

VASANTDADA PATIL PRATISHTHAN'S COLLEGE OF ENGINEERING & VISUAL ARTS

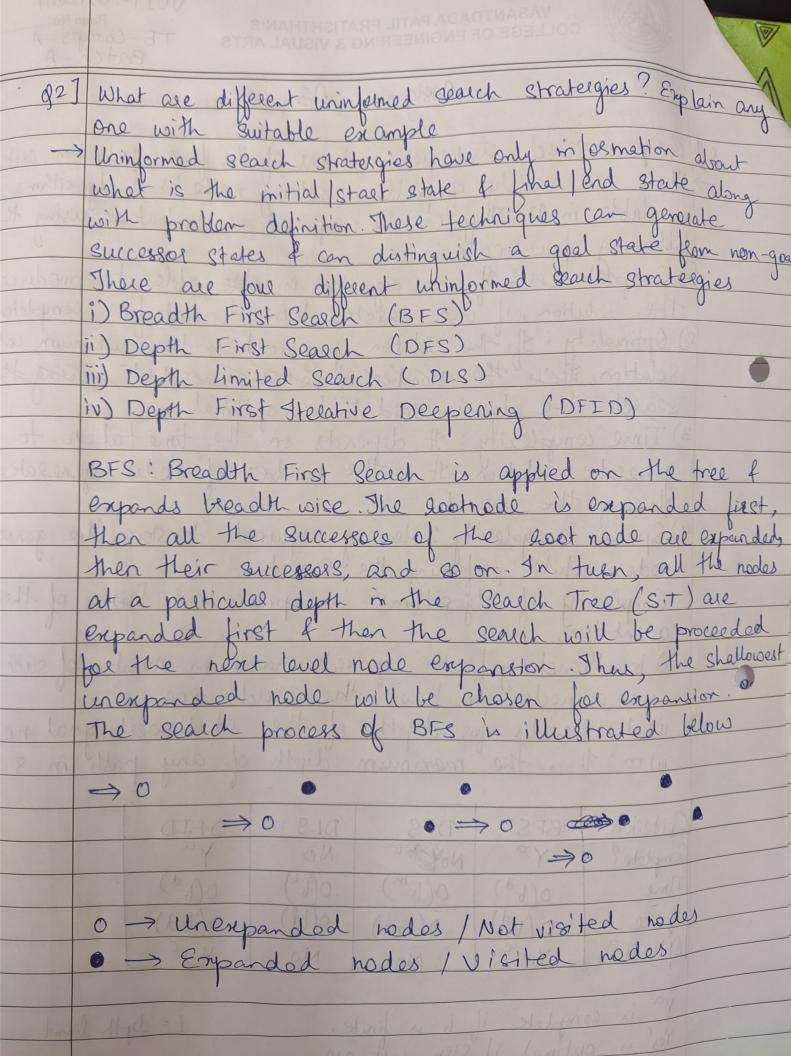
Anuag. Choèghe Page No. TE-Comps-A Batch-A

	Assignment 03
DIT	
97	These are moint algeria in teems of the four evaluation exiteria
	There are variety of problem solving methods of algorithms available in AI. They can be compared or evaluated using the following four criteria
34	available in AI they can la consider and my and lated wing the
-114	following four criteria
2000	Dempteteners: 3 an algorithm is also able to produce
	the solution, if one exists, then it satisfies completeness
	2) Optimality: 41 the solution I and and I the constant
	Solution, Steis the solution produced to the ninimum cost
	start the also is said to be optimal
	(3) Time complexity: It deleged on the the taken to
346	3) Time complexity: It depends on the time taken to generate solution. It is the number of nodes generated during the search
1	dusing the search number of nodes generated
110	4) Space complexity: Hemosy required to store the generated
134	nodes while performing the search.
	Complexity of algorithms is expressed in terms of three quantities
	and the
300	1) L. Bracking Lator vonneak marinum number of an
12	i) to: Branching factor represents maximum number of successors a node can have the search tree
100	ii) d: Stands for depth of the shallowest goal nede
10939	ii) d: Stands for depth of the shallowest goal node
	ii) m: It is the meximum depth of any path in 8.T
	C. Logina BFS DFS DL8 DFID
- 64	(TITEXION) LA
	Complete: (O(km) O(km) O(kd)
	Time Of the Court
6	Space Block Colons
	Optimal? Yes No No Yes
	va is complete if to is finite. I = depth limit

are

is complete if to is finite.

is optimal if step cost



94] What is meant by search algorithm completeness toptime Discuss for BFS skratelgies

Tompleteness: If an ordgorithm is complete, it means that
if at least one solution exists then the algorithm is
qualanteed find a solution in a finite amount of time
Optimality: If a search algorithm is optimal, then when
it finds a solution It finds the 'best' solution.

In BFS we use a FIFO queue for the finge.

Because of which the nuclei inserted nodes in the Because of which the newly inserted nodes in the floringe will automatically be placed after their parameters, children nodes, which are deeper than their parents, go to the back of the givene, fold nodes, which are shallower get expanded first. This makes BFS more complete of optimal when compared to DFS. If there are more than one sol for a given problem, then the BFS will provide sol" that Requires minimal cost of steps. g5] Consider the search problem sepresented in Figure 1, where wa' is the start node & 'f' is the goal node. Would you prefer DFS or BFS for this problem? Why?



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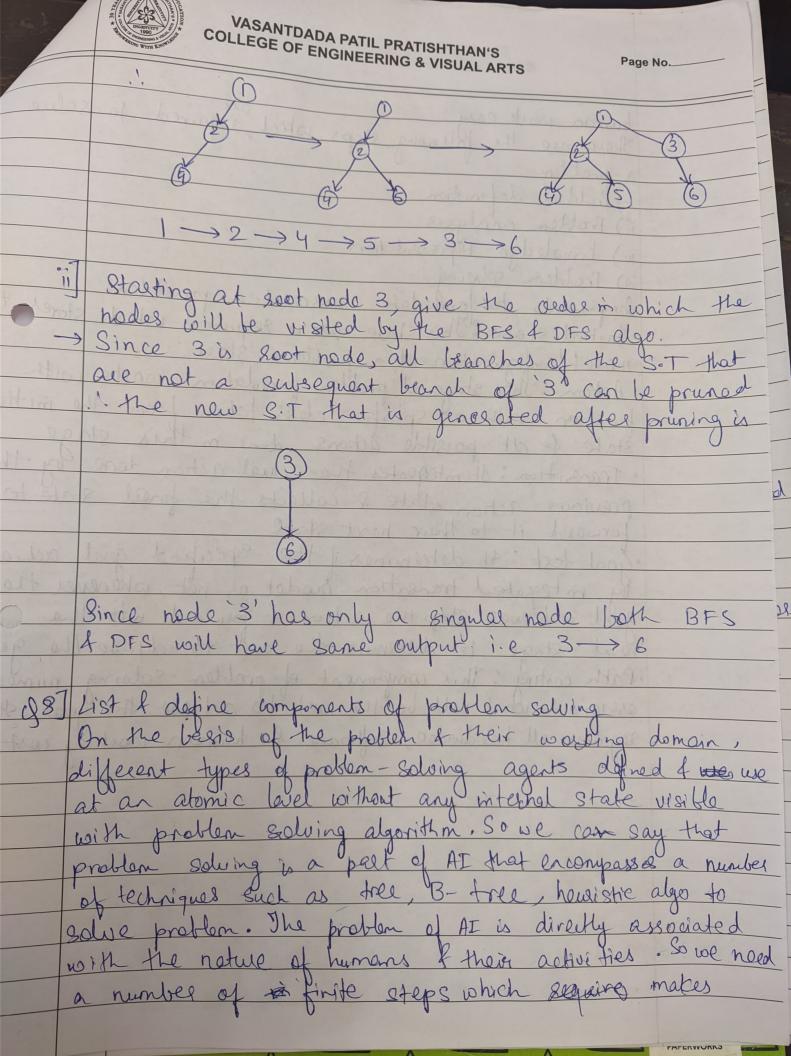
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the we would prefer BFS, because DFS could get stuck in an infinite loop abovered it it employees to the oldering of nodes of it employees to the left first it will get stuck in the loop, wholeas if employees employees a by a feel freeps oscillating between the a wheleas BFS first adds a to FIFO guere the 'abc' to FIFO queue Then path are is expanded adding a act cannot be expanded the here alaching goal state.

BFS, Breadth First Search, is a vestex based technique for finding the Shortest path in the graph. It was a gruene data structure that follows first in first out. In BFS, one vertex nade is selected at a time when it is visited & marked then its adjacent are visited & stored in the grown of is complete, provided the shallowest goal node is at some first depth. It is optimal, as it always finds the shallowest complexity is offer of space complexity is also offer). Nowever it is shower

DFS, Depth First Search, is an edge-based technique i.e.
the deepest node in the custent branch of the SIT is
expanded. It uses LIFO fainge i.e. stack data structured
The most severtly generated node, which is on the top
in the flinge is chosen first for expansion. As the node is
expanded, it is deapped from the flinge of its successor

are added . So when there are no more successors to as deepest node until that is employed. It can be implemented using iteration of necursion. It is only complete if m is finite. Since it does not quarented a solution it is not aptimal. A DFS may generate all of the o(6m) nodes in the S.T. where in is the merximum depth of any node; this can be much greated than size of State space. For a ST with branching factor of max nodes, as at a time only one branch gets employed 97] Consider the following geaph il Storting at soot node I, give the older in which the nodes will be visited by the BFS & DFS BFS: ->0-A $\bigcirc \longrightarrow \bigcirc \longrightarrow \bigcirc \longrightarrow \bigcirc \bigcirc$ B -> 0 -> 0 1. 1-2-3-3-5-6 DFS: Since it explais all children nodes that are expendable & not visited in a little stack



human work eary There are the plusing steps which required to solve a problem i) Problem definition ii) Problem analysis iv) Problem solving Components to formulate associated problem · Initial State of This state see is required which storet in the model/configuration of a problems start.

Action "This stage of problem formulation works with function with a specific class taken from the mittal state 4 all possible actions done in this stage · Transition: It integrates the actual action done by the previous action etate & collects the final state to forward it to their heart state · Goal test: It determines if the specified goal achieved by integrated transition model or not, whenever the goal achieves stop the action of forward into the next stage to determine the cost to achieve the goal · Path costing: This component of problem-solving numerical assigned what will be the cost to achieve the goal. It requires all hardware, software of human touting cost.