

Exercises 1: QFT, phase estimation and Shor's algorithm

Lecturer: Simon Apers (apers@irif.fr)

Exercise 1 (Oracles). For accessing a function $f : \{0, 1\}^n \rightarrow \{0, 1\}$ with a quantum circuit, we use a *bit oracle* O_b or a *phase oracle* O_p . For $z \in \{0, 1\}^n$ and $w \in \{0, 1\}$, these are defined as follows:

$$\begin{array}{c} |z\rangle \\ |w\rangle \end{array} \begin{array}{c} \boxed{O_b} \\ \end{array} \begin{array}{c} |z\rangle \\ |w \oplus f(z)\rangle \end{array} \quad \quad \quad |z\rangle \begin{array}{c} \boxed{O_p} \\ \end{array} (-1)^{f(z)} |z\rangle$$

We can show that both oracles are equivalent in a sense.

- Show that the phase oracle can simulate the bit oracle:

$$\begin{array}{c} |z\rangle \\ |w\rangle \end{array} \begin{array}{c} \boxed{O_p} \\ \downarrow \\ \boxed{H} \bullet \boxed{H} \end{array} \equiv \begin{array}{c} |z\rangle \\ |w\rangle \end{array} \begin{array}{c} \boxed{O_b} \\ \end{array}$$

- Show that the bit oracle can simulate the phase oracle:

$$\begin{array}{c} |z\rangle \\ |1\rangle \end{array} \begin{array}{c} \boxed{O_b} \\ \uparrow \downarrow \\ \boxed{H} \quad \boxed{H} \end{array} \equiv \begin{array}{c} |z\rangle \\ |1\rangle \end{array} \begin{array}{c} \boxed{O_p} \\ \end{array}$$

Exercise 2 (Controlled unitary). Recall the controlled unitary gate:

$$\begin{array}{c} |k\rangle \\ |\psi\rangle \end{array} \begin{array}{c} \bullet \\ \downarrow \\ \boxed{cU} \end{array} \begin{array}{c} |k\rangle \\ U^k |\psi\rangle \end{array}$$

where $k = \sum_{j=0}^{n-1} k_j 2^j$ is an n -bit integer. Expand this gate into more elementary gates of the form

$$\begin{array}{c} |k_s\rangle \\ |\psi\rangle \end{array} \begin{array}{c} \bullet \\ \downarrow \\ \boxed{U^{2^s}} \end{array} \begin{array}{c} |k_s\rangle \\ U^{k_s 2^s} |\psi\rangle \end{array}$$

for $k_s \in \{0, 1\}$ and $s \in \{0, 1, \dots, n-1\}$.

Exercise 3 (Hadamard transform). A variation on the quantum Fourier transform is the Hadamard transform H_N for $N = 2^n$. It is defined by $H_N = H^{\otimes n}$, which corresponds to the circuit

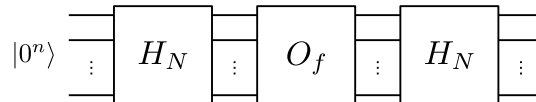
$$\begin{array}{c} \vdots \\ \vdots \end{array} \begin{array}{c} \boxed{H_N} \\ \end{array} \begin{array}{c} \vdots \\ \vdots \end{array} \equiv \begin{array}{c} \boxed{H} \\ \boxed{H} \\ \vdots \\ \boxed{H} \end{array}$$

- What is $H_N |0^n\rangle$ equal to?
- Let $k = \sum_{j=0}^{n-1} k_j 2^j$. What is $H_N |k\rangle = H_N |k_0 \dots k_{n-1}\rangle$ equal to?
(Hint: Use the inner product $x \cdot k = \sum_{\ell} x_{\ell} k_{\ell}$, and use that $H |k_{\ell}\rangle = \frac{1}{\sqrt{2}} \sum_{x_{\ell}=0}^1 (-1)^{x_{\ell} k_{\ell}} |x_{\ell}\rangle$.)

Exercise 4 (Bernstein-Vazirani algorithm). Let $N = 2^n$. Consider a function $f : \{0, 1\}^n \rightarrow \{0, 1\}$ that is determined by some hidden string $a \in \{0, 1\}^n$ in the following way:

$$f(x) = (x \cdot a) \pmod{2}.$$

We can access the function through the phase oracle $O_f |x\rangle = (-1)^{f(x)} |x\rangle$. What is the output of the following circuit?



Exercise 5 (Fourier analysis (extra)). Consider natural numbers q, m, r, s such that $q = mr$ and $s < r$. Prove the following critical identity in Shor's algorithm for period finding:

$$\frac{1}{\sqrt{m}} \sum_{j=0}^{m-1} |s + jr\rangle \xrightarrow{F_q^\dagger} \frac{1}{\sqrt{r}} \sum_{\ell=0}^{r-1} \omega_q^{-s\ell m} |\ell m\rangle,$$

$$\begin{array}{c} 0 \quad \uparrow \quad \uparrow \quad \dots \quad \uparrow \quad \dots \quad \uparrow \quad q-1 \\ \quad s \quad s+r \quad \quad \quad s+jr \quad \quad \quad s+(m-1)r \end{array} \xrightarrow{F_q^\dagger} \begin{array}{c} \omega_q^{s\ell m} / \sqrt{r} \\ \uparrow \quad \uparrow \quad \uparrow \quad \dots \quad \uparrow \quad \dots \quad \uparrow \quad q-1 \\ 0 \quad m \quad 2m \quad \quad \quad \ell m \quad \quad \quad (r-1)m \end{array}$$