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# 1 Combinatorial optimization

## 1.1 Dinic's

```

struct Dinic {
    struct Edge {
        int u, v;
        long long cap, flow;
    };
    Edge() {}
    Edge(int u, int v, long long cap): u(u), v(v), cap(
        cap), flow(0) {}
};
int N;
vector<Edge> E;
vector<vector<int>> g;
vector<int> d, pt;

Dinic(int N): N(N), E(0), g(N), d(N), pt(N) {}

void AddEdge(int u, int v, long long cap) {
    if (u != v) {
        E.emplace_back(Edge(u, v, cap));
        g[u].emplace_back(E.size() - 1);
        E.emplace_back(Edge(v, u, 0));
        g[v].emplace_back(E.size() - 1);
    }
}

bool BFS(int S, int T) {
    queue<int> q({S});
    fill(d.begin(), d.end(), N + 1);
    d[S] = 0;
    while(!q.empty()) {
        int u = q.front(); q.pop();
        if (u == T) break;
        for (int k: g[u]) {
            Edge &e = E[k];
            if (e.flow < e.cap && d[e.v] > d[e.u] + 1) {
                d[e.v] = d[e.u] + 1;
                q.emplace(e.v);
            }
        }
    }
    return d[T] != N + 1;
}

long long DFS(int u, int T, long long flow = -1) {
    if (u == T || flow == 0) return flow;

```

```

for (int &i = pt[u]; i < g[u].size(); ++i) {
    Edge &e = E[g[u][i]];
    Edge &oe = E[g[u][i]^1];
    if (d[e.v] == d[e.u] + 1) {
        long long amt = e.cap - e.flow;
        if (flow != -1 && amt > flow) amt = flow;
        if (long long pushed = DFS(e.v, T, amt)) {
            e.flow += pushed;
            oe.flow -= pushed;
            return pushed;
        }
    }
}
return 0;
}
long long MaxFlow(int S, int T) {
    long long total = 0;
    while (BFS(S, T)) {
        fill(pt.begin(), pt.end(), 0);
        while (long long flow = DFS(S, T))
            total += flow;
    }
    return total;
}
};

```

## 1.2 Min-cost Circulation

*// Runs in  $O(\text{max\_flow} * \log(V * \text{max\_edge\_cost})) = O(V^3 * \log(V * C))$  Really fast in practice,  $3e4$  edges are fine. Operates on integers, // costs are multiplied by  $N!$*

```

template<typename flow_t = int, typename cost_t = int>
struct mcSFlow {
    struct Edge {
        cost_t c;
        flow_t f;
        int to, rev;
        Edge(int _to, cost_t _c, flow_t _f, int _rev):c(_c),
            f(_f), to(_to), rev(_rev){}
    };
    static constexpr cost_t INFCOST = numeric_limits<
        cost_t>::max()/2;
    cost_t eps;
    int N, S, T;
    vector<vector<Edge>> G;
    vector<unsigned int> isq, cur;
    vector<flow_t> ex;
    vector<cost_t> h;
    mcSFlow(int _N, int _S, int _T):eps(0), N(_N), S(_S),
        T(_T), G(_N){}
    void add_edge(int a, int b, cost_t cost, flow_t cap) {
        assert(cap >= 0);
        assert(a >= 0 && a < N && b >= 0 && b < N);
        if (a == b) {assert(cost >= 0); return;}
        cost *= N;
        eps = max(eps, abs(cost));
        G[a].emplace_back(b, cost, cap, G[b].size());
        G[b].emplace_back(a, -cost, 0, G[a].size()-1);
    }
    void add_flow(Edge& e, flow_t f) {

```

```

        Edge &back = G[e.to][e.rev];
        if (!ex[e.to] && f)
            hs[h[e.to]].push_back(e.to);
        e.f -= f; ex[e.to] += f;
        back.f += f; ex[back.to] -= f;
    }
    vector<vector<int>> hs;
    vector<int> co;
    flow_t max_flow() {
        ex.assign(N, 0);
        h.assign(N, 0); hs.resize(2*N);
        co.assign(2*N, 0); cur.assign(N, 0);
        h[S] = N;
        ex[T] = 1;
        co[0] = N-1;
        for(auto &e:G[S]) add_flow(e, e.f);
        if(hs[0].size())
            for (int hi = 0; hi >= 0; ) {
                int u = hs[hi].back();
                hs[hi].pop_back();
                while (ex[u] > 0) { // discharge u
                    if (cur[u] == G[u].size()) {
                        h[u] = 1e9;
                        for(unsigned int i=0; i<G[u].size(); ++i) {
                            auto &e = G[u][i];
                            if (e.f && h[u] > h[e.to]+1) {
                                h[u] = h[e.to]+1, cur[u] = i;
                            }
                        }
                    }
                    if (++co[h[u]], !--co[hi] && hi < N)
                        for(int i=0; i<N; ++i)
                            if (hi < h[i] && h[i] < N) {
                                --co[h[i]];
                                h[i] = N + 1;
                            }
                    hi = h[u];
                }
                else if (G[u][cur[u]].f && h[u] == h[G[u][cur[u].to].to]+1)
                    add_flow(G[u][cur[u]], min(ex[u], G[u][cur[u].to].f));
                else ++cur[u];
            }
        while (hi >= 0 && hs[hi].empty()) --hi;
        return -ex[S];
    }
    void push(Edge &e, flow_t amt) {
        if(e.f < amt) amt=e.f;
        e.f -= amt; ex[e.to] += amt;
        G[e.to][e.rev].f += amt; ex[G[e.to][e.rev].to] -= amt;
    }
    void relabel(int vertex) {
        cost_t newHeight = -INFCOST;
        for(unsigned int i=0; i<G[vertex].size(); ++i) {
            Edge const&e = G[vertex][i];
            if(e.f && newHeight < h[e.to]-e.c) {
                newHeight = h[e.to] - e.c;
                cur[vertex] = i;
            }
        }
        h[vertex] = newHeight - eps;
    }

```

```

}
static constexpr int scale=2;
pair<flow_t, cost_t> minCostMaxFlow() {
    cost_t retCost = 0;
    for(int i=0; i<N; ++i)
        for(Edge &e:G[i])
            retCost += e.c*(e.f);
    //find max-flow
    flow_t retFlow = max_flow();
    h.assign(N, 0); ex.assign(N, 0);
    isq.assign(N, 0); cur.assign(N, 0);
    queue<int> q;
    for(eps; eps>>=scale) {
        //refine
        fill(cur.begin(), cur.end(), 0);
        for(int i=0; i<N; ++i)
            for(auto &e:G[i])
                if(h[i] + e.c - h[e.to] < 0 && e.f) push(e, e.f);
        for(int i=0; i<N; ++i) {
            if(ex[i]>0) {
                q.push(i);
                isq[i]=1;
            }
        }
        // make flow feasible
        while(!q.empty()) {
            int u=q.front(); q.pop();
            isq[u]=0;
            while(ex[u]>0) {
                if(cur[u] == G[u].size())
                    relabel(u);
                for(unsigned int &i=cur[u], max_i = G[u].size()
                    ; i<max_i; ++i) {
                    Edge &e=G[u][i];
                    if(h[u] + e.c - h[e.to] < 0) {
                        push(e, ex[u]);
                        if(ex[e.to]>0 && isq[e.to]==0) {
                            q.push(e.to);
                            isq[e.to]=1;
                        }
                    }
                    if(ex[u]==0) break;
                }
            }
        }
        if(eps>1 && eps>>scale==0) {
            eps = 1<<scale;
        }
    }
    for(int i=0; i<N; ++i) {
        for(Edge &e:G[i]) {
            retCost -= e.c*(e.f);
        }
    }
    return make_pair(retFlow, retCost/2/N);
}
flow_t getFlow(Edge const &e) {
    return G[e.to][e.rev].f;
}
};

```

### 1.3 Edmonds Max Matching

```

#define MAXN 505
vector<int> g[MAXN];
int pa[MAXN], match[MAXN], st[MAXN], S[MAXN], v[MAXN];
int t, n;
inline int lca(int x, int y) {
    for(++t; swap(x, y); if(x) {
        if(v[x]==t) return x;
        v[x]=t, x=st[pa[match[x]]];
    }
}
#define qpush(x) q.push(x), S[x]=0
void flower(int x, int y, int l, queue<int> &q) {
    while(st[x]!=1) {
        pa[x]=y, y=match[x];
        if(S[y]==1) qpush(y);
        st[x]=st[y]=1, x=pa[y];
    }
}
bool bfs(int x) {
    for(int i=1; i<=n; ++i) st[i]=i;
    memset(S+1, -1, sizeof(int)*n);
    queue<int> q;
    qpush(x);
    while(q.size()) {
        x=q.front(); q.pop();
        for(size_t i=0; i<g[x].size(); ++i) {
            int y=g[x][i];
            if(S[y]==-1) {
                pa[y]=x, S[y]=1;
                if(!match[y]) {
                    for(int lst=x; y=lst, x=pa[y]) {
                        lst=match[x], match[x]=y, match[y]=x;
                    }
                    return 1;
                }
                qpush(match[y]);
            } else if(!S[y] && st[y]!=st[x]) {
                int l=lca(y, x);
                flower(y, x, l, q); flower(x, y, l, q);
            }
        }
    }
    return 0;
}
int blossom() {
    int ans=0;
    for(int i=1; i<=n; ++i)
        if(!match[i] && bfs(i)) ++ans;
    return ans;
}
int main() {
    while(m--) {
        g[x].push_back(y);
        g[y].push_back(x);
    }
    printf("%d\n", blossom());
    for(int i=1; i<=n; ++i) printf("%d ", match[i]);
}

```

## 1.4 Hungarian Min-cost

```
double MinCostMatching(const VVD &cost, VI &Lmate, VI &
    Rmate) {
    int n = int(cost.size());
    // construct dual feasible solution
    VD u(n);
    VD v(n);
    for (int i = 0; i < n; i++) {
        u[i] = cost[i][0];
        for (int j = 1; j < n; j++) u[i] = min(u[i], cost[i]
            ][j]);
    }
    for (int j = 0; j < n; j++) {
        v[j] = cost[0][j] - u[0];
        for (int i = 1; i < n; i++) v[j] = min(v[j], cost[i]
            ][j] - u[i]);
    }
    // construct primal solution satisfying complementary
    // slackness
    Lmate = VI(n, -1);
    Rmate = VI(n, -1);
    int mated = 0;
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < n; j++) {
            if (Rmate[j] != -1) continue;
            if (fabs(cost[i][j] - u[i] - v[j]) < 1e-10) {
                Lmate[i] = j;
                Rmate[j] = i;
                mated++;
                break;
            }
        }
    }
    VD dist(n);
    VI dad(n);
    VI seen(n);
    // repeat until primal solution is feasible
    while (mated < n) {
        // find an unmatched left node
        int s = 0;
        while (Lmate[s] != -1) s++;
        // initialize Dijkstra
        fill(dad.begin(), dad.end(), -1);
        fill(seen.begin(), seen.end(), 0);
        for (int k = 0; k < n; k++)
            dist[k] = cost[s][k] - u[s] - v[k];
        int j = 0;
        while (true) {
            // find closest
            j = -1;
            for (int k = 0; k < n; k++) {
                if (seen[k]) continue;
                if (j == -1 || dist[k] < dist[j]) j = k;
            }
            seen[j] = 1;
            // termination condition
            if (Rmate[j] == -1) break;
            // relax neighbors
            const int i = Rmate[j];
```

```
        for (int k = 0; k < n; k++) {
            if (seen[k]) continue;
            const double new_dist = dist[j] + cost[i][k] - u
                [i] - v[k];
            if (dist[k] > new_dist) {
                dist[k] = new_dist;
                dad[k] = j;
            }
        }
        // update dual variables
        for (int k = 0; k < n; k++) {
            if (k == j || !seen[k]) continue;
            const int i = Rmate[k];
            v[k] += dist[k] - dist[j];
            u[i] -= dist[k] - dist[j];
        }
        u[s] += dist[j];
        // augment along path
        while (dad[j] >= 0) {
            const int d = dad[j];
            Rmate[j] = Rmate[d];
            Lmate[Rmate[j]] = j;
            j = d;
        }
        Rmate[j] = s;
        Lmate[s] = j;
        mated++;
    }
    double value = 0;
    for (int i = 0; i < n; i++)
        value += cost[i][Lmate[i]];
    return value;
}
```

## 1.5 König's Theorem (Text)

In any bipartite graph, the number of edges in a maximum matching equals the number of vertices in a minimum vertex cover. To exhibit the vertex cover:

1. Find a maximum matching
2. Change each edge **used** in the matching into a directed edge from **right to left**
3. Change each edge **not used** in the matching into a directed edge from **left to right**
4. Compute the set  $T$  of all vertices reachable from unmatched vertices on the left (including themselves)
5. The vertex cover consists of all vertices on the right that are **in**  $T$ , and all vertices on the left that are **not in**  $T$

## 1.6 Minimum Edge Cover (Text)

If a minimum edge cover contains  $C$  edges, and a maximum matching contains  $M$  edges, then  $C + M = |V|$ . To obtain the edge cover, start with a maximum matching, and then, for every vertex not matched, just select some edge incident upon it and add it to the edge set.

## 2 Geometry

### 2.1 Miscellaneous geometry

```
double INF = 1e100, EPS = 1e-12;
struct PT {
    double x, y;
    PT() {}
    PT(double x, double y) : x(x), y(y) {}
    PT(const PT &p) : x(p.x), y(p.y) {}
    PT operator + (const PT &p) const { return PT(x+p.x,
        y+p.y); }
    PT operator - (const PT &p) const { return PT(x-p.x,
        y-p.y); }
    PT operator * (double c) const { return PT(x*c,
        y*c); }
    PT operator / (double c) const { return PT(x/c,
        y/c); }
};

double dot(PT p, PT q) { return p.x*q.x+p.y*q.y; }
double dist2(PT p, PT q) { return dot(p-q, p-q); }
double cross(PT p, PT q) { return p.x*q.y-p.y*q.x; }
ostream &operator<<(ostream &os, const PT &p) {
    os << "(" << p.x << ", " << p.y << ")";
}

// rotate a point CCW or CW around the origin
PT RotateCCW90(PT p) { return PT(-p.y, p.x); }
PT RotateCW90(PT p) { return PT(p.y, -p.x); }
PT RotateCCW(PT p, double t) {
    return PT(p.x*cos(t)-p.y*sin(t), p.x*sin(t)+p.y*cos(t));
}

// project point c onto line through a and b
// assuming a != b
PT ProjectPointLine(PT a, PT b, PT c) {
    return a + (b-a)*dot(c-a, b-a)/dot(b-a, b-a);
}

// project point c onto line segment through a and b
PT ProjectPointSegment(PT a, PT b, PT c) {
    double r = dot(b-a, b-a);
    if (fabs(r) < EPS) return a;
    r = dot(c-a, b-a)/r;
    if (r < 0) return a;
    if (r > 1) return b;
    return a + (b-a)*r;
}

// compute distance from c to segment between a and b
double DistancePointSegment(PT a, PT b, PT c) {
    return sqrt(dist2(c, ProjectPointSegment(a, b, c)));
}

// compute distance between point (x,y,z) and plane ax+
// by+cz=d
double DistancePointPlane(double x, double y, double z,
    double a, double b, double c,
    double d)
{
    return fabs(a*x+b*y+c*z-d)/sqrt(a*a+b*b+c*c);
}
```

```
// determine if lines from a to b and c to d are
// parallel or collinear
bool LinesParallel(PT a, PT b, PT c, PT d) {
    return fabs(cross(b-a, c-d)) < EPS;
}

bool LinesCollinear(PT a, PT b, PT c, PT d) {
    return LinesParallel(a, b, c, d)
        && fabs(cross(a-b, a-c)) < EPS
        && fabs(cross(c-d, c-a)) < EPS;
}

// determine if line segment from a to b intersects with
// line segment from c to d
bool SegmentsIntersect(PT a, PT b, PT c, PT d) {
    if (LinesCollinear(a, b, c, d)) {
        if (dist2(a, c) < EPS || dist2(a, d) < EPS ||
            dist2(b, c) < EPS || dist2(b, d) < EPS) return
            true;
        if (dot(c-a, c-b) > 0 && dot(d-a, d-b) > 0 && dot(c-
            b, d-b) > 0)
            return false;
        return true;
    }
    if (cross(d-a, b-a) * cross(c-a, b-a) > 0) return
        false;
    if (cross(a-c, d-c) * cross(b-c, d-c) > 0) return
        false;
    return true;
}

// compute intersection of line passing through a and b
// with line passing through c and d, assuming that
// unique
// intersection exists; for segment intersection, check
// if
// segments intersect first
PT ComputeLineIntersection(PT a, PT b, PT c, PT d) {
    b=b-a; d=c-d; c=c-a;
    assert(dot(b, b) > EPS && dot(d, d) > EPS);
    return a + b*cross(c, d)/cross(b, d);
}

// compute center of circle given three points
PT ComputeCircleCenter(PT a, PT b, PT c) {
    b=(a+b)/2;
    c=(a+c)/2;
    return ComputeLineIntersection(b, b+RotateCW90(a-b), c
        , c+RotateCW90(a-c));
}

// determine if point is in a possibly non-convex
// polygon (by William
// Randolph Franklin); returns 1 for strictly interior
// points, 0 for
// strictly exterior points, and 0 or 1 for the
// remaining points.
// Note that it is possible to convert this into an *
// exact* test using
// integer arithmetic by taking care of the division
// appropriately
// (making sure to deal with signs properly) and then by
// writing exact
```

```

// tests for checking point on polygon boundary
bool PointInPolygon(const vector<PT> &p, PT q) {
    bool c = 0;
    for (int i = 0; i < p.size(); i++) {
        int j = (i+1)%p.size();
        if ((p[i].y <= q.y && q.y < p[j].y ||
            p[j].y <= q.y && q.y < p[i].y) &&
            q.x < p[i].x + (p[j].x - p[i].x) * (q.y - p[i].y) /
                (p[j].y - p[i].y))
            c = !c;
    }
    return c;
}

// determine if point is on the boundary of a polygon
bool PointOnPolygon(const vector<PT> &p, PT q) {
    for (int i = 0; i < p.size(); i++)
        if (dist2(ProjectPointSegment(p[i], p[(i+1)%p.size()],
            q), q) < EPS)
            return true;
    return false;
}

// compute intersection of line through points a and b
// with
// circle centered at c with radius r > 0
vector<PT> CircleLineIntersection(PT a, PT b, PT c,
    double r) {
    vector<PT> ret;
    b = b-a;
    a = a-c;
    double A = dot(b, b);
    double B = dot(a, b);
    double C = dot(a, a) - r*r;
    double D = B*B - A*C;
    if (D < -EPS) return ret;
    ret.push_back(c+a+b*(-B+sqrt(D+EPS))/A);
    if (D > EPS)
        ret.push_back(c+a+b*(-B-sqrt(D))/A);
    return ret;
}

// compute intersection of circle centered at a with
// radius r
// with circle centered at b with radius R
vector<PT> CircleCircleIntersection(PT a, PT b, double r,
    double R) {
    vector<PT> ret;
    double d = sqrt(dist2(a, b));
    if (d > r+R || d+min(r, R) < max(r, R)) return ret;
    double x = (d*d-R*R+r*r)/(2*d);
    double y = sqrt(r*r-x*x);
    PT v = (b-a)/d;
    ret.push_back(a+v*x + RotateCCW90(v)*y);
    if (y > 0)
        ret.push_back(a+v*x - RotateCCW90(v)*y);
    return ret;
}

// This code computes the area or centroid of a (
// possibly nonconvex)
// polygon, assuming that the coordinates are listed in
// a clockwise or

```

```

// counterclockwise fashion. Note that the centroid is
// often known as
// the "center of gravity" or "center of mass".
double ComputeSignedArea(const vector<PT> &p) {
    double area = 0;
    for (int i = 0; i < p.size(); i++) {
        int j = (i+1) % p.size();
        area += p[i].x*p[j].y - p[j].x*p[i].y;
    }
    return area / 2.0;
}

double ComputeArea(const vector<PT> &p) {
    return fabs(ComputeSignedArea(p));
}

PT ComputeCentroid(const vector<PT> &p) {
    PT c(0,0);
    double scale = 6.0 * ComputeSignedArea(p);
    for (int i = 0; i < p.size(); i++) {
        int j = (i+1) % p.size();
        c = c + (p[i]+p[j])*(p[i].x*p[j].y - p[j].x*p[i].y);
    }
    return c / scale;
}

// tests whether or not a given polygon (in CW or CCW
// order) is simple
bool IsSimple(const vector<PT> &p) {
    for (int i = 0; i < p.size(); i++) {
        for (int k = i+1; k < p.size(); k++) {
            int j = (i+1) % p.size();
            int l = (k+1) % p.size();
            if (i == l || j == k) continue;
            if (SegmentsIntersect(p[i], p[j], p[k], p[l]))
                return false;
        }
    }
    return true;
}

```

## 2.2 3D geometry

```

public class Geom3D {
    // distance from point (x, y, z) to plane aX + bY + cZ
    // + d = 0
    public static double ptPlaneDist(double x, double y,
        double z,
        double a, double b, double c, double d) {
        return Math.abs(a*x + b*y + c*z + d) / Math.sqrt(a*a
            + b*b + c*c);
    }

    // distance between parallel planes aX + bY + cZ + d1
    // = 0 and
    // aX + bY + cZ + d2 = 0
    public static double planePlaneDist(double a, double b,
        double c,
        double d1, double d2) {
        return Math.abs(d1 - d2) / Math.sqrt(a*a + b*b + c*c);
    }
}

```



```
// distance from point (px, py, pz) to line (x1, y1,
// z1)-(x2, y2, z2)
// (or ray, or segment; in the case of the ray, the
// endpoint is the
// first point)
public static final int LINE = 0;
public static final int SEGMENT = 1;
public static final int RAY = 2;
public static double ptLineDistSq(double x1, double y1
, double z1,
double x2, double y2, double z2, double px, double
py, double pz,
int type) {
double pd2 = (x1-x2)*(x1-x2) + (y1-y2)*(y1-y2) + (z1
-z2)*(z1-z2);
double x, y, z;
if (pd2 == 0) {
x = x1;
y = y1;
z = z1;
} else {
double u = ((px-x1)*(x2-x1) + (py-y1)*(y2-y1) + (
pz-z1)*(z2-z1)) / pd2;
x = x1 + u * (x2 - x1);
y = y1 + u * (y2 - y1);
z = z1 + u * (z2 - z1);
if (type != LINE && u < 0) {
x = x1;
y = y1;
z = z1;
}
if (type == SEGMENT && u > 1.0) {
x = x2;
y = y2;
z = z2;
}
}
return (x-px)*(x-px) + (y-py)*(y-py) + (z-pz)*(z-pz)
;
}

public static double ptLineDist(double x1, double y1,
double z1,
double x2, double y2, double z2, double px, double
py, double pz,
int type) {
return Math.sqrt(ptLineDistSq(x1, y1, z1, x2, y2, z2
, px, py, pz, type));
}
}
```

## 2.3 Convex hull

```
typedef pair<long long, long long> PT;
inline long long cross(PT o, PT a, PT b) {
PT OA = {a.first-o.first, a.second-o.second};
PT OB = {b.first-o.first, b.second-o.second};
return OA.first*OB.second - OA.second*OB.first;
}
inline long double dist(PT a, PT b) {
```

```
return sqrt(pow(a.first-b.first,2)+pow(a.second-b.
second,2));
}
vector<PT> convexhull(vector<PT> a) {
vector<PT> hull;
sort(a.begin(), a.end(), [](PT i, PT j) {
if(i.second!=j.second)
return i.second < j.second;
return i.first < j.first;
});
for(int i=0; i<a.size(); ++i) {
while(hull.size()>1 && cross(hull[hull.size()-2],
hull.back(), a[i])<=0)
hull.pop_back();
hull.push_back(a[i]);
}
for(int i=a.size()-1, siz = hull.size(); i--;) {
while(hull.size()>siz && cross(hull[hull.size()-2],
hull.back(), a[i])<=0)
hull.pop_back();
hull.push_back(a[i]);
}
return hull;
}
```

## 2.4 Min Enclosing Circle

```
// Minimum enclosing circle, Welzl's algorithm
// Expected linear time.
// If there are any duplicate points in the input, be
// sure to remove them first.
struct point {
double x;
double y;
};
struct circle {
double x;
double y;
double r;
circle() {}
circle(double x, double y, double r): x(x), y(y)
, r(r) {}
};
circle b_md(vector<point> R) {
if (R.size() == 0) {
return circle(0, 0, -1);
} else if (R.size() == 1) {
return circle(R[0].x, R[0].y, 0);
} else if (R.size() == 2) {
return circle((R[0].x+R[1].x)/2.0,
(R[0].y+R[1].y)/2.0,
hypot(R[0].x-R
[1].x, R
[0].y-R[1].
y)/2.0);
} else {
double D = (R[0].x - R[2].x)*(R[1].y - R
[2].y) - (R[1].x - R[2].x)*(R[0].y -
R[2].y);
double p0 = ((R[0].x - R[2].x)*(R[0].x
+ R[2].x) + (R[0].y - R[2].y)*(R[0].y
```

```

        + R[2].y)) / 2 * (R[1].y - R[2].y) -
        ((R[1].x - R[2].x)*(R[1].x + R[2].x)
        + (R[1].y - R[2].y)*(R[1].y + R[2].y)
        )) / 2 * (R[0].y - R[2].y))/D;
    double p1 = (((R[1].x - R[2].x)*(R[1].x
    + R[2].x) + (R[1].y - R[2].y)*(R[1].y
    + R[2].y)) / 2 * (R[0].x - R[2].x) -
    ((R[0].x - R[2].x)*(R[0].x + R[2].x)
    + (R[0].y - R[2].y)*(R[0].y + R[2].y)
    )) / 2 * (R[1].x - R[2].x))/D;
    return circle(p0, p1, hypot(R[0].x - p0,
    R[0].y - p1));
}
}
circle b_minidisk(vector<point>& P, int i, vector<point>
R) {
    if (i == P.size() || R.size() == 3) {
        return b_md(R);
    } else {
        circle D = b_minidisk(P, i+1, R);
        if (hypot(P[i].x-D.x, P[i].y-D.y) > D.r)
        {
            R.push_back(P[i]);
            D = b_minidisk(P, i+1, R);
        }
        return D;
    }
}
// Call this function.
circle minidisk(vector<point> P) {
    random_shuffle(P.begin(), P.end());
    return b_minidisk(P, 0, vector<point>());
}

```

## 2.5 Pick's Theorem (Text)

For a polygon with all vertices on lattice points,  $A = i + b/2 - 1$ , where  $A$  is the area,  $i$  is the number of lattice points strictly within the polygon, and  $b$  is the number of lattice points on the boundary of the polygon. (Note, there is no generalization to higher dimensions)

## 2.6 Slow Delaunay triangulation

```

// Slow but simple Delaunay triangulation. Does not
// handle degenerate cases (from O'Rourke)
//
// Running time: O(n^4)
// INPUT:      x[] = x-coordinates
//            y[] = y-coordinates
// OUTPUT:     triples = a vector containing m triples of
//            indices corresponding to triangle vertices
typedef double T;
struct triple {
    int i, j, k;
    triple() {}
    triple(int i, int j, int k) : i(i), j(j), k(k) {}
};

```

```

vector<triple> delaunayTriangulation(vector<T>& x,
vector<T>& y) {
    int n = x.size();
    vector<T> z(n);
    vector<triple> ret;
    for (int i = 0; i < n; i++)
        z[i] = x[i] * x[i] + y[i] * y[i];
    for (int i = 0; i < n-2; i++) {
        for (int j = i+1; j < n; j++) {
            for (int k = i+1; k < n; k++) {
                if (j == k) continue;
                double xn = (y[j]-y[i])*(z[k]-z[i]) - (y[k]-y[i]
                )*(z[j]-z[i]);
                double yn = (x[k]-x[i])*(z[j]-z[i]) - (x[j]-x[i]
                )*(z[k]-z[i]);
                double zn = (x[j]-x[i])*(y[k]-y[i]) - (x[k]-x[i]
                )*(y[j]-y[i]);
                bool flag = zn < 0;
                for (int m = 0; flag && m < n; m++)
                    flag = flag && ((x[m]-x[i])*xn + (y[m]-y[i])*
                    yn + (z[m]-z[i])*zn <= 0);
                if (flag) ret.push_back(triple(i, j, k));
            }
        }
    }
    return ret;
}

int main() {
    T xs[]={0, 0, 1, 0.9};
    T ys[]={0, 1, 0, 0.9};
    vector<T> x(&xs[0], &xs[4]), y(&ys[0], &ys[4]);
    vector<triple> tri = delaunayTriangulation(x, y);
    //expected: 0 1 3
    //           0 3 2
    for(int i = 0; i < tri.size(); i++)
        printf("%d %d %d\n", tri[i].i, tri[i].j, tri[i].k);
}

```

## 3 Numerical algorithms

### 3.1 Pollard Rho

```

typedef long long unsigned int llui;
typedef long long int lli;
typedef long double float64;
llui mul_mod(llui a, llui b, llui m){
    llui y = (llui) ((float64) a*(float64) b/m+(float64) 1/2);
    y = y * m;
    llui x = a * b;
    llui r = x - y;
    if ( (lli) r < 0 ) {
        r = r + m; y = y - 1;
    }
    return r;
}
llui C,a,b;
llui gcd(){
    llui c;
    if(a>b){
        c = a; a = b; b = c;
    }
}

```



```

    }
    while(1){
        if(a == 1LL) return 1LL;
        if(a == 0 || a == b) return b;
        c = a; a = b*a;
        b = c;
    }
}
llui f(llui a, llui b){
    llui tmp;
    tmp = mul_mod(a,a,b);
    tmp+=C; tmp%=b;
    return tmp;
}
llui pollard(llui n){
    if(!(n&1)) return 2;
    C=0;
    llui iteracoes = 0;
    while(iteracoes <= 1000){
        llui x,y,d;
        x = y = 2; d = 1;
        while(d == 1){
            x = f(x,n);
            y = f(f(y,n),n);
            llui m = (x>y)?(x-y):(y-x);
            a = m; b = n; d = gcd();
        }
        if(d != n)
            return d;
        iteracoes++; C = rand();
    }
    return 1;
}
llui pot(llui a, llui b, llui c){
    if(b == 0) return 1;
    if(b == 1) return a%c;
    llui resp = pot(a,b>>1,c);
    resp = mul_mod(resp,resp,c);
    if(b&1)
        resp = mul_mod(resp,a,c);
    return resp;
}
// Rabin-Miller primality testing algorithm
bool isPrime(llui n){
    llui d = n-1;
    llui s = 0;
    if(n <=3 || n == 5) return true;
    if(!(n&1)) return false;
    while(!(d&1)){ s++; d>>=1; }
    for(llui i = 0; i<32; i++){
        llui a = rand();
        a <<=32;
        a+=rand();
        a%=(n-3); a+=2;
        llui x = pot(a,d,n);
        if(x == 1 || x == n-1) continue;
        for(llui j = 1; j<= s-1; j++){
            x = mul_mod(x,x,n);
            if(x == 1) return false;
            if(x == n-1) break;
        }
    }
}

```

```

        if(x != n-1) return false;
    }
    return true;
}
map<llui,int> factors;
// Precondition: factors is an empty map, n is a
// positive integer
// Postcondition: factors[p] is the exponent of p in
// prime factorization of n
void fact(llui n){
    if(!isPrime(n)){
        llui fac = pollard(n);
        fact(n/fac); fact(fac);
    }else{
        map<llui,int>::iterator it;
        it = factors.find(n);
        if(it != factors.end()){
            (*it).second++;
        }else{
            factors[n] = 1;
        }
    }
}
}
}

```

## 3.2 Simplex algorithm

```

// Two-phase simplex algorithm for solving linear
// programs of the form
//      maximize      c^T x
//      subject to    Ax <= b ; x >= 0
// INPUT: A -- an m x n matrix
//         b -- an m-dimensional vector
//         c -- an n-dimensional vector
//         x -- a vector where the optimal solution will
//         be stored
// OUTPUT: value of the optimal solution (infinity if
//         unbounded above, nan if infeasible)
// To use this code, create an LPSolver object with
// A, b, and c as arguments. Then, call Solve(x).
typedef long double DOUBLE;
typedef vector<DOUBLE> VD;
typedef vector<VD> VVD;
typedef vector<int> VI;
const DOUBLE EPS = 1e-9;
struct LPSolver {
    int m, n;
    VI B, N;
    VVD D;
    LPSolver(const VVD &A, const VD &b, const VD &c) :
        m(b.size()), n(c.size()), N(n + 1), B(m), D(m + 2,
            VD(n + 2)) {
        for (int i = 0; i < m; i++) for (int j = 0; j < n; j
            ++) D[i][j] = A[i][j];
        for (int i = 0; i < m; i++) { B[i] = n + i; D[i][n]
            = -1; D[i][n + 1] = b[i]; }
        for (int j = 0; j < n; j++) { N[j] = j; D[m][j] = -c
            [j]; }
        N[n] = -1; D[m + 1][n] = 1;
    }
    void Pivot(int r, int s) {
        double inv = 1.0 / D[r][s];
    }
}

```

```

    for (int i = 0; i < m + 2; i++) if (i != r)
        for (int j = 0; j < n + 2; j++) if (j != s)
            D[i][j] -= D[r][j] * D[i][s] * inv;
    for (int j = 0; j < n + 2; j++) if (j != s) D[r][j]
        *= inv;
    for (int i = 0; i < m + 2; i++) if (i != r) D[i][s]
        *= -inv;
    D[r][s] = inv;
    swap(B[r], N[s]);
}
bool Simplex(int phase) {
    int x = phase == 1 ? m + 1 : m;
    while (true) {
        int s = -1;
        for (int j = 0; j <= n; j++) {
            if (phase == 2 && N[j] == -1) continue;
            if (s == -1 || D[x][j] < D[x][s] || D[x][j] == D
                [x][s] && N[j] < N[s]) s = j;
        }
        if (D[x][s] > -EPS) return true;
        int r = -1;
        for (int i = 0; i < m; i++) {
            if (D[i][s] < EPS) continue;
            if (r == -1 || D[i][n + 1] / D[i][s] < D[r][n +
                1] / D[r][s] ||
                (D[i][n + 1] / D[i][s]) == (D[r][n + 1] / D[r
                ][s]) && B[i] < B[r]) r = i;
        }
        if (r == -1) return false;
        Pivot(r, s);
    }
}
DOUBLE Solve(VD &x) {
    int r = 0;
    for (int i = 1; i < m; i++) if (D[i][n + 1] < D[r][n
        + 1]) r = i;
    if (D[r][n + 1] < -EPS) {
        Pivot(r, n);
        if (!Simplex(1) || D[m + 1][n + 1] < -EPS) return
            -numeric_limits<DOUBLE>::infinity();
        for (int i = 0; i < m; i++) if (B[i] == -1) {
            int s = -1;
            for (int j = 0; j <= n; j++)
                if (s == -1 || D[i][j] < D[i][s] || D[i][j] ==
                    D[i][s] && N[j] < N[s]) s = j;
            Pivot(i, s);
        }
    }
    if (!Simplex(2)) return numeric_limits<DOUBLE>::
        infinity();
    x = VD(n);
    for (int i = 0; i < m; i++) if (B[i] < n) x[B[i]] =
        D[i][n + 1];
    return D[m][n + 1];
}
};
int main() {
    const int m = 4, n = 3;
    DOUBLE _A[m][n] = {
        { 6, -1, 0 },

```

```

        { -1, -5, 0 },
        { 1, 5, 1 },
        { -1, -5, -1 }
    };
    DOUBLE _b[m] = { 10, -4, 5, -5 }, _c[n] = { 1, -1, 0
    };
    VVD A(m);
    VD b(_b, _b + m), c(_c, _c + n);
    for (int i = 0; i < m; i++) A[i] = VD(_A[i], _A[i] + n
    );
    LPSolver solver(A, b, c);
    VD x;
    DOUBLE value = solver.Solve(x);
    cerr << "VALUE: " << value << endl; // VALUE: 1.29032
    cerr << "SOLUTION:"; // SOLUTION: 1.74194 0.451613 1
    for (size_t i = 0; i < x.size(); i++) cerr << " " << x
        [i];
}

```

### 3.3 Reduced row echelon form

*// Reduced row echelon form via Gauss-Jordan elimination  
 // with partial pivoting. This can be used for  
 // computing the rank of a matrix.  
 // Running time:  $O(n^3)$   
 // INPUT:  $a[][]$  = an  $n \times m$  matrix  
 // OUTPUT:  $rref[][]$  = an  $n \times m$  matrix (stored in  $a[][]$ )  
 // returns rank of  $a[][]$*

```

const double EPSILON = 1e-10;
typedef double T;
typedef vector<T> VT;
typedef vector<VT> VVT;
int rref(VVT &a) {
    int n = a.size(), m = a[0].size(), r = 0;
    for (int c = 0; c < m && r < n; c++) {
        int j = r;
        for (int i = r + 1; i < n; i++)
            if (fabs(a[i][c]) > fabs(a[j][c])) j = i;
        if (fabs(a[j][c]) < EPSILON) continue;
        swap(a[j], a[r]);
        T s = 1.0 / a[r][c];
        for (int j = 0; j < m; j++) a[r][j] *= s;
        for (int i = 0; i < n; i++) if (i != r) {
            T t = a[i][c];
            for (int j = 0; j < m; j++) a[i][j] -= t * a[r][j
            ];
        }
        r++;
    }
    return r;
}

```

### 3.4 Fast Fourier transform

```

template<typename fpt>
struct fft_wrap {
    using cpx_t = complex<fpt>;
    const fpt two_pi = 4 * acos(0);
    vector<cpx_t> roots; //stores the N-th roots of unity.
    int N;

```

```

fft_wrap(int N) : roots(N), N(N) {
    for (int i = 0; i < N; ++i) {
        roots[i] = EXP(two_pi * i / fpt(N));
    }
}

cpx_t EXP(fpt theta) {
    return {cos(theta), sin(theta)};
}

void fft(cpx_t *in, cpx_t *out, int size, int dir) {
    bit_reverse(in, out, size);
    for (int s = 0; (1 << s) < size; ++s) {
        int s_ = s + 1;
        for (int k = 0; k < size; k += (1 << s_)) {
            for (int j = 0; j < (1 << s); ++j) {
                int id = (N + dir * (N >> s_) * j) & (N - 1);
                cpx_t w = roots[id];
                cpx_t t = w * out[k + j + (1 << s)];
                cpx_t u = out[k + j];
                out[k + j] = u + t;
                out[k + j + (1 << s)] = u - t;
            }
        }
    }

    void bit_reverse(cpx_t *in, cpx_t *out, int size) {
        for (int i = 0; i < size; ++i) {
            int rev = 0, i_copy = i;
            for (int j = 0; (1 << j) < size; ++j) {
                rev = (rev << 1) + (i_copy & 1);
                i_copy >>= 1;
            }
            out[rev] = in[i];
        }
    }
};

int main() {
    typedef complex<double> cpx_t;
    fft_wrap<double> fft_wrapper(2048);

    vector<cpx_t> in = {1.0, 1.0, 1.0, 1.0, 0.0, 0.0, 0.0,
                       0.0}, out(8);

    fft_wrapper.fft(&in[0], &out[0], 8, 1);
    fft_wrapper.fft(&out[0], &in[0], 8, -1);
    for (int i = 0; i < 8; ++i) {
        cout << in[i].real() << ' ' << in[i].imag() << endl;
    }
}

```

### 3.5 Number Theoretic transform

```

// gen should be 5^((p - 1) / fft_wrapper_size)

template<int P>
struct field_t { ... };

template<typename fpt, int gen>
struct fft_wrap {
    vector<fpt> roots; //stores the N-th roots of unity.
    int N;

```

```

fft_wrap(int _N) : roots(_N), N(_N) {
    roots[0] = 1;
    for (int i = 1; i < N; ++i) {
        roots[i] = roots[i - 1] * gen;
        if (i != N - 1) {
            assert(roots[i].v != 1);
        }
    }
}

void fft(fpt *in, fpt *out, int size, int dir) {
    bit_reverse(in, out, size);
    for (int s = 0; (1 << s) < size; ++s) {
        int s_ = s + 1;
        for (int k = 0; k < size; k += (1 << s_)) {
            for (int j = 0; j < (1 << s); ++j) {
                int id = (N + dir * (N >> s_) * j) & (N - 1);
                fpt w = roots[id];
                fpt t = w * out[k + j + (1 << s)];
                fpt u = out[k + j];
                out[k + j] = u + t;
                out[k + j + (1 << s)] = u - t;
            }
        }
    }

    void bit_reverse(fpt *in, fpt *out, int size) {
        for (int i = 0; i < size; ++i) {
            int rev = 0, i_copy = i;
            for (int j = 0; (1 << j) < size; ++j) {
                rev = (rev << 1) + (i_copy & 1);
                i_copy >>= 1;
            }
            out[rev] = in[i];
        }
    }
};

using fpt = field_t<mod>;
fft_wrap<fpt, gen> fft_wrapper(1 << 18);

vector<fpt> polymul(vector<fpt> P1, vector<fpt> P2) {
    int fsize = P1.size() + P2.size() - 1;
    int N = max(P1.size(), P2.size());
    while(N & (N - 1)) {
        ++N;
    }
    N *= 2;
    P1.resize(N); P2.resize(N);
    vector<fpt> temp(N);
    fft_wrapper.fft(&P1[0], &temp[0], N, 1);
    fft_wrapper.fft(&P2[0], &P1[0], N, 1);
    for (int i = 0; i < N; ++i) {
        P1[i] *= temp[i];
    }
    fft_wrapper.fft(&P1[0], &temp[0], N, -1);
    field_t<mod> inv(N);
    inv = inv.modexp(mod - 2);
    for (int i = 0; i < N; ++i) {
        temp[i] *= inv;
    }
    temp.resize(fsize);
    return temp;
}

```

}

### 3.6 Discrete Logarithm

```
// Calculates x such that g^x % md == h
int baby_giant(int g, int h, int md) {
    unordered_map<int,int> mp;
    int sq = ceil(sqrtl(md));
    for(int i=0, now=1; i<sq; ++i)
        mp[now] = i, now = (long long) now*g % md;
    for(int i=0, jmp=power(g, md-1-sq); i<sq; ++i) {
        if(mp.find(h) != mp.end()) return i*sq+mp[h];
        h = (long long) h*jmp % md;
    }
    return -1;
}
```

### 3.7 Mobius Inversion (Text)

$$\mu(n) = \begin{cases} 0 & n \text{ not squarefree} \\ 1 & n \text{ squarefree w/ even no. of prime factors} \\ -1 & n \text{ squarefree w/ odd no. of prime factors} \end{cases}$$

Note that  $\mu(a)\mu(b) = \mu(ab)$  for  $a, b$  relatively prime

$$\text{Also } \sum_{d|n} \mu(d) = \begin{cases} 1 & \text{if } n = 1 \\ 0 & \text{otherwise} \end{cases}$$

**Möbius Inversion** If  $g(n) = \sum_{d|n} f(d)$  for all  $n \geq 1$ , then  $f(n) = \sum_{d|n} \mu(d)g(n/d)$  for all  $n \geq 1$ .

### 3.8 Burnside Lemma (Text)

The number of orbits of a set  $X$  under the group action  $G$  equals the average number of elements of  $X$  fixed by the elements of  $G$ .

Here's an example. Consider a square of  $2n$  times  $2n$  cells. How many ways are there to color it into  $X$  colors, up to rotations and/or reflections? Here, the group has only 8 elements (rotations by 0, 90, 180 and 270 degrees, reflections over two diagonals, over a vertical line and over a horizontal line). Every coloring stays itself after rotating by 0 degrees, so that rotation has  $X^{4n^2}$  fixed points. Rotation by 180 degrees and reflections over a horizontal/vertical line split all cells in pairs that must be of the same color for a coloring to be unaffected by such rotation/reflection, thus there exist  $X^{2n^2}$  such colorings for each of them. Rotations by 90 and 270 degrees split cells in groups of four, thus yielding  $X^{n^2}$  fixed colorings. Reflections over diagonals split cells into  $2n$  groups of 1 (the diagonal itself) and  $2n^2 - n$  groups of 2 (all remaining cells), thus yielding  $X^{2n^2-n+2n} = X^{2n^2+n}$  unaffected colorings. So, the answer is  $(X^{4n^2} + 3X^{2n^2} + 2X^{n^2} + 2X^{2n^2+n})/8$ .

### 3.9 Number Theory (Modular, CRT, Linear Diophantine)

```
// All algorithms described here work on nonnegative integers.
int mod(int a, int b) {
    return ((a%b) + b) % b;
}
int lcm(int a, int b) {
    return a / __gcd(a, b) * b;
}
int powermod(int a, int b, int m) {
    return b?powermod(a*a%m, b/2, m) * (b%2?a:1)%m:1;
}
// returns g = gcd(a, b); finds x, y such that d = ax + by
int extended_euclid(int a, int b, int &x, int &y) {
    int xx = y = 0;
    int yy = x = 1;
    while (b) {
        int q = a / b;
        int t = b; b = a%b; a = t;
        t = xx; xx = x - q*xx; x = t;
        t = yy; yy = y - q*yy; y = t;
    }
    return a;
}
// finds all solutions to ax = b (mod n)
VI modular_linear_equation_solver(int a, int b, int n) {
    int x, y;
    VI ret;
    int g = extended_euclid(a, n, x, y);
    if (!(b%g)) {
        x = mod(x*(b/g), n);
        for (int i = 0; i < g; i++)
            ret.push_back(mod(x + i*(n/g), n));
    }
    return ret;
}
// computes b such that ab = 1 (mod n), returns -1 on failure
int mod_inverse(int a, int n) {
    int x, y;
    int g = extended_euclid(a, n, x, y);
    if (g > 1) return -1;
    return mod(x, n);
}
// Chinese remainder theorem (special case): find z such
// that z % m1 = r1, z % m2 = r2.
// Here, z is unique modulo M = lcm(m1, m2).
// Return (z, M). On failure, M = -1.
PII chinese_remainder_theorem(int m1, int r1, int m2,
    int r2) {
    int s, t;
    int g = extended_euclid(m1, m2, s, t);
    if (r1%g != r2%g) return make_pair(0, -1);
    return make_pair(mod(s*r2*m1 + t*r1*m2, m1*m2) / g, m1
        *m2 / g);
}
// Find z such that z % m[i] = r[i] for all i.
// The solution is unique modulo M = lcm_i (m[i]).
// Return (z, M). On failure, M = -1.
// We don't require a[i]'s to be relatively prime.
PII chinese_remainder_theorem(const VI &m, const VI &r) {
```

```

PII ret = make_pair(r[0], m[0]);
for (int i = 1; i < m.size(); i++) {
    ret = chinese_remainder_theorem(ret.second, ret.first, m[i], r[i]);
    if (ret.second == -1) break;
}
return ret;
}
// computes x and y such that ax + by = c
// returns whether the solution exists
bool linear_diophantine(int a, int b, int c, int &x, int &y) {
    if (!a && !b) {
        if (c) return false;
        x = 0; y = 0;
        return true;
    }
    if (!a) {
        if (c % b) return false;
        x = 0; y = c / b;
        return true;
    }
    if (!b) {
        if (c % a) return false;
        x = c / a; y = 0;
        return true;
    }
    int g = __gcd(a, b);
    if (c % g) return false;
    x = c / g * mod_inverse(a / g, b / g);
    y = (c - a*x) / b;
    return true;
}

```

## 4 Graph algorithms

### 4.1 Dynamic Connectivity

```

struct UnionFind {
    int n, comp;
    vector<int> uf, si, c;
    UnionFind(int n=0): n(n), comp(n), uf(n), si(n, 1) {
        for (int i=0; i<n; ++i)
            uf[i]=i;
    }
    int find(int x) { return x==uf[x]?x:find(uf[x]); }
    bool join(int x, int y) {
        if ((x=find(x))==(y=find(y))) return false;
        if (si[x]<si[y]) swap(x, y);
        si[x]+=si[y]; uf[y]=x; comp--;
        c.push_back(y);
        return true;
    }
    int snap() { return c.size(); }
    void rollback(int snap) {
        while (c.size()>snap) {
            int x=c.back(); c.pop_back();
            si[uf[x]]-=si[x]; uf[x]=x; comp++;
        }
    }
};

```

```

enum {ADD, DEL, QUERY};
struct Query {int type, x, y;};
struct DynCon {
    vector<Query> q;
    UnionFind dsu;
    vector<int> mt;
    map<pair<int, int>, int> last;
    DynCon(int n): dsu(n) {}
    void add(int x, int y) {
        if (x>y) swap(x, y);
        q.push_back((Query){ADD, x, y});
        last[make_pair(x, y)] = q.size()-1;
        mt[make_pair(x, y)] = q.size()-1;
    }
    void remove(int x, int y) {
        if (x>y) swap(x, y);
        q.push_back((Query){DEL, x, y});
        int pr = last[make_pair(x, y)];
        mt[pr] = q.size()-1;
        mt.push_back(pr);
    }
    void query(int x, int y) {
        q.push_back((Query){QUERY, x, y});
        mt.push_back(-1);
    }
    void process() { // answers all queries in order
        if (!q.size()) return;
        for (int i=0; i<q.size(); ++i)
            if (q[i].type==ADD && mt[i]<0)
                mt[i] = q.size();
        go(0, q.size());
    }
    void go(int l, int r) {
        if (l+1==r) {
            if (q[l].type==QUERY) // answer query using DSU
                puts(dsu.find(q[l].x)==dsu.find(q[l].y)? "YES":"NO");
            return;
        }
        int s=dsu.snap(), m=(l+r)/2;
        for (int i=r-1; i>=m; --i)
            if (mt[i]>=0 && mt[i]<l)
                dsu.join(q[i].x, q[i].y);
        go(l, m);
        dsu.rollback(s);
        for (int i=m-1; i>=l; --i) if (mt[i]>=r) dsu.join(q[i].x, q[i].y);
        go(m, r);
        dsu.rollback(s);
    }
};

```

### 4.2 Bridges

```

// Finds bridges and cut vertices
//
// Receives:
// N: number of vertices
// l: adjacency list
//
// Gives:
// vis, seen, par (used to find cut vertices)

```

```
// ap - 1 if it is a cut vertex, 0 otherwise
// brid - vector of pairs containing the bridges
```

```
typedef pair<int, int> PII;

int N;
vector<int> l[MAX];
vector<PII> brid;
int vis[MAX], seen[MAX], par[MAX], ap[MAX];
int cnt, root;

void dfs(int x) {
    if(vis[x] != -1)
        return;
    vis[x] = seen[x] = cnt++;

    int adj = 0;
    for(int i = 0; i < (int)l[x].size(); i++) {
        int v = l[x][i];
        if(par[x] == v)
            continue;
        if(vis[v] == -1) {
            adj++;
            par[v] = x;
            dfs(v);
            seen[x] = min(seen[x], seen[v]);
            if(seen[v] >= vis[x] && x != root)
                ap[x] = 1;
            if(seen[v] == vis[v])
                brid.push_back(make_pair(v, x));
        }
        else {
            seen[x] = min(seen[x], vis[v]);
            seen[v] = min(seen[x], seen[v]);
        }
    }
    if(x == root) ap[x] = (adj > 1);
}

void bridges() {
    brid.clear();
    for(int i = 0; i < N; i++) {
        vis[i] = seen[i] = par[i] = -1;
        ap[i] = 0;
    }
    cnt = 0;
    for(int i = 0; i < N; i++)
        if(vis[i] == -1) {
            root = i;
            dfs(i);
        }
}
```

### 4.3 Strongly connected components

```
struct SCC {
    int V, group_cnt;
    vector<vector<int>> adj, radj;
    vector<int> group_num, vis;
    stack<int> stk;
    // V = number of vertices
    SCC(int V): V(V), group_cnt(0), group_num(V), vis(V),
        adj(V), radj(V) {}
}
```

```
// Call this to add an edge (0-based)
void add_edge(int v1, int v2) {
    adj[v1].push_back(v2);
    radj[v2].push_back(v1);
}

void fill_forward(int x) {
    vis[x] = true;
    for(int i = 0; i < adj[x].size(); i++) {
        if(!vis[adj[x][i]]) {
            fill_forward(adj[x][i]);
        }
    }
    stk.push(x);
}

void fill_backward(int x) {
    vis[x] = false;
    group_num[x] = group_cnt;
    for(int i = 0; i < radj[x].size(); i++) {
        if(vis[radj[x][i]]) {
            fill_backward(radj[x][i]);
        }
    }
}

// Returns number of strongly connected components.
// After this is called, group_num contains component
// assignments (0-based)
int get_scc() {
    for(int i = 0; i < V; i++) {
        if(!vis[i]) fill_forward(i);
    }
    group_cnt = 0;
    while(!stk.empty()) {
        if(vis[stk.top()]) {
            fill_backward(stk.top());
            group_cnt++;
        }
        stk.pop();
    }
    return group_cnt;
};
```

## 5 String Stuff

### 5.1 Suffix Automaton

```
struct SuffixAutomaton {
    vector<map<char, int>> edges; // edges[i]: the labeled
    // edges from node i
    vector<int> link; // link[i] : the parent of i
    vector<int> length; // length[i]: length of longest
    // string in ith class
    vector<int> cnt; // No. of times substring occurs
    int last; // index of equivalence class of whole
    // string
    SuffixAutomaton(string const& s) {
        // add the initial node
        edges.push_back(map<char, int>());
        link.push_back(-1);
        length.push_back(0);
    }
}
```



```

cnt.push_back(0);
last = 0;
for(int i=0; i<s.size(); i++) {
    // construct r
    edges.push_back(map<char,int>());
    length.push_back(i+1);
    link.push_back(0);
    cnt.push_back(1);
    int r = edges.size() - 1;
    // add edges to r and find p with link to q
    int p = last;
    while(p >= 0 && !edges[p].count(s[i])) {
        edges[p][s[i]] = r;
        p = link[p];
    }
    if(p != -1) {
        int q = edges[p][s[i]];
        if(length[p] + 1 == length[q]) {
            // we do not have to split q, just set the
            // correct suffix link
            link[r] = q;
        } else {
            // we have to split, add q'
            edges.push_back(edges[q]); // copy edges of q
            length.push_back(length[p] + 1);
            link.push_back(link[q]); // copy parent of q
            cnt.push_back(0);
            int qq = edges.size()-1;
            // add qq as the new parent of q and r
            link[q] = link[r] = qq; cnt[r] = 1;
            // move short classes pointing to q to point
            // to q'
            while(p >= 0 && edges[p][s[i]] == q) {
                edges[p][s[i]] = qq;
                p = link[p];
            }
        }
    }
    last = r;
}

vector<int> ind(length.size());
iota(ind.begin(), ind.end(), 0);
sort(ind.begin(), ind.end(), [&](int i, int j){
    return length[i] > length[j];
});
for(auto i:ind) if(link[i] >= 0)
    cnt[link[i]] += cnt[i];
};

```

## 5.2 Suffix array

```

vector<int> suffix_array(string &A) {
    int n=A.size(), i=n, *M=new int[5*n];
    int *B=M, *C=M+n, *F=M+2*n, *G=M+3*n, *S=M+4*n;
    for(; i-->0; S[i]=n-i-1) B[i]=A[i];
    stable_sort(S, S+n, [&](int i, int j) {return A[i]<A[j];});
    for(int L=1, p; L<n; L*=2) {
        for(++i<n; F[i]=B[S[i]], G[i]=B[S[i]+L/2]);
    }
}

```

```

for(; --i; F[i]=F[i]==F[i-1]&&G[i]==G[i-1]&&S[i-1]<n-L);
for(p=B[*S]=0; ++i<n; B[S[i]]=p=F[i]?p:i);
for(fill_n(G, n, 0); i-->0; F[i]=S[i]<L?-1:B[S[i]-L]);
for(iota(C, C+n, 0); ++i<n; ~F[i]?G[i]=C[F[i]]++:0);
for(copy_n(S, n, F); i-->0; F[i]<L?0:S[G[i]]=F[i]-L);
}
vector<int> res(S, S+n);
delete[] M;
return res;
}

vector<int> kasai(string &s, vector<int> &sa) {
    int n = s.size();
    vector<int> lcp(n), inv(n);
    for(int i=0; i<n; ++i) inv[sa[i]] = i;
    for(int i=0, k=0; i<n; ++i) {
        if(k<0) k = 0;
        if(inv[i]==n-1) { k=0; continue; }
        for(int j=sa[inv[i]+1]; max(i, j)+k<n&&s[i+k]==s[j+k]; ++k);
        lcp[inv[i]] = k--;
    }
    return lcp;
}

```

## 5.3 Z Algorithm

```

vector<int> compute_Z(string s) {
    int n = s.length();
    vector<int> z(n, 0);
    z[0] = n;
    for (int i = 1, l = 0, r = 0; i < n; ++i) {
        if (r >= i) {
            z[i] = min(z[i - l], r - i + 1);
        }
        while (i+z[i] < n and s[i+z[i]] == s[z[i]]) {
            ++z[i];
        }
        if (i + z[i] - 1 > r) {
            r = i + z[i] - 1;
            l = i;
        }
    }
    return z;
}

```

## 5.4 KMP

```

vector<int> find_prefix(const vector<int> &P) {
    int M = P.size();
    vector<int> pi(M);
    /* pi[i] <- largest prefix P[0..pi[i]] which is a
       suffix of P[0..i]
       * (but not equal to it) */
    pi[0] = -1;
    for (int i = 1, k = -1; i < M; ++i) {
        while(k > -1 && P[k + 1] != P[i])
            k = pi[k];
        if (P[k + 1] == P[i]) ++k;
        pi[i] = k;
    }
}

```

```

    return pi;
}
int kmp_matcher(const vector<int> &T, const vector<int>
&P){
    int M = P.size(), N = T.size();
    vector<int> pi = find_prefix(P);
    int q = -1, matches = 0;
    for (int i = 0; i < N; ++i) {
        while(q > -1 && P[q + 1] != T[i])
            q = pi[q];
        if (P[q + 1] == T[i]) ++q;
        if (q == M - 1)
            ++matches, q = pi[q];
    }
    return matches;
}

```

## 5.5 String Hashing

```

struct hasher{
    int hashes[MAXN+5];
    int *pow, *inv;
    int mod;
    int n;
    void init(string &str, int *p, int *i, int m){
        pow = p;
        inv = i;
        mod = m;
        n = str.size();
        int last = 0;
        for (int i = 0; i < n; i++){
            int c = str[i] - 'a' + 1;
            last = (last + 1ll*c*pow[i]) % mod;
            hashes[i] = last;
        }
    }
    int getHash(int l, int r){
        if (r >= n || l < 0)
            return -1;
        int curr = hashes[r] - (l-1 >= 0 ? hashes[l-1] : 0);
        curr = ((curr % mod) + mod) % mod;
        curr = (1ll*curr*inv[l]) % mod;
        return curr;
    }
} A, B, C;

```

## 5.6 Palindrome DSU

```

// given an unknown string s and Q ranges that
// are known to be palindromes, this computes the
// characters
// that have to be equal in O(Q + n log n)
struct Palindrome{
    Palindrome(){}
    Palindrome(int n_):n(n_), m(3+__lg(n)), qs(m), p(2*n){
        iota(p.begin(), p.end(), 0);
    }
    int f(int i){
        return p[i] == i ? i : p[i] = f(p[i]);
    }
    void u(int a, int b){

```

```

assert(0 <= a && a < 2*n);
assert(0 <= b && b < 2*n);
// union with splicing is a bit faster than
// just path compression also guarantees p[i] <= i
while(p[a] != p[b]){
    if(p[a] < p[b]) swap(a, b);
    if(p[a] == a){
        p[a] = b;
        return;
    }
    int tmp = p[a];
    p[a] = p[b];
    a = p[tmp];
}
}
int components(){
    int ret = 0;
    for(int i=0; i<(int)p.size(); ++i)
        if(p[i] == i) ++ret;
    return ret;
}
// call this after adding all queries
void compute(){
    vector<int> p2(2*n);
    for(int l=m-1; l>=0; --l){
        const int s = 1<<l;
        for(int i=0; i<2*n; ++i) p2[i] = f(i);
        for(int i=0; i+s<2*n; ++i){
            const int j = p2[i];
            if(j+s < 2*n) u(i+s, j+s);
        }
        for(auto const&e:qs[l])
            u(e.first, e.second);
    }
    // link point with mirror-image
    for(int i=0; i<n; ++i)
        u(i, 2*n-1-i);
}
// force [l, r] to be a palindrome
void add_q(int l, int r){
    assert(0 <= l && l <= r && r < n);
    if(l==r) return;
    const int range = r-l+1;
    const int k = __lg(range);
    qs[k].emplace_back(l, 2*n-1-r);
}
int n, m;
vector<vector<pair<int, int>>> qs;
vector<int> p;
};

```

## 5.7 Eertree

```

/*
Palindrome tree. Useful structure to deal with
palindromes in strings. O(N)
This code counts no. of palindrome substrings of string.
Based on problem 1750 from informatics.mccme.ru:
http://informatics.mccme.ru/moodle/mod/statements/view.
php?chapterid=1750
*/

```

```

const int MAXN = 105000;
struct node {
    int next[26];
    int len;
    int sufflink;
    int num;
};
int len;
char s[MAXN];
node tree[MAXN];
int num;           // node 1 - root with len -1, node 2
                  // - root with len 0
int suff;          // max suffix palindrome
long long ans;

bool addLetter(int pos) {
    int cur = suff, curlen = 0;
    int let = s[pos] - 'a';
    while (true) {
        curlen = tree[cur].len;
        if (pos - 1 - curlen >= 0 && s[pos - 1 - curlen]
            == s[pos])
            break;
        cur = tree[cur].sufflink;
    }
    if (tree[cur].next[let]) {
        suff = tree[cur].next[let];
        return false;
    }
    num++;
    suff = num;
    tree[num].len = tree[cur].len + 2;
    tree[cur].next[let] = num;
    if (tree[num].len == 1) {
        tree[num].sufflink = 2;
        tree[num].num = 1;
        return true;
    }
    while (true) {
        cur = tree[cur].sufflink;
        curlen = tree[cur].len;
        if (pos - 1 - curlen >= 0 && s[pos - 1 - curlen]
            == s[pos]) {
            tree[num].sufflink = tree[cur].next[let];
            break;
        }
    }
    tree[num].num = 1 + tree[tree[num].sufflink].num;
    return true;
}

void initTree() {
    num = 2; suff = 2;
    tree[1].len = -1; tree[1].sufflink = 1;
    tree[2].len = 0; tree[2].sufflink = 1;
}

int main() {
    gets(s);
    len = strlen(s);
    initTree();
    for (int i = 0; i < len; i++) {
        addLetter(i);

```

```

        ans += tree[suff].num;
    }
    cout << ans << endl;
}

```

## 6 Data structures

### 6.1 BIT Range Queries

```

struct BIT {
    int n;
    vector<int> slope;
    vector<int> intercept;

    // BIT can be thought of as having entries f[1], ...,
    // f[n] which are 0-initialized
    BIT(int n): n(n), slope(n+1), intercept(n+1) {}
    // returns f[1] + ... + f[idx-1]
    // precondition idx <= n+1

    int query(int idx) {
        int m = 0, b = 0;
        for (int i = idx-1; i > 0; i -= i&-i) {
            m += slope[i];
            b += intercept[i];
        }
        return m*idx + b;
    }

    // adds amt to f[i] for i in [idx1, idx2)
    // precondition 1 <= idx1 <= idx2 <= n+1 (you can't
    // update element 0)
    void update(int idx1, int idx2, int amt) {
        for (int i = idx1; i <= n; i += i&-i) {
            slope[i] += amt;
            intercept[i] -= idx1*amt;
        }
        for (int i = idx2; i <= n; i += i&-i) {
            slope[i] -= amt;
            intercept[i] += idx2*amt;
        }
    }
};

```

### 6.2 Treaps

```

typedef struct node {
    int prior, size;
    int val; //value stored in the array
    int sum; //whatever info you want to maintain in
             //segtree for each node
    int lazy; //whatever lazy update you want to do
    struct node *l, *r;
} node;

struct Treap {
    typedef node* pnode;
    int sz(pnode t) {
        return t?t->size:0;
    }
}

```

```

void upd_sz(pnode t) {
    if(t) t->size=sz(t->l)+1+sz(t->r);
}
void lazy(pnode t) {
    if(!t || !t->lazy) return;
    t->val+=t->lazy; //operation of lazy
    t->sum+=t->lazy*sz(t);
    if(t->l) t->l->lazy+=t->lazy; //propagate lazy
    if(t->r) t->r->lazy+=t->lazy;
    t->lazy=0;
}
void reset(pnode t) {
    if(t) t->sum = t->val; //no need to reset lazy coz
        when we call this lazy would itself be propagated
}
void combine(pnode& t, pnode l, pnode r) { //combining two
    ranges of segtree
    if(!l || !r) return void(t = l?l:r);
    t->sum = l->sum + r->sum;
}
void operation(pnode t) { //operation of segtree
    if(!t) return;
    reset(t); //reset the value of current node assuming
        it now represents a single element of the array
    lazy(t->l); lazy(t->r); //imp: propagate lazy before
        combining t->l, t->r;
    combine(t, t->l, t);
    combine(t, t, t->r);
}
void split(pnode t, pnode &l, pnode &r, int pos, int add
    =0) {
    if(!t) return void(l=r=NULL);
    lazy(t);
    int curr_pos = add + sz(t->l);
    if(curr_pos<=pos) //element at pos goes to left
        subtree(l)
        split(t->r, t->r, r, pos, curr_pos+1), l=t;
    else
        split(t->l, l, t->l, pos, add), r=t;
    upd_sz(t);
    operation(t);
}
void merge(pnode &t, pnode l, pnode r) { //l->leftarray, r
    ->rightarray, t->resulting array
    lazy(l); lazy(r);
    if(!l || !r) t = l?l:r;
    else if(l->prior>r->prior) merge(l->r, l->r, r), t=l;
    else merge(r->l, l, r->l), t=r;
    upd_sz(t);
    operation(t);
}
pnode init(int val) {
    pnode ret = (pnode) malloc(sizeof(node));
    ret->prior=rand(); ret->size=1;
    ret->val=val;
    ret->sum=val; ret->lazy=0;
    return ret;
}
int range_query(pnode t, int l, int r) { //[l, r]
    pnode L, mid, R;

```

```

    split(t, L, mid, l-1);
    split(mid, t, R, r-1); //note: r-1!!
    int ans = t->sum;
    merge(mid, L, t);
    merge(t, mid, R);
    return ans;
}
void range_update(pnode t, int l, int r, int val) { //[l, r]
    pnode L, mid, R;
    split(t, L, mid, l-1);
    split(mid, t, R, r-1); //note: r-1!!
    t->lazy+=val; //lazy_update
    merge(mid, L, t);
    merge(t, mid, R);
}
};

```

### 6.3 Link-Cut Tree

```

const int MXN = 100005, MEM = 100005;
struct Splay {
    static Splay nil, mem[MEM], *pmem;
    Splay *ch[2], *f;
    int val, rev, size;
    Splay (int _val=-1) : val(_val), rev(0), size(1)
    { f = ch[0] = ch[1] = &nil; }
    bool isr()
    { return f->ch[0] != this && f->ch[1] != this; }
    int dir()
    { return f->ch[0] == this ? 0 : 1; }
    void setCh(Splay *c, int d) {
        ch[d] = c;
        if (c != &nil) c->f = this;
        pull();
    }
    void push() {
        if (!rev) return;
        swap(ch[0], ch[1]);
        if (ch[0] != &nil) ch[0]->rev ^= 1;
        if (ch[1] != &nil) ch[1]->rev ^= 1;
        rev=0;
    }
    void pull() {
        size = ch[0]->size + ch[1]->size + 1;
        if (ch[0] != &nil) ch[0]->f = this;
        if (ch[1] != &nil) ch[1]->f = this;
    }
} Splay::nil, Splay::mem[MEM], *Splay::pmem = Splay::mem;

Splay *nil = &Splay::nil;
void rotate(Splay *x) {
    Splay *p = x->f;
    int d = x->dir();
    if (!p->isr()) p->f->setCh(x, p->dir());
    else x->f = p->f;
    p->setCh(x->ch[!d], d); x->setCh(p, !d); p->pull();
    x->pull();
}
vector<Splay*> splayVec;
void splay(Splay *x) {
    splayVec.clear();

```

```

for(Splay *q=x;; q=q->f) {
    splayVec.push_back(q);
    if (q->isr()) break;
}
reverse(begin(splayVec), end(splayVec));
for (auto it : splayVec) it->push();
while (!x->isr())
    if (x->f->isr()) rotate(x);
    else if (x->dir() == x->f->dir()) rotate(x->f), rotate(x);
    else rotate(x), rotate(x);
}
int id(Splay *x) { return x - Splay::mem + 1; }
Splay* access(Splay *x) {
    Splay *q = nil;
    for (; x != nil; x = x->f) { splay(x); x->setCh(q, 1); q = x; }
    return q;
}
void chroot(Splay *x) {
    access(x); splay(x); x->rev ^= 1; x->push(); x->pull();
}
void link(Splay *x, Splay *y) {
    access(x); splay(x);
    chroot(y); x->setCh(y, 1);
}
void cut_p(Splay *y) {
    access(y); splay(y); y->push(); y->ch[0] = y->ch[0]->f
    = nil;
}
void cut(Splay *x, Splay *y) {
    chroot(x); cut_p(y);
}
Splay* get_root(Splay *x) {
    access(x); splay(x);
    for (; x->ch[0] != nil; x = x->ch[0]) x->push();
    splay(x); return x;
}
bool conn(Splay *x, Splay *y) {
    return (x = get_root(x)) == (y = get_root(y));
}
Splay* lca(Splay *x, Splay *y) {
    access(x); access(y); splay(x);
    if (x->f == nil) return x;
    else return x->f;
}

```

```

Two_Sat(int N_) : N(N_) { // create a formula over N
    variables (numbered 1 to N)
    G.resize(2*N);
}
int add_variable() {
    G.emplace_back(); G.emplace_back();
    return N++;
}
private:
    // converts a signed variable index to its position in
    // val[] and G[]
    int to_ind(int x) {
        return 2*(abs(x)-1) + (x<0);
    }
    // Add a directed edge to the graph.
    // You most likely do not want to call this yourself!
    void add_edge(int a, int b) {
        G[to_ind(a)].push_back(to_ind(b));
    }
    int time() {
        return valid.size()-1;
    }
    bool dfs(int x) {
        if (valid[abs(val[x])]) return val[x]>0;
        val[x] = time();
        val[x^1] = -time();
        for (int e : G[x])
            if (!dfs(e))
                return false;
        return true;
    }
public:
    // Add the or-clause: (a or b)
    void add_or(int a, int b) {
        add_edge(-a, b); add_edge(-b, a);
    }
    // Add the implication: a -> b
    void add_implication(int a, int b) {
        add_or(-a, b);
    }
    // Add condition: x is true
    void add_true(int x) {
        add_or(x, x);
    }
    // At most one with linear number of clauses
    template<typename T>
    void add_at_most_one(T vars) {
        if (vars.begin() == vars.end()) return;
        int last = *vars.begin(), cur = 0;
        for (int const& e : vars) {
            if (e == last) continue;
            if (cur == 0) cur = e;
            else {
                add_or(-cur, -e);
                int new_cur = add_variable();
                cur = add_implication(cur, new_cur);
                add_implication(e, new_cur);
                cur = new_cur;
            }
        }
        if (cur != 0) add_or(-cur, -last);
    }
}

```

## 7 Miscellaneous

### 7.1 2-SAT

```

class Two_Sat {
    int N; // number of variables
    vector<int> val; // assignment of x is at val[2x] and
    // -x at val[2x+1]
    vector<char> valid; // changes made at time i are kept
    // iff valid[i]
    vector<vector<int>> G; // graph of implications G[x][
    // i] = y means (x -> y)
}

```

```

    }
    bool solve() {
        val.assign(2*n, 0); valid.assign(1, 0);
        for(int i=0; i<val.size(); i+=2) {
            if(!valid[abs(val[i])]) {
                valid.push_back(1);
                if(!dfs(i)) {
                    valid.back()=0;
                    valid.push_back(1);
                    if(!dfs(i+1)) return false;
                }
            }
        }
        return true;
    }
};

// Taken from https://github.com/dacin21/
// dacin21_codebook/blob/master/dfs_stuff/2sat.cpp
// 2-sat in linear time via backtracking.
class Two_Sat {
    int N; // number of variables
    vector<int> val; // assignment of x is at val[2x]
    // and -x at val[2x+1]
    vector<char> valid; // changes made at time i are
    // kept iff valid[i]
    vector<vector<int>> G; // graph of implications G[x
    ][i] = y means (x -> y)

    Two_Sat(int N_) : N(N_) { // create a formula over N
        // variables (numbered 1 to N)
        G.resize(2*N);
    }

    int add_variable() {
        G.emplace_back();
        G.emplace_back();
        return N++;
    }

private:
    // converts a signed variable index to its position
    // in val[] and G[]
    int to_ind(int x) {
        return 2*(abs(x)-1) + (x<0);
    }

    // Add a directed edge to the graph.
    // You most likely do not want to call this yourself
    // !
    void add_edge(int a, int b) {
        G[to_ind(a)].push_back(to_ind(b));
    }

    int time() {
        return valid.size()-1;
    }

    bool dfs(int x) {
        if(valid[abs(val[x])]) return val[x]>0;
        val[x] = time();
        val[x^1] = -time();
        for(int e:G[x])
            if(!dfs(e))
                return false;
    }
};

```

```

        return true;
    }

public:
    // Add the or-clause: (a or b)
    void add_or(int a, int b) {
        add_edge(-a,b);
        add_edge(-b,a);
    }

    // Add the implication: a -> b
    void add_implication(int a, int b) {
        add_or(-a, b);
    }

    // Add condition: x is true
    void add_true(int x) {
        add_or(x,x);
    }

    // At most one with linear number of clauses
    template<typename T>
    void add_at_most_one(T vars) {
        if(vars.begin() == vars.end()) return;
        int last = *vars.begin();
        int cur = 0;
        for(int const&e:vars){
            if(e == last) continue;
            if(cur == 0) cur = e;
            else {
                add_or(-cur, -e);
                int new_cur = add_variable();
                cur = add_implication(cur, new_cur);
                add_implication(e, new_cur);
                cur = new_cur;
            }
        }
        if(cur != 0){
            add_or(-cur, -last);
        }
    }

    bool solve() {
        val.assign(2*n, 0);
        valid.assign(1, 0);
        for(int i=0; i<val.size(); i+=2) {
            if(!valid[abs(val[i])]) {
                valid.push_back(1);
                if(!dfs(i)) {
                    valid.back()=0;
                    valid.push_back(1);
                    if(!dfs(i+1)) return false;
                }
            }
        }
        return true;
    }
};

```

## 7.2 Merge Insertion

// Sorting in  $O(n^2)$  time with near-optimal number of comparisons



```

// Number of comparisons used is:  $n \lg n - 1.415 n$ 
// The lower bound is:  $\lg(n!) = n \lg n - 1.443 n$ 
// Binary search insertion sort would need:  $n \log n - n$ 
struct Merge_Insertion_Sort{
    template<typename F>
    static void apply_permutation(vector<int> const&p, F
        get){
        const int n = p.size();
        vector<int> q(n);
        for(int i=0; i<n; ++i) q[p[i]] = i;
        for(int i=0; i<n; ++i){
            while(q[i] != i){
                swap(get(i), get(q[i]));
                swap(q[i], q[q[i]]);
            }
        }
    }
    // ret.first is the sorted vector ret.second[i] is the
    // index of the i-th smallest element in the original
    // vector
    // i.e. ret.second is the permutation that was applied
    // to sort
    template<typename T, typename F>
    static pair<vector<T>, vector<int>> sort(vector<T> v,
        F comp){
        const int n = v.size();
        if(n <= 1) return {move(v), {{0}}};
        vector<int> preperm(n);
        iota(preperm.begin(), preperm.end(), 0);
        const int M = n-n/2;
        for(int i=0; i<n/2; ++i){
            if(comp(v[M+i], v[i])){
                swap(v[M+i], v[i]);
                swap(preperm[M+i], preperm[i]);
            }
        }
        auto ret=sort(vector<T>(v.begin(), v.begin()+n/2),
            comp);
        apply_permutation(ret.second, [&preperm, M](int const&
            i)->int&{return preperm[M+i];});
        apply_permutation(ret.second, [&preperm](int const&i)
            ->int&{return preperm[i];});
        apply_permutation(ret.second, [&v, M](int const&i)->T
            &{return v[M+i];});
        iota(ret.second.begin(), ret.second.end(), 0);
        // insert one element without comparisons
        ret.first.push_back(v.back());
        ret.second.push_back(n-1);
        // now insert the rest in blocks that optimize the
        // binary search sizes
        for(int it=1, r=n-1, s=2; r>n/2; ++it, r-=s, s=(1<<it)
            -s){
            for(int i=r-s; i<r; ++i){
                if(i>=n/2){
                    int a = find(ret.second.begin(), ret.second.
                        end(), i-M) - ret.second.begin();
                    int b = ret.first.size();
                    if(a==b){
                        assert(i==M-1);
                        a = -1;
                    }
                }
            }
        }
    }
};

```

```

        while(a+1 < b){
            const int m = a+(b-a)/2;
            if(comp(ret.first[m], v[i])) a = m;
            else b = m;
        }
        ret.first.insert(ret.first.begin()+b, v[i]);
        ret.second.insert(ret.second.begin()+b, i);
    }
}
// compose permutations
apply_permutation(ret.second, [&preperm](int const&i
    )->int&{return preperm[i];});
ret.second.swap(preperm);
return ret;
}
template<typename T>
static pair<vector<T>, vector<int>> sort(vector<T> v)
{
    return sort(move(v), std::less<T>{});
}
};

```

## 7.3 DP Optimizations

$A[i][j]$  : The smallest  $k$  that gives optimal answer

Divide and Conquer:

$dp[i][j] = \min(k < j) \{dp[i-1][k] + C[k][j]\}$   
 $O(kn^2) \rightarrow O(kn \log(n))$

Conditions:

$A[i][j] \leq A[i][j+1]$  OR  
 $C[a][d] + C[b][c] \geq C[a][c] + C[b][d]$  where  $a < b < c < d$

Short Description:

$A[i][1] \leq A[i][2] \leq \dots \leq A[i][n]$

Knuth Optimization:

$dp[i][j] = \min(i < k < j) \{dp[i][k] + dp[k][j]\} + C[i][j]$   
 $O(n^3) \rightarrow O(n^2)$

Conditions:

$A[i, j-1] \leq A[i, j] \leq A[i+1, j]$  OR  
 $C[a][d] + C[b][c] \geq C[a][c] + C[b][d]$  AND  
 $C[b][c] \leq C[a][d]$  where  $a \leq b \leq c \leq d$

Short Description:

For  $dp[i][j]$ , loop  $k$  from  $A[i][j-1]$  to  $A[i+1][j]$

## 7.4 Convex Hull Trick (Dynamic)

```

struct Line {
    long long m, b;
    mutable function<const Line*> succ;
    bool operator<(const Line& rhs) const{
        if(rhs.b!==(1ll<<62)) return m>rhs.m; // < for max
        const Line* s = succ();
        if(!s) return 0;
        return b-s->b > (s->m -m)*rhs.m; // < for max
    }
};

```

```

struct HullDynamic : public multiset<Line> {
    bool bad(iterator y) {
        auto z = next(y);
        if(y==begin()){
            if(z==end()) return 0;
            return y->m == z->m && y->b >= z->b; // <= for max
        }
        auto x = prev(y);
        if (z==end()) return y->m == x->m && y->b >= x->b;
        // <= for max
        return (x->b - y->b)*1.0*(z->m - y->m) >= (y->b - z->b)*1.0*(y->m - x->m);
    }
    void insert_line(long long m, long long b) {
        auto y = insert({ m, b });
        y->succ = [=] { return next(y)==end()? 0:&*next(y); };
        if(bad(y)) { erase(y); return; }
        while(next(y)!=end() && bad(next(y))) erase(next(y));
        while(y!=begin() && bad(prev(y))) erase(prev(y));
    }
    long long eval(long long x) {
        auto l = *lower_bound((Line){x,-(1ll<<62)});
        return l.m * x + l.b;
    }
};

```

## 7.5 Convex Hull Trick (Static)

```

struct ConvexHullTrick {
    typedef long long LL;
    vector<LL> M;
    vector<LL> B;
    vector<double> left;
    ConvexHullTrick() {}
    bool bad(LL m1, LL b1, LL m2, LL b2, LL m3, LL b3) {
        // Careful, this may overflow
        return (b3-b1)*(m1-m2) < (b2-b1)*(m1-m3);
    }
    // Add a new line to the structure, y = mx + b.
    // Lines must be added in decreasing order of slope.
    void add(LL m, LL b) {
        while (M.size() >= 2 && bad(M[M.size()-2], B[B.size()-2], M.back(), B.back(), m, b)) {
            M.pop_back(); B.pop_back(); left.pop_back();
        }
        if (M.size() && M.back() == m) {
            if (B.back() > b) {
                M.pop_back(); B.pop_back(); left.pop_back();
            } else {
                return;
            }
        }
        if (M.size() == 0) {
            left.push_back(-numeric_limits<double>::infinity());
        } else {
            left.push_back((double) (b - B.back()) / (M.back() - m));
        }
        M.push_back(m);
        B.push_back(b);
    }
};

```

```

}
// Get the minimum value of mx + b among all lines in
// the structure.
// There must be at least one line.
LL query(LL x) {
    int i = upper_bound(left.begin(), left.end(), x) -
        left.begin();
    return M[i-1]*x + B[i-1];
};

```

## 7.6 BigInt library

```

struct bignum {
    typedef unsigned int uint;
    vector<uint> digits;
    static const uint RADIX = 1000000000;
    bignum(): digits(1, 0) {}
    bignum(const bignum& x): digits(x.digits) {}
    bignum(unsigned long long x) { *this = x; }
    bignum(const char* x) { *this = x; }
    bignum(const string& s) { *this = s; }
    bignum& operator=(const bignum& y) {
        digits = y.digits; return *this;
    }
    bignum& operator=(unsigned long long x) {
        digits.assign(1, x%RADIX);
        if (x >= RADIX)
            digits.push_back(x/RADIX);
        return *this;
    }
    bignum& operator=(const char* s) {
        int slen=strlen(s), i, l;
        digits.resize((slen+8)/9);
        for (l=0; slen>0; l++, slen-=9) {
            digits[l]=0;
            for (i=slen>9?slen-9:0; i<slen; i++)
                digits[l]=10*digits[l]+s[i]-'0';
        }
        while (digits.size() > 1 && !digits.back()) digits.pop_back();
        return *this;
    }
    bignum& operator=(const string& s) {
        return *this = s.c_str();
    }
    void add(const bignum& x) {
        int l = max(digits.size(), x.digits.size());
        digits.resize(l+1);
        for (int d=0, carry=0; d<=l; d++) {
            uint sum=carry;
            if (d<digits.size()) sum+=digits[d];
            if (d<x.digits.size()) sum+=x.digits[d];
            digits[d]=sum;
            if (digits[d]>=RADIX)
                digits[d]-=RADIX, carry=1;
            else
                carry=0;
        }
        if (!digits.back()) digits.pop_back();
    }
    void sub(const bignum& x) {

```

```

// if ((*this)<x) throw; //negative numbers not yet
// supported
for (int d=0, borrow=0; d<digits.size(); d++) {
    digits[d]-=borrow;
    if (d<x.digits.size()) digits[d]-=x.digits[d];
    if (digits[d]>>31) { digits[d]+=RADIX; borrow=1; }
    else borrow=0;
}
while (digits.size() > 1 && !digits.back()) digits.
    pop_back();
}
void mult(const bignum& x) {
    vector<uint> res(digits.size() + x.digits.size());
    unsigned long long y,z;
    for (int i=0; i<digits.size(); i++) {
        for (int j=0; j<x.digits.size(); j++) {
            unsigned long long y=digits[i]; y*=x.digits[j];
            unsigned long long z=y/RADIX;
            res[i+j+1]+=z; res[i+j]+=y-RADIX*z; //mod is
            slow
            if (res[i+j] >= RADIX) { res[i+j] -= RADIX; res[
                i+j+1]++; }
            for (int k = i+j+1; res[k] >= RADIX; res[k] -=
                RADIX, res[+k]++);
        }
    }
    digits = res;
    while (digits.size() > 1 && !digits.back()) digits.
        pop_back();
}
// returns the remainder
bignum div(const bignum& x) {
    bignum dividend(*this);
    bignum divisor(x);
    fill(digits.begin(), digits.end(), 0);
    // shift divisor up
    int pwr = dividend.digits.size() - divisor.digits.
        size();
    if (pwr > 0) {
        divisor.digits.insert(divisor.digits.begin(), pwr,
            0);
    }
    while (pwr >= 0) {
        if (dividend.digits.size() > divisor.digits.size()
            ) {
            unsigned long long q = dividend.digits.back();
            q *= RADIX; q += dividend.digits[dividend.digits
                .size()-2];
            q /= 1+divisor.digits.back();
            dividend -= divisor*q; digits[pwr] = q;
            if (dividend >= divisor) { digits[pwr]++;
                dividend -= divisor; }
            assert(dividend.digits.size() <= divisor.digits.
                size()); continue;
        }
        while (dividend.digits.size() == divisor.digits.
            size()) {
            uint q = dividend.digits.back() / (1+divisor.
                digits.back());
            if (q == 0) break;
            digits[pwr] += q; dividend -= divisor*q;

```

```

        }
        if (dividend >= divisor) { dividend -= divisor;
            digits[pwr]++; }
        pwr--; divisor.digits.erase(divisor.digits.begin()
            );
    }
    while (digits.size() > 1 && !digits.back()) digits.
        pop_back();
    return dividend;
}
string to_string() const {
    ostringstream oss;
    oss << digits.back();
    for (int i = digits.size() - 2; i >= 0; i--) {
        oss << setfill('0') << setw(9) << digits[i];
    }
    return oss.str();
}
bignum operator+(const bignum& y) const
    {bignum res(*this); res.add(y); return res;}
bignum operator-(const bignum& y) const
    {bignum res(*this); res.sub(y); return res;}
bignum operator*(const bignum& y) const
    {bignum res(*this); res.mult(y); return res;}
bignum operator/(const bignum& y) const
    {bignum res(*this); res.div(y); return res;}
bignum operator%(const bignum& y) const
    {bignum res(*this); return res.div(y);}
bignum& operator+=(const bignum& y)
    {add(y); return *this;}
bignum& operator-=(const bignum& y)
    {sub(y); return *this;}
bignum& operator*=(const bignum& y)
    {mult(y); return *this;}
bignum& operator/=(const bignum& y)
    {div(y); return *this;}
bignum& operator%=(const bignum& y)
    {*this = div(y);}
bool operator==(const bignum& y)
    {return digits == y.digits;}
bool operator<(const bignum& y) const {
    if (digits.size() < y.digits.size()) return true;
    if (digits.size() > y.digits.size()) return false;
    for (int i = digits.size()-1; i >= 0; i--)
        if (digits[i] < y.digits[i])
            return true;
        else if (digits[i] > y.digits[i])
            return false;
    return false;
}
bool operator>(const bignum& y) const
    {return y<*this;}
bool operator<=(const bignum& y) const
    {return !(y<*this);}
bool operator>=(const bignum& y) const
    {return !(*this<y);}
};

```

## 7.7 Manachers algorithm

```
// Ex: "opposes" -> [0,1,0,1,4,1,0,1,0,1,0,3,0,1,0]
vector<int> fastLongestPalindromes(string str) {
    int i=0,j,d,s,e,lLen,pallen=0;
    vector<int> res;
    while (i < str.length()) {
        if (i > pallen && str[i-pallen-1] == str[i]) {
            pallen += 2; i++; continue;
        }
        res.push_back(pallen);
        s = res.size()-2;
        e = s-pallen;
        bool b = true;
        for (j=s; j>e; j--) {
            d = j-e-1;
            if (res[j] == d) { pallen = d; b = false; break; }
            res.push_back(min(d, res[j]));
        }
        if (b) { pallen = 1; i++; }
    }
    res.push_back(pallen);
    lLen = res.size();
    s = lLen-2;
    e = s-(2*str.length()+1-lLen);
    for (i=s; i>e; i--) { d = i-e-1; res.push_back(min(d,
        res[i])); }
    return res;
}
```

## 7.8 Dates

```
// Months are expressed as integers from 1 to 12, Days
// are expressed as integers from 1 to 31, and Years are
// expressed as 4-digit integers.
string dayOfWeek[] = {"Mon", "Tue", "Wed", "Thu", "Fri",
    "Sat", "Sun"};

//converts Gregorian date to integer(Julian day number)
int dateToInt (int m, int d, int y){
    return
        1461 * (y + 4800 + (m - 14) / 12) / 4 +
        367 * (m - 2 - (m - 14) / 12 * 12) / 12 -
        3 * ((y + 4900 + (m - 14) / 12) / 100) / 4 +
        d - 32075;
}

// converts integer (Julian day number) to Gregorian
// date: month/day/year
void intToDate (int jd, int &m, int &d, int &y){
    int x, n, i, j;
    x = jd + 68569;
    n = 4 * x / 146097;
    x -= (146097 * n + 3) / 4;
    i = (4000 * (x + 1)) / 1461001;
    x -= 1461 * i / 4 - 31;
    j = 80 * x / 2447;
    d = x - 2447 * j / 80;
    x = j / 11;
    m = j + 2 - 12 * x;
    y = 100 * (n - 49) + i + x;
}

// converts integer (Julian day number) to day of week
string intToDay (int jd){
```

```
    return dayOfWeek[jd % 7];
}
```

## 7.9 Bitset (Text)

Remember `_Find_first()` and `_Find_next()` exist, and run in  $O(N/W)$ , where  $W$  is word size of machine.

## 7.10 Template

```
g++ -std=c++17 -DLOCAL -O2 -Wall -Wshadow -Wextra -
    pedantic -Wfloat-equal -Wlogical-op

#pragma comment(linker, "/stack:200000000")
#pragma GCC optimize("Ofast")
#pragma GCC optimize("unroll-loops")
#pragma GCC target("sse,sse2,sse3,ssse3,sse4,popcnt,abm,
    mmx,avx,tune=native") // codeforces
//#pragma GCC target("avx,avx2,fma")
//#pragma GCC target("sse,sse2,sse3,ssse3,sse4,popcnt,
    abm,mmx,tune=native") // yandex

#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
#include <bits/stdc++.h>

using namespace __gnu_pbds;
using namespace std;

typedef long double ld;
typedef long long ll;
typedef pair<int, int> pii;
typedef pair<ll, ll> pll;

mt19937 rng(std::chrono::duration_cast<std::chrono::
    nanoseconds>(chrono::high_resolution_clock::now().
    time_since_epoch()).count());

template<typename has_less>
using ordered_set =
    tree<has_less,
        null_type,
        less<has_less>,
        rb_tree_tag,
        tree_order_statistics_node_update>;

//insert using pref_trie.insert
//get range for prefix using pref_trie.prefix_range
//use iterator from range.first until != range.second
typedef
    trie<string,
        null_type,
        trie_string_access_traits<>,
        pat_trie_tag,
        trie_prefix_search_node_update>
    pref_trie;

struct chash {
    int operator()(int x) const { return x ^ RANDOM; }
};
gp_hash_table<key, int, chash> table;

int main() {}
```

## 7.11 Numerical Integration

```
// different schemes for numerical integration
// approximatively ordered by accuracy
// do NOT use integer types for integration range!
struct Integration_Midpoint{
    template<typename Func, typename S>
    static typename result_of<Func(S)>::type
        integrate_step(Func f, S l, S r){
        S m = (l+r)/2;
        return f(m) * (r-l);
    }
};

struct Integration_Simpson{
    template<typename Func, typename S>
    static typename result_of<Func(S)>::type
        integrate_step(Func f, S l, S r){
        S m = (l+r)/2;
        return (f(l) + 4*f(m) + f(r))/6 * (r-l);
    }
};

struct Integration_Gauss_2{
    static constexpr long double A = 1.01/sqrtl(3)/2, x1
        =0.51-A, x2 = 0.51+A;
    template<typename Func, typename S>
    static typename result_of<Func(S)>::type
        integrate_step(Func f, S l, S r){
        return (f(l*x1 + r*x2) + f(l*x2+r*x1))/2 * (r-l);
    }
};

struct Integration_NCotes_Open_4{
    template<typename Func, typename S>
    static typename result_of<Func(S)>::type
        integrate_step(Func f, S l, S r){
        S h = (r-l)/5;
        return (11*f(l+h) + f(l+2*h) + f(r-2*h) + 11*f(r-h))
            /24 * (r-l);
    }
};

struct Integration_Gauss_3{
    static constexpr long double A = sqrtl(3.01/5.01)/2,
        x1=0.5-A, x2 = 0.5+A;
    template<typename Func, typename S>
    static typename result_of<Func(S)>::type
        integrate_step(Func f, S l, S r){
        return (5*f(l*x1 + r*x2) + 8*f((l+r)/2) + 5*f(l*x2+r
            *x1))/18 * (r-l);
    }
};

template<typename Integration_Method>
struct Integrator_Fixedstep{
    template<typename Func, typename S>
```

```
static typename result_of<Func(S)>::type integrate(
    Func f, S const l, S const r, size_t const steps){
    assert(steps>0);
    typename result_of<Func(S)>::type ret(0);
    S cur_l = l, cur_r;
    for(size_t i=0;i<steps;++i){
        cur_r = (l*(steps-i-1) + r*(i+1))/steps;
        ret+=Integration_Method::integrate_step(f, cur_l,
            cur_r);
        cur_l = cur_r;
    }
    return ret;
};

template<typename Integration_Method>
class Integrator_Adaptive{
private:
    template<size_t depth_limit, typename Func, typename S>
    static typename result_of<Func(S)>::type integrate(
        Func f, S const l, S const r, typename result_of<
        Func(S)>::type const val, typename result_of<Func(S)
        >::type const eps, const size_t depth){
        if(depth>=depth_limit){
            return val;
        }
        S const m = (l+r)/2;
        typename result_of<Func(S)>::type val_l =
            Integration_Method::integrate_step(f, l, m);
        typename result_of<Func(S)>::type val_r =
            Integration_Method::integrate_step(f, m, r);
        typename result_of<Func(S)>::type error = abs(val -
            val_l - val_r);
        if(error < eps){
            return val_l + val_r;
        }
        return integrate<depth_limit>(f, l, m, val_l, eps/2,
            depth+1)
            + integrate<depth_limit>(f, m, r, val_r, eps/2,
            depth+1);
    }
public:
    template<size_t depth_limit, typename Func, typename S>
    static typename result_of<Func(S)>::type integrate(
        Func f, S const l, S const r, typename result_of<
        Func(S)>::type const eps){
        return integrate<depth_limit>(f, l, r,
            Integration_Method::integrate_step(f, l, r), eps,
            0);
    }
};
```