

Problem Solving Workshop #28

Tech Interviews and Competitive Programming Meetup

September 2, 2017

<https://www.meetup.com/tech-interviews-and-competitive-programming/>

Instructor: Eugene Yarovoi (can be [contacted](#) through the group Meetup page above under Organizers)

More practice questions: leetcode.com, glassdoor.com, geeksforgeeks.org

Books: Elements of Programming Interviews, Cracking the Coding Interview

Have questions you want answered? [Contact the instructor](#), or ask on [Quora](#). You can post questions and [follow the instructor](#) and other people who write about algorithms.

Try to find optimized solutions, and provide a time and space complexity analysis with every solution.

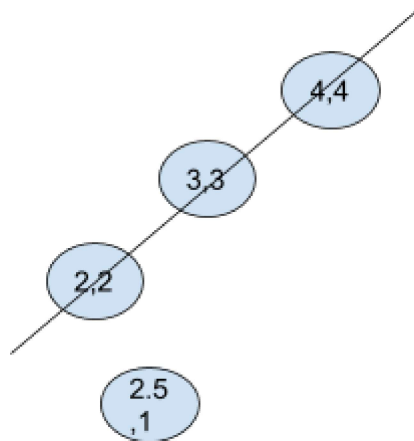
Problem #1, “Max Points on a Line” (Easy / Medium)

Given N points in 2D space (assume integer coordinates), find the largest set of points that are collinear (lie on the same line).

Example: [(2,2), (3, 3), (4, 4), (2.5, 1)]

Output: [(2,2), (3, 3), (4, 4)]

Explanation:

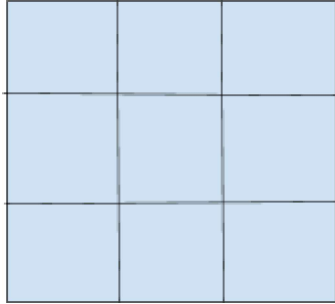


It is not possible to select a larger set of points so that they’re all on the same line.

Complexity Guidance: **[Very Easy]** $O(N^3)$ time **[Easy]** $O(N^2)$ time / $O(N^2)$ space, **[Medium]** $O(N^2)$ time / $O(N)$ space.

Problem #2, “Square Counting” (Medium)

You’re given an $N \times N$ grid, that starts out looking like this. For example;



Your goal is to count how many total **squares** there are in the diagram. We count something as a square if it is an arrangement of 4 lines in the way you would normally expect square to be defined, without regard as to whether there are any extra lines or what is inside the square. A line segment may be part of multiple different squares. For example, the above diagram has 1 3x3 square, 4 2x2 squares, and 9 1x1 squares, for 14 total squares.

(a) **[Easy]** Provide an algorithm for counting the number of squares in an $N \times N$ grid.

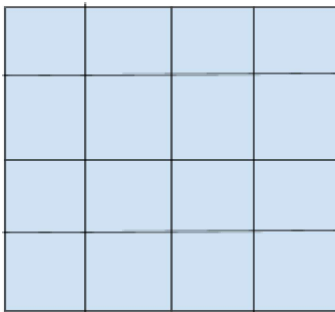
Example Input: 3 (meaning 3x3 grid)

Output: 14 (as explained above)

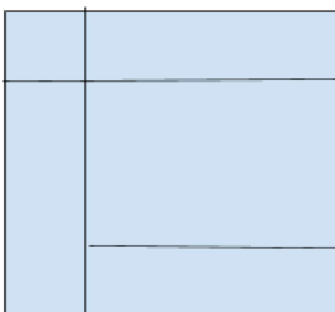
Example Input: 4

Output: 30

(b). Now, some grid lines may be erased. For example, a 4x4 grid starts out like this:



And then may have some lines erased to look like this:



Line segments can only be erased between intersection points of lines in the complete grid. In other words, the above drawing can be specified as a $(N+1) \times N$ matrix of booleans denoting which horizontal line segments (of length 1) to keep and a $N \times (N+1)$ matrix denoting which vertical line segments to keep. The above drawing would be represented as:

Horizontal lines:

```
[T T T T
T T T T
F F F F
F T T T
T T T T]
```

Vertical lines:

```
[T T F F T
T T F F T
T T F F T
T T F F T]
```

(b) Count the number of **squares** in such an arrangement of lines.

The above example has 1 1×1 square (top left), 2 3×3 squares (overlapping, on the right side), and 1 4×4 square (the overall outline). So the answer would be 4.

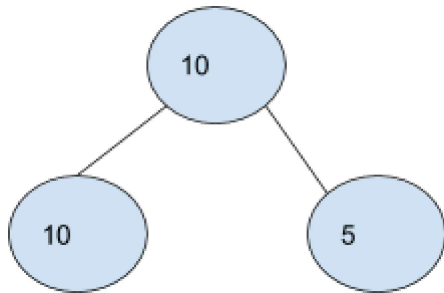
Complexity Guidance: **[Easy]** $O(N^4)$ naive solution, **[Medium]** $O(N^3)$ time, **[Hard, don't get too stuck on it]** $O(N^2 \log N)$ time.

Problem #3, "Graph Sums" (Hard)

Consider a graph where every vertex has a value associated with it. We want to create a data structure on this graph that will support the following operations:

- `Update(N, V)`: changes the value of node N to V .
- `SumNeighbors(N)`: returns the sum of the values of all of N 's neighbors.

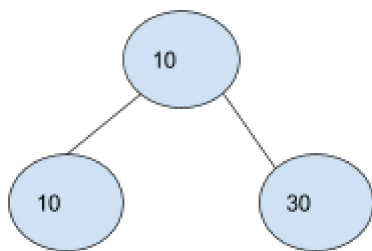
Example:



SumNeighbors(top node) -> 15 (Explanation: 5+10)

SumNeighbors(right node) -> 10 (Explanation: only connected to the top node)

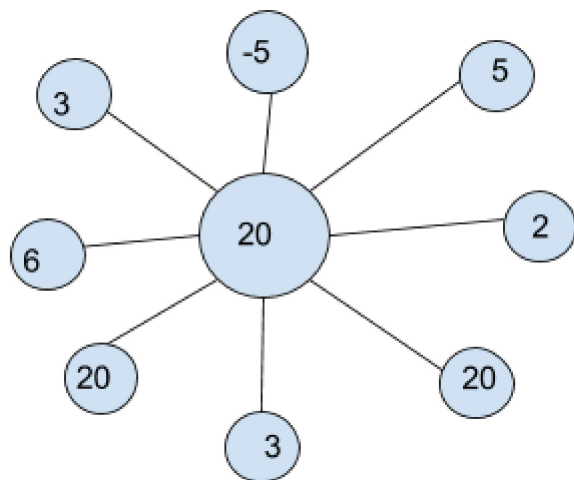
Update(right node, 30)



SumNeighbors(top node) -> 40

SumNeighbors(right node) -> 10 (top node didn't change)

The obvious approach is to just store each node's value in the node itself so that Update can run in $O(1)$, but then SumNeighbors runs in $O(\text{neighbor count of node } N)$ time. The problem with this is, for example, with graphs that are structured like this:



Here the center node could take a long time to sum up (imagine an extension of this concept where the center node has a million neighbors).

Can you design an approach that better balances the running time of Update against the running time of SumNeighbors?