

CSE-233 : Week 1  
Summer 2020

# Introduction to Finite Automata

Reference:

Book1 Chapter 1.5,  
Book2 Chapter 1.1

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# Key Concepts

1. Alphabet  $\Sigma$
2. String,  $w$
3. Language,  $L$

# Alphabet

Alphabet of English Language  $\Sigma_{\text{Eng}} = \{a, b, c, \dots\}$

Alphabet of Bengali Language  $\Sigma_{\text{Beng}} = \{\text{অ, আ, ক}, \dots\}$

Alphabet of Arabic Language  $\Sigma_{\text{Arabic}} = \{\text{ج, ب, أ}, \dots\}$

# Strings

Alphabet of English Language  $\Sigma_{\text{Eng}} = \{a, b, c, \dots\}$

Strings =  $\{\epsilon, a, bd, cat, pq, \dots\}$

Alphabet of Bengali Language  $\Sigma_{\text{Beng}} = \{\text{অ}, \text{আ}, \text{ক}, \dots\}$

Strings =  $\{\epsilon, \text{কখ}, \text{আ}, \text{গাড়ি}, \dots\}$

Empty String is  $\epsilon$

Do you remember what is  $\Sigma^*$ ?

What is  $\Sigma_{\text{eng}}^*$ ? What is  $\Sigma_{\text{Beng}}^*$ ?

# Language

Are all strings meaningful?

Alphabet of English Language  $\Sigma_{\text{Eng}} = \{a, b, c, \dots\}$

Strings =  $\{a, bd, cat, pq, \dots\}$

Language = set of strings from  $\Sigma^*$  that are meaningful  
=  $\{cat, dog, car, \dots\}$

# Mapping Problems as Language

Languages can be used to describe problems with “yes/no” answers, for example:

- $L_1 =$  The set of all strings over  $S_1$  that contain the substring “food”
- $L_2 =$  The set of all strings over  $S_2$  that are divisible by 7  
 $= \{7, 14, 21, \dots\}$
- $L_3 =$  The set of all strings of the form  $w\#w$  where  $w$  is any string over  $\{a, b, \dots, z\}$
- $L_4 =$  The set of all strings over  $S_4$  where every ( can be matched with a subsequent )

# More Examples of Language

1. The language of all strings consisting of  $n$  0s followed by  $n$  1s ( $n \geq 0$ ):

$$\{\epsilon, 01, 0011, 000111, \dots\}$$

2. The set of strings of 0s and 1s with an equal number of each:

$$\{\epsilon, 01, 10, 0011, 0101, 1001, \dots\}$$

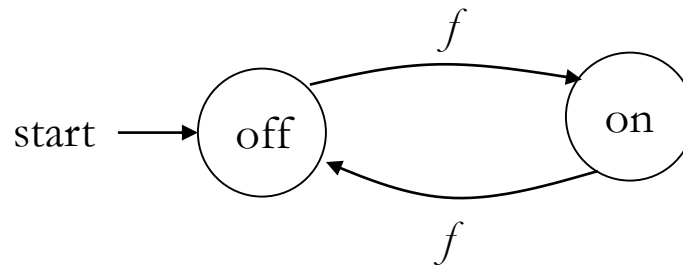
3.  $\Sigma^*$  is a language for any alphabet  $\Sigma$
4.  $\emptyset$ , the empty language, is a language over any alphabet
5.  $\{\epsilon\}$ , the language consisting of only the empty string, is also a language over any alphabet

NOTE:  $\emptyset \neq \{\epsilon\}$  since  $\emptyset$  has no strings and  $\{\epsilon\}$  has one

6.  $\{w \mid w \text{ consists of an equal number of 0 and 1}\}$
7.  $\{0^n 1^n \mid n \geq 1\}$
8.  $\{0^i 1^j \mid 0 \leq i \leq j\}$

# Finite Automata

Small program that can recognize language (Meaningful Inputs)



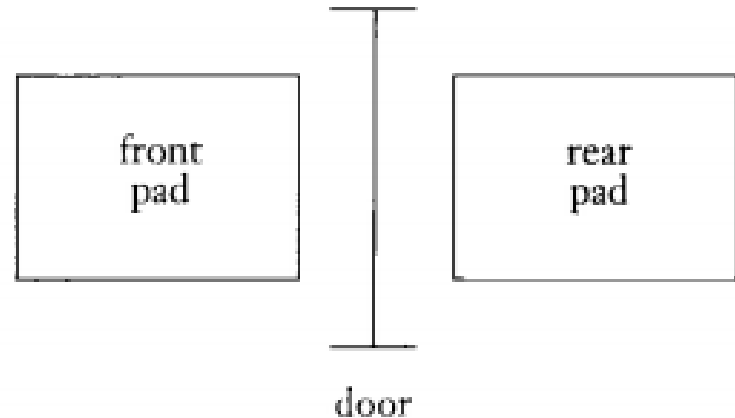
bulb is on if and only if  
there was an **odd** number  
of flips



# Deterministic Finite Automata

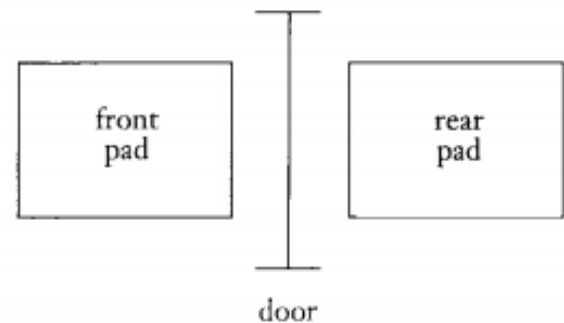
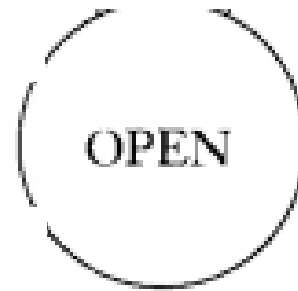
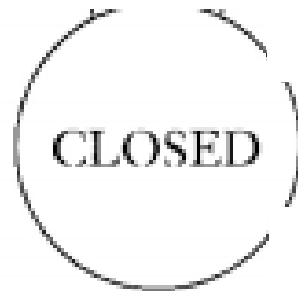
Small program that can recognize language (Meaningful Inputs) and takes one step at a time

Example: Shopping Mall Glass Door



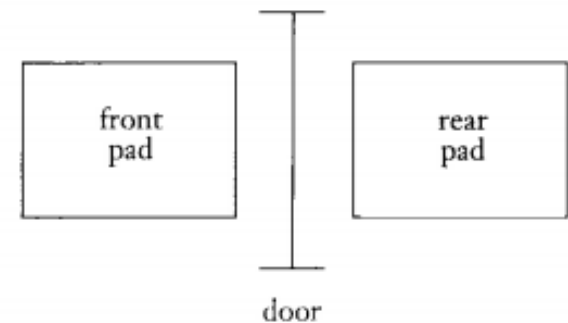
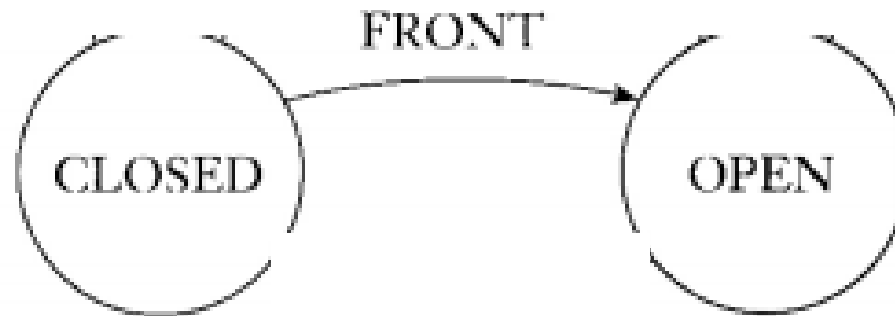
# DFA Example

Shopping Mall Glass Door



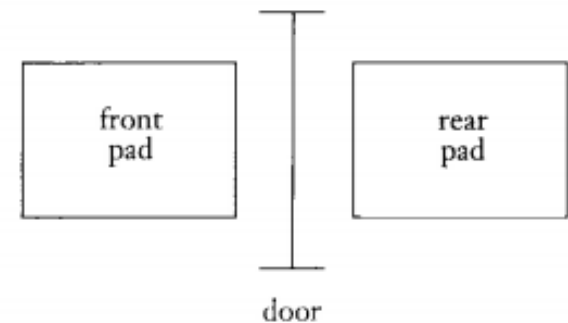
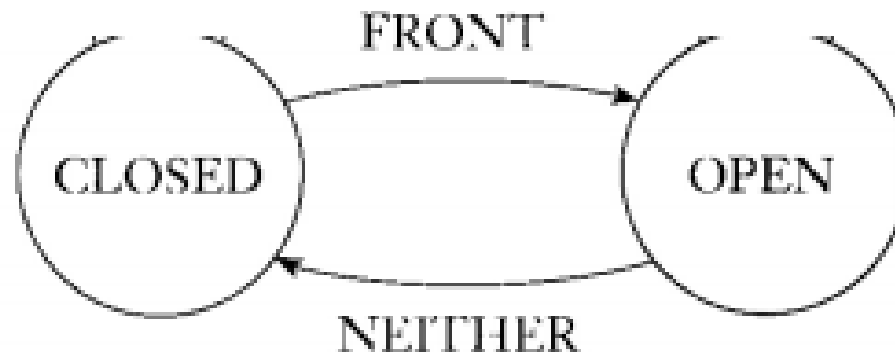
# DFA Example

Shopping Mall Glass Door



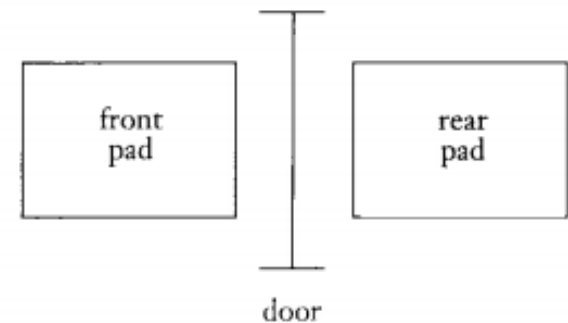
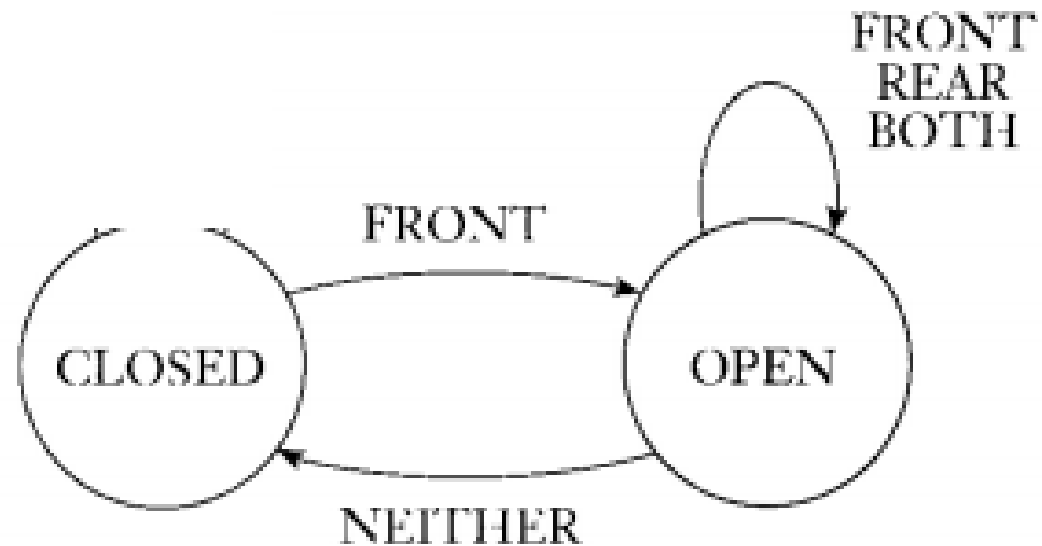
# DFA Example

Shopping Mall Glass Door



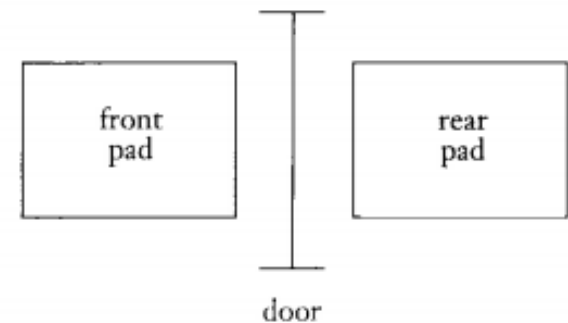
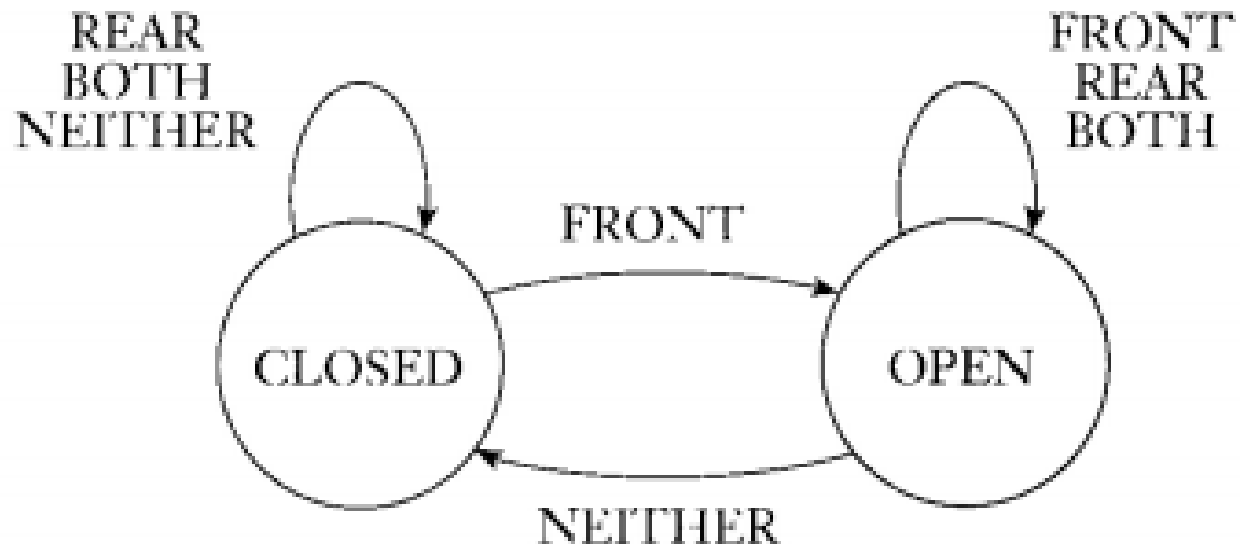
# DFA Example

Shopping Mall Glass Door



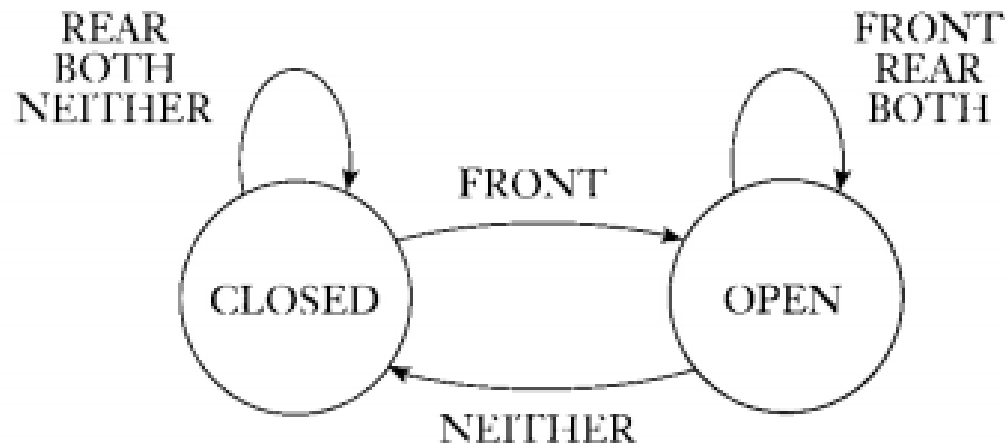
# DFA Example

Shopping Mall Glass Door



# Transition Function

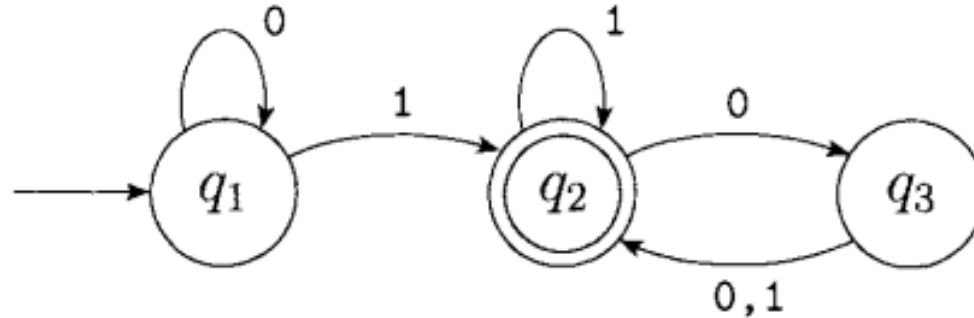
Shopping Mall Glass Door



input signal

state				
	NEITHER	FRONT	REAR	BOTH
CLOSED	CLOSED	OPEN	CLOSED	CLOSED
OPEN	CLOSED	OPEN	OPEN	OPEN

# Another Example



**FIGURE 1.4**

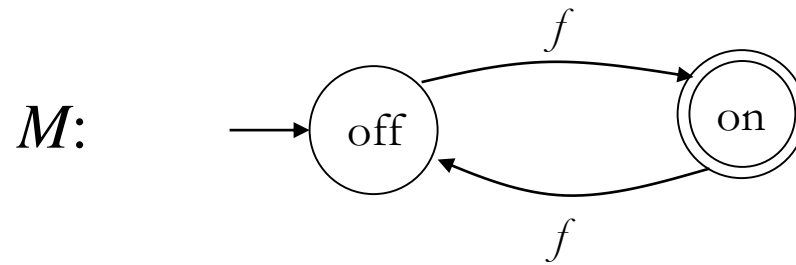
A finite automaton called  $M_1$  that has three states

1. States = ?
2. Accepted Inputs (Alphabets)?
3. Initial/Starting State = ?
4. Transition Function = ?
5. Accept State = ?



# Language of DFA

All inputs that are accepted by the FDA



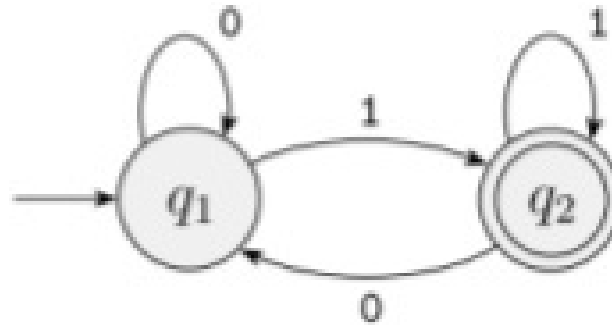
- Language of  $M$  is  $\{f, fff, fffff, \dots\} = \{f^n: n \text{ is odd}\}$

# Formal Definition of DFA

A *finite automaton* is a 5-tuple  $(Q, \Sigma, \delta, q_0, F)$ , where

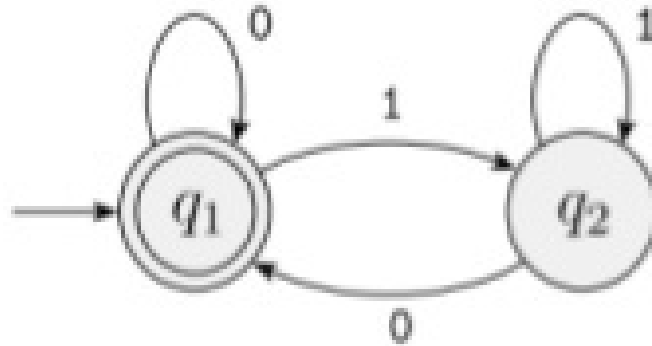
1.  $Q$  is a finite set called the *states*,
2.  $\Sigma$  is a finite set called the *alphabet*,
3.  $\delta: Q \times \Sigma \longrightarrow Q$  is the *transition function*,<sup>1</sup>
4.  $q_0 \in Q$  is the *start state*, and
5.  $F \subseteq Q$  is the *set of accept states*.<sup>2</sup>

# More Examples



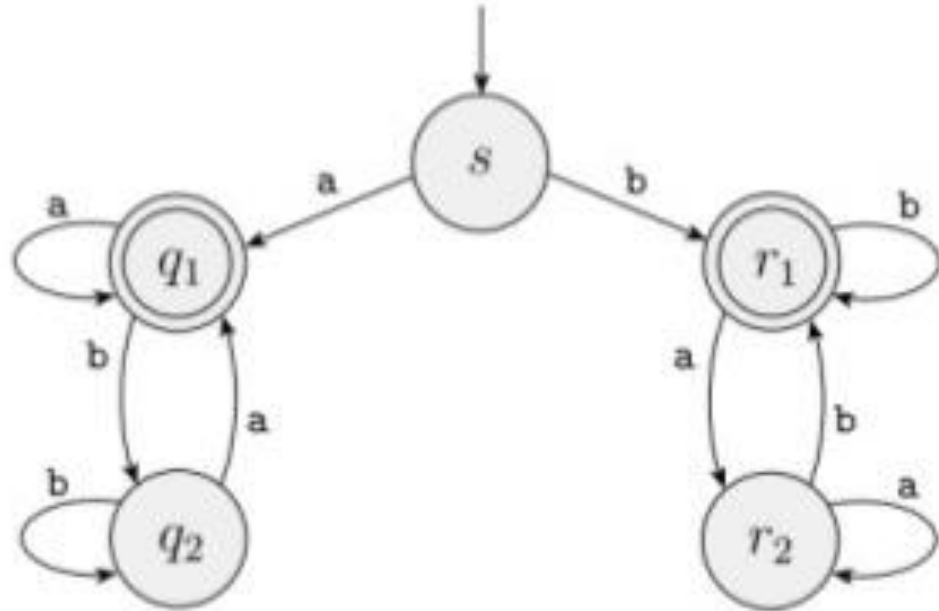
1. States = ?
2. Accepted Inputs (Alphabets)?
3. Initial/Starting State = ?
4. Transition Function = ?
5. Final State = ?
6. Language = ?

# More Examples



1. States = ?
2. Accepted Inputs (Alphabets)?
3. Initial/Starting State = ?
4. Transition Function = ?
5. Final State = ?
6. Language = ?

# More Examples



1. States = ?
2. Accepted Inputs (Alphabets)?
3. Initial/Starting State = ?
4. Transition Function = ?
5. Final State = ?
6. Language = ?