1. Introduction to the judge environment

2. Partial sums/precomputation trick

Sample Problem: Even Pairs

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- lacktriangleright first line: a positive integer n
- second line: a sequence $x_0, \ldots, x_{n-1} \in \{0, 1\}$

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- lacktriangle first line: a positive integer n
- second line: a sequence $x_0, \ldots, x_{n-1} \in \{0, 1\}$

Output: a single line containing the number of pairs $0 \leq i \leq j < n$ such that

$$x_i + \cdots + x_j$$

is even.

- (1) for all pairs $i \leq j$, compute the sum $x_i + \cdots + x_j$
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- this means that our algorithm is too slow
- we have three nested loops: two for going over all pairs $i \leq j$, and one for summing up the $x_i + \cdots + x_j$
- running time is $\Theta(n^3)$.
- this type of analysis is very important in this course.

Second approach

Observation: if we know the parity of the sum

$$x_i + \cdots + x_j$$

then based on the parity of x_{j+1} we also know the parity of

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Running time: $\Theta(n^2)$

Third approach

Observation:

$$x_i + \dots + x_j = \sum_{a=1}^{j} x_a - \sum_{b=1}^{i-1} x_b$$

= $S_j - S_{i-1}$

Third approach

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- (1) calculate partial sums $S_i = \sum_{a=1}^i x_a$ in one iteration
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Fourth approach

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- (1) calculate partial sums $S_i = \sum_{a=1}^i x_a$ in one iteration
- (2) for every j: $E_j = \#$ of S_i $(i \le j)$ that are even
- (3) for every $j: O_j = \#$ of S_i $(i \leq j)$ that are odd

Fourth approach

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- (4) if S_j is even: increase the counter by $E_{j-1}+1$ if S_j is odd: increase the counter by O_{j-1}

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Running time: $\Theta(n)$

Even pairs - conclusion

Tricks/techniques: Partial sums/Precomputing

- Precomputing partial sums allows computing the sum of the elements in an interval in constant time.
- More generally, precomputing certain values can speed up the running time of an algorithm.

Judge Results

Besides correct, timelimit and wrong-answer, the judge can give the following results.

```
assertion-failure SIGABRT: memory screwup or assertion failure segmentation-fault SIGSEGV: memory screwup run-error nonzero exit status forbidden bad syscall or other safety
```

Forum

The forum is your main tool for discussing ideas and getting help. Use it.

Of course, you will only learn if you first try to solve the problems on your own.

Apply spoiler warnings

Example

SPOILER <<<

Set this text to have a white foreground. It will then be invisible unless marked. The <<< . . . >>> exploit a bug in the email plugin to also remove the text in plain-text email.

>>>

- Apply spoiler warnings
- 2 Describe the problem, not your guesses or summaries

Example

Bad When I compile, it tells me it cannot find it.

Good When I run g++ -o foo foo.cpp, I get bash: \$'g++\302\240-o': command not found

- Apply spoiler warnings
- Describe the problem, not your guesses or summaries
- 3 Code: describe what fails and what you expect instead

Example

Bad The code below doesn't work. Help?

Good I am trying to solve problem 1. I tried strategy <u>blah blah</u>. My code is below. For some reason, when running it on the provided testcase it emits no solution instead of 1. What am I doing wrong?

- Apply spoiler warnings
- Describe the problem, not your guesses or summaries
- 3 Code: describe what fails and what you expect instead
- 4 Code: post minimal examples

Example

Bad When I call .foo() on a vector, it segfaults. Bug!

Good I am trying to <u>blah</u>. The code is below. I get a segfault in the line that calls .foo(), but if I remove that line the program continues. What am I doing wrong?

- Apply spoiler warnings
- 2 Describe the problem, not your guesses or summaries
- 3 Code: describe what fails and what you expect instead
- 4 Code: post minimal examples
- 5 Don't rush to claim that you have found a bug

Further reading: How To Ask Questions The Smart Way

Learning C++ is beyond the scope of this course.

- True beginners should probably read a book
 - Andrew Koenig and Barbara E. Moo: Accelerated C++, Addison-Wesley, 2000
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- Useful in any case: C++ FAQ Lite (link)

STL is part of the C++ standard library. It is important that you know

- how and when to use the classes std::vector, std::priority_queue, std::set and std::map,
- how to do I/O using <iostream>, and
- how to use the sort function from the <algorithm> header file.

The most important things to remember:

You can find all that you need at our website: http://www.cadmo.ethz.ch/education/lectures/HS16/ algolab/index.html

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- Today, we will publish the first week's exercises on the webpage.
- Next Monday at 17:00, we have the first problem of the week.