

**Homework 5**

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**1. Introduction**

In this paper, I will introduce a simple dynamic memory allocation library and its new version 2.0. In version 1.0, basic memory allocation and deallocation functions are implemented. Furthermore, some utility function `print_sm_containers` is implemented. At the new version of `smalloc`, new features are available. First, the first feature is a change from first fit to best fit. The old `smalloc` was implemented with the first fit algorithm. In contrast, the new `smalloc` function now uses the best fit algorithm. Also, the new `sfree` function does the extra work of merging unused nearby memory. New functions have been added. The `print_mem_uses` function tells you about the memory usage. The `srealloc` function is a function that extends existing memory. The `srealloc` function extends the existing memory and reallocates it. The `sshink` function is a function that compresses memory by lowering the upper limit of the heap if there is unused memory at the end of the heap. I will introduce some tests to show working properly and discuss with further works at the version 3.0.

**2. Approach**

First of all, to implement best fit in `smalloc`, I have slightly modified the function to find the smallest allocable memory. Additional merging unused neighbor memory have added at `sfree` routine. Additional merging unused neighbor memory have added at `sfree` routine. If there is unused memory on the left, the function tries to merge it to the left. Conversely, if it is right, the function tries to merge it to the right. The `print_mem_uses` prints three memory allocation data: sum of all allocated memory, sum of busy memory size and sum of unused memory size. Iteration and adding each type by type are very simple solution to acquire these memory data. The `srealloc` is implemented in this way. First it finds current memory node. After that, it checks next node. If next node is unused and enough size to allocate, it splits next node and merges next node to current node. Otherwise, it calls `smalloc` to find another available memory and calls `memcpy` to copy existing data. Finally, `sshink` function is implemented to adjust top of the heap point to start of the unused node. While iterating in reverse direction, iteration stops at first busy node appearing. After that, `brk` function is called to resize heap size with the node address.

**3. Evaluation**

In this section, I will show the test #4 to compare first fit and best fit.

Best fit (Left)		First fit (Right)	
<code>smalloc(1000):0x1140a04</code>	<pre> 0:0x113f020: Busy: 2000:00 00 00 00 00 00 1:0x113f810: Unused: 2032:00 00 00 00 00 00 2:0x1140020: Busy: 2500:00 00 00 00 00 00 3:0x1140a04: Busy: 1000:00 00 00 00 00 00 4:0x1140e0c: Unused: 500:00 00 00 00 00 00 </pre>	<code>smalloc(1000):0x82b810</code>	<pre> 0:0x82b020: Busy: 2000:00 00 00 00 00 00 1:0x82b810: Busy: 1000:00 00 00 00 00 00 2:0x82bc18: Unused: 1000:00 00 00 00 00 00 3:0x82c020: Busy: 2500:00 00 00 00 00 00 4:0x82ca04: Unused: 1532:00 00 00 00 00 00 </pre>
<code>smalloc(500):0x1140e0c</code>	<pre> 0:0x113f020: Busy: 2000:00 00 00 00 00 00 1:0x113f810: Unused: 2032:00 00 00 00 00 00 2:0x1140020: Busy: 2500:00 00 00 00 00 00 3:0x1140a04: Busy: 1000:00 00 00 00 00 00 4:0x1140e0c: Busy: 500:00 00 00 00 00 00 </pre>	<code>smalloc(500):0x82bc18</code>	<pre> 0:0x82b020: Busy: 2000:00 00 00 00 00 00 1:0x82b810: Busy: 1000:00 00 00 00 00 00 2:0x82bc18: Busy: 500:00 00 00 00 00 00 3:0x82bc2c: Unused: 468:00 00 00 00 00 00 4:0x82c020: Busy: 2500:00 00 00 00 00 00 5:0x82ca04: Unused: 1532:00 00 00 00 00 00 </pre>
<code>smalloc(2000):0x1140e0c</code>	<pre> 0:0x113f020: Busy: 2000:00 00 00 00 00 00 1:0x113f810: Unused: 2032:00 00 00 00 00 00 2:0x1140020: Busy: 2500:00 00 00 00 00 00 3:0x1140a04: Busy: 1000:00 00 00 00 00 00 4:0x1140e0c: Busy: 500:00 00 00 00 00 00 5:0x1162020: Busy: 2000:00 00 00 00 00 00 6:0x1162810: Unused: 2032:00 00 00 00 00 00 </pre>	<code>smalloc(2000):0x82bc18</code>	<pre> 0:0x82b020: Busy: 2000:00 00 00 00 00 00 1:0x82b810: Busy: 1000:00 00 00 00 00 00 2:0x82bc18: Busy: 500:00 00 00 00 00 00 3:0x82bc2c: Unused: 468:00 00 00 00 00 00 4:0x82c020: Busy: 2500:00 00 00 00 00 00 5:0x82ca04: Unused: 1532:00 00 00 00 00 00 6:0x82d020: Busy: 2000:00 00 00 00 00 00 7:0x82d810: Unused: 2032:00 00 00 00 00 00 </pre>

Figure 1. Test4.c (Left: Best fit, Right: First fit)

As you can see in Figure 1, Best fit memory fragmentation is less than first fit one. Memory allocation bytes sequence is following: <2000, 2500, 1000, 500, 2000>

When allocating 1000 bytes, best fit allocated 1000 bytes in index 3 node. But first fit allocated 1000 bytes in index 1 node. Allocating 500 bytes and last 2000 bytes also show difference.

**4. Discussion**

In this section, I will discuss further version of this `smalloc` library 3.0. First thing is to implement worst fit algorithm. Implementing the worst-fit algorithm can give users one way to manage heap memory more efficiently. Furthermore, making an API to switch the algorithms in the runtime can provide more powerful options to user to manage heap memory. And it will be good to implement heap memory dump feature. More future version of `smalloc` can provide simple implementation of reference

**5. Conclusion**

I discussed `smalloc` library 2.0 version implementations and also showed the future works. Best fit algorithm and printing mem uses is simply done with iterating all nodes in sequence. And `sfree` is done with merging adjacency memories. I showed you the test #4 to compare algorithms and talked about future works.