

I Silas Davis declare that I have completed this assignment completely and entirely on my own, without any consultation with others. I understand that any breach of the UAB Academic Honor Code may result in severe penalties.

Silas Davis

1.  $D \rightarrow DCA$ .  $CE \rightarrow CEAD$ .  $AE \rightarrow AEDC$ .

No closure determines all attributes, no key found.

Since attributes B and E are not present in the RHS of any functional dependencies, they will be apart of candidate key.

$BDE \rightarrow BDECA$ .  $BCE \rightarrow BCEAD$ .  $ABE \rightarrow ABECD$ . These will be the candidate keys.

2. Firstly we can conclude that any functional dependencies which contain E attribute will be ignored since the given relation S does not have the attribute E present. That means at  $AB \rightarrow C$  and  $BC \rightarrow D$  are the FDs that remain, and they will hold.  $AB \rightarrow C$  can translate to  $A \rightarrow C$  and  $B \rightarrow C$ .  $BC \rightarrow D$  can translate to  $B \rightarrow D$  and  $C \rightarrow D$ . So using transitive dependency,  $A \rightarrow D$  is true. FDs like  $ABC \rightarrow D$ ,  $BD \rightarrow C$ , and  $AC \rightarrow B$  will hold in S.

3. A. The candidate key will be AB, AC, BD.

$BD \rightarrow C = BDCA$ .  $AB \rightarrow D = ABDC$ .  $AC \rightarrow B = ACBD$ .

$A \rightarrow B$  will be in BCNF if A is super key, therefore a is in BCNF.

- B. The candidate key will be AC, D.

$AC \rightarrow D$ , AC super key.  $D \rightarrow A$ , D super key.  $D \rightarrow C$ , D super key.

$D \rightarrow B$ , D super key.

Therefore, relation is in BCNF.

- C. The candidate key will be A, C, D.

$C \rightarrow D = CBDA$ .  $D \rightarrow A = DACB$ .  $C \rightarrow D = CBDA$ .  $A \rightarrow C = ACDB$ .

This relation is already in BCNF.

- D. The candidate key is AD, CD, BD.

$BC \rightarrow A = BCA$ , not super key.  $AD \rightarrow C = ADCB$ , super key.  $CD \rightarrow B = CDBA$ , super key.  $BD \rightarrow C = BDCA$ . Loss less join, relation not in BCNF.

4. Since all table are not dependent on one key, which is the primary key, then it is not in 3NF.

$A \rightarrow BCDE$

A->EF

A->C

Since all table are dependent on A, the primary key, we can fetch all table data with “A”.

5. A.

```
sdavis08=> CREATE TABLE Employee(  
sdavis08(> eid INT PRIMARY KEY NOT NULL,  
sdavis08(> name TEXT NOT NULL,  
sdavis08(> age INT NOT NULL,  
sdavis08(> sex CHAR(50),  
sdavis08(> salary INT NOT NULL  
sdavis08(> );  
CREATE TABLE  
sdavis08=> INSERT INTO Employee VALUES(  
sdavis08(> 1, 'Silas1', 45, 'MALE', 45000);  
INSERT 0 1  
sdavis08=> INSERT INTO Employee VALUES(  
2, 'Silas2', 32, 'MALE', 32000);  
INSERT 0 1  
sdavis08=> INSERT INTO Employee VALUES(  
3, 'Silas3', 30, 'MALE', 30000);  
INSERT 0 1  
sdavis08=> SELECT * FROM Employee;  
  eid |  name  | age | sex | salary  
-----+-----+-----+-----+-----  
    1 | Silas1 | 45 | MALE | 45000  
    2 | Silas2 | 32 | MALE | 32000  
    3 | Silas3 | 30 | MALE | 30000  
(3 rows)
```

B.

```

sdavis08=> CREATE TABLE Dependant(
sdavis08(> did INT PRIMARY KEY NOT NULL,
sdavis08(> name TEXT NOT NULL,
sdavis08(> age INT NOT NULL,
sdavis08(> sex CHAR(50)
sdavis08(> );
CREATE TABLE
sdavis08=> INSERT INTO Dependant VALUES(1, 'Davis', 99, 'MALE');
INSERT 0 1
sdavis08=> INSERT INTO Dependant VALUES(2, 'Mavis', 50, 'MALE');
INSERT 0 1
sdavis08=> INSERT INTO Dependant VALUES(3, 'Vavis', 1, 'MALE');
INSERT 0 1
sdavis08=> SELECT * FROM Dependant;
 did | name  | age | sex
-----+-----+-----+-----
    1 | Davis |  99 | MALE
    2 | Mavis |  50 | MALE
    3 | Vavis |   1 | MALE
(3 rows)

```

C.

D. In the Employee table, the eid column is the primary key and the did column is the primary key for the Dependant table.

```

sdavis08=> SELECT * FROM Dependant;
 did | name  | age | sex
-----+-----+-----+-----
    1 | Davis |  99 | MALE
    2 | Mavis |  50 | MALE
    3 | Vavis |   1 | MALE
(3 rows)

sdavis08=> SELECT * FROM Employee;
 eid | name   | age | sex | salary
-----+-----+-----+-----+-----
    1 | Silas1 |  45 | MALE |  45000
    2 | Silas2 |  32 | MALE |  32000
    3 | Silas3 |  30 | MALE |  30000
(3 rows)

```

E.

```
sdavis08=> CREATE INDEX in_emp ON Employee(eid);  
CREATE INDEX  
sdavis08=> CREATE INDEX in_emp2 ON Employee(eid, age);  
CREATE INDEX
```

F.

```
sdavis08=> CREATE TABLE Department(  
sdavis08(> depid INT PRIMARY KEY NOT NULL,  
sdavis08(> depname CHAR(50) NOT NULL,  
sdavis08(> eid INT references EMPLOYEE(eid)  
sdavis08(> );  
CREATE TABLE  
sdavis08=> SELECT * FROM Department;  
  depid | depname | eid  
-----+-----+-----  
(0 rows)
```