

## CSC115 Lecture 12



### Chapters 7 & 8

## Stacks & Queues

- Concepts
- ADT
- Implementation ideas

Next Class: implementation details

## New Concept: Stack

- A stack
  - Last-in, first-out (LIFO) property
    - The last item placed on the stack will be the first item removed
  - Analogy
    - A stack of dishes in a cafeteria

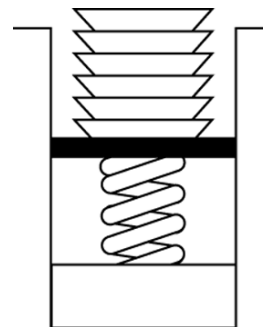


Figure 7-1  
Stack of cafeteria dishes

## New Concept: Queue

- A queue
  - First in, first out (FIFO) property
    - The first item added is the first item to be removed



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## Recall: Abstract Data Types (from Lecture 4)

- An ADT is composed of
  - A collection of data
  - A set of operations on that data
- Specifications of an ADT indicate
  - What the ADT operations do,  
NOT: how to implement them
- Implementation of an ADT
  - Includes choosing a particular data structure

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## Recall: Abstract Data Types

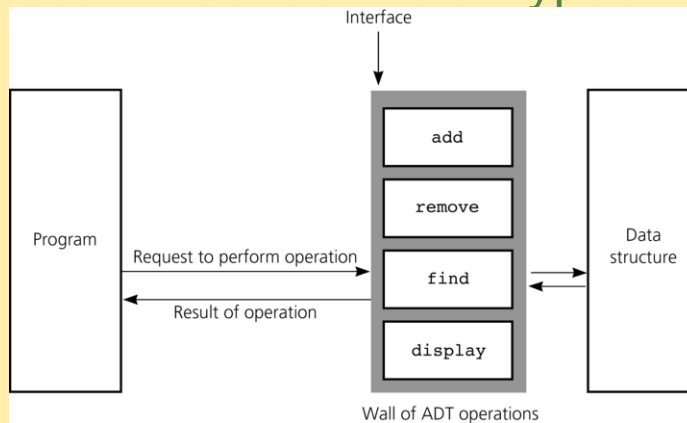


Figure 4-4

A wall of ADT operations

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So . . . .

Lets consider ADTs for Stack & Queue

## ADT Stack

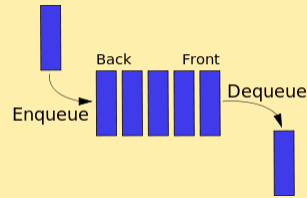
- ADT stack operations
  - Determine whether a stack is empty
  - Add a new item to the stack
  - Remove from the stack the item that was added most recently
  - Remove all the items from the stack
  - Retrieve from the stack the item that was added most recently



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## The Abstract Data Type Queue



Graphic from: wikipedia.org

- ADT queue operations
  - Determine whether a queue is empty
  - Add a new item to the queue
  - Remove from the queue the item that was added earliest
  - Remove all the items from the queue
  - Retrieve from the queue the item that was added earliest

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## The ADT Stack → Java's Interface

- Pseudocode for the ADT stack operations

```
isEmpty()
// Determines whether a stack is empty.

push(newItem) throws StackException
// Adds newItem to the top of the stack.
// Throws StackException if the insertion is
// not successful.
```

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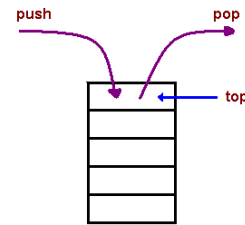
## The ADT Stack → Java's Interface

- Pseudocode for the ADT stack operations (Continued)

```
pop() throws StackException
// Retrieves and then removes the top of the stack.
// Throws StackException if the deletion is not
// successful.
```

```
popAll()
// Removes all items from the stack.
```

```
peek() throws StackException
// Retrieves the top of the stack. Throw
// StackException if the retrieval is not successful
```



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## The Abstract Data Type Queue

- Pseudocode for the ADT queue operations

```
isEmpty()
// Determines whether a queue is empty
```

```
enqueue(newItem) throws QueueException
// Adds newItem at the back of a queue. Throws
// QueueException if the operation is not
// successful
```

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## The Abstract Data Type Queue

- Pseudocode for the ADT queue operations (Continued)

```
dequeue() throws QueueException
// Retrieves and removes the front of a queue.
// Throws QueueException if the operation is
// not successful.
```

```
dequeueAll()
// Removes all items from a queue
```

```
peek() throws QueueException
// Retrieves the front of a queue. Throws
// QueueException if the retrieval is not
// successful
```

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## Example Application: the ADT Stack: - Checking for Balanced Braces

- A stack can be used to verify whether a program contains balanced braces
  - An example of balanced braces  
`abc{defg{ijk}}{l{mn}}op}qr`
  - An example of unbalanced braces  
`abc{def}}{ghij{kl}m`

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## Checking for Balanced Braces

- Requirements for balanced braces
  - Each time you encounter a “}”, it matches an already encountered “{”
  - When you reach the end of the string, you have matched each “{”

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## Checking for Balanced Braces

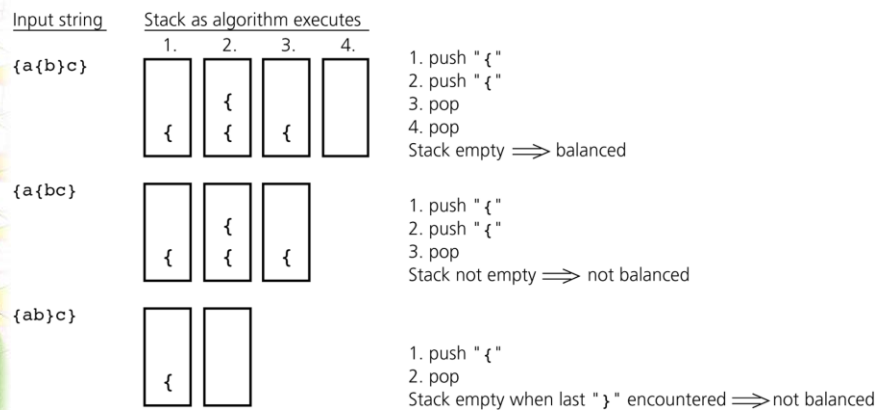


Figure 7-3

Traces of the algorithm that checks for balanced braces

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## Example: Balanced Braces: ie, javac

```
public class TestQuest4 {
    public static int theSum(int num) {
        if (num < 2) return 2*num;
        else {
            return theSum(num/2) + theSum(num-4);
        }
    }

    public static void main(String[] args) {
        int answer=theSum(6);
        System.out.println("The Sum = " + answer);
    }
}
```

Imagine: This brace is missing

## Checking for Balanced Braces

- The exception `StackException`
  - A Java method that implements the balanced-braces algorithm should do one of the following
    - Take precautions to avoid an exception
    - Provide `try` and `catch` blocks to handle a possible exception



## Implementations of the ADT Stack

- The ADT stack can be implemented using

- An array
- A linked list

Implementation of methods discussed in class:

- see Lecture 12 Code for result
- also, some implementation pictures next 2 slides

- StackInterface

- Provides a common specification for the implementations

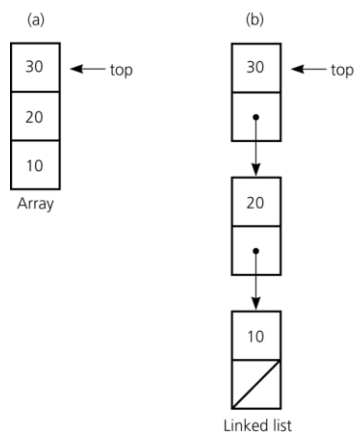
- StackException

- Used by StackInterface
- Extends `java.lang.RuntimeException`

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## Implementations of the ADT Stack



**Figure 7-4**

Implementation of the ADT stack that use a) an array; b) a linked list;

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## An Array-Based Implementation of the ADT Stack

- `StackArrayBased` class
  - Implements `StackInterface`
  - Instances
    - Stacks
  - Private data fields
    - An array of Objects called `items`
    - The index `top`

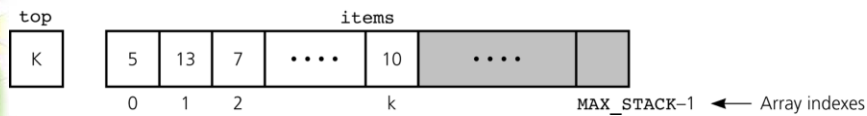


Figure 7-5

An array-based implementation

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## A Reference-Based Implementation of the ADT Stack

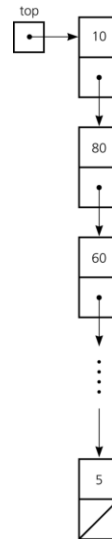
- A reference-based implementation
  - Required when the stack needs to grow and shrink dynamically
- `StackReferenceBased`
  - Implements `StackInterface`
  - `top` is a reference to the head of a linked list of items

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## A Reference-Based Implementation of the ADT Stack

**Figure 7-6**  
A reference-based  
implementation



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## The Abstract Data Type Queue

- Queues
  - Are appropriate for many real-world situations
    - Example: A line to buy a movie ticket
  - Have applications in computer science
    - Example: A request to print a document

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## The Abstract Data Type Queue

Operation	Queue after operation
<code>queue.createQueue()</code>	↓ Front
<code>queue.enqueue(5)</code>	5
<code>queue.enqueue(2)</code>	5 2
<code>queue.enqueue(7)</code>	5 2 7
<code>queueFront = queue.peek()</code>	5 2 7 (queueFront is 5)
<code>queueFront = queue.dequeue()</code>	5 2 7 (queueFront is 5)
<code>queueFront = queue.dequeue()</code>	2 7 (queueFront is 2)

Figure 8-2

Some queue operations

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## Simple Applications of the ADT Queue: Reading a String of Characters

- A queue can retain characters in the order in which they are typed

```
queue.createQueue()
while (not end of line) {
    Read a new character ch
    queue.enqueue(ch)
}
```

- Once the characters are in a queue, the system can process them as necessary

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## Recognizing Palindromes

- A palindrome
  - A string of characters that reads the same from left to right as it does from right to left
- To recognize a palindrome, a queue can be used in conjunction with a stack
  - A stack can be used to reverse the order of occurrences
  - A queue can be used to preserve the order of occurrences

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## Recognizing Palindromes

- A nonrecursive recognition algorithm for palindromes
  - As you traverse the character string from left to right, insert each character into both a queue and a stack
  - Compare the characters at the front of the queue and the top of the stack

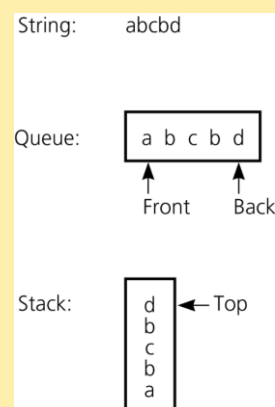


Figure 8-3

The results of inserting a string into both a queue and a stack

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## Implementations of the ADT Queue

- A queue can have either
  - An array-based implementation
  - A reference-based implementation

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## A Reference-Based Implementation

- Possible implementations of a queue
  - A linear linked list with two external references
    - A reference to the front
    - A reference to the back

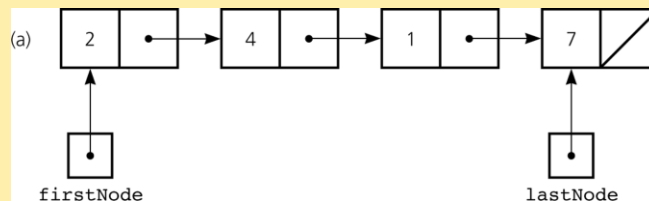


Figure 8-4a

A reference-based implementation of a queue: a) a linear linked list with two external references

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## A Reference-Based Implementation

- Possible implementations of a queue (Continued)
  - A circular linked list with one external reference
    - A reference to the back

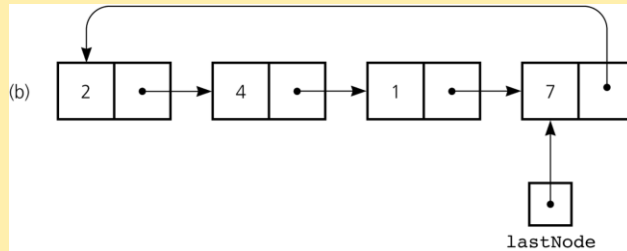


Figure 8-4b

A reference-based implementation of a queue: b) a circular linear linked list with one external reference

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## A Reference-Based Implementation

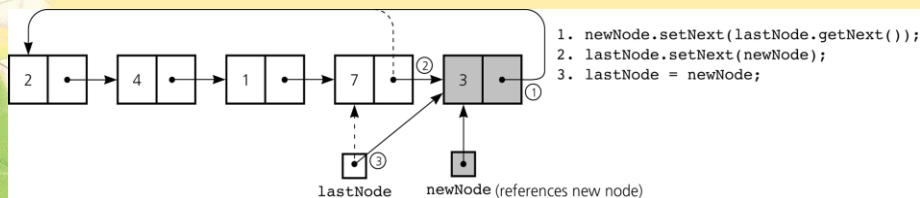


Figure 8-5

Inserting an item into a nonempty queue

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## A Reference-Based Implementation

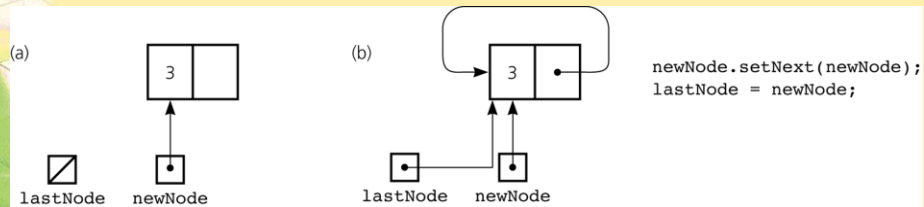


Figure 8-6

Inserting an item into an empty queue: a) before insertion; b) after insertion

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## A Reference-Based Implementation

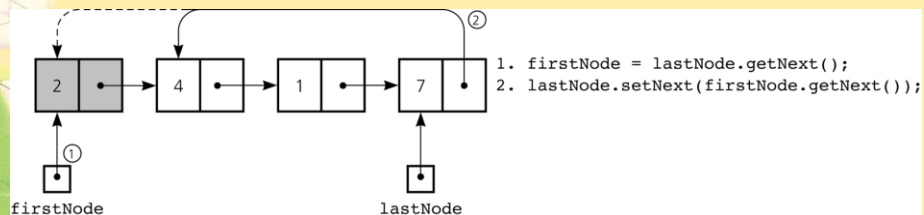


Figure 8-7

Deleting an item from a queue of more than one item

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## An Array-Based Implementation

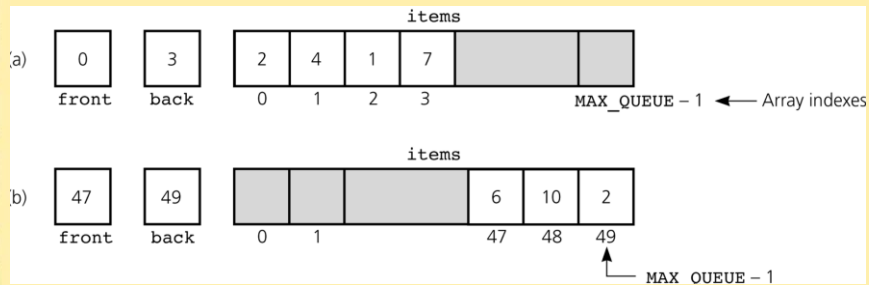


Figure 8-8

a) A naive array-based implementation of a queue; b) rightward drift can cause the queue to appear full

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## An Array-Based Implementation

- A circular array eliminates the problem of rightward drift

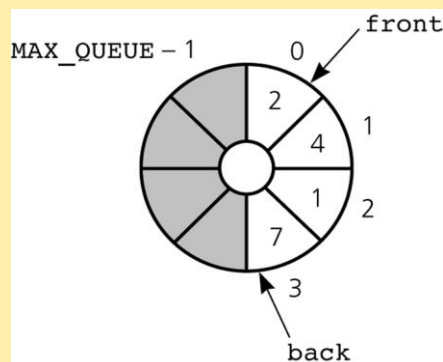


Figure 8-9

A circular implementation of a queue

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## An Array-Based Implementation

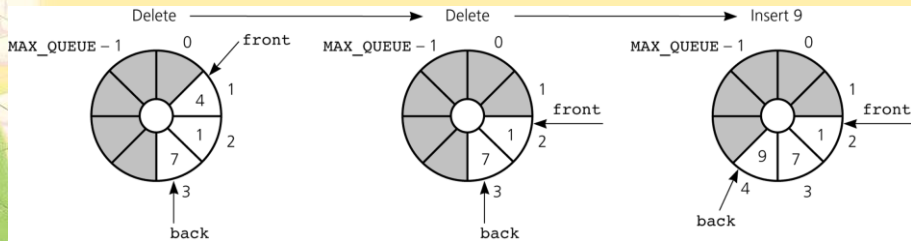


Figure 8-10

The effect of some operations of the queue in Figure 8-8

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## An Array-Based Implementation

- A problem with the circular array implementation
  - $front$  and  $back$  cannot be used to distinguish between queue-full and queue-empty conditions

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## An Array-Based Implementation

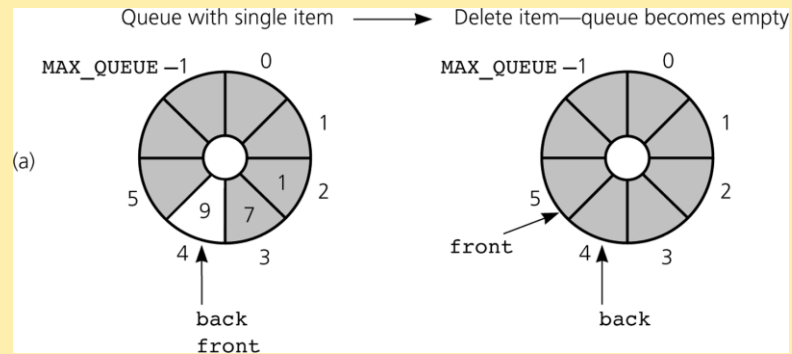


Figure 8-11a

a) **front** passes **back** when the queue becomes empty

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## An Array-Based Implementation

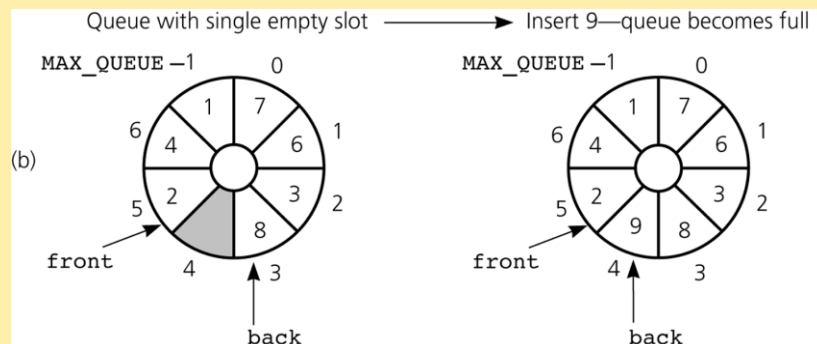


Figure 8-11b

b) **back** catches up to **front** when the queue becomes full

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## An Array-Based Implementation

- Inserting into a queue
 

```
back = (back+1) % MAX_QUEUE;
items[back] = newItem;
++count;
```
- Deleting from a queue
 

```
front = (front+1) % MAX_QUEUE;
--count;
```

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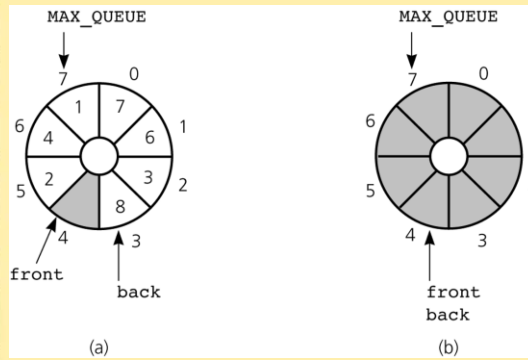
## An Array-Based Implementation

- Variations of the array-based implementation
  - Use a flag `full` to distinguish between the full and empty conditions
  - Declare `MAX_QUEUE + 1` locations for the array items, but use only `MAX_QUEUE` of them for queue items

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## An Array-Based Implementation



**Figure 8-12**

A more efficient circular implementation: a) a full queue; b) an empty queue