when ever we are putting nols into
character fun's (length, upper, lower, initcap) then Oracle Server
automatically converts no.type into character type.

Ex:- Select length(1111) from ideal;

O.P:- 4

Note 3 :- In oracle when ever we are passing date string into date fun's then oracle server automatically converts date string into date type but here passed parameter must be in default date format.

Ex:- Select last_day ('17-dec-05') from dual.

Out:- 31-Dec-05.

Implicit Conversion Table.

| From | To | Assignment | Evaluation of Expression. |
|--------------------------|--------------------------|------------|---------------------------|
| Varchar2 (or) char | number | yes | yes |
| Varchar2 (or) char | Date | yes | yes |
| number | Varchar2 (or) char | yes | no |
| Date | Varchar2 (or) char | yes | no |

Explicit Conversions:-

In this case developers only explicitly convert 1 d.t.y. into another d.t.y. using Sole explicit conversion fun's.

wing

value

pe but

to format

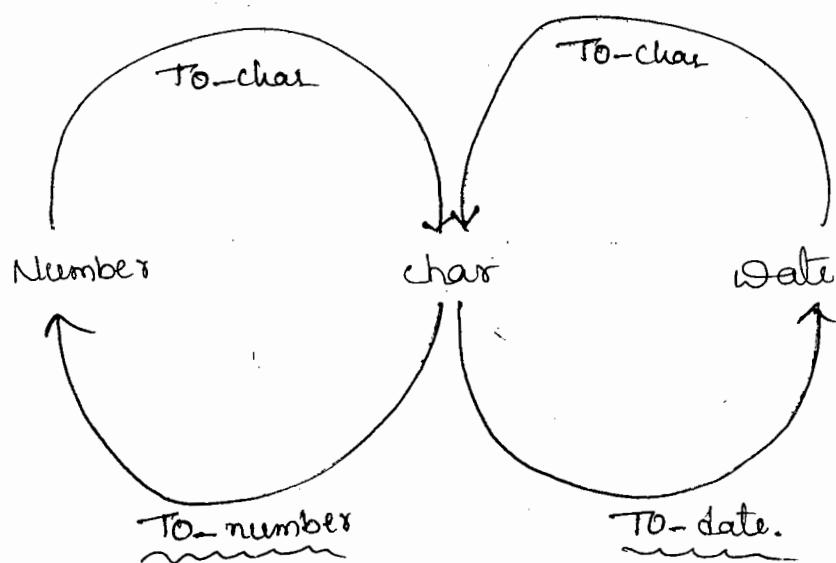
I

→ There are

(1) to-number()

(2) to-char()

(3) to-date().



① to-number-

Converting a String representing ('\$') a numeric value (45.6) with format (\$) into a string representing ('') a numeric value without format (45.6)

Ex:-
only

Select '\$45.6'+3 from dual;

error: Invalid number.

sing
using
Solution :-

Select to-number('\$45.6')+3 from dual

Output: 488, error: invalid no.

➤ Select

to-number ('\$45.6', '\$99.9') + 3 from dual;

O.P:- 48.6.

On oracle:

- ① \$ - \$ (or) L (local currency)
- ② : - . (d). (only special symbols)

③ , - , (or) g.

\$ - L

- ④ 123 - 999
- ⑤ ab - x Invalid Format model.

TO-char :- to-char is an overloading fun. which is used to convert no. type into character type and also convert date type into date string.

format :-

(.) g → group separator.

(.) d → decimal indicator.

→ Select to-char(2000.00, '9999') from dual.

O.P:- ######

→ Select to-char(20000, '99999') from dual

20000.

→ Select to-char(20000, '999.99g99d9') from dual

20000.00

2,00,00.0.

To-

→ Select to-char(20000, '\$99999d99') from dual

20000.00

→ Select

Sal, to_char (sal, 'L99G,999'), ('NLS_CURRENCY=RS')

from emp

Sal

800

1600

1250

TO-CHAR

RS 800

RS 1,600

RS 1,250.

→ Select

Sal, to_char (sal, 'L99G,999') from emp.

SAL

TO-CHAR.

800

\$ 800

}
L-local currency

1600

\$ 1,600.

1250

\$ 1,250

→ Select

ename, nvl (to_char (comm), 'no commission')

O.P:

Comm

2000

from emp

no commission (null replaced by)

no commission

\$ 00
200

→ Select

ename, nvl (to_char (mgr), 'no manager') from emp;

→ Select

to_date ('sysdate', 'DD/MM/YYYY') from deal

28/08/2014.

To_date: It is used to convert date stored into date type

→ Select

to_date ('16/June/05') from deal;

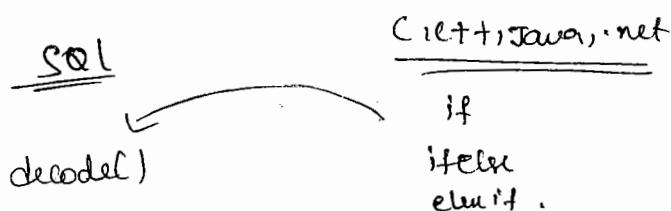
16-JUN-05.

decode(): (Conversion fun) :-

decode() is a conversion fun which is used to convert 1 d.t.y. into another d.t.y.

→ It is used to convert m.type into String type & also converts single character into string.

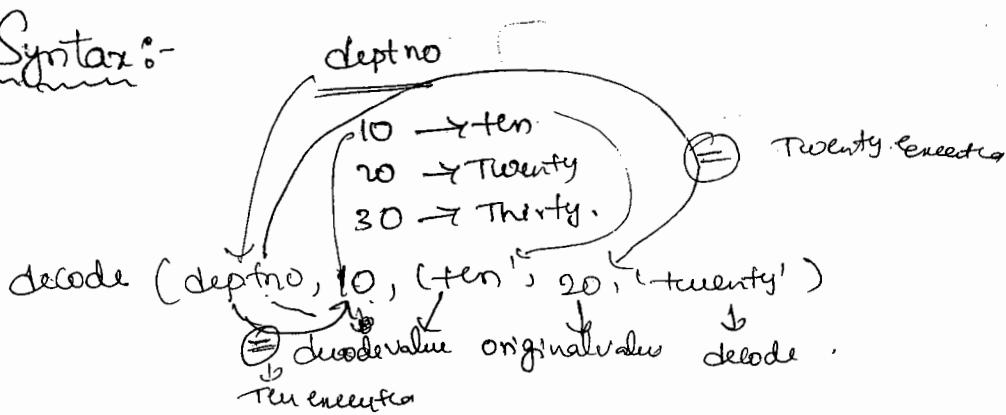
$$'a' \rightarrow 'ab' \quad 123 \rightarrow '(123)'.$$



It is also same as if -- then -- elif construct in PL/SQL.

→ decode() internally uses equality operator (=)

Syntax:-



decode (Columnname, value1, Stmt1, value2, Stmt2, ... Stmtn)

> Select ename, sal, deptno, decode (deptno, 10, 'ten', 20, 'Twenty', 'Others') from emp

29-8-14.

to connect.

Note:- We can also use decode() for modifying data in a table conditionally. In this case, we are using

type &

decode() in update stmt.

(Q) wAQ to modify commission of the employees from emp table based on following conditions.

- (1) if job = 'CLERK' then Comm \rightarrow 10% of sal.
- (2) if job = 'SALESMAN' then Comm \rightarrow 20% of sal.
- (3) if job = 'ANALYST' then Comm \rightarrow 30% of salary.
- (4) else comm \rightarrow 50% of salary.

\rightarrow update emp set Comm = decode(job, "CLERK", sal*0.1, "SALESMAN", sal*0.2, "ANALYST", sal*0.3, sal*0.5);

NOTE:- We can also use decode() within group by clause to display summarize data conditionally.

for eg:- \rightarrow select job, sum(decode(deptno, 10, sal)) "deptno10", sum(decode(deptno, 20, sal)) "deptno20", sum(decode(deptno, 30, sal)) "deptno30" from emp
group by job;

| O.P:- | <u>JOB</u> | <u>Deptno 10</u> | <u>Deptno 20</u> | <u>Deptno 30</u> |
|-------|------------|------------------|------------------|------------------|
| 20, | Analyst | | 6000 | |
| | Clerk | 1300 | 1900 | 980 |
| | Manager | 2450 | 2975 | 2880 |
| | President | 5000 | | |

whenever we want to display multiple values for single columns, then we use decode(): decode() is also used for counting purpose. for eg Counting no. of clerks in sales department.

$\text{Sum}(\text{decode}(\text{Job}, 'SALESMAN', 1)) \rightarrow$ Here, for every one Salesman, it count 1 and sum means add 1 to previous value.

Eg:- select dname, sum(decode(Job, 'CLERK', 1)) "Clerk",
sum(decode(Job, 'SALESMAN', 1)) "Salesman",
sum(decode(Job, 'ANALYST', 1)) "Analyst",
where e.deptno = d.deptno group by dname;

O.P:-

| DNAME | Clerks | Salesman | Analyst |
|------------|--------|----------|---------|
| Accounting | 1 | 0 | 0 |
| Research | 2 | 0 | 2 |
| Sales | 1 | 4 | 0 |

{ Now days, decode is replaced by case i.e, Oracle 8.0 onwards}

CASE Statement :-

→ Oracle 8.0 introduced Case Stmt which is used to decoding the value.

→ Oracle 8i introduced Case Conditional Stmt.

→ Case Stmt performance is very high Compared to decode()

decode() internally uses '=' operator implicitly.

Case Explicitly uses all SQL relational operator from 8i which is advantages

Note:- decode() internally uses equality (=) operator where as in Case Stmt we can also use all SQL operators explicitly.

for
also
of
one
to

"clerk",

Method 1 (B.O) :- Syntax:-

case column name
when value1 then stmt1
when value2 then stmt2

else stmt end

Ex:- > Select ename, sal, deptno
case deptno
when 10 then ('ten')
when 20 then ('twenty')
else 'Others' end
from emp;

Method(2) :- Case Conditional Stmt (Oracle 8i) :-

Syntax:- Case

when condition1 then stmt1
when cond2 then stmt2

else stmts end

Ex:- Select ename, sal, case when sal
1000 then 'low salary' when sal
bw 1000 and 2000 then 'medium sal'
when sal in (2300, 2500, 2800) then
'special sal' else 'other salary' end
from emp;

.. Oracle Server - architecture 80-8-14.

Oracle Server mainly consists of 2 parts.

(1) Storage Area.

(2) Instance.

(1) Storage Area:- when ever we are installing

to Oracle Server then automatically 3 files are created
in disk. These are

(1) Data files (.dbf).

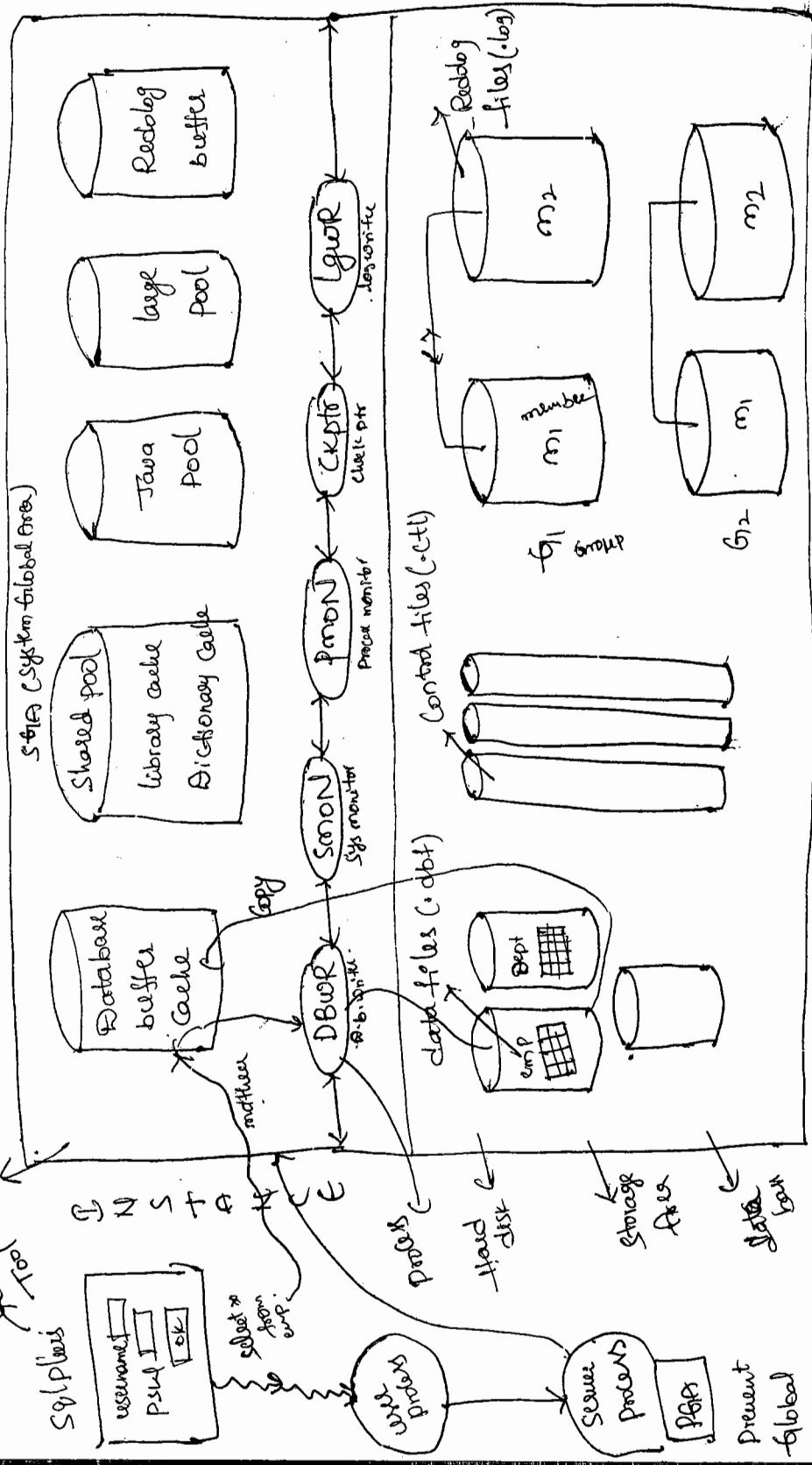
(2) Control files (.ctl).

(3) Redolog files (.log).

These 3 types of files are also called as "Storage Area"
or physical database.

Oracle Server Architecture

Oracle Server



Holes

Ex:

Recov

Alter

Ex:

Ex:

Ex:

foos

J

J

J

- Data files stores all database objects physically
- Data files extension is .dbf
- In oracle if u want to view path of the data files then we are using "dba_data_files" date dictionary

Ex:

> Conn sys as sysdba;

Enter psw: sys.

> desc dba_data_files;

> Select FILE_NAME from dba_data_files;

→ Control files stores databases information these files extension is .ctl. these files are used by DBA's in recovery process.

→ All Control files information stored under "V\$controlfile" data dictionary.

Ex: > Conn sys as sysdba;

Enter psw: sys.

> desc V\$controlfile;

> Select NAME from V\$controlfile;

→ Redolog files stores Committed (Saved) transactions from Redolog buffer.

→ These files also used by DBA in recovery process

→ All Non log files information stored under "V\$logfile" date dictionary

Ex:- > desc V\$logfile;

> Select member from v\$logfile;

Instance:-

when ever we are connecting to the Server using SQL plus tool then automatically an instance is created in "RAM" memory area.

This instance having 2 parts.

(1) SGA

(2) Back ground process.

(1) SGA:- (Shared / System Global Area)

SGA is also called as shared (or) system-GA.

→ SGA Contains Set of buffers. These are database buffer cache, shared pool, Java pool, large pool, Redolog buffer.

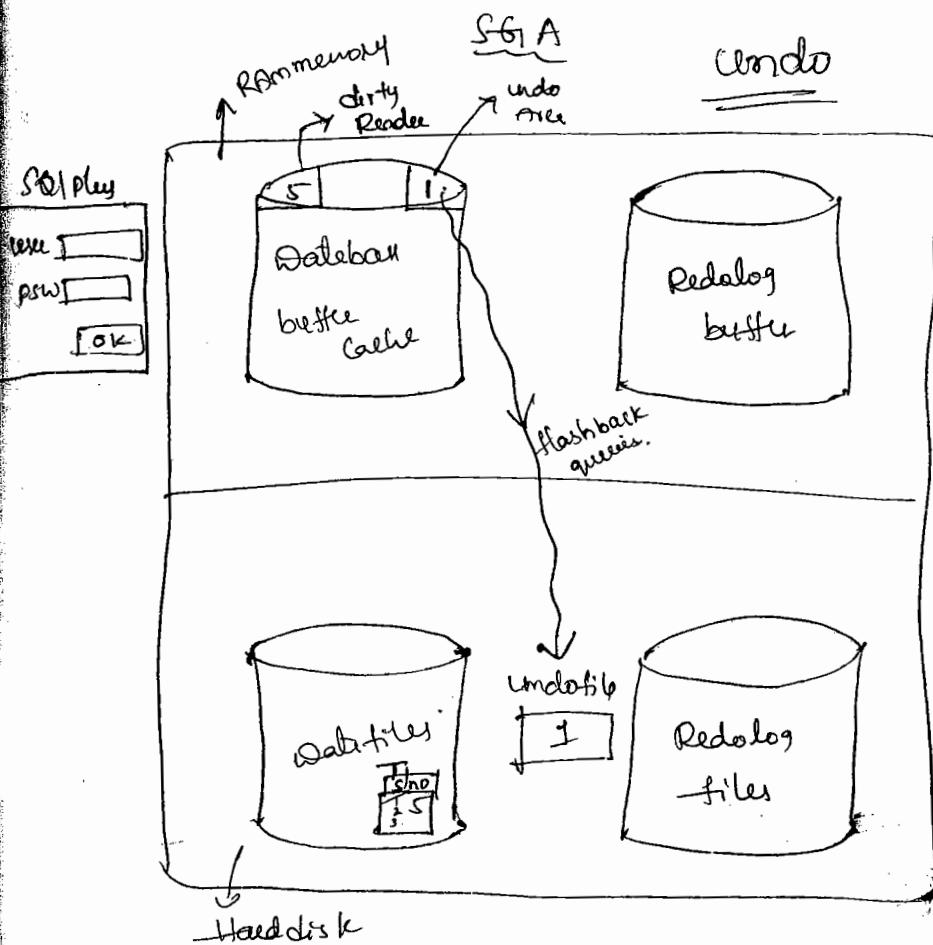
→ when ever we are submitting SQL, PLSQL Code that code is stored in library cache within shared pool.

→ Library Cache always checks Syntaxes Semantics

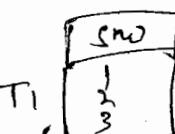
→ Library Cache reduces parsing by using parsing algorithms.

→ Library Cache also stores Cache values defined in sequence object.

- Shared pool also contains dictionary cache, which executes DCL related objects.
- Java pool executes java related objects. Large pool used by DBA in recovery process.
- Redolog buffer stores new information for the transaction.
- Server process also contains PGA which identify no. of clients uniquely & also PGA stores all collections informations from the PL/SQL.



- Update T_1 set $Sno=3$ where $Sno=1$,
- Commit;

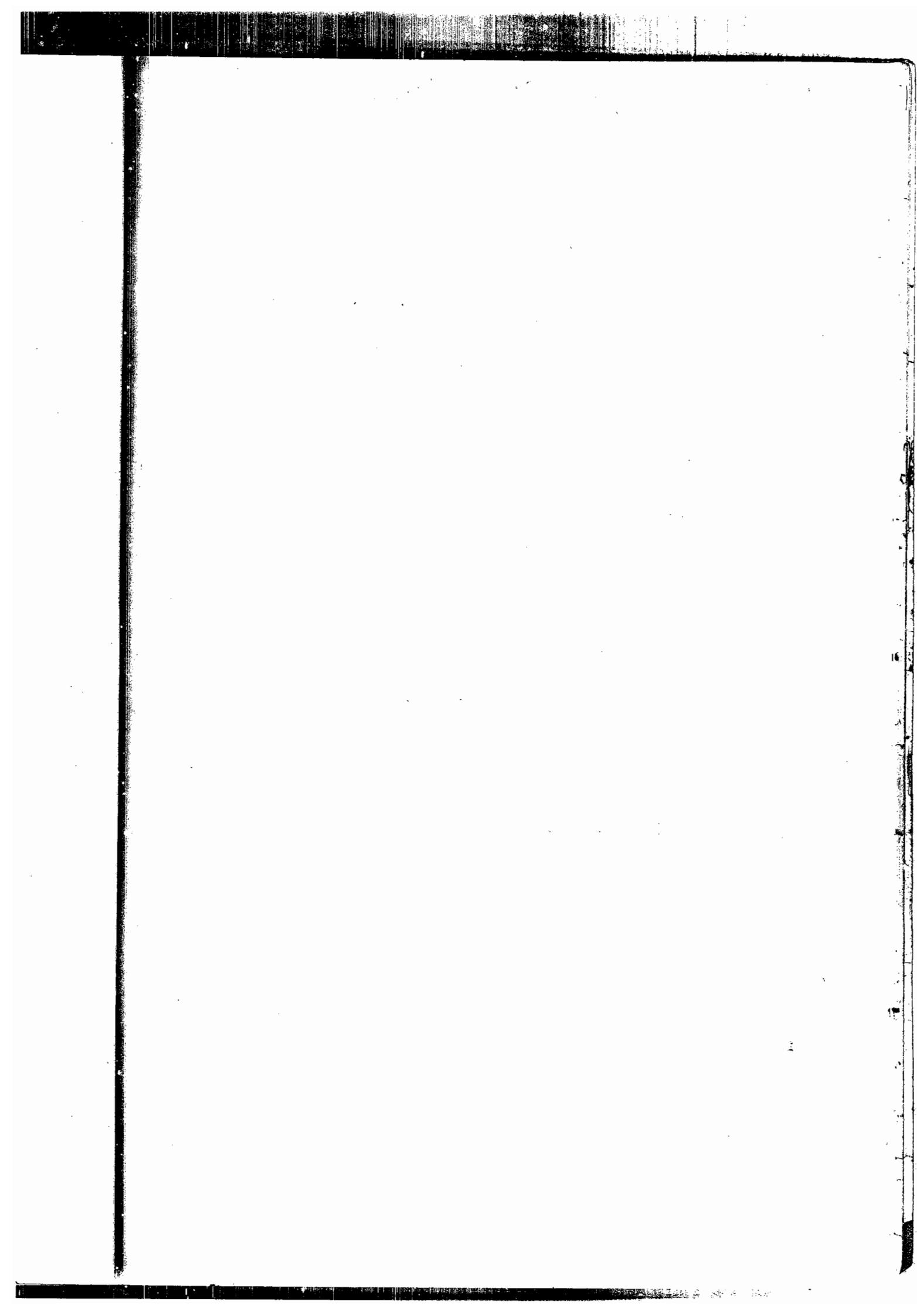


→ In Oracle whenever we are performing DML transaction it updates new values for the transaction automatically stores in dirty buffer & old value for the transaction stored in undo area within d.b. buffer cache.

→ Whenever ever user gives Commit dirty buffer value automatically permanently stored in data files and also undo area data automatically stored in undo file within database.

→ Oracle 9i introduced flashback queries to retrieve old data after committing the transaction also.

→ These flashback queries all works based on undotables.



Indexes:-

Indexes (use-ind-columns)

- Index is a d.b. object which is used to retrieve the data very fastly from the d.b. This process automatically improves performance of the application.
 - Generally indexes are created on table columns.
 - These indexes are created by d.b. sys.
 - In all D.B. Systems we are creating indexes in 2 ways
 - (1) Automatically
 - (2) manually.
 - In every D.B. Systems whenever we are requesting data using where clause or order by clause then only D.B. Servers searching for indexes.
- Note :- In oracle whenever where clause contains not equal to ($<$, $>$), is null, is not null operators then oracle server does not search for indexes if those columns already having indexes also. (group by clause)
2-8-14.

- (1) Automatically :- when ever we are creating p.key or unique key constraints then oracle servers internally automatically creates "btree indexes" on those cols
- (2) Manually :- we can also create indexes explicitly by using create index privilege

Syntax:-

[Create index indexname on tablename(columnname);]

→ Oracle having 2 types of indexes

(1) btree indexes.

(2) bitmap indexes.

→ generally when ever high Cardinality Cols are there
then only we are using "btree indexes".

where as in low cardinality Cols only we are creating "bitmap indexes".

Formulae:-

$$\text{Cardinality} = \frac{\text{distinct no. of records}}{\text{total no. of records}}$$

Ex: find empno. Cardinality.

→ Select distinct(empno) from emp;

⇒ 14.

$$\text{empno Cardinality} = \frac{14}{14} = 1 \text{ (high cardinality)} \text{ (btree)}$$

② Job cardinality.

→ Select distinct(job) from emp;

— 5.

$$\text{Job Cardinality} = \frac{5}{14} = 0.357 \text{ (low cardinality)}$$

bitmap index

btree indexes:- In oracle by default indexes are

"btree indexes" (unique, non unique)
default when index
or indexable via index.



→ when ever we are requesting data using where clause or order by clause then Oracle Server Searching for indexes on those columns in appropriate data dictionary.

→ If btree index is available then oracle server automatically creates btree structure on those columns.

→ In btree structure always "leaf blocks" stores

actual data along with rowids. Based on the

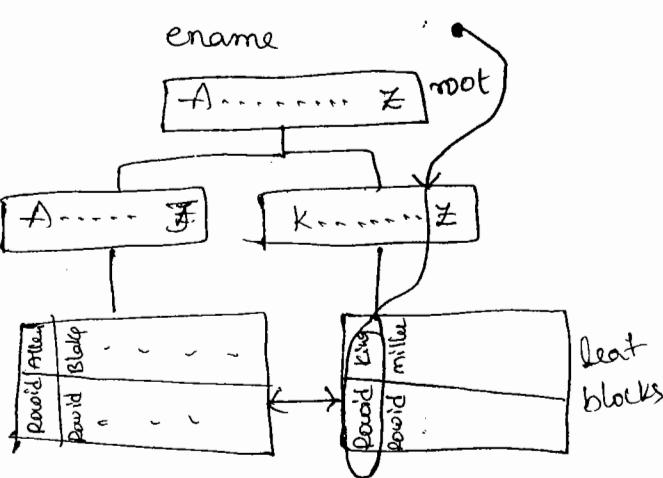
indexes scan mechanism oracle Server Search data

very fastly in btree structure & also retrieve data very fastly from the leaf blocks.

→ This process automatically improves performance of the application.

Developer

>Select * from emp (ename created as index to where ename='KING'; btreeformed)

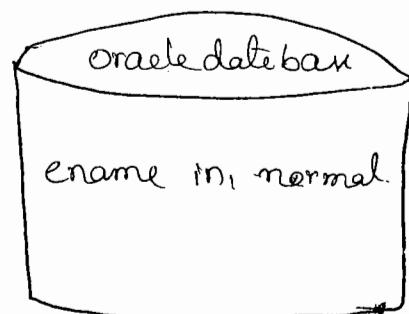


(Indexes Scan mechanism)

DBA

Create index in,

on emp(ename);



btree k.w → normal.

Calculate performance of the Query:-

Step 1:- Syntax:-

- explain plan for select * from tablename where Condition;

Step 2:-

After planning Select query. Then we must use display method from "dbms_xplan" package using following

Syntax.

➤ Select * from table(dbms_xplan.display());

Ex:- without using indexes

- explain plan for

Select * from emp where ename = 'KING';

- Select * from

table(dbms_xplan.display());

O.P:- Cost (-, CPU)
3.

Using indexes :-

- Create Index int on emp(ename);

- explain plan for select * from emp where ename = 'KING';

➤ Select * from table(dbms_xplan.display());

O.P:- Cost
1 (high performance)

→ In oracle all indexes information stored under "User_Indexes" dictionay

> desc user_indexes;

> Select INDEX_NAME, INDEX_TYPE from user_indexes
where TABLE_NAME = 'EMP';

| INDEX-NAME | INDEX-TYPE |
|--|---------------|
| (automatically when P.K. created than this index created) | PK-EMP (auto) |
| (manually we all created index) | IN1 (manu) |

* * *
* * * User_IND_COLUMNS * * *

→ In oracle if u want to view column names along with index names then we are using "User_IND_COLUMNS" data dictionary.

Ex:- > desc user_IND_COLUMNS;

> Select INDEX_NAME, COLUMN_NAME from

User_IND_COLUMNS

where TABLE_NAME = 'EMP';

O.P:

| INDEX-NAME | COLUMN-NAME |
|------------|-------------|
| PK-EMP | EMPNO. |
| IN1 | ENAME. |

to read

Function Based indexes:- (by default it is
normal index)
Btree

Oracle 8i introduced F.B.I's.

→ Generally in oracle when ever where clause contains fun's (or) expressions then oracle server does not search for indexes. if those columns already having indexes also.

→ To overcome this problem oracle 8i introduced extension of the btree indexes called function based indexes which creates indexes on fun's along with columns.

→ Syntax:-

create index index_name on
table_name (functionname (columnname));
or
expression

Ex:-

Create index idx on

emp(upper(ename));

- Select * from emp where upper(ename)='KINTI';
- Select index-name, index-type from user_indexes where
table-name = 'emp';

INDEX-NAME

IN1

INDEX-TYPE

Function-Based Normal

~~After~~ Virtual Columns:-

3-~~8~~-14.

Oracle 11g introduce Virtual Columns, which stores stored expressions directly in database.

→ generally before oracle 11g we can also stored stored expressions in oracle d.b. indirectly through Views, function based indexes.

→ oracle 11g having Virtual Columns which stores stored expression through "generated always as clause"

Syntax :-

Columnname datatype(size) generated always as

(function-name (columnname)) [virtual].
(or)
expression

- > Create table t1 (a number(10), b number(10), c number(10) generated always as ($a+b$) virtual);
↳ stored expression (previous col values)
- > desc t1;
- > insert into t1(a,b) values (10,20);
- > Select * from t1;

| a | b | c |
|----|----|----|
| 10 | 20 | 30 |

Note:- If we want to view virtual column expressions then we are using "data-default" property from "user-tab-columns" data dictionary.

Stores

Ex: > desc user-tab-columns;

> select column-name, data-default from

user-tab-columns where table-name = 'T1';

| <u>COLUMN-NAME</u> | <u>DATA-DEFAULT</u> |
|--------------------|---------------------|
| C | "A+B" |

h Stores

Bitmap Indexes

class

Oracle 7.3 introduced Bitmap Indexes

→ Generally Bitmap indexes are created on "low cardinality" columns

→ Bitmap indexes are created on low cardinality.

and in "duplicating applications,"

→ Generally we are not allowed to create unique btree indexes on "duplicate value columns" (job)

→ In place of this one we can also create bitmap indexes.

Syntax:-

Create bitmap index indexname on tablename (columnname);

Create unique index in_1 on emp(ename);

Index created

Create unique index in_2 on emp(job);

error: Can not create unique syntax Index: duplicate value taken.
80

Solutions:-

→ Create bitmap index in_2 on emp(job);

→ Index created.

⇒ whenever we are creating bitmap index then oracle server internally automatically creates a bitmap table.

→ whenever user requesting data using logical operators then oracle server directly operates on the bits within bitmap table.

→ Then resultant bitmap automatically converted into rowid using an internal bitmap fun.

→ Create bitmap index in_2 on emp (job).

| Job | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 | 11 | 12 | 13 | 14 |
|-----------|---|---|---|---|---|---|---|---|---|----|----|----|----|----|
| Clerk | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 | 0 | 1 |
| Salesman | 0 | 1 | 1 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 |
| Manager | 0 | 0 | 0 | 1 | 0 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |
| Analyst | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 1 | 0 |
| President | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 |

Table can also drop Index using "drop index"
index name"

(reference name) Synonyms :- (view(security purpose) (by default private)
dropping synonym to table etc to do

edit
See below.

- Synonym is a d.b. object
- which hides another Schema, username, object name, d.b. link
- Synonym is an alias name or reference name for the table, view, procedure, function etc
- Synonyms are created by d.b. administrator
- All d.b. systems having 2 types of Synonyms
 - (1) public Synonyms,
 - (2) private Synonyms.
- In Oracle db. by default Synonyms are "private"
- If we want to create public Synonyms then we must give "Create public Synonym Privileges" to user Syntax

Grant Create public Synonym to username;

Private Synonym Syntax:-

- Create Synonym Synonym name for username • object name @ database link;

Public Synonym:-

Create public Synonym Synonymname

for username.objectname @ database link;

| Scott/Tiger | murali/murali° | ar/ar |
|--|---|------------------------------|
| > Grant all on emp to murali ar; > Select * from Scott.emp ; (v) > Create synonym abc _{synonym} for Scott.emp ; > Select * from abc; (v) > Create public synonym xyz for Scott.emp ; > Conn sys as sysdba; Enter psw: sys In sufficient privileges > Conn murali° / murali°; > Create public synonym xyz for Scott.emp ; > Select * from xyz; (v) | > Select * from scott.emp ; (v) > Select * from abc; <u>error</u> : Table or view does not exist > Select * from xyz; (v) | 2 1 → + (1) too |
| | > Select * from xyz; (v) | (2) |

21-9-14

→ All synonym information stored under user-synonyms data dictionary.

> desc user-synonyms;

TCL:- (Transaction Control Language) :-

Transaction- A logical unit of work in between

2 points is called "transaction".

→ All d.b. systems having 2 transactional Commands

(1) Commit.

(2) Savepoint.

(1) Commit- This command is used to save the transaction permanently into d.b. (All DDL cmds automatically committed)

Syntax:-

Commit ;

(2) Savepoint- It is a logical mark b/w transactions

Syntax:-

Savepoint Savepointname ;

rollback- This command is used to undo the transactions from memory.

Syntax:-

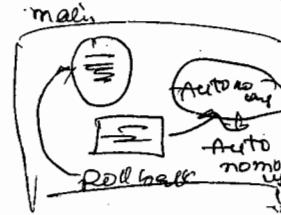
rollback ;

Particular

Rollback to Save Point:-

Syntax:

`rollback to Savepointname;`



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Rollback to Save Point:-

Syntax:

`rollback to Savepointname;`

Eg:- delete from emp where empno=1;

> update emp set sal = sal - 200 where ename='SMITH';

> Samepoint S1;

> update emp set sal = sal - 100 where ename='SMITH';

> Insert into emp (empno) values (5);

> Samepoint S2;

> update emp set sal = sal + 100 where ename='SMITH';

> Rollback to S1; (S1, S2 Rollback) above S1 are committed.

> Commit;

— Commit Completed.

(Create or add type) [undefined dty] Instead Table :- (Index by Table ^{overcome problem} extentions is Not) :- Not here in db.

→ Oracle 8.0 introduced instead table basically instead table is a user defined type which is used to store no. of data items in a "single unit".

→ generally if we want to store large amount of data into single unit then we are using index by table in PL/SQL.

→ But we are not allowed to store index by tables permanently into database.

→ To overcome this problem Oracle 8.0 introduced extension of the index by tables called nested tables which stores

permanently into d.b. using SQL..

Creating Nested Tables

Step 1:- Create an object type.

Step 2:- Create a nested table type

Step 3:- Create actual table.

Create an object type

This is a user-defined type which is used to represent different date types into single unit.

It is also same as structures in C language.

Syntax:-

create or replace type typename

as object (attributename1 datatype(size), attributename2 datatype(size) ...)

Create an nested table type:-

create or replace type typename

as table of objectname type;

Step 3:- Create actual table:-

Syntax:- create table tablename (col1 datatype (size),

... coln nested table type) nested table column

Stores as any name

Ex: Create or replace type e1 as object (book ID number (10), book name Varchar2(10), price number (10));

> Create or replace type xy as table of e1;

/ - complex.

> Create table Student (Stno number(10), sname Varchar2(10),

Col3 xy) instead table Col3 Same Stored as PQX;
Kind address Copy name
its content)

> desc student;

| name | type. |
|-------|-------------------|
| Stno | Number (10) |
| sname | Varchar2 (10) |
| Col3 | xy , type name |

Note: If u want to store data into instead table type then we must use Constructor.

→ here Constructor name is same as type name.

> insert into student values (1, 'murali') as (e1 (101, Java, 900))

900); e1(102, 'PLSQL', 400));

> Select * from student;

O.P.: Col3 (BOOKID, Bookname, price)

murali

xy (e1 (101, 'Java', 900), e1 (102, 'PLSQL', 400)));

Now created

mes

BOOK Pd

- (10);

address (10),

(Copy name
its works)

| Stno | Sname | Col3 |
|------|--------|------|
| 1 | murali | |

Indirectly Storing
Nested table

| Bookid | Bookname | price |
|--------|----------|-------|
| 101 | java | 900 |
| 102 | PLSQL | 400 |

> Select * from table (Select Col3 from student);

| Bookid | Bookname | price |
|--------|----------|-------|
| 101 | Java | 900 |
| 102 | PLSQL | 400 |

> update table (Select Col3 from Student) Set price=1000

where BookId = 101;

ad table

> Select * from table (Select Col3 from student);

| | | |
|-----|-------|------|
| 101 | Java | 1000 |
| 102 | PLSQL | 400 |

(101, Java)

→ If u want to view nested table date or modified nested table date then we must use 'Table' operator.

→ & also along with table operator we must select nested table columns within subquery

Syntax:

- Select * from table (select col3 from student);
- Select * from table (Select nested table colname from actual table name);

D.P. :-

| | | |
|-----|-------|------|
| 101 | Java | 900 |
| 102 | PLSQL | 400. |

Partitions:

- partition tables are created by database administrator partition tables are used in back up and recovery process which is used to improves performance of
- Partitions tables are used in data warehousing applications.
- generally partitions are created based on a key column
- This Col is also called as partition key
- when database contains large amount of data then only we are creating partition table.
- Oracle having 3 types of partition

- (1) Range Partitions.
- (2) List Partitions.
- (3) Hash Partitions.

Range Partitions:- In this method database admin +);
-strator create partitions based on range of values.
from

Syntax:-

create table tablename (col, datatype (size), ...)
partition by range (key column name)

(Partition partitionname values less than (actual value)),

partition partitionname2 values less than (actual value),

partition partitionname value less than (maxvalue)

retrieve particular partitions:-

Syntax:-

Select * from tablename

partition (partitionname1, partitionname2, ...),

Ex:-

Create table test (empno number(10), ename
varchar2(10), sal number(10))

partition by range (sal)

partition P1 values less than (1000),

partition P2 values less than (2000),

partition Others values less than (maxvalue));

> insert into test values (...);
> Select * from test;
> Select * from test;

Partition (P1);

| EMPNO | ENAME | SAL |
|-------|-------|------|
| 1 | a | 500. |

Note:- We can also add partition by using alter with add into the existing partition table

Syntax:

```
alter table tablename add  
partition partitionname values  
less than (actual value);
```

List Partitions

Oracle 9i introduce list partition in this method data base administrator partitioning the table by using list of values.

Syntax: Create table tablename (col ,daty (size), ...)

partition by list (key columnname)

(partition partition name1 values (value1, value2, ...),

partition partition name2 values (value1, value2, ...),

Particular particular name values (Default));

Ex:- Create table test1 (sno number(10), name varchar(10))
Partition by list (name)
(partition P1 values ('india', 'Pakistan'),
partition P2
values ('us', 'uk', 'canada'),
partition others
values (default));

alter
table
 > insert into test1 values (...);
 > select * from test1 partition (others);

| <u>SNO</u> | <u>Name</u> |
|------------|-------------|
| 1 | Japan |
| 2 | australia |

Hash Partition- In this method oracle server
only create partition based on specified column.
by using an internal "hash algorithm".

Syntax:-

Create table tablename (col1 datatype(size), ...)

Partition by hash(key colname)

Partitions any number;

Ex:- Create table test3 (sno number(10), sal number(10))

Partition by hash(sal)

Partitions 5;

→ all partition information stored user-tab-partition date dictionary.

Ex: SQL> desc user_tab_partitions;

→ Select partition-name from user-tab-partitions
where table-name = 'Test 3';

Hierarchical queries- In relational date bases we can also stored hierarchical data

→ In oracle if we want to retrieve hierarchical data then we are using level pseudo column along with following clauses.

(1) Start with

(2) Connect by.

→ If we want to store hierarchical data in relational table minimum 2 columns are required and also these 2 column must belong to same datatype & also having some relation.

→ level is a Pseudo Column which assigns no's to each level in Tree Structure

→ Start with:-

This clause specify searching condition with in hierarchy

ab-

Syntax:- Start with Condition

Connect by :- This clause is used to specify deletion
row or

ship b/w parent, child Column.

- In this clause we must used prior operator

Syntax's Select Col1, Col2

from tablename

where Condition

Start with Condition

Connect by prior

Parent column name = child column name;

> Write a hierarchical query to display employee who
are working under another employees root node
onwards from emp table using employee number,
mfr Columns.

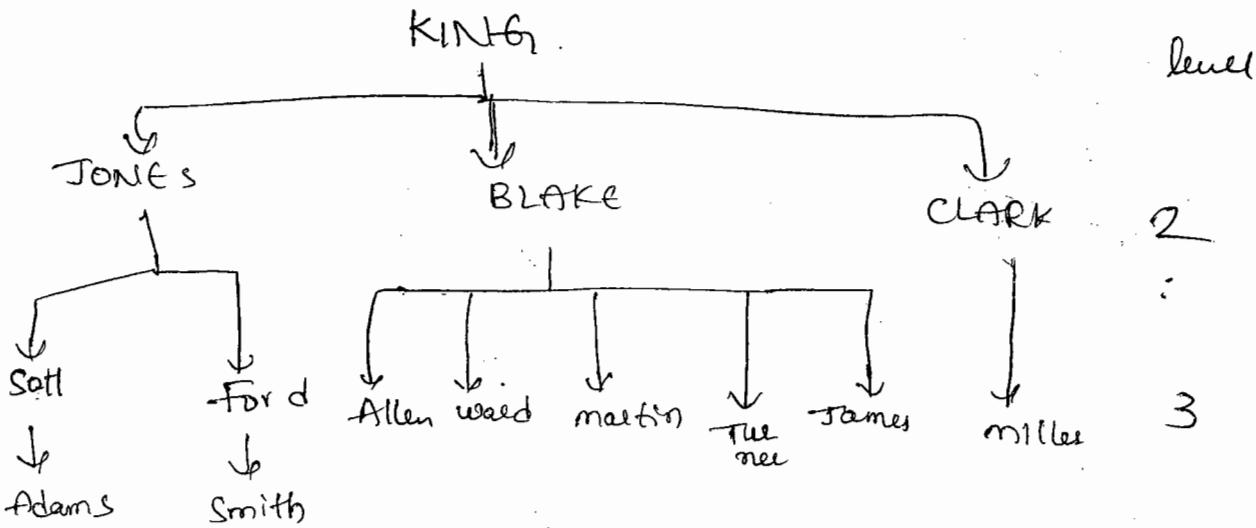
SELECT level, fname

from emp

level

Start with mfr is null

Connect By prior Empno=mfr.



Note:- Oracle 9i introduced sys-connect-by-path() function which returns path of hierarchy entry structure.

This fun accepts 2 Parameters

Syntax:-

sys-connect-by-path

(Columnname, 'delimitername')

Ex:-

Select

level : sys-connect-by-path (ename, '→')

From emp

Start with mgr is null

Connect by Prior empno = mgr

1.

level

Q) write a hierarchical query to display the employee who are working under Blake from emp table using emp no, mgr columns.

2

Select level,

3

Sys - Connect by - path(ename, '→')

From emp

Start with ename = 'Blake'

Connect by empno = MGR

/

Select level : Sys - Connect - by - path(ename, '→')

From emp

Start with ename = 'miller'

Connect by empno = prior mgr.

/

level ename

1 → miller

2 → miller → clark

3 → miller → clark → king

Plus is an unary operator and along with
Connect by clause always prior operator between
Parent column in hierarchy

Using prior operator we can also travel top to bottom (or) Bottom to top in the tree.

Note: oracle 10g introduce **Connect - By - IsLeaf**

Pseudo column which returns 1 to the leaf node and '0' to the all other nodes in the tree.

Syntax:

```
Select  
level, ename, Connect - by - IsLeaf  
from emp
```

Start with mgr is null

Connect by prior empno = mgr

Note: oracle 10g introduced **Connect - By - Root**

operation which returns root node in the hierarchy

Syntax:

Connect - by - Root Columnname,

Ex:

```
Select level, ename,  
Connect - by - root ename  
from emp
```

Start with mgr is null

Connect by prior empno = mgr

Up to

Interview questions Flash back Queries.

- Oracle 9i introduced flashback queries using flashback queries data base administrator to retrieve accidental date
- flashback queries always allows content of a table can be queried with respect to a specific point of time using As of clause
- Flash back queries works based on undo file using flashback query we can retrieve previous data after committing date also
- To retrieve accidental date then we are using 2 methods
 - (1) using timestamp datatype
 - (2) using SCN (System Changed Number)

Method - 1:- using timestamp datatype

Syntax:

Select * from tablename

as Of timestamp(specific point of time);

- Ex > Create table t1 (Sno number (10));
 > Insert into t1 values (.....);
 > Commit;
 > Select * from t1;
- | | |
|------|-----|
| O.P. | SNO |
| 1 | |
| 2 | |
| 3 | |
| 4 | |
- > delete from t1;
 > Commit;
 — Commit Complete.
 > Select * from t1;
 — no rows selected.
 > Select * from t10 as of
 timestamp (Systimestamp - interval '2' minute);

| Sno |
|-----|
| 1 |
| 2 |
| 3 |
| 4 |

Method 2 (using seq)

→ In oracle whenever we are creating transaction then automatically a unique identification no is created in the date base.

→ This unique identification number is also called as system changed number and on this number also data base administrator can do accidental date

Syntax :-

Select * from tablename

as of scn System changenumber;

~ ~ ~ ~ ~ > conn sys as sysdba;

Enter password : sys -

> Create table m1 (Sno number(10)),

> desc v\$database;

> Select Current_scn from v\$database;

CURRENT_SCN :-

2306347.

> Insert into m1 values(5);

> Commit;

> Select * from m1;

SNO
5.

> Select Count(*) from m1 as of scn 2306347;

Count(*)
0.

Question

10 → 25

