

PUSHDOWN AUTOMATA (PDA)

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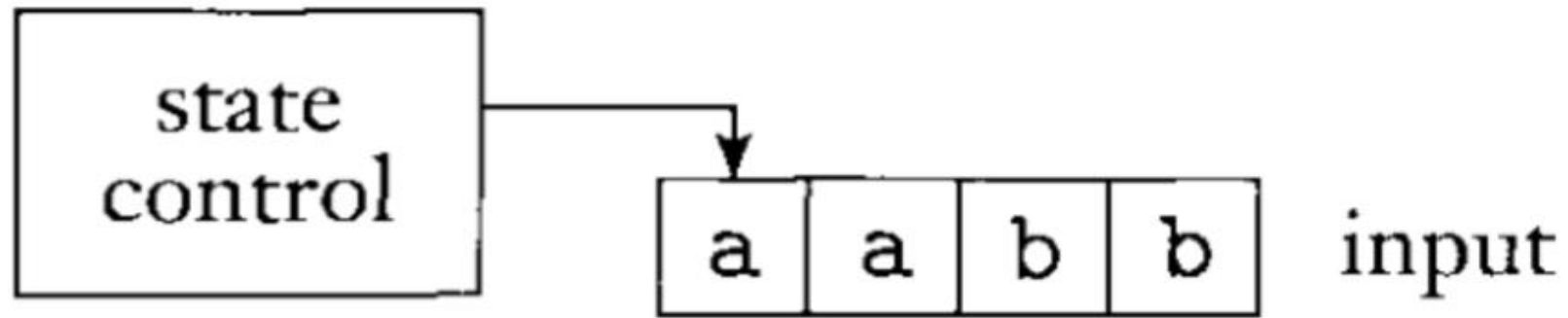


FIGURE 2.11

Schematic of a finite automaton

Pushdown Automata (Continuation...)

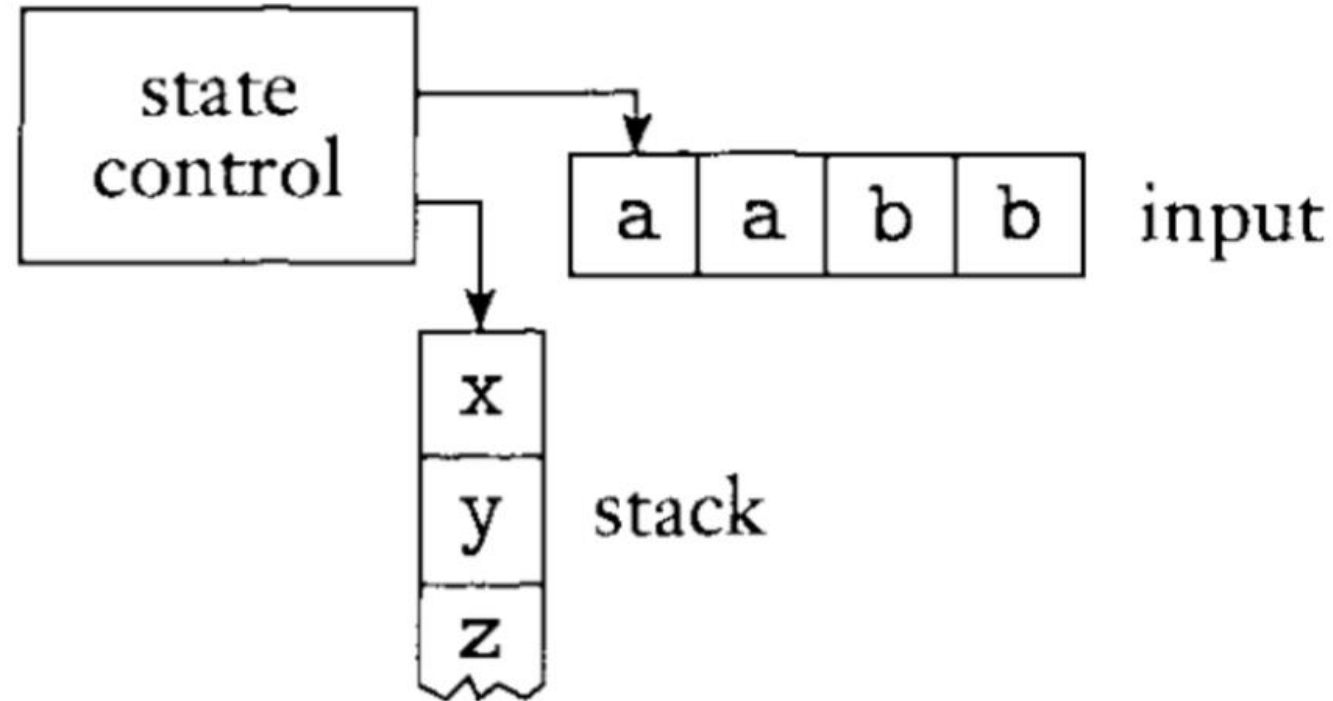


FIGURE 2.12
Schematic of a pushdown automaton

Formal Definition of a Pushdown Automaton

DEFINITION 2.13

A *pushdown automaton* is a 6-tuple $(Q, \Sigma, \Gamma, \delta, q_0, F)$, where Q , Σ , Γ , and F are all finite sets, and

1. Q is the set of states,
2. Σ is the input alphabet,
3. Γ is the stack alphabet,
4. $\delta: Q \times \Sigma_{\epsilon} \times \Gamma_{\epsilon} \longrightarrow \mathcal{P}(Q \times \Gamma_{\epsilon})$ is the transition function,
5. $q_0 \in Q$ is the start state, and
6. $F \subseteq Q$ is the set of accept states.

Examples of Pushdown Automaton

EXAMPLE 2.14

The following is the formal description of the PDA (page 110) that recognizes the language $\{0^n 1^n \mid n \geq 0\}$. Let M_1 be $(Q, \Sigma, \Gamma, \delta, q_1, F)$, where

$$Q = \{q_1, q_2, q_3, q_4\},$$

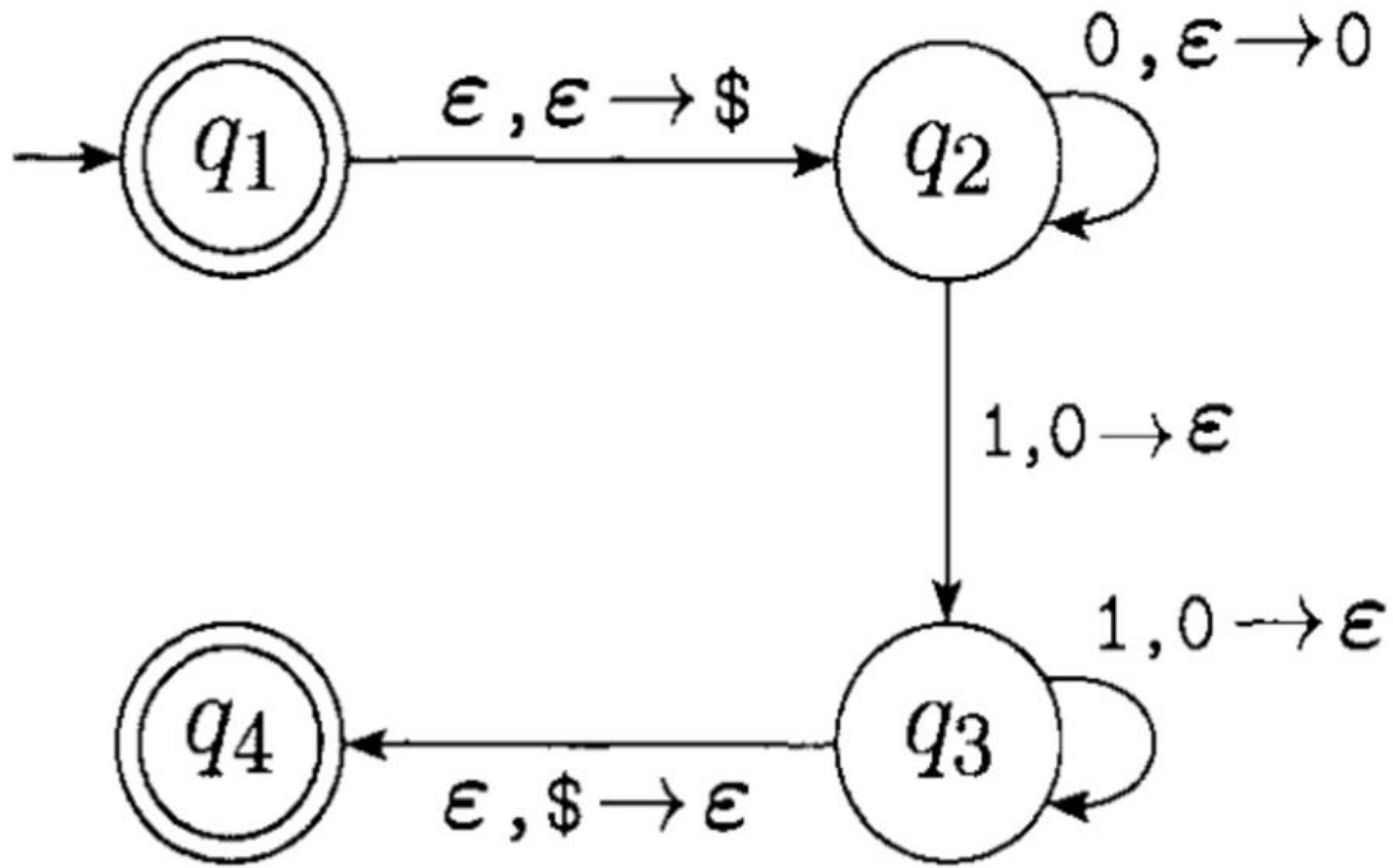
$$\Sigma = \{0,1\},$$

$$\Gamma = \{0, \$\},$$

$$F = \{q_1, q_4\}, \text{ and}$$

δ is given by the following table, wherein blank entries signify \emptyset .

Input:	0			1			ϵ		
Stack:	0	\$	ϵ	0	\$	ϵ	0	\$	ϵ
q_1									$\{(q_2, \$)\}$
q_2			$\{(q_2, 0)\}$		$\{(q_3, \epsilon)\}$				
q_3					$\{(q_3, \epsilon)\}$			$\{(q_4, \epsilon)\}$	
q_4									

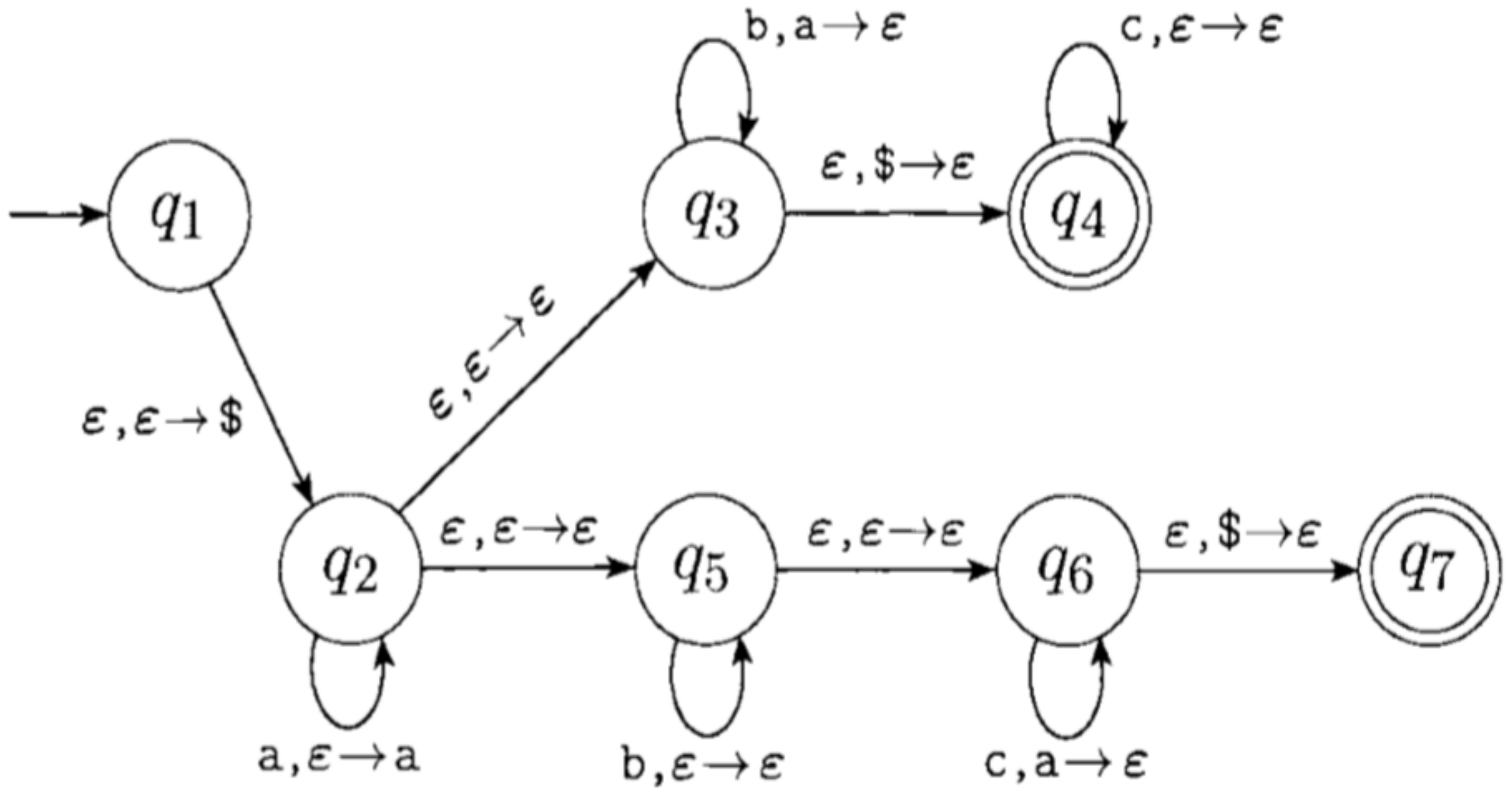


Examples of Pushdown Automaton : 2

EXAMPLE 2.16

This example illustrates a pushdown automaton that recognizes the language

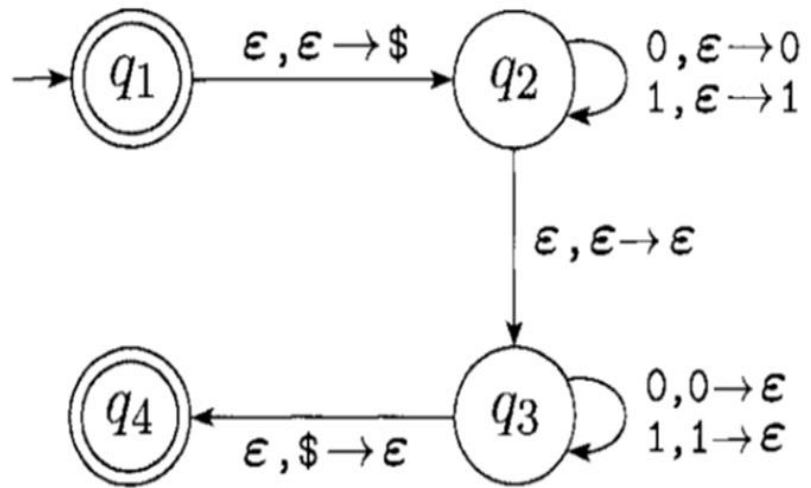
$$\{a^i b^j c^k \mid i, j, k \geq 0 \text{ and } i = j \text{ or } i = k\}.$$



Examples of Pushdown Automaton: 3

EXAMPLE 2.18

In this example we give a PDA M_3 recognizing the language $\{ww^R \mid w \in \{0,1\}^*\}$.



Practice Problems of Pushdown Automaton

- Design a NPDA for the following languages:
 - $\{0^n 1^{2n} \mid \text{where } n \geq 1\}$
 - $\{0^{2n} 1^n \mid \text{where } n \geq 1\}$
 - $\{a^n b^m c^r \mid \text{where } m, n, r \geq 0 \text{ and } r = n + m\}$
 - $\{a^n b^m c^r \mid \text{where } m, n, r \geq 0 \text{ and } r = n - m\}$