

Worksheet 6 (Solved)

HoTTEST Summer School 2022

The HoTTEST TAs 25 July 2022

1 (*)

Prove that $\neg isContr(\emptyset)$.

Recall that, for any type A,

$$isContr(A) \doteq \sum_{c:A} \prod_{a:A} c = a.$$

So we have the projection map

$$\sum_{c:\emptyset} \left(\prod_{a:\emptyset} c = a \right) \xrightarrow{\mathbf{pr}_1} \emptyset$$

Which is a map $\mathsf{isContr}(\emptyset) \to \emptyset$.

2 (**)

Recall the observational equality of natural numbers $Eq-\mathbb{N}$: $\mathbb{N} \to \mathbb{N} \to \mathcal{U}$

Eq-
$$\mathbb{N} \ 0 \qquad \qquad \dot{=} \ \mathbb{1}$$

$$\mathsf{Eq}\text{-}\mathbb{N}\;(\mathsf{suc}\;m)\;0\qquad \ \, \dot{=}\;\emptyset$$

Eq-
$$\mathbb{N} \ 0 \qquad (\operatorname{suc} \ n) \doteq \emptyset$$

$$\mathsf{Eq}\text{-}\mathbb{N}\ (\mathsf{suc}\ m)\ (\mathsf{suc}\ n) \doteq \mathsf{Eq}\text{-}\mathbb{N}\ m\ n$$

Prove that, for every $n : \mathbb{N}$,

$$\mathsf{Eq}\text{-}\mathbb{N}\; n\; (\mathsf{suc}\; n) = \emptyset$$

Proceed by induction on n. For $n \doteq 0$, we have that Eq- \mathbb{N} 0 1 is definitionally equal to \emptyset , so we use refl_{\emptyset}. Then, for some n, assuming

$$\mathbf{ih} \colon (\mathbf{Eq}\text{-}\mathbb{N}\ n\ (\mathbf{suc}\ n)) = \emptyset$$

we have that Eq- $\mathbb N$ (suc n) (suc(suc n)) = \emptyset by the recursive definition of Eq- $\mathbb N$.

In Lecture 4, we did most of the proof that

$$\mathsf{Eq}\text{-}\mathbb{N}\ m\ n \quad \simeq \quad m =_{\mathbb{N}} n.$$

Use this (and the fact proved above) to prove that \neg (isContr \mathbb{N}).

Suppose

$$(c,\varphi)$$
: isContr \mathbb{N} .

So $c: \mathbb{N}$ and $\varphi: \prod_{n:\mathbb{N}} c = n$. So $\varphi(\operatorname{suc} c): c = (\operatorname{suc} c)$. But we have

$$c = (\operatorname{suc} c) \simeq \operatorname{Eq-N} c (\operatorname{suc} c)$$

We showed that the right-hand side is \emptyset , and thus isContr(\mathbb{N}) implies \emptyset .

$3 \quad (\star \star \star)$

Show that if A is contractible, then for any x, y : A, the identity type x = y is also contractible.

Let c:A be the center of contraction. It will suffice to show that $(x=y)\simeq \mathbb{1}$. For the map $f:(x=y)\to \mathbb{1}$, we can just use the constant \star map. For the other direction, $g:\mathbb{1}\to (x=y)$, we proceed by singleton induction and can therefore take $x \doteq c \doteq y$ and put

$$g(\star) \doteq \mathbf{refl}_c$$
.

To see that g(f(p)) = p for any p: x = y, we can proceed by path induction and assume $x \doteq y$ and $p \doteq \mathbf{refl}_x$. By singleton induction again, we can assume $x \doteq c$. So $p \doteq \mathbf{refl}_c$, which is the same as g(f(p)). The other direction is simpler: for any $t: \mathbb{1}$, we prove that f(g(t)) = t by singleton induction, i.e. assuming $t \doteq \star$. But $f(g(t)) \doteq \star$ by definition, so $\mathbf{refl}_\star: f(g(t)) = t$.

4 $(\star\star\star)$

Recall the first projection function

$$\operatorname{pr}_1 \quad : \quad \sum_{x:A} B(x) \to A$$

Show that $\operatorname{\mathsf{pr}}_1$ is an equivalence iff each B(a) is contractible. Hint: Use the results about identity types of Σ types we proved in a previous lecture.

First the 'only if' direction: if pr_1 is an equivalence, that means it's a contractible map, i.e.

$$isContr(fib_{pr_1}(a))$$

for each a:A. Then, given any a:A and b,b':B(a), we have

$$((a,b),\mathsf{refl}_a): \mathsf{fib}_{\mathsf{pr}_1}(a) \qquad \text{ and } \qquad ((a,b'),\mathsf{refl}_a): \mathsf{fib}_{\mathsf{pr}_1}(a).$$

Since $fib_{pr_1}(a)$ is contractible, we have

$$((a,b), \mathbf{refl}_a) = ((a,b'), \mathbf{refl}_a).$$

Using our result characterizing the identity type of Σ types (Lecture 5), this gives us a proof of

$$(a,b) = (a,b')$$

and, applying it again,

$$b = b'$$
.

Thus we have that B(a) is contractible.

In the other direction, suppose each B(a) is contractible with center c_a . Then define an inverse $q: A \to \sum_{x:A} B(x)$ by

$$q(a) \doteq (a, c_a).$$

Then $\operatorname{pr}_1(q(a)) \doteq \operatorname{pr}_1(a, c_a) \doteq a$, so q is a section of pr_1 . Furthermore, for any $(a,b): \sum_{x:A} B(x)$, we have a proof that $b=c_a$, so, again using that identities of pairs are equivalent to pairs of identities, we get

$$(a, c_a) = (a, b).$$

The left-hand side is exactly $q(\mathbf{pr}_1(a,b))$, as desired.

Show that for any a:A, the map

$$\lambda((x,y),p).\mathsf{tr}_B(p,y)\colon \mathsf{fib}_{\mathsf{pr}_1}(a)\to B(a)$$

is an equivalence.

Call this map k. Now, pick some b:B(a) and observe that $((a,b), \mathbf{refl}_a): \mathbf{fib}_{\mathbf{pr}_1}(a)$ and that

$$k((a,b), refl_a) \doteq tr_B(refl_a, b) \doteq b$$

So now pick some ((x,y),p): $fib_{pr_1}(b)$ and

$$q: \mathbf{tr}_B(p, y) = b.$$

Recall (again) from Lecture 5 that

$$(x,y)=(a,b)$$
 \simeq $\sum_{p:x=a} \mathbf{tr}_B(p,y)=b$

So, since we have p: x = a and $q: \operatorname{tr}_B(p, y) = b$, we get a proof r: (x, y) = (a, b). Finally, by the observations we made in lecture, to prove

$$((x,y),p)=((a,b),\mathbf{refl}_a)$$

we just need to have our proof r:(x,y)=(a,b) and then show

$$p = \mathsf{ap}_{\mathsf{pr}_1} r \cdot \mathsf{refl}_a.$$

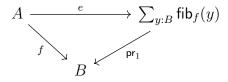
But we can check that $ap_{pr_1}r$ gives us p, and $p \cdot refl_a = p$.

5 (**)

Construct for any map $f: A \to B$ an equivalence

$$e:A\simeq \sum_{y:B}\operatorname{fib}_f(y)$$

with a homotopy $H: f \sim \mathsf{pr}_1 \circ e$ witnessing that the triangle



commutes.

The natural definition for e is

$$e(a) \doteq (f(a), (a, \mathbf{refl}_{f(a)}))$$

The inverse for e is given by

$$e^{-1}(y,(x,p)) \doteq x.$$

Then $e^{-1}(e(a)) \doteq a$, so we have $e^{-1} \circ e \sim id_A$. For the other composition, take any x : A. To prove that for any y : B and p : f(x) = y

$$(y,(x,p)) = (f(x),(x,refl_{f(x)}))$$

we use path induction. Therefore, we can take $y \doteq f(x)$ and $p \doteq \text{refl}_{f(a)}$, and then we have that

$$e(e^{-1}(y,(x,p))) = (y,(x,p))$$

as desired.

The homotopy H is given by

$$H(x) \stackrel{.}{=} \operatorname{refl}_{f(x)} : f(x) = \operatorname{pr}_{1}(e(x))$$