

# Horn Minimization

An overview of some existing algorithms

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# Let's discuss 1.1 - Toy example



#### List of film genres:

► sci-fi, fantastic, fantasy, cyber-punk, post-apo, documentary, drama, romance, . . .

### We can *derive* genres from others, e.g :

- cyber-punk leads probably to sci-fi,
- ▶ a documentary may not be fantastic,
- **...**



# Implications at first sight

1.1 - Toy example



#### Rewritting of previous examples :

- ► cyber-punk → sci-fi,
- ▶ documentary, war → historical,

# Closure operator and systems

2.1 - Elements of set theory



Set  $\Sigma$  of attributes. A map  $\varphi: 2^{\Sigma} \longrightarrow 2^{\Sigma}$  is a *closure operator* if,  $\forall X, Y, Z \subseteq \Sigma$ :

- $\triangleright$   $X \subseteq \varphi(X)$  (increasing)
- $\triangleright$   $X \subseteq Y \longrightarrow \varphi(X) \subseteq \varphi(Y)$  (isotone)
- $\triangleright \varphi(\varphi(X)) = \varphi(X)$  (idempotent)

#### Some details:

- ightharpoonup X is *closed* if  $X = \varphi(X)$ ,
- $\triangleright$   $\Sigma^{\varphi}$  set of closed sets : *closure system*.
- $\triangleright \Sigma^{\varphi}$  is closed under *intersection*, contains  $\Sigma$ .



Pouet.

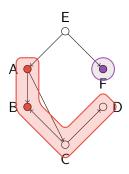
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# Closure Example

#### 2.1 - Elements of set theory



- ▶ Directed graph G = (V, E).
- ► Closure  $\varphi(X)$  of  $X \subseteq V$ : every vertices reachable from X.
- $ightharpoonup \varphi(\{A,B\}) = \{A,B,C,D\}.$
- $\varphi(\{F\}) = \{F\}, \{F\}$  is *closed*.



 $\label{eq:Figure} \begin{aligned} &\operatorname{Figure} - \operatorname{Closure} \text{ of a vertex in a} \\ &\operatorname{directed} \text{ graph} \end{aligned}$ 

# **Implications**

2.1 - Elements of set theory



## $A, B \subseteq \Sigma$ . An *implication* is :

- $\triangleright$  A  $\longrightarrow$  B, A premise, B conclusion.
- ▶ relation : "If we have A, we have B". (different from causality)
- $M \subseteq \Sigma$  model of  $A \longrightarrow B$  if

$$B \subseteq M \lor A \nsubseteq M$$

denoted  $M \models A \longrightarrow B$ ,  $A \longrightarrow B$  follows from M.

Set of implications  $\mathcal{L}$ : *implication system*.

- $ightharpoonup M \models \mathcal{L}$  if each element of  $\mathcal{L}$  follows from M,
- $\triangleright$   $\mathcal{L} \models A \longrightarrow B$ : all models of  $\mathcal{L}$  are models of  $A \longrightarrow B$ .



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# Implications and closure

2.1 - Elements of set theory



### Given $\mathcal{L}$ an implication system :

- closure operator  $\mathcal{L}(X)$ : smallest model (inclusion wise) of  $\mathcal{L}$ containing  $X, X \subseteq \Sigma$ .
- ightharpoonup Models of  $\mathcal L$  form a *closure system* :

$$\Sigma^{\mathcal{L}} = \{ M \subseteq \Sigma \mid M = \mathcal{L}(M) \}$$

### Important property:

$$\blacktriangleright \ \mathcal{L} \models A \longrightarrow B \text{ iff } B \subseteq \mathcal{L}(A)$$



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# Small Example

#### 2.1 - Elements of set theory



#### Let $\mathcal L$ be an implication system :

- $\triangleright \Sigma = \{a, b, c, d, e\},\$
- $\blacktriangleright \ \mathcal{L} = \{ab \longrightarrow c, \ bd \longrightarrow a, ce \longrightarrow abd\}$

#### We have :

- $\triangleright$   $\mathcal{L}(b) = b$ , b is closed, hence a model of  $\mathcal{L}$ ,
- $\triangleright$   $\mathcal{L}(bd) = abcd$ , bd is not a model (abcd is).
- $\triangleright$   $\emptyset$  is also a model.



# **Applications**

2.1 - Elements of set theory



- ► Relational Databases,
- ► Formal Concept analysis,
- Conceptual exploration,
- Linguistics?

# title

3.1 - early 80s



# title

3.1 - early 80s



## title 4.1 - Pouf



Pouet

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## title 4.1 - Pouf

