

CSCI15 Lecture 4

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Intro to Abstract Data Types

Chapter 4: Data Abstraction: The Walls

Recall Program Modularity

- Modularity
 - Keeps the complexity of a large program manageable by systematically controlling the interaction of its components
 - Isolates errors
 - Eliminates redundancies
 - A modular program is
 - · Easier to write
 - Easier to read
 - · Easier to modify

Procedural Abstraction

- Procedural abstraction
 - Separates the purpose and use of a module from its implementation
 - A module's specifications should
 - Detail how the module behaves
 - · Identify details that can be hidden within the module
- Information hiding
 - Hides certain implementation details within a module
 - Makes these details inaccessible from outside the module

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Abstract Data Types

- The isolation of modules is not total
 - Methods' specifications, or contracts, govern how they interact with each other

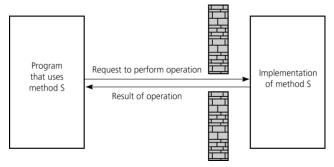


Figure 4-2

A slit in the wall © 2006 Pearson Addison-Wesley. All rights reserved

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Data Abstraction

- Typical operations on data
 - Add data to a data collection
 - Remove data from a data collection
 - Ask questions about the data in a data collection
- Data abstraction
 - Asks you to think what you can do to a collection of data independently of how you do it
 - Allows you to develop each data structure in relative isolation from the rest of the solution
 - A natural extension of procedural abstraction

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Abstract Data Types

- Abstract data type (ADT)
 - An ADT is composed of
 - A collection of data
 - A set of operations on that data
 - Specifications of an ADT indicate
 - What the ADT operations do, not how to implement them
 - Implementation of an ADT
 - Includes choosing a particular data structure

Abstract Data Types

- Data structure
 - A construct that is defined within a programming language to store a collection of data
 - Example: arrays
- ADTs and data structures are not the same
- Data abstraction
 - Results in a wall of ADT operations between data structures and the program that accesses the data within these data structures

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Abstract Data Types Interface add remove Data Program Request to perform operation structure find Result of operation display Wall of ADT operations Figure 4-4 A wall of ADT operations isolates a data structure from the program that uses it © 2006 Pearson Addison-Wesley. All rights reserved

An ADT Example: A List



Figure 4-5
list A grocery

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- In a list
 - Except for the first and last items, each item has
 - A unique predecessor
 - · A unique successor
 - Head or front
 - · Does not have a predecessor
 - Tail or end
 - · Does not have a successor

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The ADT List

- ADT List operations
 - Determine whether a list is empty
 - Determine the number of items in a list
 - · Add an item at a given position in the list
 - Remove the item at a given position in the list
 - Remove all the items from the list
 - Retrieve (get) the item at a given position in the list
- Items are referenced by their position within the list



The ADT List

- Specifications of the ADT operations
 - Define the contract for the ADT list
 - Do not specify how to store the list or how to perform the operations
- ADT operations can be used in an application without the knowledge of how the operations will be implemented

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An Array-Based Implementation of the ADT List

- An array-based implementation
 - A list's items are stored in an array items
 - A natural choice
 - Both an array and a list identify their items by number
 - A list's kth item will be stored in items [k-1]
 - Determine whether a list is empty
 - Determine the number of items in a list
 - > Add an item at a given position in the list
 - Remove the item at a given position in the list
 - Remove all the items from the list
 - Retrieve (get) the item at a given position in the list

An Array-Based Implementation of the ADT List

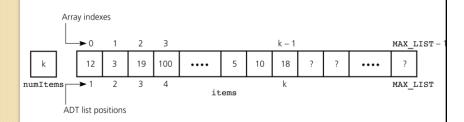


Figure 4-11

An array-based implementation of the ADT list

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A Simpler Example: The ADT IntegerList

- The ADT simplified: IntegerList
 - · adds items to the beginning or end of the list
 - determines the size of the list, and
 - · determines the location of a specified item in the list
- The Java structure for an ADT: an interface

```
public interface IntegerList
  public void addFront (int x);
  public void addBack (int x);
  public int size();
  public int get (int pos);
  public String toString();
}
```

We will implement this ADT in class

Designing an ADT

- The design of an ADT should evolve naturally during the problem-solving process
- Questions to ask when designing an ADT
 - What data does a problem require?
 - What operations does a problem require?

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Implementing ADTs

- Choosing the data structure to represent the ADT's data is a part of implementation
 - Choice of a data structure depends on
 - Details of the ADT's operations
 - · Context in which the operations will be used
- Implementation details should be hidden behind a wall of ADT operations
 - A program would only be able to access the data structure using the ADT operations

Summary

- A client should only be able to access the data structure by using the ADT operations
- An object encapsulates both data and operations on that data
 - In Java, objects are instances of a class, which is a programmer-defined data type
- A Java class contains at least one constructor, which is an initialization method
- Typically, you should make the data fields of a class private and provide public methods to access some or all of the data fields