### Advanced Computer Networks

**Overlay Networks** 

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# Paper reading summaries

- Starting from this week
  - H reading list is on xc (will be further updated)
  - H summary submission schedule to be posted soon
  - H 1-page reading summary on assigned papers
  - H what's the problem? important then? and now?
  - H main ideas? previous work? follow-on work?
  - H strengths then? and now? why?
  - H weaknesses then? and now? why?
  - H how to improve? how to do better or differently?

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<sup>\*</sup> reading guideline at: http://www.cs.uvic.ca/~pan/csc466/reading.txt (template to be posted soon)

#### What do we "have" so far?

- Internet design and architecture
  - H store-and-forward packet switching
  - H end-to-end arguments
    - smart end-hosts vs dumb networks
  - H best-effort services, client-server applications
- Initially, the Internet was an "overlay"
  - H over telephone networks
- By design, the Internet is a "peer-to-peer"
  - H for all end-hosts

# Reality check

- A network of service-provider's networks
  H still mostly packet switching, end-to-end, best-effort
- But hierarchical structures almost everywhere
  - H tiered service provider networks
  - H hierarchies in naming, addressing, routing, service provisioning, content delivery etc
  - H the (only) way to deal with scalability
- Two sides of the story
  - H a lot of details/redundancy invisible to externals

### Examples

- Internet routing
  - H routing pathologies
    - a considerable percentage of routes is affected
  - H delayed convergence
    - after a fault, it takes tens of minutes to converge
  - **H** extended recovery
    - some faults take hours to recover
- Dependable Internet?
  - H not yet

# Adding ??? into the network?

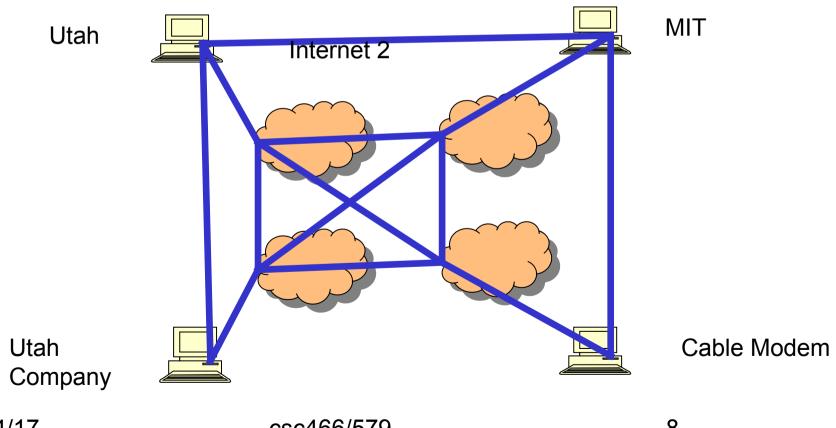
- Changing the infrastructure is difficult
  - <sup>⊬</sup> in a competitive ISP market
  - H only end-to-end counts
  - H and not all applications need perfect ???
- Alternatives
  - **H** application overlays
    - e.g., virtual private networks (VPN)
    - content delivery networks (CDN)
  - H end-to-end or edge-to-edge

### Resilient Overlay Networks

- http://nms.lcs.mit.edu/ron
  - н́ [ABKM01] D. Anderson, H. Balakrishnan, F. Kaashoek, R. Morris, Resilient Overlay Networks, In Proc. of SOSP '01. [RON]
- Design goals
  - H fast failure detection and recovery
    - active probing, re-routing
  - H tighter integration with applications
    - application-specific, e.g., video conferencing
  - **H** expressive policy routing
    - e.g., "no commercial traffic on Internet2" 7

#### Observations

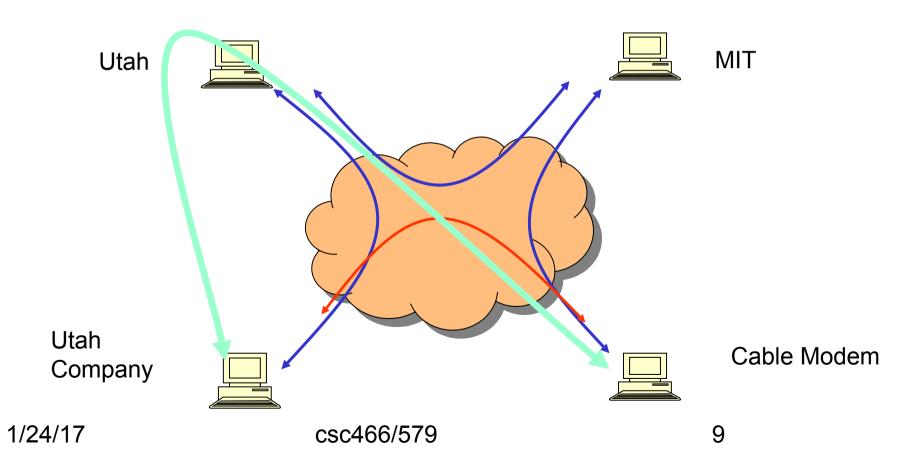
Network redundancy, invisible to applications



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#### Ideas

Route around failures



#### Approaches

- Characterize "links" between nodes
  - H active probing: delay, loss
- Disseminate link characteristics
  - H "link-state" advertisement
- Choose the "best" route
  - H only at the entry node
  - H with possibly one intermediate node
- Forward the packets
- H RON encapsulation

### Design details

Path selection

<sup>Ĥ</sup> delay

- exponentially weighted moving average (EWMA)
- delay<sub>i+1</sub> = a \* delay<sub>i</sub> + (1-a) \* last\_rtt, a = 0.9

H loss: moving window average

window size: 100

**H** throughput

TCP-like, proportional to 1/(rtt \* sqrt(p))

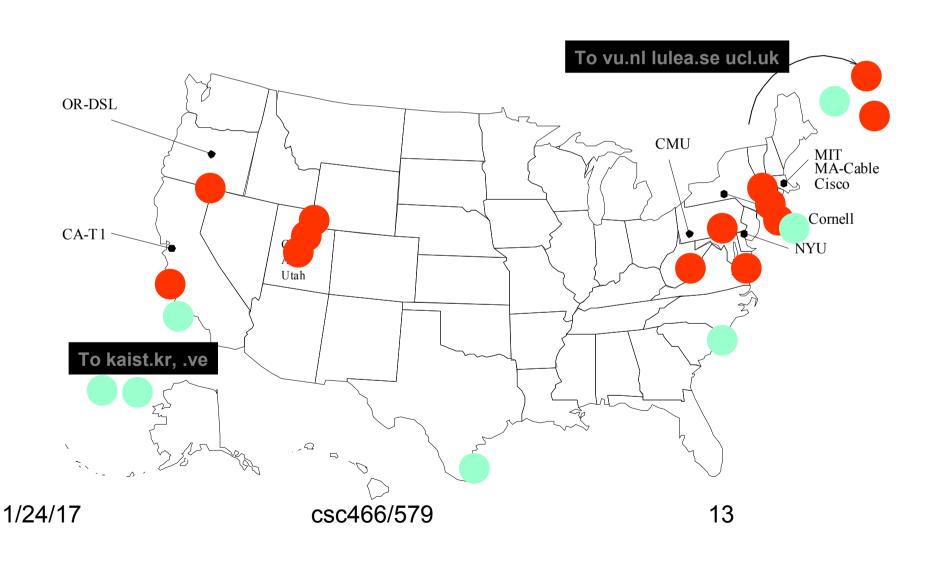
H application-specific

• priority among delay, loss, throughput, etc

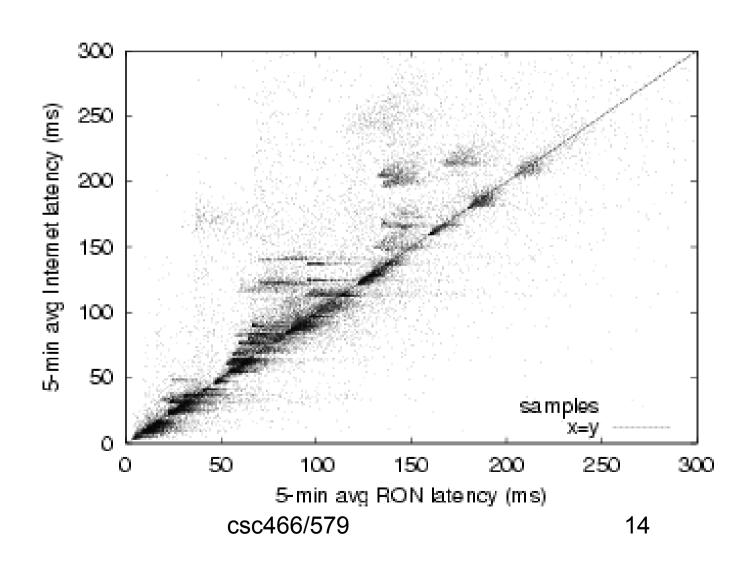
# Membership management

- Static membership
  - H load other peer nodes from a configuration file
- Announcement-based, soft-state membership
  - H know at least one peer node
  - H announce its existence by broadcast
  - **H** soft-state
    - flood peer node list every 5 minutes
    - if a node is not heard for 60 minutes, the node has left
- Search?

#### Performance evaluation

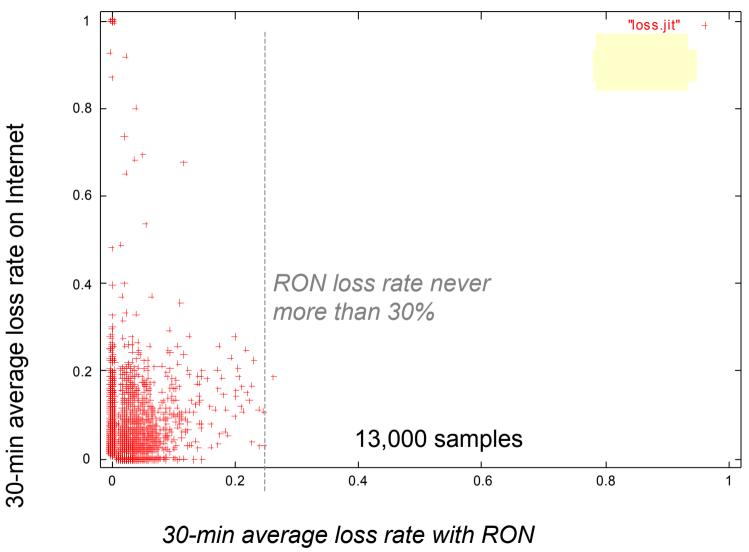


# Reduced delay



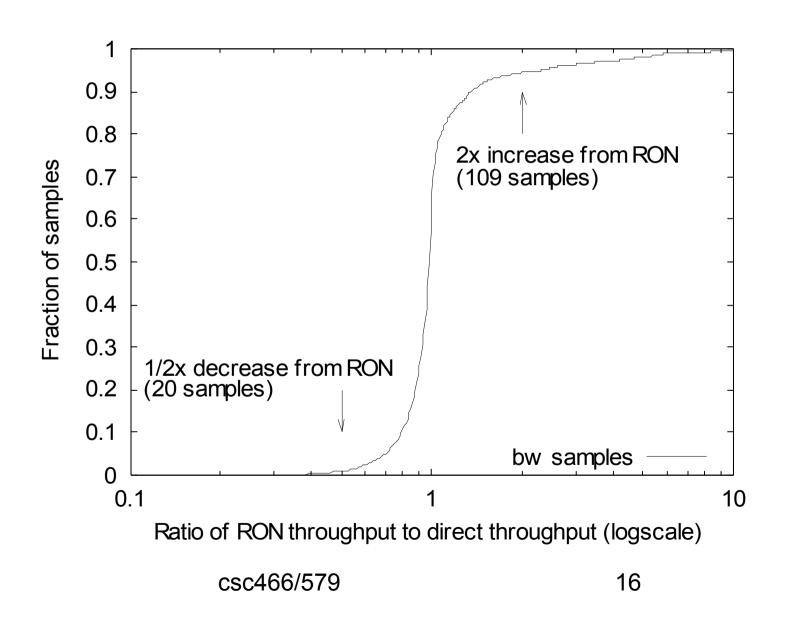
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#### Reduced loss



30-min average loss rate with RON csc466/579

### Improved throughput



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#### Overhead

Link probing

H size: 69 bytes; interval: 12 seconds

Link advertisement

H size: 60+20\*(N-1); interval: 14 seconds

• Recovery time: 12~25 seconds (N=50)

10 nodes	20 nodes	30 nodes	40 nodes	50 nodes
1.8 Kbps	5.9 Kbps	12 Kbps	21 Kbps	32 Kbps

#### More discussion

- One hop?
- Route stability
  - **H** hysteresis
- Path selection
  - H tradeoff between delay, loss, etc
- Routing policy
- Scalability
- NAT (network address translator)

### More overlay networks

- Planet-lab network testbed and GENI
- Peer-to-peer applications
  - H Napster: with centralized directory server
  - H Gnutella: distributed flooding search (ERS)
  - H KaZaA: hierarchy introduced; supernode
  - H BitTorrent: trackers; files in chunks; tit-for-tat
  - <sup>⊬</sup> Skype
  - **H** Structured P2P
    - Distributed Hash Table (HDT): Chord, CAN, Pastry, etc

#### Next lectures

#### DHT

- H Required reading
  - Chord
  - [RFHKS01] S. Ratnasamy, P. Francis, M. Handley, R. Karp, and S. Shenker, "A scalable content-addressable network. In SIGCOMM," Aug. 2001. [CAN]
  - [RD01] Rowstron and P. Druschel, "Pastry: Scalable, distributed object location and routing for largescale peerto-peer systems," Proc. 18th IFIP/ACM Int'l. Conf. Distributed Systems Platforms (Middleware), 2001.
     [Pastry]
- Gnutella, BitTorrent, Skype