

CANDIDATE

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**TEST** 

## Quiz 6

Subject code	
Evaluation type	<del></del>
Test opening time	27.03.2024 07:00
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Question	Status	Marks	Question type
1.1	Correct	1/1	Multiple Response
1.2	Correct	1/1	True / False
1.3	Correct	1/1	Multiple Response
1.4	Correct	1/1	Multiple Response
1.5	Correct	1/1	Multiple Response
2.1	Correct	1/1	True / False
2.2	Correct	1/1	True / False
2.3	Correct	1/1	Multiple Response
2.4	Correct	1/1	Multiple Response
2.5	Correct	1/1	Multiple Response

**1.1** Suppose you are trying to prove that the square of every natural number has remainder 0 or 1 on division by 4. If P(n) is the proposition:

$$P(n): \qquad n^2 =_{(4)} 0 \quad \text{or} \quad n^2 =_{(4)} 1$$

which of the following would be sufficient to prove the result by induction?

Select all that apply:

 $\square$  P(0); and for all  $k \in \mathbb{N}$  P( $k^2$ )  $\Rightarrow$  P( $(k+1)^2$ )

|--|

- None of the other options
- $\square$  P(0); P(1); and for all  $k \in \mathbb{N}$  P(k)  $\Rightarrow$  P(k+4)
- $\square$  P(0); for all  $k \in \mathbb{N}$  P(k)  $\Rightarrow$  P(k+2); and for all  $k \in \mathbb{N}$  P(k+1)  $\Rightarrow$  P(k+3)
- $\square$  P(0); and for all  $k \in \mathbb{N}$  P( $k^2$ )  $\Rightarrow$  P( $k^2+1$ )

- **1.2** Suppose  $f: \mathbb{N} \times \mathbb{N} \to \mathbb{N}$  is defined recursively as follows for all  $m, n \in \mathbb{N}$ :
  - f(0,m) = m
  - f(n+1,m) = 1 + f(n,m)

True or false:

f(m,n) = f(n,m) for all  $m,n \in \mathbb{N}$ 





False

**1.3** Let C(n) be the maximum number of elementary operations required to compute A+B where A and B are n×n matrices.

What can be said about the asymptotic behaviour of C(n)? Select all that apply

 $\square$  C(n)  $\subseteq$  O(log n)

$C(n) \subseteq O(\log n)$	
$C(n) \subseteq O(n^3)$	
■ None of these options	
$\square$ $C(n) \subseteq O(1)$	
$\square$ $C(n) \in O(n)$	
$ ightharpoonup C(n) \in O(n^2)$	<b>⊘</b>

**1.4** Consider the following code snippet:

myFunc(n):	myFunc(n):	
i = 1		
while i < n:		
j = n		
while j > 0:		
print('*')		
j = j/3		
i = i+2		

Which of the following hold with regard to the running time T(n) of this code (assume / is integer division)

Select all that apply

$ ightharpoonup T(n) \in O(n^2)$	•
$ ule{transformation} T(n) \in O(n \log n)$	•
$  T(n) \in O(\log n) $	
■ None of these options	
$\Box$ T(n) $\in$ O(n)	

**1.5** Consider the following code snippet:

myFunc(n):	
if n==0:	
return 1	
i = 0	
x = 0	
while i < n:	
x = x + myFunc(0)	
i = i+1	
return x	

Which of the following hold with regard to the running time  $\mathsf{T}(\mathsf{n})$  of this code Select all that apply

$ ightharpoonup T(n) \in O(2^n)$	•
$ ule{transformation} T(n) \in O(n \log n)$	•
$T(n) \in O(n)$	•
■ None of these options	
$ ightharpoonup T(n) \in O(n^2)$	•

- 2.1 Let EXP be as defined in Quiz 6, Question M3. That is:
  - $\varnothing$  and  $\mathcal{U}$  are elements of EXP
  - X,Y,Z ∈ EXP
  - If  $E \in EXP$ , then:
    - ∘ (E) ∈ EXP
    - Ec ∈ EXP
  - If  $E_1, E_2 \in EXP$ , then:
    - $\circ$  (E<sub>1</sub>  $\cap$  E<sub>2</sub>)  $\in$  EXP
    - $\circ$  (E<sub>1</sub>  $\cup$  E<sub>2</sub>)  $\in$  EXP

Suppose we can show that the following holds for a proposition P:

- P(∅) holds
- $P(\mathcal{U})$  holds
- P(X), P(Y), and P(Z) hold
- If P(E) holds then P((E)) holds
- If P(E) holds then P(Ec) holds
- If P(E₁) and P(E₂) hold, then P((E₁ ∩ E₂)) holds
- If  $P(E_1)$  or  $P(E_2)$  hold, then  $P((E_1 \cup E_2))$  holds

True or false: P(E) holds for all  $E \in EXP$ 

Select one alternative:

False





- **2.2** Suppose  $f: \mathbb{N} \times \mathbb{N} \to \mathbb{N}$  is defined recursively as follows for all  $m, n \in \mathbb{N}$ :
  - f(0,m) = 0
  - f(n+1,m) = m + f(n,m)

True or false:

f(m,n) = f(n,m) for all  $m,n \in \mathbb{N}$ 

False





2.3 Which of the following are true for all functions  $f: \mathbb{N} \to \mathbb{N}$  with  $f(n) \in \Theta(n)$ ? Select all that apply

$ ightharpoonup f \circ f \in O(n^2)$	•
${\color{red} leve{ }}$ $f \circ f \in \Omega(n)$	•
✓ f∘f ∈ O(n)	•

**2.4** Consider the following code fragment that works on a (non-empty) array:

myFunc(A):
if length(A) <= 5:
return A[0]
i = 0
while i < length(A):
B = A[i, i+5) # Take the i - (i+4)th elements of A
C[i/5] = myFunc(B)
i = i+5
return <b>myFunc</b> (C)

If n = length(A), which of the following hold for T(n), the running time of the above code?

## Select all that apply

- $\square$  T(n)  $\subseteq$  O( log n)
- $\Box$  T(n)  $\in$  O(1)
- None of these options

✓ T(	$(n) \in O(n^2)$	•
<b>✓</b> T(	$(n) \in O(n \log n)$	•
✓ T(	$(n) \in O(n)$	•

**2.5** Consider the following code snippet:

myFunc(n):
if n==0:
return 1
else:
return myFunc(n-1) + myFunc2(n-1)
myFunc2(n):
if n==0:
return 0
else:
return myFunc(n-1)

Which of the following hold with regard to the running time T(n) of this code **Select all that apply** 

- $\square$  T(n)  $\in$  O(n)
- $\square$  T(n)  $\subseteq$  O(n<sup>2</sup>)
- $Arr T(n) \in O(2^n)$



- None of these options
- $\square$  T(n)  $\subseteq$  O(log n)