## INTERVIEW CNRS (SECTION 6)

· E(w(T15,517) - w(T15,5977))

(x,y) = Tx T =

45. learn ar(TIS,S)

WITH ENTINE

Ely(TRO-u(Troca) -pattay condo

 $= \sum_{i,j} V_i \, v_j \, \, \mathbb{E} \Big[ \, \underline{1}(x_1) \, \, \underline{1}(\hat{p}_i x_i) \Big] = \sum_i V_i \, \, \, \mathbb{E} \big[ \, \underline{x}_i^i x_i \big] + \sum_{(i,j)} \, \underline{x}_i \, v_j^i \, \mathbb{E} \big[ \, \underline{x}_i^i (x_i^i) \big]$ 

 $= \frac{1}{K} \cdot \|\mathbf{v}\|_{L^{\infty}}^{L} + \frac{1}{K} \frac{1}{L^{\infty}} \cdot \|\mathbf{v}\|_{L^{\infty}}^{L} + \frac{1}{K} \frac{1}{L^{\infty}} \cdot \|\mathbf{v}\|_{L^{\infty}}^{L} - \frac{1}{L^{\infty}} \|\mathbf{v}\|_{L^{\infty}}^{L} + \frac{1}{L^{\infty}}$ 

 $E[X] = \frac{k}{n} \|V\|_{1}$ ,  $E[I_{(k1)}] = \frac{k}{n}$ ,  $E[I_{(k1)}]_{(k1)}$ 

Simon Apers (CWI, ULB)

April 2021

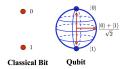
X= |Vs| . Ty 1/(x) rordon voriable

 $E[X_t] = E[(\sum_{i=1}^{k} k_i)_i] = E[\sum_{i=1}^{k} k_i k_i]$ 

- E [ = 7, v, 2(in) 2(is)

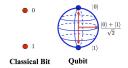
## QUANTUM COMPUTERS

fundamentally new way of computing, built on quantum theory

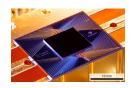


## QUANTUM COMPUTERS

fundamentally new way of computing, built on quantum theory



on cusp of revolution: first quantum computing devices being built



(Google Sycamore quantum processor)

ı

# critical to understand: what are quantum computers good for?

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 cryptography and communication (network of quantum computers)



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 cryptography and communication (network of guantum computers)



• algorithms and simulation



### BACKGROUND

2014 - 2018: PhD student at **Ghent University**⊃ long term visit to **University of Calgary** 

2019 - 2020: postdoc at **INRIA** (Paris) and **CWI** (Amsterdam)

2020 - ... : postdoc at **CWI** (Amsterdam) and **ULB** (Brussels)

## RESEARCH THEMES

- quantum algorithms and quantum walks
  - graph theory and random walks

### **PUBLICATIONS**

- 5 conference publications (FOCS, STACS, ESA, ...)
  - 6 journal publications (PRL, Quantum, PRA, ...)
    - among those, 2 single author publications
      - 3 preprints (2 submitted)

## CONTRIBUTION

$$qLx = b$$

or

"quantum speedups in modern optimization"

#### starting point is

## "Quantum Speedup for Graph Sparsification, Cut Approximation and Laplacian Solving"

by Simon Apers and Ronald de Wolf

in FOCS'20

presented at QIP'21

invited talks at Simons institute (Berkeley), University of Cambridge, University of Bristol, . . .

## MODERN OPTIMIZATION

combinatorics  $\leftrightarrow$  linear algebra

graph 
$$G = igcup_{G} \longleftrightarrow Laplacian L_G = egin{bmatrix} 3 & -1 & -1 & -1 \ -1 & 2 & -1 & 0 \ -1 & -1 & 3 & -1 \ -1 & 0 & -1 & 2 \ \end{bmatrix}$$

electrical flow on  $G \leftrightarrow \text{Laplacian system } L_{G}x = b$ 

### MODERN OPTIMIZATION

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electrical flow on  $G \leftrightarrow \text{Laplacian system } L_{G}x = b$ 

landmark result by Spielman and Teng ('04): solve Laplacian system  $L_{G^X}=b$  in near-linear time  $\widetilde{O}(m)$ 



# faster algorithms for many fundamental problems in optimization and learning

flow problems

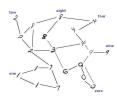
(max flow, min cost flow, ...)

learning on graphs

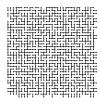
(clustering, semi-supervised learning, ...)

combinatorial optimization

(graph matching, matrix scaling, random spanning trees, ...)



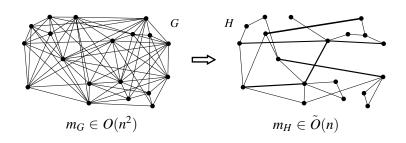
(semi-supervised learning)



(random spanning tree)

## cornerstone of fast Laplacian solvers:

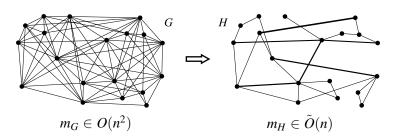
## graph sparsification



such that  $G \approx H$ 

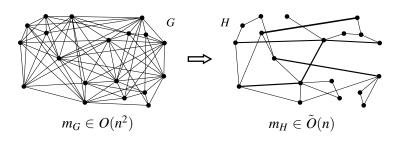
## classically: sparsification in time

$$\widetilde{O}(m) \in \widetilde{O}(n^2)$$



## classically: sparsification in time

$$\widetilde{O}(m) \in \widetilde{O}(n^2)$$



Apers - de Wolf '19: (optimal) quantum speedup for graph sparsification

$$\widetilde{O}(\sqrt{mn}) \in \widetilde{O}(n^{3/2})$$

### "LOW-HANGING FRUIT":

first quantum speedups for approximating max cut, min cut, sparsest cut, ... √ first quantum speedups for approximately solving Laplacian systems, electrical flows, . . .

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first quantum speedups for approximating max cut, min cut, sparsest cut, ... √ first quantum speedups for approximately solving Laplacian systems, electrical flows, . . .

↓ "HIGHER-HANGING FRUIT":

1

✓ first quantum speedup for exact min cut (Apers-Lee '20, Apers-Gawrychowski-Lee '21) ? quantum speedups for modern combinatorial optimization algorithms (IPMs, max flow, learning, ...)

# EXISTING WORK: QUANTUM ALGORITHMS FOR LINEAR SYSTEMS $Ax = b \label{eq:algorithm}$

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• quantum solution  $|b\rangle$ : speedup for sparse, well-conditioned systems = key primitive in quantum algorithms (Harrow-Hassidim-Lloyd '08,...)

# EXISTING WORK: QUANTUM ALGORITHMS FOR LINEAR SYSTEMS $Ax = b \label{eq:algorithms}$

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explicit solution b:no speedups known

# EXISTING WORK: QUANTUM ALGORITHMS FOR LINEAR SYSTEMS Ax=b

quantum solution |b⟩:
 speedup for sparse, well-conditioned systems
 = key primitive in quantum algorithms
 (Harrow-Hassidim-Lloyd '08,...)

explicit solution b:no speedups known



our work: quantum speedup for explicit solution of general Laplacian system

 quantum speedup over "classic" algorithms for spanning tree, max flow, . . .

(Dürr (CNRS, LIP6) et al. '06, Ambainis-Špalek '06, ...)

 quantum speedup over "classic" algorithms for spanning tree, max flow, . . .

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quantum SDP solversopen question: quantum speedup for max cut?

(Brandão-Svore '17, van Apeldoorn et al. '17, ...)

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quantum SDP solvers open question: quantum speedup for max cut?

(Brandão-Svore '17, van Apeldoorn et al. '17, ...)

 $\downarrow$ 

### our work:

quantum speedup over "modern" algorithms, quantum speedup for max cut

## **TECHNIQUES:**

- sequential quantum algorithm built from distributed classical algorithm
  - 2. quantum speedup for graph spanners
- 3. simulate randomness using hash functions

connection to distributed algorithms & streaming algorithms



= important trend in modern quantum algorithms!

# connection to distributed algorithms & streaming algorithms



= important trend in modern quantum algorithms!

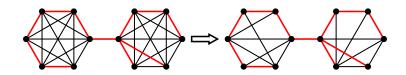


relevant expertise in IRIF/LIP6 on distributed/streaming algorithms

based on distributed sparsification algorithm (Koutis-Xu '16):

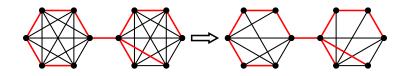
based on distributed sparsification algorithm (Koutis-Xu '16):

- a) detect "important" edges by constructing spanner (see next slide)
- b) sample remaining edges with constant probability



based on distributed sparsification algorithm (Koutis-Xu '16):

- a) detect "important" edges by constructing spanner (see next slide)
- b) sample remaining edges with constant probability



c) repeat a)-b) until  $\widetilde{O}(n)$  edges

## 2. QUANTUM SPEEDUP FOR GRAPH SPANNERS

"graph spanner" = sparse subgraph that approximately preserves distance

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quantum speedup:

 $\widetilde{O}(m)$  classical  $\longrightarrow \widetilde{O}(\sqrt{mn})$  quantum

buildings blocks:

Thorup-Zwick ('01) spanner algorithm & Dürr et al. ('06) quantum algorithm for shortest-path tree

#### 2. QUANTUM SPEEDUP FOR GRAPH SPANNERS

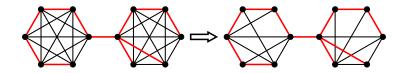
## follow-up work:

build on distributed spanner algorithm by Baswana et al. ('12)

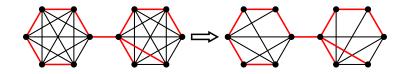


faster quantum algorithms for spanners and sparsification of *unweighted* graphs (Apers-Lee-Le Gall, in preparation)

- a) detect "important" edges by constructing spanner
- b) sample remaining edges with constant probability



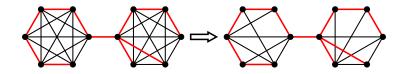
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in sublinear time,

can explicitly write down spanners

- a) detect "important" edges by constructing spanner
- b) sample remaining edges with constant probability



in sublinear time,

can explicitly write down spanners

cannot explicitly sample  $\Omega(n^2)$  edges  $(=\Omega(n^2)$  coin tosses)

 $n^2$  coin tosses

1 1 0 1 0 0 1 1 0 0 0 ...

 $n^2$  coin tosses

↑ access?

sublinear  $k \in o(n^2)$ -time quantum algorithm

 $n^2$  coin tosses

1 1 0 1 0 0 1 1 0 0 0 ...

↑ access?

sublinear  $k \in o(n^2)$ -time quantum algorithm

### solution:



### FOLLOW-UP/RELATED WORK:

- "Quantum complexity of minimum cut"
   by Apers and Lee (submitted to CCC '21)
- "Faster quantum algorithms for unweighted spanners and sparsification" by Apers, Lee and Le Gall \*
  - "Approximate minimum cut partition in near-linear time" by Apers, Gawrychowski and Lee \*
    - "Quantum cut query complexity of minimum cut"
       by Apers, Efron, Lee, Mukhopadhyay and Nanongkai \*

\* in preparation

= collaborations with
University of Technology Sydney (Australia),
Nagoya University (Japan),
University of Wrocław (Poland),
University of Toronto (Canada),
KTH Royal Institute of Technology (Sweden),
University of Copenhagen (Denmark)



### RESEARCH PROGRAMME:

## **QUANTUM ALGORITHMS**

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### **QUANTUM ALGORITHMS**

**FOR** 

# NEAR-TERM QUANTUM COMPUTERS:

 $\sim$  applications

focus on sampling problems

### RESEARCH PROGRAMME:

### **QUANTUM ALGORITHMS**

**FOR** 

NEAR-TERM
QUANTUM COMPUTERS:

 $\sim$  applications

focus on sampling problems

LONG-TERM
QUANTUM COMPUTERS:

 $\sim$  foundations

focus on optimization

# NEAR-TERM QUANTUM COMPUTERS and their applications

'19: Google claims "quantum supremacy" =

**artificial** sampling problem solved exponentially faster on quantum device

# NEAR-TERM QUANTUM COMPUTERS and their applications

'19: Google claims "quantum supremacy"

artificial sampling problem solved exponentially faster on quantum device



? ∃ natural sampling problem solved exponentially faster on quantum device

# e.g., near-term quantum devices for **Monte Carlo method?**

## e.g., near-term quantum devices for Monte Carlo method?

=

use of sampling to approximate intractable problems

## e.g., near-term quantum devices for Monte Carlo method?

=

use of sampling to approximate intractable problems

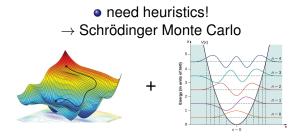
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key algorithmic technique in

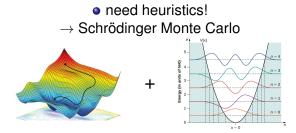
- machine learning and statistics
  - combinatorial optimization
    - finance, chemistry, . . .

Monte Carlo method on near-term quantum devices?

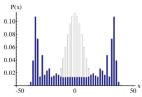
## Monte Carlo method on near-term quantum devices?



### Monte Carlo method on near-term quantum devices?



- quantum walk speedup?
- ightarrow build on *quantum fast-forwarding* (Apers-Sarlette QIC'19)
  - ightarrow ongoing work with Laurent Miclo (CNRS, Toulouse)





use cases of large error-corrected quantum computers?

•

use cases of large error-corrected quantum computers?

•

fundamental limits of quantum computation?

•

use cases of large error-corrected quantum computers?

•

fundamental limits of quantum computation?

•

new algorithmic primitives?

## • qLx=b or quantum Laplacian paradigm

quantum algorithms for optimization, machine learning, ...

## qLx=b or quantum Laplacian paradigm

quantum algorithms for optimization, machine learning, ...

theory of quantum walks

connection with electric networks algorithms for graph problems

## qLx=b or quantum Laplacian paradigm

quantum algorithms for optimization, machine learning, ...

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• exponential quantum speedups?

candidate: submodular function optimization

## Paris region:

## Paris region:

quantum hub

PARIS CENTRE FOR QUANTUM COMPUTING



## Paris region:

quantum hub



• bright prospects thanks to plan quantique

## Paris region:

quantum hub



- bright prospects thanks to plan quantique
- connections with INRIA: postdoc at COSMIQ INRIA-Microsoft joint project, partially based on QFF (Apers-Sarlette QIC'19)

### IRIF:

- most overlap with foundations of quantum algorithms
- collaboration:
  - quantum algorithms/complexity (Laplante, Magniez, Santha)
  - quantum machine learning (Kerenidis)
  - combinatorial optimization (Mathieu, Vladu)
  - distributed/streaming algorithms (Fraigniaud, Rosén)
  - graph theory/combinatorics (Charbit, Habib)

### LIP6:

- most overlap with near-term applications of quantum computers
- collaboration:
  - quantum computation (Diamanti, Grilo, Grosshans, Kashefi, Markham)
  - operations research (Dürr, Escoffier, Cohen-Addad)
  - learning theory (Gallinari, Piwowarski)

### SIMON APERS

'14-'18: Ghent University
'19-'20: INRIA (Paris) and CWI (Amsterdam)
'20-'21: CWI (Amsterdam) and ULB (Brussels)

10 publications (FOCS, PRL, ...), 3 preprints

# focus on quantum algorithms

research programme on

- 1. near-term quantum computers for sampling
- 2. long-term quantum computers for optimization

 <sup>\*</sup> application updates: ● awarded Marie Skłodowska-Curie Seal of Excellence
 ● invited to European Symposium on Algorithms '21 PC

#### NETWORK

(active collaboration or paper under review in past year)

Sweden: Danupon Nanongkai (KTH), Sagnik Mukhopadhyay (KTH), Florian

Adriaens (KTH)

France: Anthony Leverrier (INRIA), Christophe Vuillot (INRIA), Alain Sarlette

(INRIA), Laurent Miclo (Toulouse)

Belgium: Jérémie Roland (ULB), Shantanav Chakraborty (ULB), Leonardo

Novo (ULB)

Italy: Francesco Ticozzi (Padova)

the Netherlands: Stacey Jeffery (CWI), Ronald de Wolf (CWI), Ido Niesen (CWI), Kirill

Neklyudov (CWI)

Poland: Pawel Gawrychowski (Wrocław)

USA: András Gilyén (Caltech)
Canada: Yuval Efron (Toronto)

Australia: Troy Lee (Sydney)

Japan: Francois Le Gall (Nagoya University)

quantum algorithms/complexity, combinatorial optimization, data science/machine learning, quantum control, quantum error correction, probability theory, data structures

#### IMAGE CREDITS

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- Quantum internet (slide 2): KPN (link).
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   Proceedings of the 20th International conference on Machine learning (ICML-03). 2003.
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- World map (slide 22): clker (link).
- Optimization landscape (slide 26): Medium (link).
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- Quantum walks plot (slide 26): University of Paderborn (link).

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