English documentation for mcq_solver

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August 20, 2023

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1 Introduction

mcq_solver is a library written in Rust and compiled, with maturin and pyo3, into a shared library accessible by a Python "CPython" interpreter, version ≥ 3.10 (tested on 3.10.6). The installation and uninstallation processes are explained in the repository's README.md file.

Its functions allow, from the answer sheets of a MCQ and their grades, to determine the correction of said MCQ. The algorithm is based on the "simulated annealing" statistical method in order to converge to a correction that would explain the best way the grades observed.

An example for each function can be found in the repository's examples directory.

2 Available functions

The library mcq_solver proposes two functions. Their arguments are given below in order with their original names.

2.1 solve

Description Determines the correction of a MCQ from raw Python data. The output is a couple (result, success). If success is True, then result approaches or is the MCQ's correction. Otherwise, result is an error message to print out.

Mandatory arguments (On the left lies the name, on the right lies description)

sheets A list of answer sheets, each sheet being a list of answer tokens

_answer_tokens A string or a list of <u>distincts</u> tokens (characters) corresponding to the different

possible answers

grades A list of integers ≤ 255 corresponding the the grades of the sheets

Optional arguments (From left to right, the name, the default value, the description)

starting_beta 0.1 β initial value in the annealing algorithm

max_beta 0.5 The value of β for which the algorithm stops

lambda_inv 1.01 The value of $1/\lambda$ in the annealing algorithm

2.2 solve_from_files

Description Determines the correction of an MCQ based on data written in text files. The output is the same as for solve. This function can help shipping and using the data of hundreds of sheets, each containing more than a hundred answers.

Mandatory arguments (On the left lies the name, on the right lies description)

sheets_path A list of paths to the answer sheets

grades_path A path leading to a file containing the grades, integers ≤ 255

grades_separator A token meant to separate the grades in the file at the end of grades_path

(typically comma ', ' or a newline '\n')

number_of_questions The number of questions of the MCQ

Optional arguments The same as solve's

3 The algorithm used

3.1 Potential to minimize

Suppose we have a problem consisting in $q \in \mathbb{N}^*$ questions and $N \in \mathbb{N}^*$ answer sheets, denoted $F_1, ..., F_N$. Each sheet F_n has a grade $g_n \in [0, q]$ and a list of answers $a_{n,1}...a_{n,q}$. We also suppose that there is only one correction to the MCQ $a = a_1...a_q$ that can give the resulting grades (otherwise, the problem is unsolvable).

Therefore,

$$\forall n, g_n = \#\{k \in [0, q] \mid a_{n,k} = a_k\}$$

The goal is to determine the list of answers satisfying this condition that is a characteristic of the correction. To this end, we represent a candidate correction by a list $c = c_1...c_q$ of answers. We then define the *disagreement* between c and a over F_n by

$$d_n(a,c) = \#\{k \in [0,q] \mid a_{n,k} = a_k\} - \#\{k \in [0,q] \mid a_{n,k} = c_k\}$$

And we introduce the following function, that we will call c's potential:

$$\phi(c) = \sum_{n=1}^{N} d_n(a, c)^2$$

It is therefore clear that, under the chosen conditions, $a = c \Leftrightarrow \phi(c) = 0$. And so we will construct a sequence of candidates (c_i) that musts make ϕ converge to its minimum. The exact method is described in the next subsection.

3.2 Simulated annealing algorithl

The simulated annealing is a general statistical method to search for the extremum of a function that is hard to analyze. Here is how we use it in our case:

• Setup

- 1. Choose an initial value for the temperature $T_0 > 1$.
- 2. Choose a stopping value for the temperature $T_0 > T^* > 1$. The algorithm will iterate as long as $T > T^*$.
- 3. Choose a constant value for a coefficient $0 < \lambda < 1$. The usual value picked is $\lambda = 0.99$ but here we will rather take $1/\lambda = 1.01$.
- 4. Choose the first candidate c_0 to be the answer sheet with the best grade (or randomly picked between the best sheets if there are several)

• Main loop (while $T > T^*$ and $\phi(c) \neq 0$)

- 1. Randomly choose a question $k \in [0, q]$ and change randomly c's answer to this question to construct another candidate c'.
- 2. Note the change in potential $\Delta \phi = \phi(c') \phi(c)$.
- 3. If $\Delta \phi < 0$, replace c by c'.
- 4. Otherwise, replace c by c' with a probability $\exp\left(-\frac{\Delta\phi}{T}\right)$ and, if the change happens, replace T by $\lambda T < T$.
- 5. Memoize the candidate with the lowest potential along the execution.

• Output: lowest potential candidate

3.3 Optimizations

The first optimization lies in remplacing T by the Boltzmann coefficient $\beta = \frac{1}{T}$ (the Boltzmann constant is taken equal to 1 in this representation) and λ by $1/\lambda$ so that there is no division to make in the algorithm, and these are more expansive than multiplications.

A second one, to compute fast a value of $d_n(c,a)$, if to vectorize the main loop of the calculation (the loop that counts matching answers) by using the SIMDs. This optimization was removed from the code as Rust's compiler is powered by the LLVM toolchain which applies a similar optimization automatically.

Finally, initially, every structure was statically allocated on the stack, which allowed in fine to compute a value of $d_n(c, a)$ in about 20 nanoseconds. In order to allow the user to choose easily the number of questions per answer sheets and other parameters, the structures are now dynamically allocated on the heap, which reduced this performance a bit.

4 Performances

Several dozens of thousands of randomly¹ generated MCQs with hundreds of answer sheets for each where given to the algorithm. The number of possible answers per question ranged from 2 to 8 and the number of questions was 120.

4.1 Correction

When 2 choices per question are given, 150 copies were often enough to solve the problem. For 8 choices, this number had to be two or threefolded. What was generally observed is that, out of 5000 MCQs, the algorithm failed solving between 1 and 5 of them. The grade of the candidate was still much higher than the best grade of the input sheets.

4.2 Speed

In the same conditions, it took the algorithm at most a few milliseconds to solve an MCQ. 5000 of them were solved in 2.5 seconds generally.

¹Every random input followed a uniform law

5 Organisation du dossier src

This folder contains the Rust source code of the library. The functions defined here are slightly modified and re-exported by the script <code>__init__.py</code> in the directory <code>python/mcq_solver</code>.

- The file lib.rs declares the other modules and exports the Rust functions _solve and _solve_from_files called by the library's Python layer.
- vec_util.rs defines a function that is not used anymore to initialize a Rust dynamic array from a generator function.
- sheet.rs defines the structure Sheet for the answer sheets and the type Answer for the answers.
- parameters.rs defines the structure Annealing Parameters encapsulating the parameters of the annealing algorithm.
- annealing.rs defines in particular the structure AnnealingSolver and the function solve_mcq allowing to use the annealing algorithm to correct the MCQ.
- parsing.rs defines the backend of solve_from_files.
- basic.rs defines the backend of solve.