MPS¹ Protocol – BAN logic

Sign Up Real Protocol

$$M_1: A \rightarrow S: K_a, E_{k_s}(A, N_a)$$

$$M_2: S \to A: E_{k_a}(A, N_a, N_s, K_{as}) | |E_{k_s}(S(E_{k_a}(A, N_a, N_s, K_{as})))$$

$$M_3$$
: A \rightarrow S: $E_{k_{as}}(A, N_s, g_a, p_a)$

$$M_4: S \rightarrow A: E_{k_{as}}(A, N_s + 1, g_a, p_a)$$

Sign Up Idealized Protocol

$$M_1$$
: A \rightarrow S: $\{A, N_a\}_{k_s}$

$$M_2: S \to A: \left\{A, N_a, \ \#(N_a), N_s, (A \overset{k_{as}}{\longleftrightarrow} S) \right\}_{k_a} | \left| \left\{S(\left\{A, N_a, \ \#(N_a), N_s, (A \overset{k_{as}}{\longleftrightarrow} S) \right\}_{k_a}\right)\right\}_{k_s^-}$$

$$M_3$$
: A \rightarrow S: $\{A, N_s, \#(N_s), g_a, p_a\}_{k_{as}}$

$$M_4: S \to A: \{A, N_S + 1, \#(N_S), g_a, p_a\}_{k_{g_S}}$$

Sign Up Protocol Analysis

Objectives

$$S \models (g_a, p_a)$$

$$A \models S \models (g_a, p_a)$$

 $A \models A \stackrel{k_{as}}{\longleftrightarrow} S$ # Symmetric Key

$$S \models A \models A \stackrel{k_{as}}{\longleftrightarrow} S$$

Assumptions

 $A \models \stackrel{k_S}{\rightarrow} S$

 $S \models^{k_a} A$

plaintext in the first message

 $S \Rightarrow A \stackrel{k_{as}}{\longleftrightarrow} S$

 $S \models A \Rightarrow g_a, p_a$

After M_1 :

$$S \triangleleft \{A, N_a\}_{k_c}$$

Hardcoded Server's Public Key

Diffie Hellman Parameters

Alice's Public Key sent in

Symmetric Key

¹ Magherini – Pochiero – Sieni (MPS)

After M_2 :

$$A \models \stackrel{k_s}{\to} S , A \triangleleft \left\{ A, N_a, \#(N_a), N_s, (A \stackrel{k_{as}}{\leftrightarrow} S) \right\}_{k_s^-}$$

$$A \models S \mid \sim (A, N_a, \#(N_a), N_s, (A \stackrel{k_{as}}{\leftrightarrow} S))$$

$$A \models S \mid \sim \left(A, N_a, \#(N_a), N_s, (A \stackrel{k_{as}}{\leftrightarrow} S) \right), A \models \#(N_a)$$

$$A \models S \models (A, N_a, \#(N_a), N_s, (A \stackrel{k_{as}}{\leftrightarrow} S))$$

$$A \models S \models \left(A, N_a, \#(N_a), N_s, (A \stackrel{k_{as}}{\leftrightarrow} S) \right), S \Rightarrow A \stackrel{k_{as}}{\leftrightarrow} S$$

$$A \models A \stackrel{k_{as}}{\longleftrightarrow} S$$

After M_3 :

$$S \models A \stackrel{k_{as}}{\longleftrightarrow} S , S \triangleleft \left\{A, N_s, \#(N_s), g_a, p_a, (A \stackrel{k_{as}}{\longleftrightarrow} S)\right\}_{k_{as}}$$

$$S \models A \mid \sim (A, N_s, \#(N_s), g_a, p_a, (A \stackrel{k_{as}}{\longleftrightarrow} S))$$

$$S \models A \mid \sim (A, N_s, \#(N_s), g_a, p_a, (A \stackrel{k_{as}}{\longleftrightarrow} S)), S \models \#(N_s)$$

$$S \models A \models (A \stackrel{k_{as}}{\longleftrightarrow} S)$$

$$S \models A \Rightarrow g_a, p_a, S \models A \models (A, N_s, \#(N_s), g_a, p_a, (A \stackrel{k_{as}}{\longleftrightarrow} S))$$

$$S \models (g_a, p_a)$$

After M_4 :

$$\frac{A \vDash A \overset{k_{as}}{\longleftrightarrow} S , A \vartriangleleft \{A, N_s, \#(N_s), g_a, p_a\}_{k_{as}}}{A \vDash S \mid \sim (A, N_s, \#(N_s), g_a, p_a)}$$

$$\frac{A \vDash S \mid \sim (A, N_s, \#(N_s), g_a, p_a), A \vDash \#(N_s)}{A \vDash S \vDash g_a, p_a}$$

Authentication Real Protocol

$$M_1$$
: A \rightarrow S: $E_{k_{as}}(A, N_a)$

$$M_2: S \rightarrow A: E_{k_{as}}(A, N_a, N_s)$$

$$M_3: A \rightarrow S: E_{k_{as}}(A, N_s)$$

Authentication Idealized Protocol

$$M_1: A \rightarrow S: \{A, N_a\}_{k_{as}}$$

$$M_2: S \rightarrow A: \{A, N_a, \#(N_a), N_s\}_{k_{as}}$$

$$M_3$$
: A \rightarrow S: $\{A, N_s, \#(N_s)\}_{k_{as}}$

Authentication Analysis

Objectives

$$A \vDash \#(N_S)$$
, $S \vDash A \vDash \#(N_S)$

Session ID Establishing

Assumptions

$$A \models A \stackrel{k_{as}}{\longleftrightarrow} S$$
 $S \models A \stackrel{k_{as}}{\longleftrightarrow} S$

$$A \vDash S \Rightarrow N_c$$
. $S \vDash \#(N_c)$

Key Registration

Nonce Authority

After M_1 :

$$\frac{S \vDash A \overset{k_{as}}{\longleftrightarrow} S, S \vartriangleleft \{A, N_a, \#(N_a), N_s\}_{k_{as}}}{S \vDash A \mid \sim (A, N_a, \#(N_a), N_s)}$$

After M_2 :

$$\frac{A \vDash A \overset{k_{as}}{\longleftrightarrow} S \ , \ A \vartriangleleft \{A, N_a, \#(N_a), N_s\}_{k_{as}}}{A \vDash S \mid \sim (A, N_a, \#(N_a), N_s)}$$

$$\frac{A \models S \mid \sim (A, N_a, \#(N_a), N_s), A \models \#(N_a)}{A \models S \models (A, N_a, \#(N_a), N_s)}$$

$$\frac{A \vDash S \vDash N_s, \ A \vDash S \Rightarrow N_s}{A \vDash N_s}$$

$$\frac{A \vDash N_s, A \vDash \#(N_a)}{A \vDash \#(N_s)}$$

After M_3 :

$$\frac{S \vDash A \overset{k_{as}}{\longleftrightarrow} S, S \vartriangleleft \{A, N_s, \#(N_s)\}_{k_{as}}}{S \vDash A \mid \sim (A, N_s, \#(N_s))}$$
$$\frac{S \vDash A \mid \sim (A, N_s, \#(N_s)), S \vDash \#(N_s)}{S \vDash A \vDash \#(N_s)}$$

Online Key Exchange Real Protocol

$$\begin{split} &M_{1} \colon \mathsf{A} \to \mathsf{S} \colon E_{k_{as}}(\ A, B, N_{sa}) \\ &M_{2} \colon \mathsf{S} \to \mathsf{A} \colon E_{k_{as}}(\ A, B, N_{sa}, g_{b}, p_{b,}k_{b,}E_{k_{bs}}(k_{a}, N_{sb}, N_{sa}) \big) \\ &M_{3} \colon \mathsf{A} \to \mathsf{B} \colon E_{k_{b}}(\ A, B, Y_{A}, E_{k_{bs}}(k_{a}, N_{sb}, N_{sa}) \big) | |E_{k_{a}^{-}}(S(\ E_{k_{b}}\left(\ A, B, Y_{A}, E_{k_{bs}}(k_{a}, N_{sb}) \right))) \\ &M_{4} \colon \mathsf{B} \to \mathsf{A} \colon E_{k_{a}}(\ A, B, Y_{B}, N_{sa}, E_{k_{ab}}(N_{b})) | |E_{k_{b}^{-}}(S(\ E_{k_{a}}\left(\ A, B, Y_{B}, E_{k_{ab}}(N_{b}) \right)))) \\ &M_{5} \colon \mathsf{A} \to \mathsf{B} \colon E_{k_{ab}}(N_{b}) \end{split}$$

 M_x : A \rightarrow B: $E_{k_{ab}}(data, SeqNum)$

Online Key Exchange Idealized Protocol

$$\begin{split} &M_{1} \colon \mathsf{A} \to \mathsf{S} \colon \{A, B, N_{sa}\}_{k_{as}} \\ &M_{2} \colon \mathsf{S} \to \mathsf{A} \colon \left\{A, B, N_{sa}, g_{b}, p_{b}, \stackrel{k_{b}}{\to} B, \left\{\stackrel{k_{a}}{\to} A, N_{sb}, N_{sa}\right\}_{k_{bs}}, \#(N_{sa})\right\}_{k_{as}} \\ &M_{3} \colon \mathsf{A} \to \mathsf{B} \colon \left\{A, B, Y_{A}, \left\{\stackrel{k_{a}}{\to} A, N_{sb}, N_{sa}\right\}_{k_{bs}}, \#(N_{sb})\right\}_{k_{b}} \mid \left\{S\left(\left\{A, B, Y_{A}, \left\{\stackrel{k_{a}}{\to} A, N_{sb}\right\}_{k_{bs}}, \#\left(\stackrel{k_{a}}{\to} A\right)\right\}_{k_{b}}\right\}_{k_{a}} \\ &M_{4} \colon \mathsf{B} \to \mathsf{A} \colon \left\{A, B, Y_{b}, N_{sa}, \{N_{b}\}_{k_{ab}}\right\}_{k_{a}} \mid \left\{S\left(\left\{A, B, Y_{b}, \{N_{b}\}_{k_{ab}}\right\}_{k_{a}}\right\}_{k_{b}} \right\}_{k_{b}} \\ &M_{5} \colon \mathsf{A} \to \mathsf{B} \colon \{N_{b}\}_{k_{ab}} \end{split}$$

Online Key Exchange Analysis

Objectives

$$A \models A \stackrel{k_{ab}}{\longleftrightarrow} B$$
, $B \models A \stackrel{k_{ab}}{\longleftrightarrow} B$

$$A \vDash B \vDash A \stackrel{k_{ab}}{\longleftrightarrow} B$$
, $B \vDash A \vDash A \stackrel{k_{ab}}{\longleftrightarrow} B$

$$B \vDash A \vDash A \stackrel{k_{ab}}{\longleftrightarrow} E$$

Key Confirmation

Assumptions

$$A \models A \stackrel{k_{as}}{\longleftrightarrow} S$$
, $S \models A \stackrel{k_{as}}{\longleftrightarrow} S$

Symmetric keys

$$B \models B \stackrel{k_{bs}}{\leftrightarrow} S$$
, $S \models B \stackrel{k_{bs}}{\leftrightarrow} S$

Symmetric keys

Freshness of the session ids used during the authentication protocol

$$S \models \#(N_{sa}), A \models \#(N_{sa}), B \models \#(N_{sb})$$

$$S \models (g_h, p_h)$$

Diffie Hellman's parameters

$$A \models B \Rightarrow Y_B$$
, $B \models A \Rightarrow Y_A$

Authority on Y parameters

After M_1 :

$$\frac{S \vDash A \overset{k_{as}}{\longleftrightarrow} S, S \vartriangleleft \{A, B, N_{sa}\}_{k_{as}}}{S \vDash A \mid \sim (A, B, N_{sa})}$$

$$\frac{S \models A \mid \sim (A, B, N_{sa}), S \models \#(N_{sa})}{S \models A \models (A, B, N_{sa})}$$

After M_2 :

$$\frac{A \vDash A \overset{k_{as}}{\longleftrightarrow} S, \quad A \vartriangleleft \left\{A, B, N_{sa}, g_b, p_b, \overset{k_b}{\to} B, \left\{\overset{k_a}{\to} A, N_{sb}, N_{sa}\right\}_{k_{bs}}, \#(N_{sa})\right\}_{k_{as}}}{A \vDash S \mid \sim (A, B, N_{sa}, g_b, p_b, \overset{k_b}{\to} B, \left\{\overset{k_a}{\to} A, N_{sb}, N_{sa}\right\}_{k_{bs}}, \#(N_{sa}))}$$

$$\frac{A \models S \mid \sim (X), \ A \models \#(N_{Sa})}{A \models S \models (X)}$$

$$\frac{A \models S \models \stackrel{k_b}{\rightarrow} B, \ A \ trusts \ S \ on \ k_b}{A \models \stackrel{k_b}{\rightarrow} B}$$

$$\frac{A \models S \models (g_b, p_b), \ A \ trusts \ S \ on \ (g_b, p_b)}{A \models (g_b, p_b)}$$

After M_3 :

$$\frac{B \models B \overset{k_{bs}}{\longleftrightarrow} S , B \lhd \left\{\overset{k_a}{\to} A, N_{sb}, N_{sa}\right\}_{k_{bs}}}{B \models S \mid \sim \left(\overset{k_a}{\to} A, N_{sb}, N_{sa}\right)}$$

$$\frac{B \models S \mid \sim \binom{k_a}{\to} A \ , \ B \models \#(N_{sb})}{B \models S \models \binom{k_a}{\to} A \ }$$

$$B \models S \models \stackrel{k_a}{\to} A \quad B \text{ trusts } S \text{ on } B$$

$$\frac{B \models S \models \stackrel{k_a}{\rightarrow} A, \ B \ trusts \ S \ on \ k_a}{B \models \stackrel{k_a}{\rightarrow} A}$$

$$B \models \stackrel{k_a}{\to} A , B \triangleleft \left\{ A, B, Y_A, \left\{ \stackrel{k_a}{\to} A, N_{sb}, N_{sa} \right\}_{k_{bs}}, \#(N_{sb}) \right\}_{k_a^-}$$

$$B \models A \mid \sim (A, B, Y_A, \left\{ \stackrel{k_a}{\to} A, N_{sb}, N_{sa} \right\}_{k_b}, \#(N_{sb}))$$

$$\frac{B = A \mid \sim (A, B, Y_A, \left\{ \stackrel{k_a}{\to} A, N_{sb}, N_{sa} \right\}_{k_{bs}}, \#(N_{sb})), B = \#(N_{sb})}{B = A = (A, B, Y_A, \left\{ \stackrel{k_a}{\to} A, N_{sb}, N_{sa} \right\}_{k_{bs}}, \#(N_{sb}))}$$

$$\frac{B \models A \models Y_A, B \models A \Rightarrow Y_A}{B \models Y_A}$$

$$\frac{B \models Y_A, B \Rightarrow Y_B}{B \models A \stackrel{k_{ab}}{\longleftrightarrow} B}$$

After M_4 :

$$\frac{A \vDash \stackrel{k_b}{\rightarrow} B \;,\;\; A \mathrel{\vartriangleleft} \left\{A,B,Y_b,N_{sa},\{N_b\}_{k_{ab}}\right\}_{k_b^-}}{A \vDash B \mid \sim (A,B,Y_b,N_{sa},\{N_b\}_{k_{ab}})}$$

$$\frac{A \models B \mid \sim (A, B, Y_b, N_{sa}, \{N_b\}_{k_{ab}}), A \models \#(N_{sa})}{A \models B \models (A, B, Y_b, N_{sa}, \{N_b\}_{k_{ab}})}$$

$$\frac{A \models B \models Y_B, \ A \models B \Rightarrow Y_B}{A \models Y_B}$$

$$\frac{A \models Y_B, A \Rightarrow Y_A}{A \models A \stackrel{k_{ab}}{\longleftrightarrow} B}$$

$$\frac{A \vDash A \overset{k_{ab}}{\longleftrightarrow} B, \ A \vartriangleleft \{N_b\}_{k_{ab}}}{A \vDash B \mid \sim (N_b)}$$

$$\frac{A \models B \mid \sim (N_b), \ A \models \#(N_{Sa})}{A \models B \models A \stackrel{k_{ab}}{\longleftrightarrow} B}$$

After M_{5} :

$$\frac{B \vDash A \overset{k_{ab}}{\longleftrightarrow} B, \ B \vartriangleleft \{N_b\}_{k_{ab}}}{B \vDash A \mid \sim (N_b)}$$

$$\frac{B \vDash A \mid \sim (N_b), \ B \vDash \#(N_b)}{B \vDash A \vDash A \overset{k_{ab}}{\longleftrightarrow} B}$$

Offline Communication Real Protocol

$$M_1: A \rightarrow S: E_{k_{as}}(A, B, N_{sa})$$

$$M_2: S \rightarrow A: E_{k_{as}}(A, B, N_{sa}, k_b)$$

$$M_x$$
: A \rightarrow S: E_{k_b} $\Big($ data, SeqNum, $E_{k_a^-}(S(data, SeqNum))\Big) | | E_{k_a^-}(S(data, SeqNum)))$

Offline Communication Idealized Protocol

$$M_1: A \rightarrow S: \{A, B, N_{sa}\}_{k_{as}}$$

$$M_2: S \rightarrow A: \left\{A, B, N_{sa}, \stackrel{k_b}{\rightarrow} B, \#(N_{sa})\right\}_{k_{as}}$$

Offline Communication Analysis

Objective

$$A \models \stackrel{k_b}{\rightarrow} B$$

Bob's Public Key

Assumptions

$$A \models A \stackrel{k_{as}}{\longleftrightarrow} S$$
, $S \models A \stackrel{k_{as}}{\longleftrightarrow} S$

Symmetric keys

Freshness of the session id used during the authentication protocol

$$S \models \#(N_{sa}), A \models \#(N_{sa})$$

After M_1 :

$$\frac{S \vDash A \overset{k_{as}}{\longleftrightarrow} S, S \vartriangleleft \{A, B, N_{sa}\}_{k_{as}}}{S \vDash A \upharpoonright \sim (A, B, N_{sa})}$$

$$\frac{S \models A \mid \sim (A, B, N_{sa}), \ S \models \#(N_{sa})}{S \models A \models (A, B, N_{sa})}$$

After M_2 :

$$A \models A \stackrel{k_{as}}{\longleftrightarrow} S , A \triangleleft \left\{ A, B, N_{sa}, \stackrel{k_b}{\to} B, \#(N_{sa}) \right\}_{k_{as}}$$

$$A \models S \mid \sim (A, B, N_{sa}, \stackrel{k_b}{\to} B, \#(N_{sa}))$$

$$A \models S \mid \sim (X), A \models \#(N_{sa})$$

$$A \models S \models (X)$$

$$A \models S \models \stackrel{k_b}{\to} B, A trusts S on k_b$$

$$A \models \stackrel{k_b}{\to} B$$