



COMP 412  
FALL 2010

# *The Procedure Abstraction, Part VI: Inheritance in OOLs*

## *Comp 412*

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# What About Inheritance?

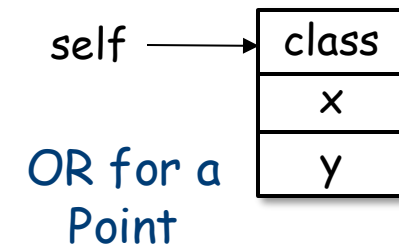
## Impact on OR Layout

Assume single inheritance

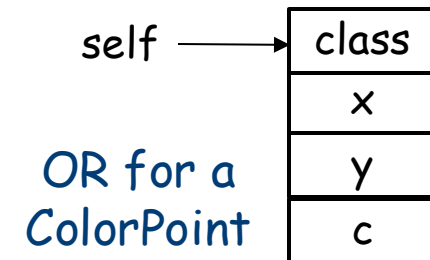
- OR needs slots for each member declared, all the way up the class hierarchy (class, superclass, super-superclass, ...)
- Can use **prefixing of storage** to lay out the OR

## Back to Our Java Example — Class Point

```
Class Point {  
    public int x, y;  
    ...  
}
```



```
Class ColorPoint extends Point {  
    Color c;  
    ...  
}
```



What happens if we cast a ColorPoint to a Point?

Take the word extends literally.



# Open World versus Closed World

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Prefixing assumes that the class structure is known when layout is performed. Two common cases occur.

## Closed-World Assumption

(Compile time)

- Class structure is known and closed prior to runtime
- Can lay out ORs in the compiler and/or the linker

## Open-World Assumption

(Interpreter or JIT)

- Class structure can change at runtime
- Cannot lay out ORs until they are allocated
  - Walk class hierarchy at allocation

C++ has a closed class structure.

Java as an open class structure.



# Open World versus Closed World

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What happens if the class structure changes for a class with active instantiated objects?

- Oops. That might be a problem.

## Changes to the structure of an instantiated class

- Languages differ on legality and advisability of such changes
- Smalltalk-80, for example, found & fixed all ORs in the class
- Python does not support such changes
  - *Programmer can, on her own, track instances and re-instantiate them ...*



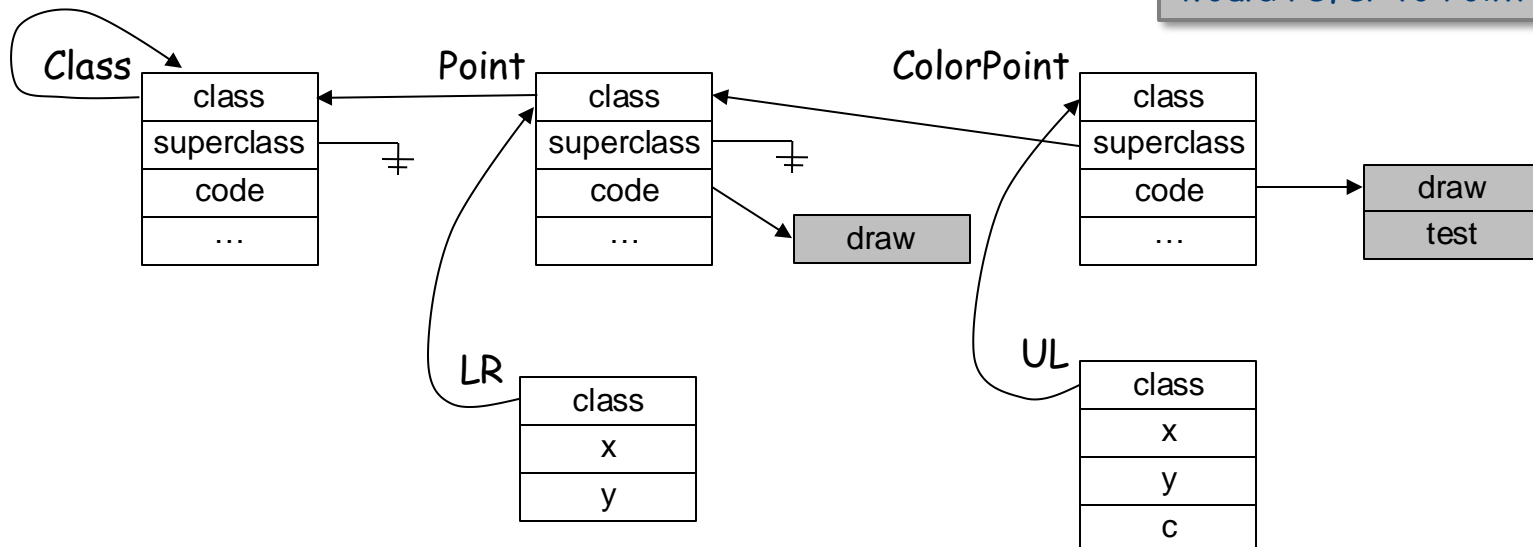
# What About Code Members?

How does the language's runtime environment find the code for a given method invocation?

## Closed Class Structure

- Mapping of names to methods is static and known (C++)
  - Fixed offsets & indirect calls
- Virtual functions force runtime resolution

If ColorPoint inherited draw from Point, its code vector would refer to Point's draw.



UL finds draw at offset 0 in ColorPoint's code vector

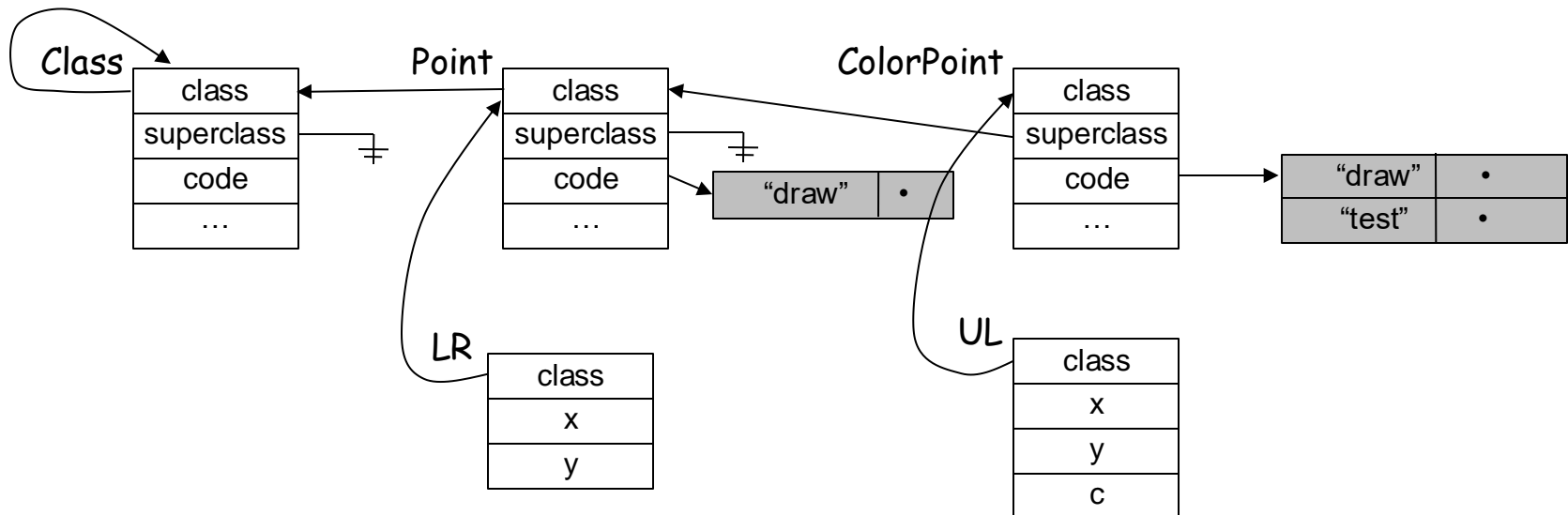


# What About Code Members?

How does the language's runtime environment find the code for a given method invocation?

## Open Class Structure

- Dynamic mapping, unknown until runtime
- In general case, need runtime representation of hierarchy
  - Lookup by textual name in class' table of methods



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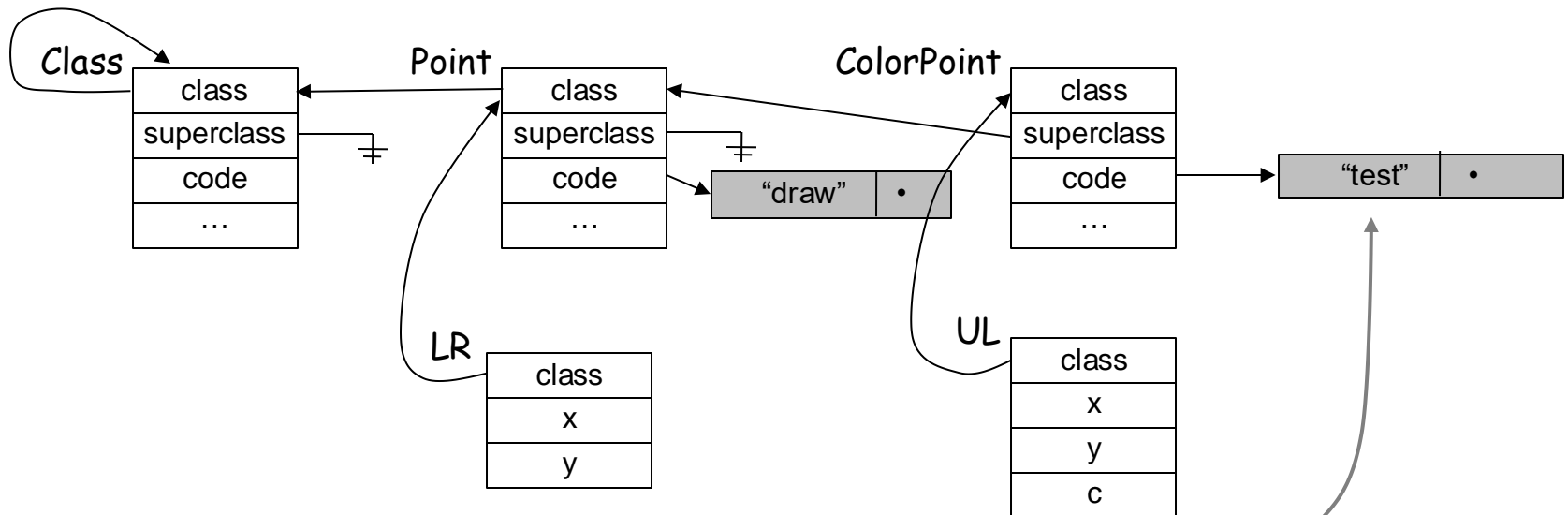


## What About Code Members?

## How does the language's runtime environment find the code for a given method invocation?

# Open Class Structure

- Dynamic mapping, unknown until runtime
- In general case, need runtime representation of hierarchy
  - Lookup by textual name in class' table of methods



If ColorPoint inherited draw from Point, its code vector would lack a draw.



# The Single Inheritance Hierarchy

## Two distinct philosophies

### Closed Class Structure (C++)

- Can map name to code at compile time
- Leads to 1-level code vector
- Copy superclass methods
- Fixed offsets & indirect calls
- Less flexible & expressive

### Open Class Structure (ST80)

- Cannot map name to code at compile time
- Multiple jump vectors (1/class)
- Must search for method
- Run-time lookup caching
- Much **more expensive** to run

## Impact of OOL on Program's Name Space

- Method can see values in the code's lexical hierarchy
- Method can see members of self, class, & superclasses
- Many different levels where a value can reside

OOL differs from ALL in the shape of its name space AND in the mechanism used to bind names to implementations





# The Single Inheritance Hierarchy

## Two distinct philosophies

### Closed Class Structure (C++)

- Can map name to code at compile time
- Leads to 1-level
- Copy superclass
- Fixed offsets
- Less flexible

### Runtime changes

- Layout OR by prefixing
- Limit changes in instantiated classes
- Use 1-level code vectors and rebuild on hierarchy change
- Fixed offsets & 1-level of indirection on calls

### Open Class Structure (ST80)

- Can map name to code at runtime
- Use 1-level code vectors (1/class) for method lookup caching
- **expensive** to run

## Impact of OOL on Program's Name Space

- Method can see values in the code's lexical hierarchy
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# What About Multiple Inheritance?

## The Idea

- Let  $C$  be a subclass of both  $A$  and  $B$
- $C$  draws some, but not all, methods from  $A$
- $C$  draws some, but not all, methods from  $B$

Need a linguistic mechanism to specify partial inheritance

## Problems arise when $C$ inherits from both $A$ & $B$

- $C$ 's OR can extend  $A$  or  $B$ , but not both
- $C$ 's code vector can extend  $A$  or  $B$ , but not both
- Some class  $D$  might inherit from either  $A$  or  $B$  or both
  - Need consistency in OR layout & code vector layout
- Both  $A$  &  $B$  might provide `fum()` — which is seen in  $C$ ?
  - C++ produces a "syntax error" when `fum()` is used

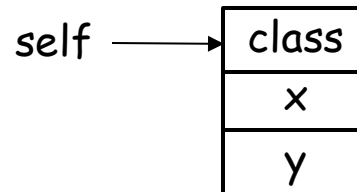
Need a better way  
to say "inherit"



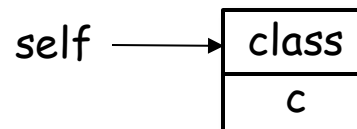
# OR Layout with Multiple Inheritance

We can try to use prefixing again

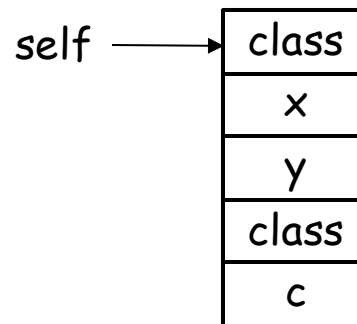
```
Class Point {
    public int x, y;
    public void draw();
    public Point add();
}
```



```
Class CThing {
    Color c;
    public void colorize();
}
```



```
Class CPoint extends
    Point and CThing {
    public void draw();
}
```



How do casts work?

- If we cast a CPoint as a Point, x & y are at the desired offsets
- If we cast a CPoint as a Cthing, c is at the wrong offset
  - Extra class field is the key
  - Cast bumps "self" to point at 2<sup>nd</sup> class field

Class that occurs out of "single inheritance prefix order" needs additional information - class field.

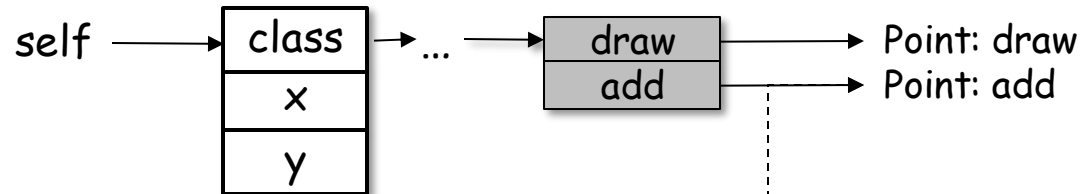


Assume 64 bit slots in the ORs

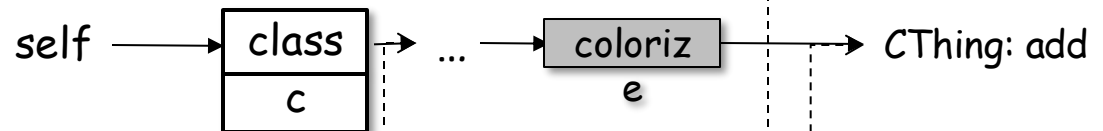
## OR Layout with Multiple Inheritance

Need to prefix both OR & code vector

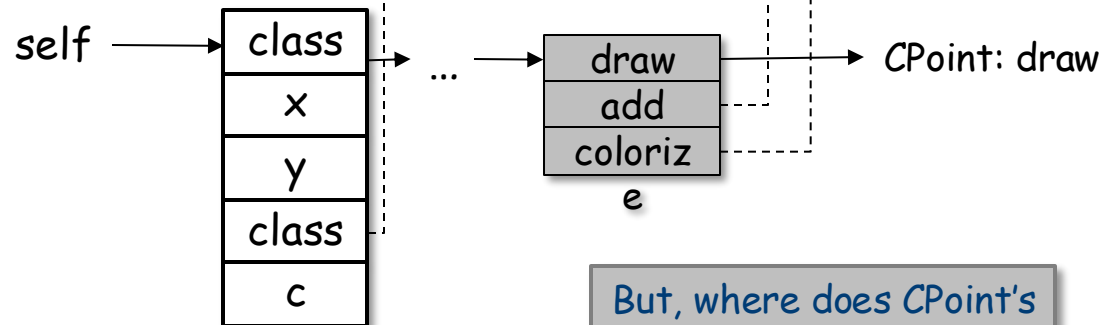
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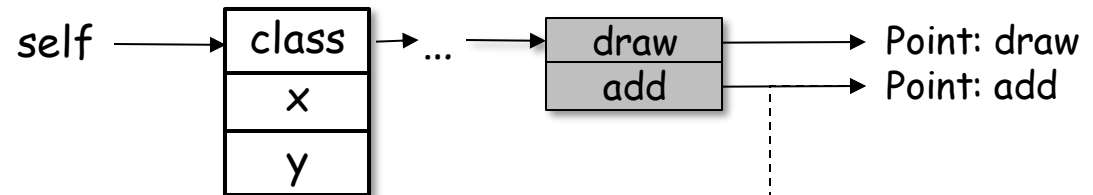
But, where does `CPoint`'s `colorize` find "`c`"?



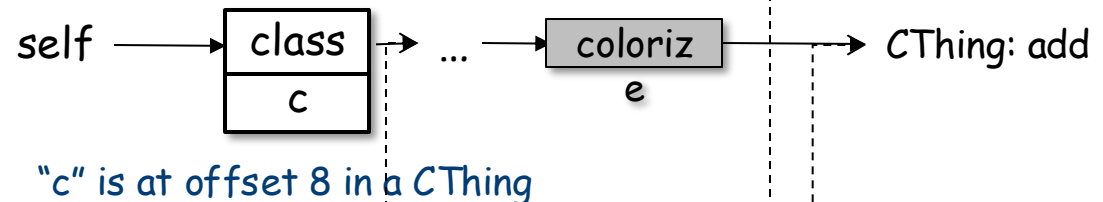
# OR Layout with Multiple Inheritance

Solution: Prefix both code storage & code vector

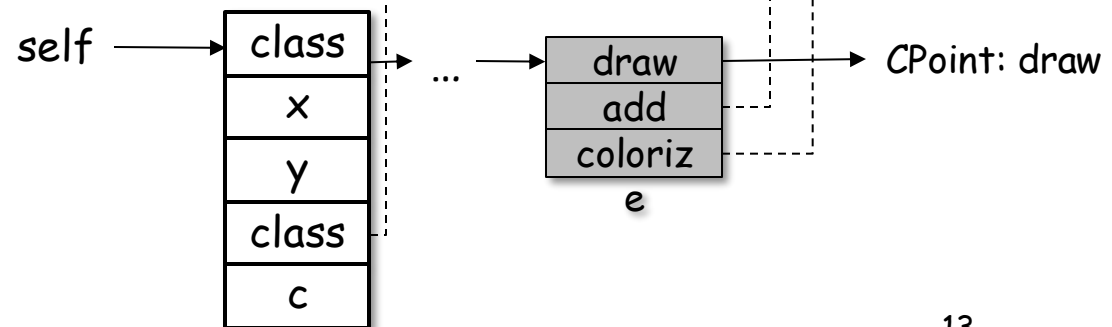
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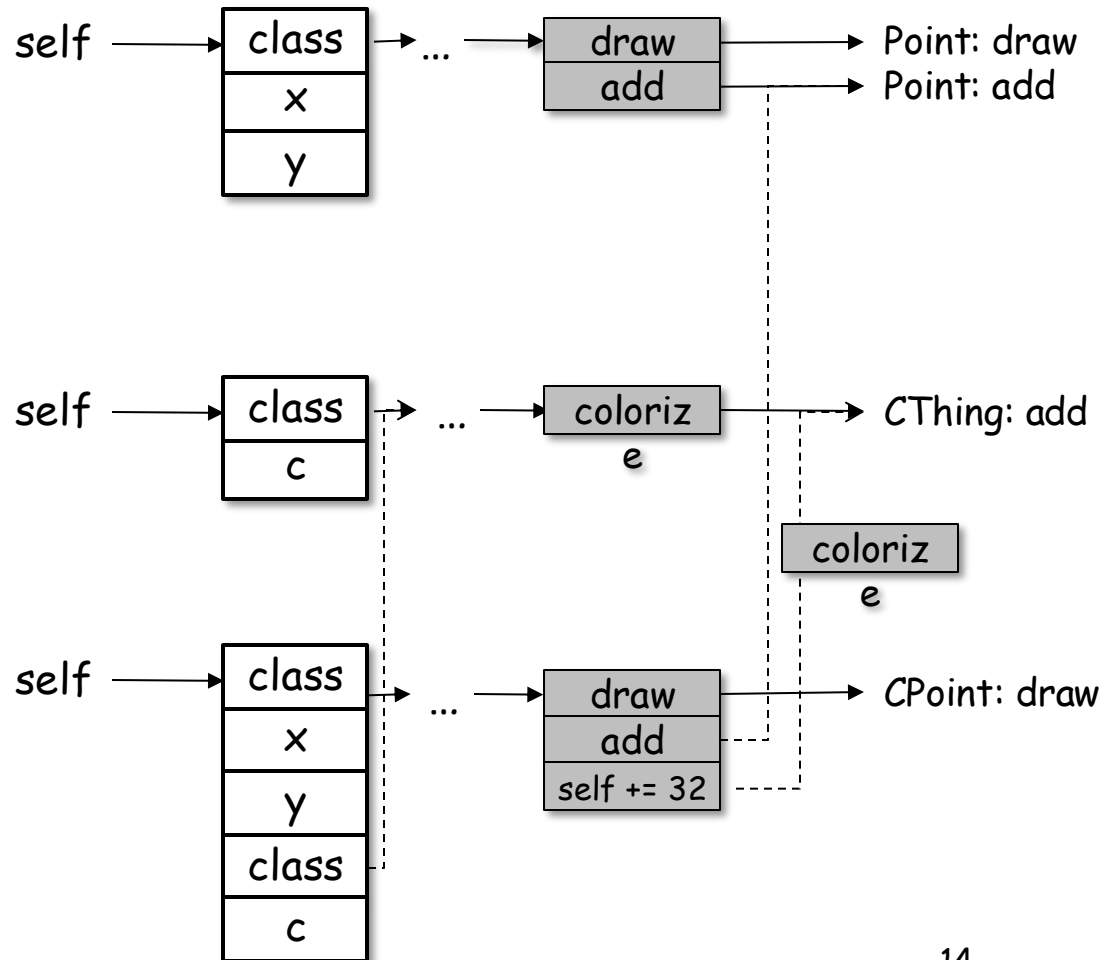
# OR Layout with Multiple Inheritance

To fix CPoint's offset for "c", use a trampoline function

```
Class Point {  
    public int x, y;  
    public void draw();  
    public Point add();  
}
```

```
Class CThing {  
    Color c;  
    public void colorize();  
}
```

```
Class CPoint extends  
    Point and CThing {  
    public void draw();  
}
```





# Open World versus Closed World

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## Closed Class Structure

- Compile time or link time layout (ORs, classes, code vectors)
- If known at compile time, can generate static code
  - No code vector, no indirection, efficient implementation
- C++ virtual functions use runtime resolution, as if open world

## Open Class Structure

- Cannot lay out ORs until hierarchy is known
- With infrequent change, can perform layout at each change
  - Single-level code vectors, fixed offsets, indirect calls
- With frequent change, may need runtime method resolution
  - Search for method in class hierarchy (w/tag) & cache result
  - Much more expensive



# Method Calls

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Given the runtime structure, how does a call work?

- Compiled code does not contain the callee's address
- Reference is relative to receiver
  - Through class to code vector

In the general case, may need dynamic dispatch

- Map code member to a search key
- Perform runtime search through hierarchy
  - This process is expensive
- Use a "method cache" to speed the search
  - Cache holds *<search key, class, method pointer>*

Smalltalk-80





# Method Calls

## Improvements are possible in special cases

- If class has no subclasses, can generate direct calls
  - Class structure must be static or class must be FINAL
- If class structure is static
  - Can generate complete method table for each class
    - Use prefixed object records and complete code vectors
  - Indirection through the class point
  - Keeps overhead low
- If class structure changes infrequently
  - Build complete method tables at initialization & when class structure changes
- If running program can create new classes
  - Well, not all things can be efficient
  - See Deutsch & Schiffman, POPL 1984

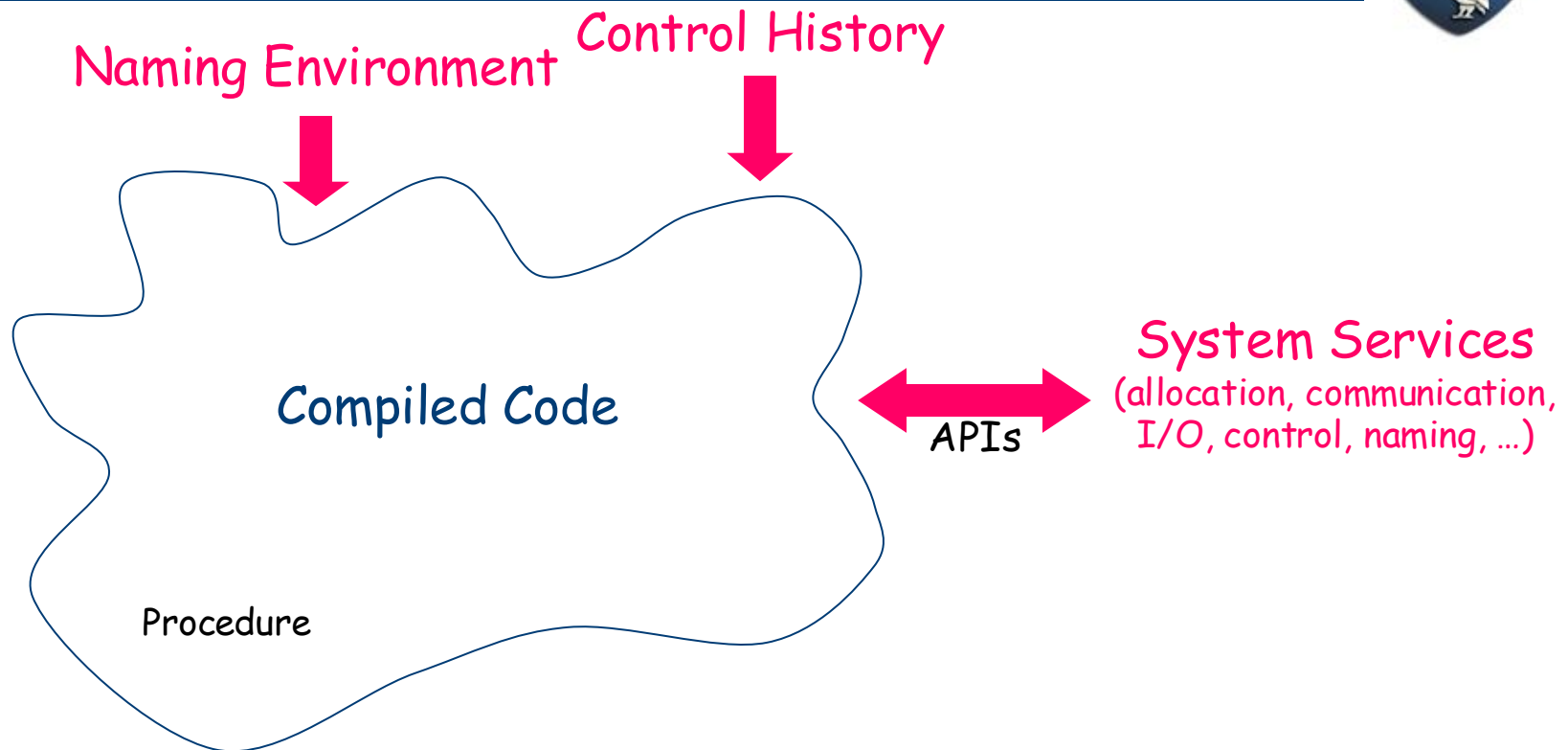
Language design

Behavior

Language design



# The Procedure & Its Three Abstractions

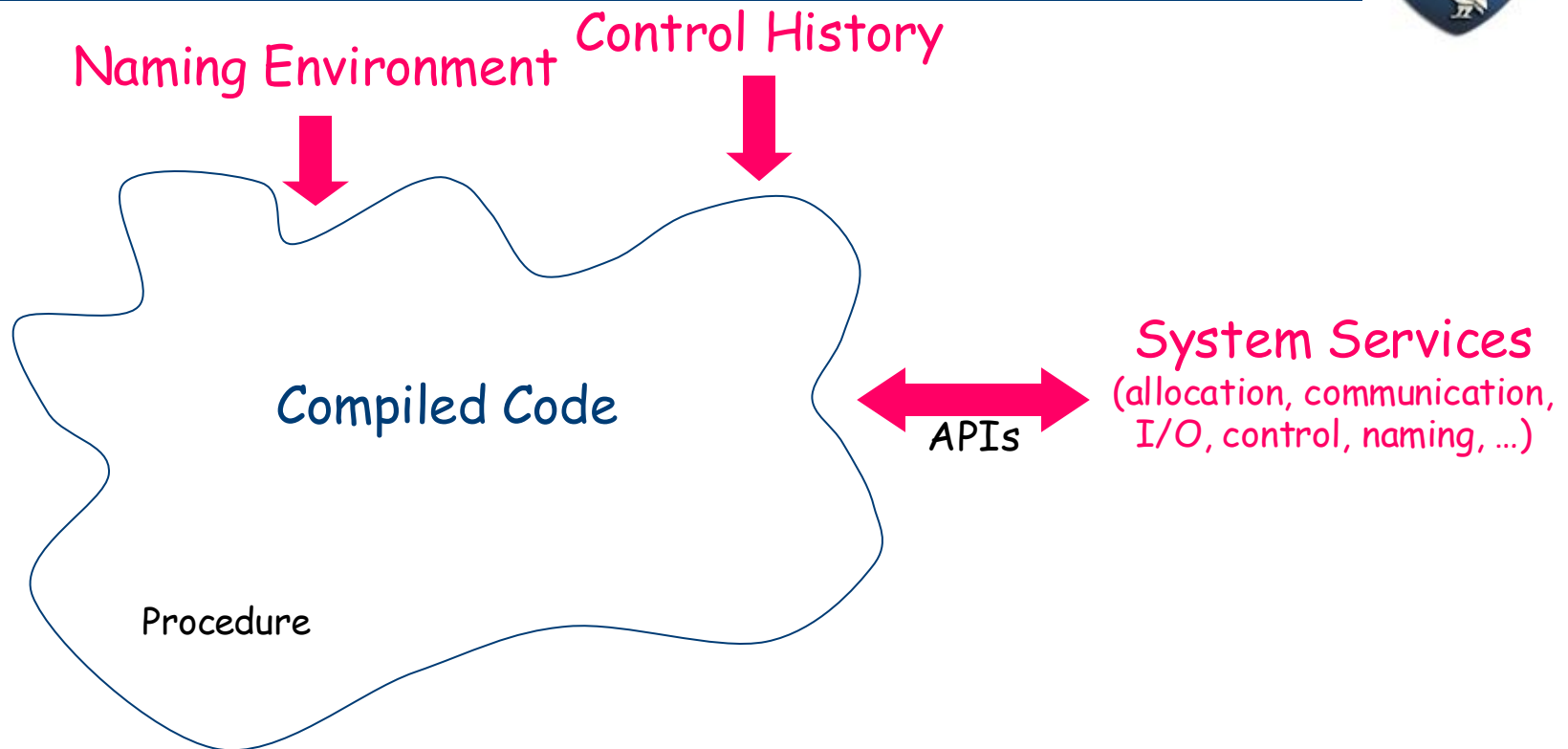


We looked at the three abstractions for ALLs and OOLs

- OOLs have more complex naming environments than ALLs
- Need both compile time and run time support for naming



# The Procedure & Its Three Abstractions



Procedure invocation reflects these differences

- OOLs use indirect calls through a code vector or jump table
- Extra loads at each call to engineer name resolution