

CSCI 15 Lecture 15

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Trees

1. Definitions & Terminology
2. *Binary Tree ADT*
3. *Tree Traversal*

But First:
A little quiz
on Queues
and Stacks!

```
public interface Stack<E>
{
    void push (E element);
    E pop() throws StackEmptyException;
    E peek() throws StackEmptyException;
    boolean isEmpty();
}
```

A Little Quiz #1

```
public interface Queue<T> {
    public void enqueue(T element);
    public T dequeue() throws QueueEmptyException;
    public boolean isEmpty();
}
```

Assume an implementation of the Stack and Queue interfaces. Draw the queue and stack the result from:

```
Queue<String> he = new linkQueue<String>;
Stack<String> she = new arrayStack<String>;
he.enqueue ( "This" );
he.enqueue ( "is" );
she.push (he.dequeue ());
she.push ( "it" );
he.enqueue ( "Wow! " );
she.push (he.dequeue ());
```

A Little Quiz #2

```
public interface Stack<E>
{
    void    push (E element);
    E       pop() throws StackEmptyException;
    E       peek() throws StackEmptyException;
    boolean isEmpty();
}
```

```
public interface Queue<T> {
    public void enqueue(T element);
    public T dequeue() throws QueueEmptyException;
    public boolean isEmpty();
}
```

Assume an implementation of the Stack and Queue interfaces. Draw the queue and stack the result from:

```
Queue<double>  act = new linkQueue<double>;
Stack<double> nit = new arrayStack<double>;
nit.push(3.14159);
act.enqueue(2.718286);
act.enqueue(43.23);
nit.push(12.43);
act.enqueue(nit.pop());
act.enqueue(nit.pop());
```

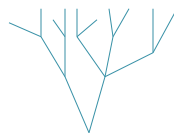
Trees

General tree

- A general tree T is a set of one or more nodes such that T is partitioned into disjoint subsets:
 - A single node r , the root
 - Sets that are general trees, called subtrees of r

Binary tree

- A binary tree is a set T of nodes such that either
 - T is empty, or
 - T is partitioned into three disjoint subsets:
 - A single node r , the root
 - Two possibly empty sets that are binary trees, called left and right subtrees of r



Terminology

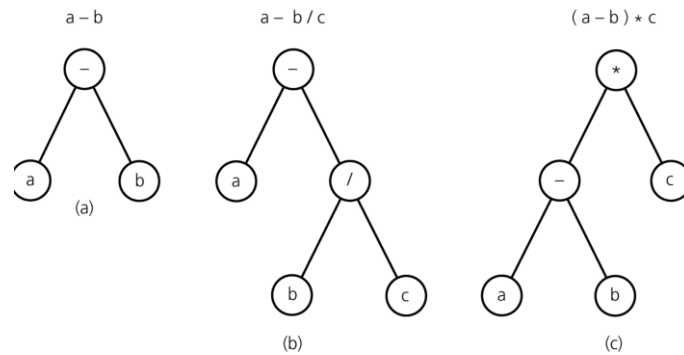


Figure 11-4

Binary trees that represent algebraic expressions

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Terminology

A binary search tree

- A binary tree that has the following properties for each node n
 - n 's value is greater than all values in its left subtree T_L
 - n 's value is less than all values in its right subtree T_R
 - Both T_L and T_R are binary search trees

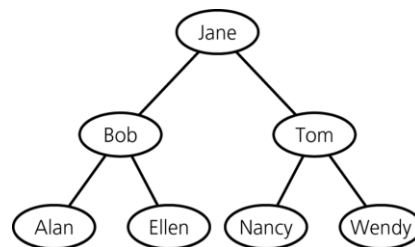


Figure 11-5

A binary search tree of names

Tree Terminology

Level of a node, n , in a Tree, T

- If n is the root of T , it is at level 1
- If n is not the root of T , its level is 1 greater than the level of its parent

Height of a tree, T

- If T is empty, its height is 0
- If T is not empty, its height is equal to the maximum level of its nodes

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Terminology

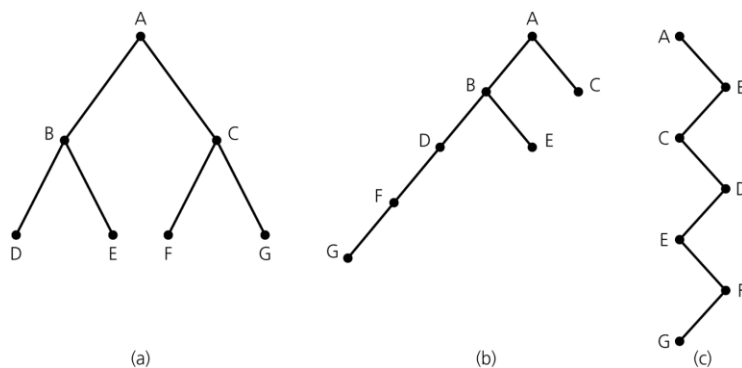


Figure 11-6

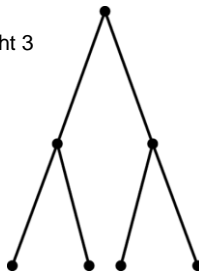
Binary trees with the same nodes but different heights

Definitions: Full, Complete, Balanced

Full Binary Trees

- If T is empty, T is a *full* binary tree of height 0
- If T is not empty and has height $h > 0$, T is a full binary tree if its root's subtrees are both full binary trees of height $h - 1$

Example:
A full binary tree, height 3



Observe: This
is a *Recursive
Definition!!*

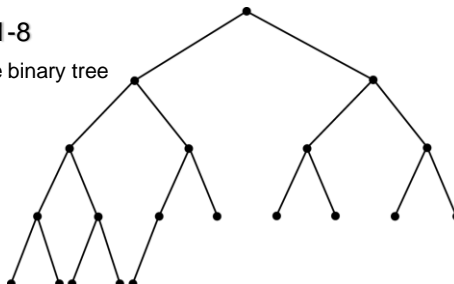
Definitions: Full, Complete, Balanced

Complete binary trees

- A binary tree T of height h is *complete* if
 - All nodes at level $h - 2$ and above have two children each, and
 - When a node at level $h - 1$ has children, all nodes to its left at the same level have two children each, and
 - When a node at level $h - 1$ has one child, it is a left child

Figure 11-8

A complete binary tree



Another
*Recursive
Definition!!*

Definitions: Full, Complete, Balanced

Balanced binary trees

- A binary tree is *balanced* if the height of any node's right subtree differs from the height of the node's left subtree by no more than 1

Full binary trees are complete

Complete binary trees are balanced

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Tree Terminology

General Tree

- A set of one or more nodes, partitioned into a root node and subsets that are general subtrees of the root

Parent of node n

- The node directly above node n in the tree

Child of node n

- A node directly below node n in the tree

Root

- The only node in the tree with no parent

Leaf

- A node with no children

Siblings

- Nodes with a common parent

Ancestor of node n

- A node on the path from the root to n

Descendant of node n

- A node on a path from n to a leaf

Subtree of node n

- A tree that consists of a child (if any) of n and the child's descendants

Tree Terminology: Summary

Height

- Number of nodes on the longest path from the root to a leaf

Binary tree

- A set of nodes that is either empty or partitioned into a root node and one or two subsets that are binary subtrees of the root
- Each node has at most two children: left child & right child

Left (right) child of node n

- A node directly below and left (right) of node n in a binary tree

Left (right) subtree of node n

- In a binary tree, left (right) child (if any) of n plus its descendants

Binary search tree

- The value in any node n is greater than every node in n 's left subtree, but less than every node in n 's right subtree

Empty binary tree

- A binary tree with no nodes

Tree Terminology: Summary

Full binary tree

- A binary tree of height h with no missing nodes
- All leaves are at level h and all other nodes each have two children

Complete binary tree

- A binary tree of height h that is full to level $h - 1$ and has level h filled in from left to right

Balanced binary tree

- A binary tree in which the left and right subtrees of any node have heights that differ by at most 1

Observe: These definitions are not given recursively!!

The ADT Binary Tree: Basic Operations of the ADT Binary Tree

- The operations available for a particular ADT binary tree depend on the type of binary tree being implemented
- Basic operations of the ADT binary tree
 - `createBinaryTree()`
 - `createBinaryTree(rootItem)`
 - `makeEmpty()`
 - `isEmpty()`
 - `getRootItem()` throws `TreeException`

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General Operations of the ADT Binary Tree

- General operations of the ADT binary tree
 - `createBinaryTree (rootItem, leftTree, rightTree)`
 - `setRootItem(newItem)`
 - `attachLeft(newItem)` throws `TreeException`
 - `attachRight(newItem)` throws `TreeException`
 - `attachLeftSubtree(leftTree)` throws `TreeException`
 - `attachRightSubtree(rightTree)` throws `TreeException`
 - `detachLeftSubtree()` throws `TreeException`
 - `detachRightSubtree()` throws `TreeException`

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Traversals of a Binary Tree

- A traversal algorithm for a binary tree visits each node in the tree
- Recursive traversal algorithms
 - Preorder traversal
 - Inorder traversal
 - Postorder traversal
- Traversal is $O(n)$

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Traversal of a Binary Tree

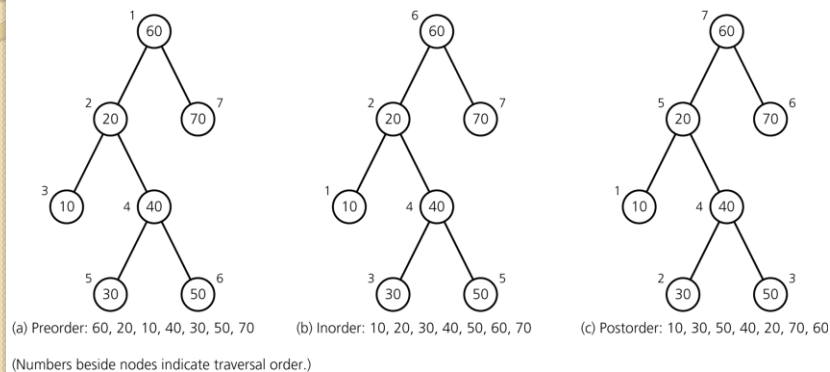


Figure 11-10

Traversals of a binary tree: a) preorder; b) inorder; c) postorder

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