

Theory of Computation

(A study of the computational devices/ machines)

– their power/ limitation with mathematical proof

Problem: There is a finite set of input and there is a finite set of desired output

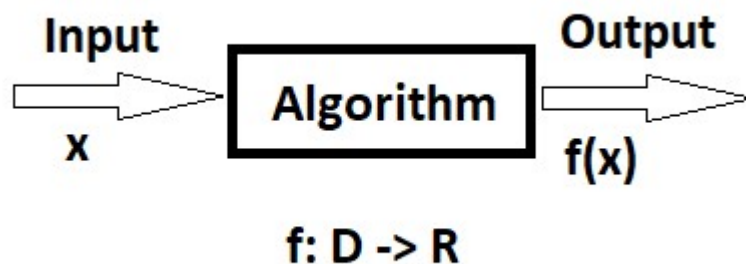
Computation: A “Program” that runs on a machine and carryout a computation

Program: “Algorithm” expressed through a programming language

Algorithm:

- is a recipe of carrying out input & output transformation?
- “finite ordered description of a process to carry out the transformation of input to output”
- it tells “how to compute a function”

Function: is a mapping of D (Domain) to some R (Range)



- in this case ‘Domain’ is all input and ‘Range’ is all possible output

Example: `is_prime.Number -> { 'yes', 'no' }`

$$is_prime(n) = \begin{cases} 'yes', & \text{if } 'n' \text{ is prime} \\ 'no', & \text{if } 'n' \text{ is not prime} \end{cases}$$

So what we have understand:

- Every 'Algorithm' computes a/ some 'Function(s)'
- 'Function' definition does not illustrate the way to achieve the right answer
- So we need an "Algorithm" to have the right results

Some situation may lead:

- for function 'f' there is an algorithm 'A' to compute 'f'
- Although having 'Function' no 'Algorithm' defines how to achieve that state

Question:

- For what 'Function' we have 'Algorithm'?
- to identify the class of 'Functions' which admit 'Algorithm' to compute them

Set Membership Problem:

Suppose S is a set

Given any 'a' to identify $a \in S$

- Assume we have a function 'f' then the $graph(f) = \{ (a,b) \mid f(a) = b \}$
- Suppose we found there is no algorithm to solve the 'Set Membership Problem' for the set $graph(f)$, then there is no algorithm to computer function 'f'

Formal Language:

A simple way (language) to present the states of the machine.

Course Outline:

- **Regular Language**

A Regular Language is a language that can be expressed with a regular expression or a deterministic or non-deterministic finite automata or state machine.

- Finite Automata
 - Simple machine, no memory
 - It uses in Model Checking, design large electronics circuit
 - Protocol design, Text Processing Application, Important component of the Compiler
- Nondeterministic Finite Automata
- Regular Expression
 - String pattern used in many systems
 - Unix command parser: `a.*b`
 - Document Type Definition (DTD) like XML:
`person(name, addr, child*)`
- Properties of Regular Languages
- Closure properties of Regular Languages

- **Context Free Language**

- Context Free Grammar (CFG)
 - Typically, the tree structure to parse/ analyze the text stream
 - Syntax of the Programming Language according to some recursive rules
- Push down automata
 - The machine model of CFG
- Decision and closure properties of CFL

- **Recursive and recursively enumerable Language**
 - Fundamental limitation our ability to compute
 - Undecidable - No program/ algorithm
 - Intractable - Solution takes exponential time with respect to input size but no proof
 - Turing Machine
 - What we can compute
- **Intractable Problem**
 - Problems that (appears to) requires exponential amount of time w.r.t. its input size

Text Book:

- Introduction to Automata Theory, Languages, and Computation by John Hopcroft, Rajee Motwani, Jeffrey D. Ullman
- Introduction to the Theory of Computation 3rd Edition by Michael Sipser