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CS 205: Assignment 1 - Railway Shunting

Integrity Statement

To complete this assignment, the following resources were consulted:

- 1) Lecture Slides by Professor Keogh
 - a) 2_Blind Search_part1
 - b) 2_Blind Search_part2
 - c) 3_Heuristic Search
- 2) Python 3.13 documentation
- 3) Project description
- 4) https://www.transum.org/software/sw/starter_of_the_day/starter_January9.asp
- 5) https://en.wikipedia.org/wiki/Train_shunting_puzzle

The following libraries and their respective documentation were consulted when appropriate:

- 1) Numpy
- 2) Seaborn
- 3) Matplotlib
- 4) Dataclasses
- 5) Queue
- 6) Typing
- 7) Os
- 8) Datetime
- 9) Pytest

All major code was originally written.

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Introduction

Railway Shunting Problem

The Railway Shunting Problem is a combinatorial optimization challenge involving the rearrangement of train cars into a specified goal configuration, subject to a constrained set of legal operations. This problem is grounded in real-world logistics, where physical limitations, such as siding capacity, directionality of movement, and limited maneuvering space add layers of complexity. Inspired by the 8-Puzzle, the Railway Shunting Problem extends the core ideas of search and reordering to a domain that mirrors practical transportation scenarios.

To solve this problem, we apply three different search strategies:

- 1) Uniform Cost Search (UCS)
- 2) A* with Misplaced Train Heuristic
- 3) A* with Manhattan Distance Heuristic

The following sections provide detailed descriptions of each search algorithm. The objective of this project is to compare these methods in terms of efficiency, optimality, and computational costs.

Background

Rail shunting, the process of sorting and assembling railcars, has evolved significantly since the birth of rail transport. In the early days of railways, shunting was performed manually or with the assistance of horses, which proved to be efficient for maneuvering wagons in yards. It wasn't until the mid-19th century that dedicated shunting locomotives were introduced, such as the Stroudley 0-6-0Ts in the 1860s, which were adaptations of mainline engines for yard work (*"When were locomotives first used for shunting? What were the first Shunting Engines?"*).

Steam-powered shunters were introduced in the 19th century, yet the last British Railways shunting horse retired in 1967. The transition to diesel and electric shunting locomotives in the 20th century marked a significant technological advancement, enhancing efficiency and safety in rail yards. Today, shunting operations are integral to railway logistics, utilizing sophisticated locomotives and automation technologies to manage the complex task of assembling and disassembling trains.

Interestingly enough, Dijkstra came up with the shunting yard algorithm, designed to convert an infix expression into a postfix expression, using the practicality of shunting in train cars (*"The Shunting Yard Algorithm"*). Just as railroad yards can temporarily hold cars on side tracks while rearranging, he realized you can sort mathematical operations by moving numbers and operators

between data structures such as stacks (*“How in the world did dijkstra come up with the shunting yards algorithm : r/computerscience”*).

Problem Description

In this version of the Railroad Shunting problem:

- The main track allows bidirectional movement of trains
- There are multiple sidings, each with a limited capacity
- The goal is to rearrange an initial sequence of train cars into a target sequence using the fewest possible moves
- Moves include pushing/pulling a train into or out of a siding or rearranging trains on the main track

Algorithm Implementation

Uniform Cost Search (UCS)

As noted in the project description, Uniform Cost Search is a simplified version of A*, where $h(n)$ (the heuristic) is hardcoded to be 0. This means that our equation in A*,

$$f_i(n) = g_i(n) + h_i(n) \text{ (path cost + heuristic)}$$

is simplified to,

$$f_i(n) = g_i(n)$$

where we expand solely based on the lowest path cost. This often results in the expansion of many irrelevant nodes, particularly in deeper search spaces.

A* with Misplaced Heuristic

This algorithm is a variant of A*, where our heuristic counts the number of train cars not in their goal position. It is simple and admissible, but not always informative, as it doesn't consider how far off a train car is from its intended destination.

$$h_{misplaced}(n) = \sum_i [car_i \neq goal_i]$$

A* with Manhattan Distance Heuristic

This algorithm calculates its heuristic by summing up the number of moves each car is away from its goal position. This heuristic is more informed than the misplaced heuristic and typically results in fewer nodes expanded.

$$h_{manhattan}(n) = \sum_i [pos_{current}(i) - pos_{goal}(i)]$$

This assumes a linear track.

Experimental Setup

Benchmarks were run across multiple puzzles of increasing difficulty:

- **Easy:** 2-3 cars, minimal rearrangement
- **Medium:** 3-4 cars, moderate sidings usage
- **Hard:** 4+ cars, significant reordering

Metrics tracked:

- Path length (solution depth)
- Number of nodes expanded
- Execution time
- Maximum queue size

Results and Analysis

The following tables summarize the performance metrics for puzzles of increasing complexity: easy, medium, and hard.

Easy

Summary Statistics for Puzzle: easy2

Initial State: $2 \rightarrow 1 \rightarrow 3$
Goal State: $1 \rightarrow 2 \rightarrow 3$
Expected Depth: 2 moves
Sidings: 2

Solution path:

Step 0:

Main Track: $2 \rightarrow 1 \rightarrow 3$

Siding 1:

Siding 2:

Step 1:

Main Track: $1 \rightarrow 3$

Siding 1:

Siding 2: 2

Step 2:

Main Track: 3

Siding 1: 1

Siding 2: 2

Step 3:

Main Track: $2 \rightarrow 3$

Siding 1: 1

Siding 2:

Step 4:

Main Track: $1 \rightarrow 2 \rightarrow 3$

Siding 1:

Siding 2:

=====

Algorithm: UCS

Path Length: 4
Nodes Expanded: 11
Max Queue Size: 10
Execution Time: 0.0012 seconds
Nodes per Second: 9100.07

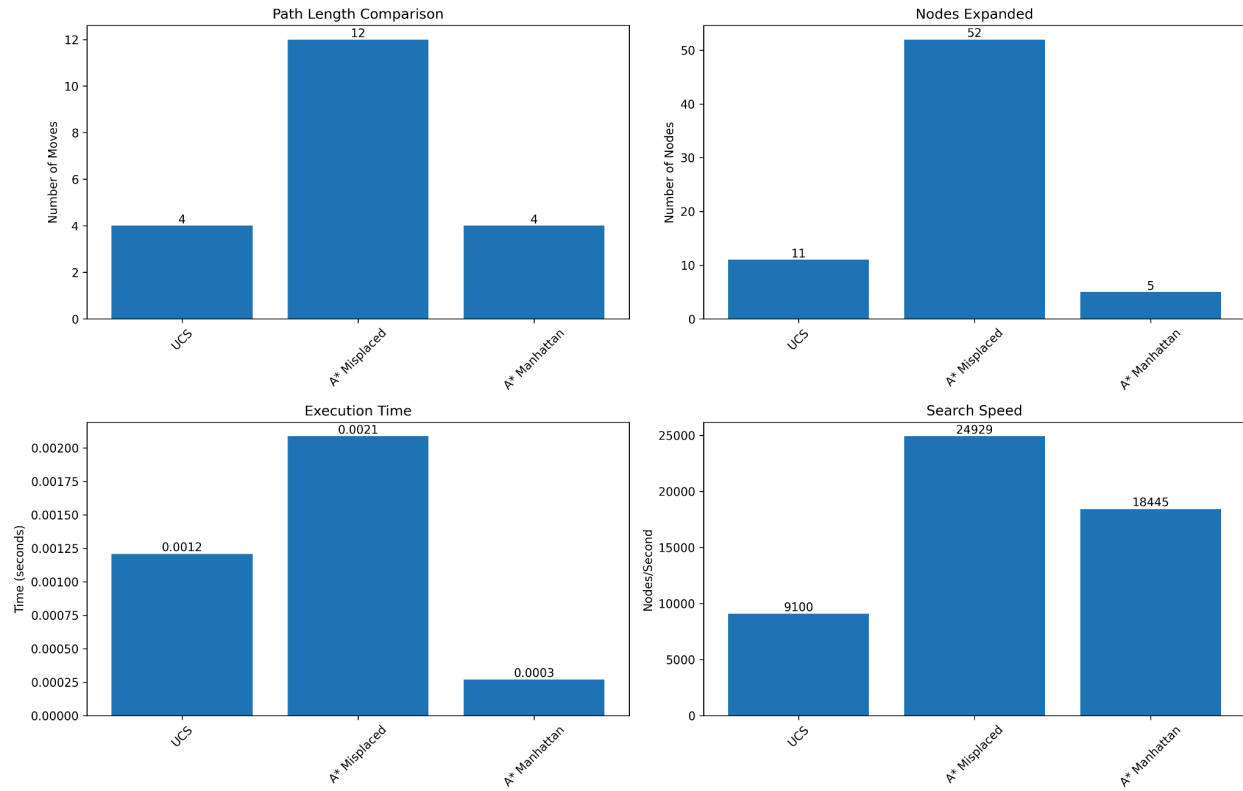
Algorithm: A* Misplaced

Path Length: 12
Nodes Expanded: 52
Max Queue Size: 19
Execution Time: 0.0021 seconds
Nodes per Second: 24929.00

Algorithm: A* Manhattan

Path Length: 4
Nodes Expanded: 5
Max Queue Size: 7
Execution Time: 0.0003 seconds
Nodes per Second: 18444.61

Performance Comparison for Puzzle: easy2



A* with Manhattan heuristic clearly outperforms the other searches by finding the optimal path with the fewest nodes expanded and fastest execution time, while A* with Misplaced heuristic drastically overestimated the solution cost.

Medium

Summary Statistics for Puzzle: medium2

Initial State: $2 \rightarrow 3 \rightarrow 1$

Goal State: $1 \rightarrow 2 \rightarrow 3$

Expected Depth: 6 moves

Sidings: 3

Solution path:

Step 0:

Main Track: $2 \rightarrow 3 \rightarrow 1$

Siding 1:

Siding 2:

Siding 3:

Step 1:

Main Track: $3 \rightarrow 1$

Siding 1: 2

Siding 2:

Siding 3:

Step 2:

Main Track: 1

Siding 1: $2 \rightarrow 3$

Siding 2:

Siding 3:

Step 3:

Main Track:

Siding 1: $2 \rightarrow 3$

Siding 2: 1

Siding 3:

Step 4:

Main Track: 3

Siding 1: 2

Siding 2: 1

Siding 3:

Step 5:

Main Track: 2 → 3

Siding 1:

Siding 2: 1

Siding 3:

Step 6:

Main Track: 1 → 2 → 3

Siding 1:

Siding 2:

Siding 3:

=====

Algorithm: UCS

Path Length: 6

Nodes Expanded: 118

Max Queue Size: 58

Execution Time: 0.0089 seconds

Nodes per Second: 13230.18

Algorithm: A* Misplaced

Path Length: 16

Nodes Expanded: 112

Max Queue Size: 56

Execution Time: 0.0058 seconds

Nodes per Second: 19357.26

Algorithm: A* Manhattan

Path Length: 6

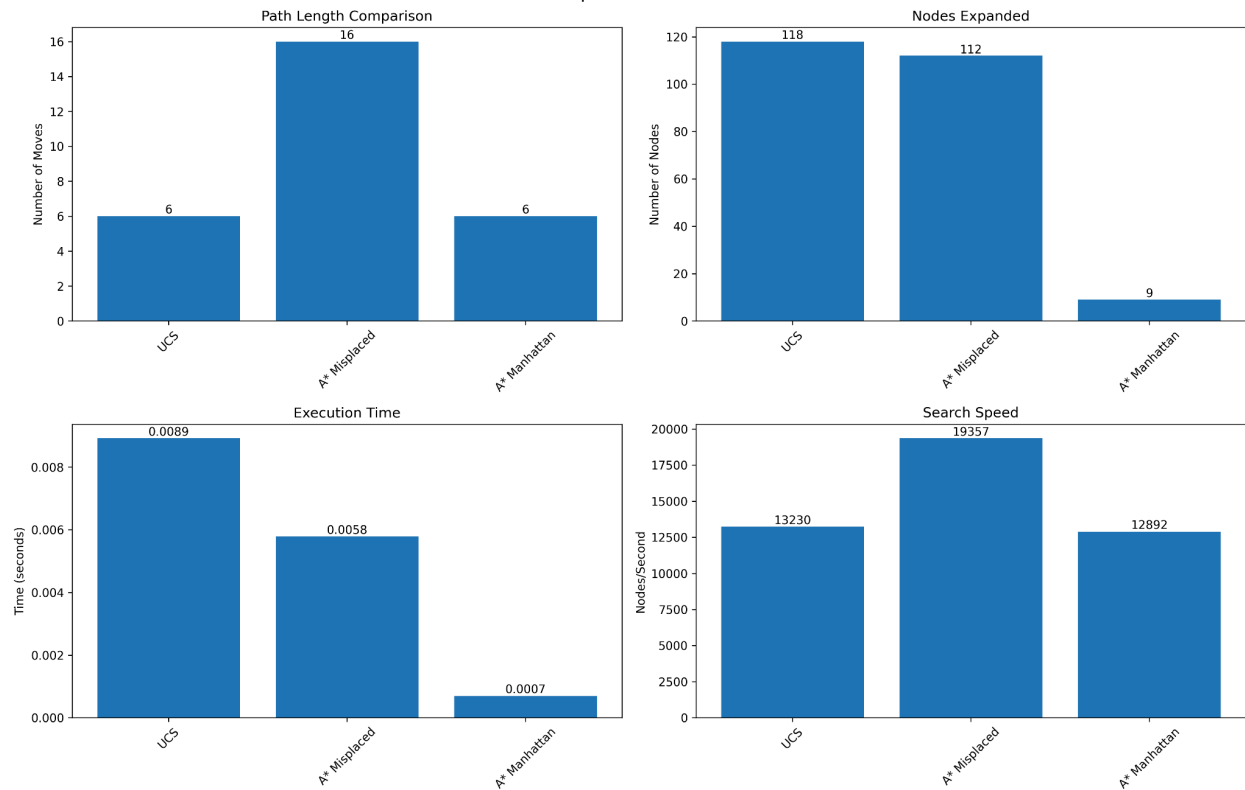
Nodes Expanded: 9

Max Queue Size: 17

Execution Time: 0.0007 seconds

Nodes per Second: 12892.33

Performance Comparison for Puzzle: medium2



A* with Manhattan heuristic again delivered the optimal solution with maximal node expansion, whereas the misplaced heuristic led A* astray, resulting in an overly long solution path, despite a fast search speed.

Hard

Summary Statistics for Puzzle: hard2

Initial State: $2 \rightarrow 4 \rightarrow 1 \rightarrow 3$

Goal State: $1 \rightarrow 2 \rightarrow 3 \rightarrow 4$

Expected Depth: 10 moves

Sidings: 3

Solution path:

Step 0:

Main Track: $2 \rightarrow 4 \rightarrow 1 \rightarrow 3$

Siding 1:

Siding 2:

Siding 3:

Step 1:

Main Track: $4 \rightarrow 1 \rightarrow 3$

Siding 1: 2

Siding 2:

Siding 3:

Step 2:

Main Track: $1 \rightarrow 3$

Siding 1: $2 \rightarrow 4$

Siding 2:

Siding 3:

Step 3:

Main Track: 3

Siding 1: $2 \rightarrow 4$

Siding 2: 1

Siding 3:

Step 4:

Main Track: $4 \rightarrow 3$

Siding 1: 2

Siding 2: 1

Siding 3:

Step 5:

Main Track: $2 \rightarrow 4 \rightarrow 3$

Siding 1:

Siding 2: 1

Siding 3:

Step 6:

Main Track: $1 \rightarrow 2 \rightarrow 4 \rightarrow 3$

Siding 1:

Siding 2:

Siding 3:

Step 7:

Main Track: $2 \rightarrow 4 \rightarrow 3$

Siding 1: 1

Siding 2:

Siding 3:

Step 8:

Main Track: $4 \rightarrow 3$

Siding 1: $1 \rightarrow 2$

Siding 2:

Siding 3:

Step 9:

Main Track: 3

Siding 1: $1 \rightarrow 2 \rightarrow 4$

Siding 2:

Siding 3:

Step 10:

Main Track:

Siding 1: $1 \rightarrow 2 \rightarrow 4$

Siding 2: 3

Siding 3:

Step 11:

Main Track: 4

Siding 1: $1 \rightarrow 2$

Siding 2: 3

Siding 3:

Step 12:

Main Track: $3 \rightarrow 4$

Siding 1: $1 \rightarrow 2$

Siding 2:

Siding 3:

Step 13:

Main Track: $2 \rightarrow 3 \rightarrow 4$

Siding 1: 1

Siding 2:

Siding 3:

Step 14:

Main Track: $1 \rightarrow 2 \rightarrow 3 \rightarrow 4$

Siding 1:

Siding 2:

Siding 3:

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Algorithm: UCS

Path Length: 10

Nodes Expanded: 705

Max Queue Size: 398

Execution Time: 0.0444 seconds

Nodes per Second: 15890.93

Algorithm: A* Misplaced

Path Length: 14

Nodes Expanded: 736

Max Queue Size: 451

Execution Time: 0.0434 seconds

Nodes per Second: 16957.76

Algorithm: A* Manhattan

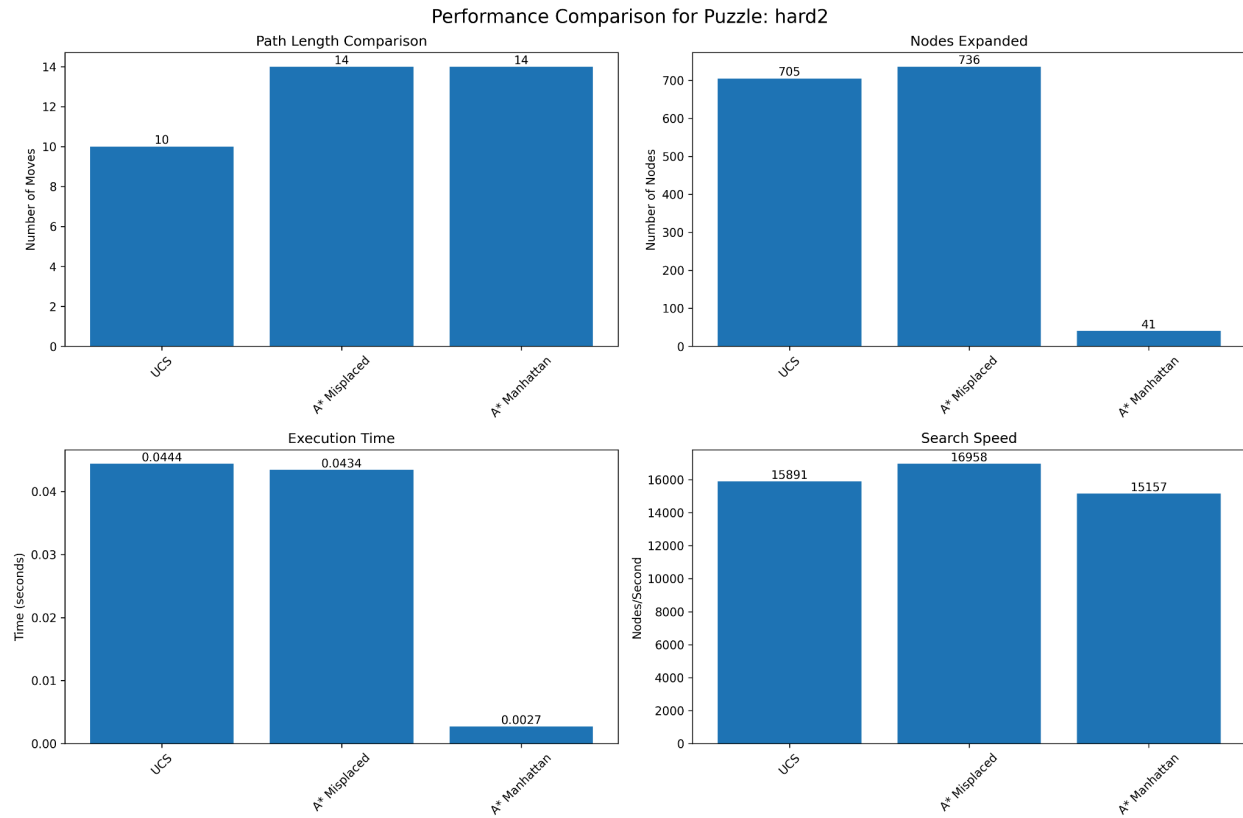
Path Length: 14

Nodes Expanded: 41

Max Queue Size: 73

Execution Time: 0.0027 seconds

Nodes per Second: 15156.57



In this more complex puzzle, A* with Manhattan heuristic maintained its advantage by achieving near-optimal results with dramatically fewer node expansions, while UCS and the Misplaced heuristic expanded significantly more nodes and produced less efficient paths.

Conclusion

In this project, we applied three different search algorithms to solve the Railway Shunting problem. After interpreting our results, we found that:

- **Uniform Cost Search**: while guaranteed to find an optimal solution, is computationally expensive
- **A*** outperforms UCS in both speed and efficiency when paired with an effective heuristic
- The **Misplaced Train Heuristic** is surprisingly effective for this problem, though the **Manhattan Distance Heuristic** offers slightly more informed search behavior

Works Cited

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