Fall 2016 CENG 355

## Assignment 5 **Due November 17, 13:59**

**NOTE:** Late submissions will **NOT** be accepted. Please put your solutions in the CENG 355 **drop-box** (ELW, second floor) – they will be collected at **14:00**.

1. [5 points] Consider the code portion of the <u>matrix-vector product</u> computation as shown below: (float) 128x128 matrix A is multiplied by (float) 128x1 vector X, producing (float) 128x1 result Y (initially all 0's).

```
for (i = 0; i < 128; i++) {
    for (j = 0; j < 128; j++) {
        Y[i] = Y[i] + A[i][j]*X[j];
    }
}</pre>
```

Storing **X**, **Y**, and **A** (each float array element is a 4-byte number) requires 128\*4 + 128\*4 + 128\*128\*4 = **65KB** of memory. If the cache (assume fully associative) is smaller than 65KB, the above code will yield many misses, considerably slowing down program execution. Alternatively, one can perform blocked computation: partition **A** into smaller blocks and perform the product computation block-by-block. If block data can fit into the cache, such blocked computation may significantly outperform the original code.

Rewrite the code fragment above using <u>blocked computation</u> and letting matrix **A**'s blocks be of size **64x64** (i.e., 4 blocks total). Assuming that such blocking yields the best performance, what can you say about the size of the cache?

**2.** [10 points] Consider a <u>C code</u> fragment below, modifying a given <u>square matrix</u> int x[n][n] (stored row by row, i.e., in the row-major order), where n = 256:

```
int average, dev;
for (i = 0; i < N; i++) {
        average = 0;
        for (j = 0; j < N; j++) {
            average = average + X[i][j]; /* sum row elements */
        }
        average = average/N; /* row average */
        for (j = 0; j < N; j++) {
            dev = X[i][j] - average; /* deviation from average */
            X[i][j] = dev*dev; /* deviation squared */
        }
}</pre>
```

Determine the x-related <u>page fault rate</u> in the following <u>two cases</u>: (1) the main memory uses **1-KB** paging with <u>four pages</u> allocated for x, and (2) the main memory uses **4-KB** paging with only <u>one page</u> allocated for x. Initially, no part of x is in the main memory.

**3.** [10 points] Consider a <u>C code</u> fragment below, working on a given <u>square matrix</u> float x[n][n] (stored row by row, i.e., in the row-major order), where n = 256:

Determine the  $\mathbf{x}$ -related <u>page fault rate</u> in the following <u>two cases</u>: (1) the main memory uses **1-KB** paging with <u>four pages</u> allocated for  $\mathbf{x}$ , and (2) the main memory uses **4-KB** paging with only <u>one page</u> allocated for  $\mathbf{x}$ . Initially, no part of  $\mathbf{x}$  is in the main memory.