CMSC 441 LCS Project 2 Report

By,

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Project Overview:

For this project our task was to design and implement a serial and parallel implementation of the longest common subsequence (LCS) algorithm. The basis of our serial design was out of the *Introduction to Algorithms 3rd. edition* textbook. After estimating the runtime the serial version would achieve, we implemented a parallel version of the LCS algorithm. We tested the parallel version against multiple CPUs with various inputs. Afterwards we implemented a memory efficient parallel version of the parallel LCS algorithm. Reducing the memory usage from (m \* n) to (m + n). Below is an analysis and look into the implementation of our groups LCS algorithms.

Design:

Calculations:

*Work Calculation:*

We define work law to be the summation of all runtimes of a serial algorithm.

Since dominates the runtime in the summation above we can ignore all the constant work that’s done with an and sufficiently large.

If we bound the parameters and such that then

*Span Calculation:*

The span calculation is based off

When analyzing the parallel LCS code, we can see that the outer most for loop will take

Within the algorithm there is constant time work we will denote as:

Replacing the span calculations within the span law we get the equation

For the equation above. Given an and sufficiently large, we can reduce ignore the constant time work and can remove any constant time addition / subtraction :

The theta above clearly dominates the runtime of the algorithm. Since any runtime greater than logarithmic runtimes will always take longer to run, then we can remove all logarithmic runtimes from the equation.

There’s only one non-logarithmic and non-constant runtime left. However defined above we bound the parameters and such that therefore

We can reduce the runtime above by the definition of the theta runtime, where

This is because the asymptotic complexity of a linear runtime is such that

Therefore the span of the parallel LCS algorithm is

*Parallelism Calculation*

The parallelism calculation is based off

*Parallelism*

Computing the *parallelism* is quite easy, we simply plug in the runtimes for work over span.

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*Linear-speed-up Estimation*

First we’ll define the term’s we’re going to be using.

= running time on P processors.

We are bounding the inputs such that

*=*

Predictions:

The expectation of the serial algorithm based off the work calculation was that it was going to take a fair amount of time for algorithm to run. Algorithms that have an (n^2) runtime get exponentially slower overtime. Therefore we should notice a steep increase in runtime the larger m and n get.

For the parallel algorithm, we are going to test the algorithm on 1, 2, 4, 8, and 16 CPUs. My expectation for a single processor is that it’s going to about the same as the serial implementation above. I expect that it might take a bit longer, since the parallel LCS algorithm is a different implementation than the serial version. Once the algorithm goes through 2 to 16 processors respectively, the expectation is to see increasingly faster runtimes the more processors we throw at it. The limitation of this however is that we can only have as many threads running on the algorithm as there are diagonal spaces allowed per iteration. So this algorithm will operate very quickly with very large inputs and with a high number of processors.

Empirical Performance:

|  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- |
| Parallel | CPUs and Runtime in Seconds | | | | |
| Input size | CPU 1 | CPU 2 | CPU 4 | CPU 8 | CPU 16 |
| 10 x 10 | 0.000028 | 0.000179 | 0.000258 | 0.000371 | 0.000422 |
| 25 x 25 | 0.000072 | 0.000381 | 0.000403 | 0.000521 | 0.000601 |
| 50 x 50 | 0.000198 | 0.000597 | 0.000518 | 0.00066 | 0.000813 |
| 75 x 75 | 0.00049 | 0.000671 | 0.000721 | 0.000785 | 0.001159 |
| 100 x 100 | 0.000842 | 0.001018 | 0.001129 | 0.001184 | 0.001668 |
| 175 x 175 | 0.002209 | 0.002181 | 0.001812 | 0.001823 | 0.002652 |
| 250 x 250 | 0.002874 | 0.003387 | 0.002618 | 0.002876 | 0.00368 |
| 500 x 500 | 0.008859 | 0.009768 | 0.006685 | 0.005299 | 0.006911 |
| 750 x 750 | 0.016808 | 0.015785 | 0.011328 | 0.007322 | 0.010195 |
| 1,000 x 1,000 | 0.030712 | 0.022057 | 0.015641 | 0.010488 | 0.014563 |
| 1,250 x 1,250 | 0.048662 | 0.035343 | 0.021419 | 0.017548 | 0.018298 |
| 1,500 x 1,500 | 0.071463 | 0.047607 | 0.027002 | 0.019819 | 0.026714 |
| 1,750 x 1,750 | 0.098836 | 0.065561 | 0.037614 | 0.026792 | 0.030848 |
| 2,000 x 2,000 | 0.130177 | 0.085987 | 0.046703 | 0.032783 | 0.029263 |
| 2,250 x 2,250 | 0.16484 | 0.103802 | 0.057866 | 0.036123 | 0.033444 |
| 2,500 x 2,500 | 0.203756 | 0.125321 | 0.067985 | 0.043438 | 0.042173 |
| 5,000 x 5,000 | 0.917295 | 0.485617 | 0.280697 | 0.149459 | 0.110557 |
| 7,500 x 7,500 | 2.138821 | 1.193619 | 0.581892 | 0.321084 | 0.198746 |
| 10,000 x 10,000 | 3.841368 | 2.140341 | 1.073969 | 0.567078 | 0.328026 |
| 12,500 x 12,500 | 6.325828 | 3.522792 | 1.762813 | 0.867938 | 0.47751 |
| 15,000 x 15,000 | 9.264918 | 5.184949 | 2.64153 | 1.275007 | 0.662077 |
| 17,500 x 17,500 | 12.879671 | 7.209504 | 3.69968 | 1.835184 | 0.892528 |
| 20,000 x 20,000 | 17.189502 | 9.504888 | 4.931689 | 2.501318 | 1.183977 |
| 22,500 x 22,500 | 22.577094 | 12.350231 | 6.430517 | 3.115865 | 1.413076 |
| 25,000 x 25,000 | 30.258455 | 15.527944 | 8.070663 | 4.017151 | 1.694669 |
| 27,500 x 27,500 | 36.305059 | 18.957842 | 9.733501 | 4.811172 | 2.061413 |
| 30,000 x 30,000 | 44.994047 | 22.409602 | 11.607725 | 5.845212 | 2.496002 |

Runtime Analysis:

Pseudocode: