

[AI2613 Lecture 6]: Concentration Inequalities, Martingale

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1 Chernoff Bounds

Recall the Markov inequality and Chebyshev's inequality we introduced before. They are used to prove that a random variable is concentrated around its expectation.

If we apply Markov inequality to

$$\Pr[f(X) \geq f(t)]$$

with $f(x) = e^{\alpha x}$ where $\alpha > 0$, then the bound amounts to bound $E[e^{\alpha X}]$ which is the *moment generating function* of X .

When the random variable X can be written as the sum of independent Bernoulli variables, its moment generating function is easy to estimate and we obtain sharp concentration bounds.

Theorem 1 (Chernoff Bound) . Let X_1, \dots, X_n be independent random variables such that $X_i \sim \text{Ber}(p_i)$ for each $i = 1, 2, \dots, n$. Let $X = \sum_{i=1}^n X_i$ and denote $\mu \triangleq E[X] = \sum_{i=1}^n p_i$, we have

$$\Pr[X \geq (1 + \delta)\mu] \leq \left(\frac{e^\delta}{(1 + \delta)^{1+\delta}} \right)^\mu$$

If $0 < \delta < 1$, then we have

$$\Pr[X \leq (1 - \delta)\mu] \leq \left(\frac{e^{-\delta}}{(1 - \delta)^{1-\delta}} \right)^\mu$$

Proof. We only prove the upper tail bound and the proof of lower tail bound is similar. For every $\alpha > 0$, we have

$$\Pr[X \geq (1 + \delta)\mu] = \Pr[e^{\alpha X} \geq e^{\alpha(1+\delta)\mu}] \leq \frac{E[e^{\alpha X}]}{e^{\alpha(1+\delta)\mu}}.$$

Therefore, we need to estimate the moment generating function $E[e^{\alpha X}]$. Since $X = \sum_{i=1}^n X_i$ is the sum of independent Bernoulli variables, we have

$$E[e^{\alpha X}] = E\left[e^{\alpha \sum_{i=1}^n X_i}\right] = E\left[\prod_{i=1}^n e^{\alpha X_i}\right] = \prod_{i=1}^n E[e^{\alpha X_i}].$$

Since $X_i \sim \text{Ber}(p_i)$, we can compute $E[e^{\alpha X_i}]$ directly: ①

$$E[e^{\alpha X_i}] = p_i e^\alpha + (1 - p_i) = 1 + (e^\alpha - 1)p_i \leq e^{((e^\alpha - 1)p_i)}.$$

Therefore, ②

$$E[e^{\alpha X}] \leq \prod_{i=1}^n e^{((e^\alpha - 1)p_i)} = e^{((e^\alpha - 1) \sum_{i=1}^n p_i)} = e^{((e^\alpha - 1)\mu)}.$$

本节讲述 Concentration Ineq 以及详细总结 Random Walk 中一些定量计算问题 (之前讨论的是收敛性、回归性等问题)

首先介绍 Chernoff bound (Zineq), Chernoff bound 用于 n 个独立 Bernoulli r.v. 之和的 n.v. 的聚集性。

证明方法与 Chebyshev 类似。

都运用 $\Pr[X > a] \Rightarrow \Pr[f(X) > f(a)]$,

其中 cheb 适用 $f(x) = (x - Ex)$.

而 Chernoff 适用类似 $f(x) = e^{\alpha x}$ ，这种方法关键在于 $E[e^{\alpha X}]$ 可解，其中 ① 运用了独立性

② 运用了 Bernoulli r.v. 条件

另外，注意 $f(x) = e^{\alpha x}$ 又称矩生成函数，用其泰勒展开为各阶矩的线性组合，权值由参数 α 控制。

因此，最后的 bound 含 α 项，可取 α 范围内的极值，取最 tight 的 bound。

方法在下一页有展示。

Therefore,

$$\Pr [X \leq (1 + \delta)\mu] \leq \frac{\mathbb{E}[e^{\alpha X}]}{e^{\alpha(1+\delta)\mu}} \leq \left(\frac{e^{(e^\alpha - 1)}}{e^{(\alpha(1+\delta))}} \right)^\mu$$

Note that above holds for any $\alpha > 0$. Therefore, we can choose α so as to minimize $\frac{e^{(e^\alpha - 1)}}{e^{(\alpha(1+\delta))}}$. To this end, we let $\left(\frac{e^{(e^\alpha - 1)}}{e^{(\alpha(1+\delta))}} \right)' = 0$. This gives $\alpha = \log(1 + \delta)$. Therefore

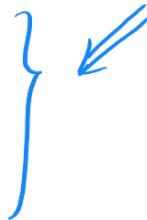
$$\Pr [X \leq (1 + \delta)\mu] \leq \left(\frac{e^{(e^\alpha - 1)}}{e^{(\alpha(1+\delta))}} \right)^\mu = \left(\frac{e^\delta}{(1 + \delta)^{(1+\delta)}} \right)^\mu.$$

□

The following form of Chernoff bound is more convenient to use (but weaker):

Corollary 2 For any $0 < \delta < 1$,

$$\begin{aligned} \Pr [X \geq (1 + \delta)\mu] &\leq \exp\left\{-\frac{\delta^2}{3}\mu\right\} \\ \Pr [X \leq (1 - \delta)\mu] &\leq \exp\left\{-\frac{\delta^2}{2}\mu\right\} \end{aligned}$$



Chernoff bound 猛然被弱化成
如左形式，便于使用。

Proof. We only prove the upper tail. It suffices to verify that for $0 < \delta < 1$, we have

$$\frac{e^\delta}{(1 + \delta)^{(1+\delta)}} \leq \exp\left\{-\frac{\delta^2}{3}\right\}$$

Taking logarithm of both sides, this is equivalent to

$$\delta - (1 + \delta) \ln(1 + \delta) \leq -\frac{\delta^2}{3}$$

Let $f(\delta) = \delta - (1 + \delta) \ln(1 + \delta) + \frac{\delta^2}{3}$ and note that

$$f'(\delta) = -\ln(1 + \delta) + \frac{2}{3}\delta, \quad f''(\delta) = -\frac{1}{1 + \delta} + \frac{2}{3}.$$

Then for $0 < \delta < 1/2$, $f''(\delta) < 0$, and for $1/2 < \delta < 1$, $f''(\delta) > 0$. Therefore, $f'(\delta)$ first decreases and then increases in $[0, 1]$. Also note that $f'(0) = 0$, $f'(1) < 0$ and $f'(\delta) \leq 0$ when $0 \leq \delta \leq 1$. Therefore $f(\delta) \leq f(0) = 0$. □

Example 1 (Tossing p -coins) Consider a p -coin that we get a head with probability p when tossing it. If we toss a p -coin n times, the average number of heads is pn . We want to determine the value δ such that with high probability (say 99%), the total number of heads is in the interval of $[(1 - \delta)pn, (1 + \delta)pn]$. We use Chernoff bound to determine δ .

Let X denote the total number of heads, and $X_i \sim \text{Ber}(p)$ be the indicator of whether the i -th toss gives a head. Then by Chernoff bound, we have

$$\Pr [|X - pn| \geq \delta \cdot pn] \leq 2 \exp\left\{-\frac{\delta^2}{3} \cdot pn\right\} \leq 0.01$$

So if p is a constant, it suffices to choose

$$\delta = \Omega\left(\frac{1}{\sqrt{n}}\right).$$

seems not
taught in
class .

2 Hoeffding's Inequality

One of annoying restrictions of Chernoff bound is that each X_i needs to be a Bernoulli random variable. We first relax this requirement by introducing Hoeffding's inequality which allows X_i to follow any distribution, provided its value is almost surely bounded.

Theorem 3 (Hoeffding's Inequality) Let X_1, \dots, X_n be independent random variables where each $X_i \in [a_i, b_i]$ for certain $a_i \leq b_i$ with probability 1. Let $X = \sum_{i=1}^n X_i$ and $\mu \triangleq \mathbb{E}[X] = \sum_{i=1}^n \mathbb{E}[X_i]$, then

①

$$\Pr [|X - \mu| \geq t] \leq 2 \exp \left(-\frac{2t^2}{\sum_{i=1}^n (b_i - a_i)^2} \right) \quad ②$$

for all $t \geq 0$.

It is instructive to compare Hoeffding and Chernoff when X_i 's are independent Bernoulli variables. Formally, let X_1, \dots, X_n be i.i.d. random variables where $X_i \sim \text{Ber}(p)$ for all $i = 1, \dots, n$. Set $X = \sum_{i=1}^n X_i$ and denote $\mathbb{E}[X] = np$ by μ . By Hoeffding's inequality, we have

$$\Pr [|X - \mu| \geq t] \leq 2 \exp \left(-\frac{2t^2}{n} \right).$$

By Chernoff Bound, we have

$$\Pr [|X - \mu| \geq t] \leq 2 \exp \left(-\frac{t^2}{3pn} \right).$$

Comparing the exponent, it is easy to see that for $p > 1/6$, Hoeffding's inequality is tighter up to a certain constant factor. However, for smaller p , Chernoff bound is significantly better than Hoeffding's inequality.

We consider the balls-in-a-bag problem. There are g green balls and r red balls in a bag. These balls are the all same except for the color. We want to estimate the ratio $\frac{r}{r+g}$ by drawing balls. There are two scenarios.

- Draw balls with replacement. Let $X_i = \mathbf{1}[\text{the } i\text{-th ball is red}]$. Let $X = \sum_{i=1}^n X_i$. Then clearly each $X_i \sim \text{Ber}\left(\frac{r}{r+g}\right)$ and $\mathbb{E}[X] = n \cdot \frac{r}{r+g}$.

Since all X_i 's are independent, we can directly apply Hoeffding's inequality and obtain

$$\Pr [|X - \mathbb{E}[X]| \geq t] \leq 2 \exp \left(-\frac{2t^2}{n} \right).$$

- Draw balls without replacement. Again we let $Y_i = \mathbf{1}[\text{the } i\text{-th ball is red}]$, then unlike the case of drawing with replacement, variables in $\{Y_i\}$ are dependent. Let $Y = \sum_{i=1}^n Y_i$. We first calculate $\mathbb{E}[Y]$.

For every $i \geq 1$, $\mathbb{E}[Y_i]$ is the probability that the i -th draw is a red ball.

Note that drawing without replacement is equivalent to first drawing a

我们可以发现 Chernoff bound 使用了蓝框(见前)的两个条件. 其中独立性很重要, 而 Bernoulli r.v. 条件可放宽.

由此诞生了 Hoeffding's Ineq. 其将条件 Bernoulli r.v. 替换为 ①. 结论变为了 ②. 证明见左.

我们额外来讨论一下 1-d RW 的一些定量结果 (定性结果见 lec4-lec5)
(w.p. $\frac{1}{2}$)

设 Z_i 为 i 时刻的位置, $\in \{-\cdots, -2, -1, 0, 1, 2, \dots\}$, $Z_0 = 0$.

一个 trick 是取 $X_t = Z_t - Z_{t-1} = \begin{cases} 1 & \text{w.p. } \frac{1}{2} \\ -1 & \text{w.p. } \frac{1}{2} \end{cases}$

1. $E[Z_i] = 0$. 因为 $Z_i = Z_0 + \sum_{t=1}^i X_t$,

取 E , 由 E 线性性,

$$E[Z_i] = 0 + \sum E[X_t],$$

$$\text{即 } E[X_t] = \frac{1}{2} \times 1 + \frac{1}{2} \times (-1) = 0.$$

$$\text{故 } E[Z_i] = 0.$$

2. $\Pr[\exists t, \text{ s.t. } Z_t = n] \geq 1 - \frac{1}{k}$ 即 kn^2 内达到 n 的概率不低于 $1 - \frac{1}{k}$

证明见 [lec5], 运用 T_n 的分类讨论, 列方程求 $E[T_n]$, 并 bound 本法论的方法.

3. $E[|Z_i|]$ (?) 和 $\Pr[|Z_i| - E[|Z_i|] \geq t]$

$$\text{证明: } |Z_i| = \underbrace{|Z_i - Z_{i-1}|}_{Y_{i-1}} + |Z_{i-1} - Z_{i-2}| + \dots + |Z_1 - Z_0|$$

$$= \sum_{i=0}^{n-1} Y_i$$

$$\text{作左右期望: } E[|Z_i|] = \sum_{i=0}^{n-1} E[Y_i]$$

* 注意 $E[Y_i]$ 与 $E[X_i]$ 不同, 其不会为 0. 特别地, 在 $Z_i = 0$ 时 $E[Y_i] = 1$
其他时刻 $E[Y_i] > 0$

$$\text{故 } E[|Z_i|] = \sum \underbrace{E[Y_i]}_{\downarrow} \Big|_{Z_i=0} + \sum \underbrace{E[Y_i]}_0 \Big|_{Z_i=1}$$

$$= \sum_{i=0}^{n-1} \Pr[Z_i = 0], \text{ 这个在 Lec 2, 3 中在 convergence 时讨论过.}$$

$$= \sum_{i=0}^{n-1} \Pr[Z_{2i} = 0]$$

$$= \sum_{i=0}^{\lfloor \frac{n}{2} \rfloor} \binom{2i}{i} / 2^{2i}$$

↓

在 $2i$ 个点中找 i 个为 '+1' move 的概率，用排列公式算。

$$\approx \sum_{i=0}^{\lfloor \frac{n}{2} \rfloor} \sqrt{\frac{2}{\pi i}} \approx C \cdot \sqrt{n} \quad C \text{ 为 constant.}$$

最后用本节的 concentration Ineq 算概率。

由 Hoeffing Ineq, $\Pr[Y_i \in \{-1, 1\}]$.

↑

$$\Pr[|X - \mu| \geq t] \leq 2 \exp\left(-\frac{2t^2}{\sum_{i=1}^n (b_i - a_i)^2}\right)$$

$$\therefore a_i = -1, b_i = 1$$

$$\text{又: } |z_i| = \sum Y_i, \mu = E[z_i] = \sum E[Y_i] = C \cdot \sqrt{n}$$

$$\text{取 } t \sim O(\sqrt{n}) = k\sqrt{n}$$

$$\therefore \text{代入 Ineq, } \Pr[|z_i| - C\sqrt{n} \geq k\sqrt{n}] \leq (\approx) \exp(-\frac{c^2}{2})$$

$$\therefore E[|z_i|] = C \cdot \sqrt{n}, \Pr \leq (\approx) \exp(-\frac{c^2}{2}) \text{ 成立.}$$

uniform permutation of $r + g$ balls and drawing each ball one by one in that order. Therefore, the probability of $Y_i = 1$ is $\frac{r \cdot (r+g-1)!}{(r+g)!} = \frac{r}{r+g}$. So we have $E[Y] = n \cdot \frac{r}{r+g}$.

However, since $\{Y_i\}$ are dependent, we cannot apply Hoeffding's inequality directly. This motivates us to generalize it by removing the requirement of independence.

3 Martingale

We develop the theory of martingale, which is a core concept in probability theory. We use martingale to get rid of the independence requirement in the concentration inequalities mentioned above.

Consider a fair gambling game in which the expected gain in each round is zero. As a result, regardless of how much one bets in each round, the money in expectation remains the same. The balances after each round form a *martingale*.

Def 1 **Definition 4 (Martingale)** Let $\{X_n\}_{n \geq 0}$ and $\{Z_n\}_{n \geq 0}$ be two sequences of random variables. Let $Z_n = \sum_{t=0}^n X_t$.¹ We say $\{Z_n\}_{n \geq 0}$ is a martingale w.r.t. $\{X_n\}_{n \geq 0}$ if

$$E[Z_{n+1} | X_0, X_1, \dots, X_n] = Z_n.$$

Sometimes we say a single sequence $\{X_n\}_{n \geq 0}$ is a martingale if it is a martingale w.r.t. itself. Formally, if for every $n \geq 0$, it holds that

$$E[X_{n+1} | X_0, \dots, X_n] = X_n.$$

For convenience, from now on we use $\bar{X}_{i,j} = (X_i, X_{i+1}, \dots, X_j)$ to simplify the notations. The conditional expectation $E[Z_{n+1} | \bar{X}_{0,n}]$ is equivalent to $E[Z_{n+1} | \sigma(\bar{X}_{0,n})]$ where $\sigma(\bar{X}_{0,n})$ is the σ -algebra generated by X_0, \dots, X_n . This motivates us to define martingale in a more general way.

Def 2

Definition 5 (Martingale (defined by filtration)) Let $\{\mathcal{F}_n\}_{n \geq 0}$ be a sequence of σ -algebras. We call such σ -algebra sequence a filtration if it satisfies

$$\mathcal{F}_0 \subseteq \mathcal{F}_1 \subseteq \dots \subseteq \mathcal{F}_n \subseteq \mathcal{F}_{n+1} \subseteq \dots$$

Given a filtration $\{\mathcal{F}_n\}_{n \geq 0}$, let $\{Z_n\}_{n \geq 0}$ be a stochastic process that Z_n is \mathcal{F}_n -measurable for every $n \geq 0$. Then we say $\{Z_n\}_{n \geq 0}$ is a martingale w.r.t. $\{\mathcal{F}_n\}_{n \geq 0}$ if for every $n \geq 0$

$$E[Z_{n+1} | \mathcal{F}_n] = Z_n.$$

Example 2 (1-D Random Walk) Consider a random walk on \mathbb{Z} starting from 0. The probability to the left and the probability to the right are both $\frac{1}{2}$ at each step. Denote the action at the n -th step by a uniform random variable

→ 在上文的 Hoeffding Ineq 中。
只有 X_i 独立才有 $X_t \in [a_i, b_i] \Rightarrow \exists X_i$ 范围的工具。
但如果引入鞅论，如 $\{X_t\}$ 是鞅，则无须独立条件也能 bound 住 $\exists X_i$

¹ Consider the problem of fair gambling where X_n is the gain of n -th round and $Z_n = \sum_{t=1}^n X_t$. $\{Z_n\}_{n \geq 0}$ is not necessarily a Markov chain. The value X_n may depend on information before round $n - 1$.

左展示了两种鞅定义，def 1 利用随机过程 $\{X_t\}$ 的变量定义，def 2 将 $X_0 \rightarrow \bar{X}_{0,1} \rightarrow \bar{X}_{0,2}$ 的过程抽象为 $\bar{X}_0 \rightarrow \bar{X}_1$ 。

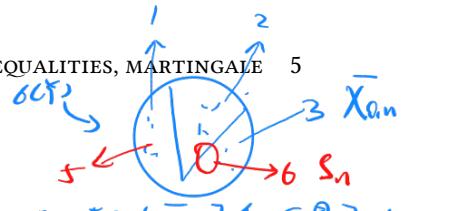
对鞅的一种直观理解在于：

将 Def 1 左一取期望，则

$E[Z_{n+1}] = E[Z_n]$ ，故鞅可以理解为一种已知 $X_0 \cup X_n$ 信息时作决策，但决策结果的期望收益是 0 ($E[Z_t] = E[Z_{t+1}]$) 的一种随机过程。

If $E[Z_{n+1} | \mathcal{F}_n] \leq Z_n$ in Definition 5, we call $\{Z_n\}_{n \geq 0}$ a supermartingale w.r.t. $\{\mathcal{F}_n\}_{n \geq 0}$. Similarly, if $E[Z_{n+1} | \mathcal{F}_n] \geq Z_n$, we call it a submartingale.

接下来我们举若干例子，证明其为 Martingale，以此加深理解。



$X_n \in \{-1, +1\}$. Let $S_n = \sum_{k=0}^n X_k$. Then we can verify $\{S_n\}_{n \geq 0}$ is a martingale

w.r.t. $\{X_n\}_{n \geq 0}$ (or w.r.t. $\{S_n\}_{n \geq 0}$) by noticing that

$$\mathbb{E}[S_{n+1} | \bar{X}_{0,n}] = \mathbb{E}[S_n + X_{n+1} | \bar{X}_{0,n}] = S_n + \mathbb{E}[X_{n+1} | \bar{X}_{0,n}] = S_n.$$

因为 S_n 是 $\bar{X}_{0,n}$ 可测，故在 $\bar{X}_{0,n}$ 条件下 $\mathbb{E}[S_n | \bar{X}_{0,n}]$ 和 S_n 是同一个 rv.

More generally, if $\mathbb{E}[X_{n+1} | \bar{X}_{0,n}] = \mu$, we define $Y_k = X_k - \mu$ and $S'_n \triangleq \sum_{k=0}^n Y_k = S_n - (n+1)\mu$. Then S'_n is a martingale w.r.t. $\{Y_n\}_{n \geq 0}$.

Example 3 Consider a sequence of random variables $\{X_n\}_{n \geq 0}$ where $\mathbb{E}[X_n | \bar{X}_{0,n-1}] = 1$ for all $n \geq 1$. Let $P_n = \prod_{k=0}^n X_k$. Then we can verify $\{P_n\}_{n \geq 0}$ is a martingale w.r.t. $\{X_n\}_{n \geq 0}$ by verifying that

$$\mathbb{E}[P_{n+1} | \bar{X}_{0,n}] = \mathbb{E}[P_n \cdot X_{n+1} | \bar{X}_{0,n}] = P_n \cdot \mathbb{E}[X_{n+1} | \bar{X}_{0,n}] = P_n.$$

Example 4 (Galton-Watson Process) Recall the Galton-Watson process we discussed in the last lecture. Suppose that all the individuals reproduce independently of each other and have the same offspring distribution. Let G_t be the number of individuals of the t -th generation. Each individual of generation t gives birth to a random number of children of generation $t+1$. Denote by $X_{t,k}$ the number of children of the k -th individual in the t -th generation. Assume that $X_{t,k}$ are i.i.d. and let $\mu \triangleq \mathbb{E}[X_{t,k}]$. Then we have $G_{t+1} = \sum_{k=1}^{G_t} X_{t,k}$. Thus,

$$\mathbb{E}[G_{t+1} | G_t] = \mathbb{E}\left[\sum_{k=1}^{G_t} X_{t,k} \middle| G_t\right] = G_t \cdot \mathbb{E}[X_{t,1}] = \mu G_t.$$

Define $M_t = \mu^{-t} G_t$. Then

$$\mathbb{E}[M_{t+1} | G_t] = \mu^{-t-1} \mathbb{E}[G_{t+1} | G_t] = \mu^{-t} G_t = M_t.$$

That is, $\{M_t\}_{t \geq 0}$ is a martingale w.r.t. $\{G_t\}_{t \geq 0}$.

Example 5 (Pólya's urn) Suppose there are some white balls and black balls in an urn. All of these balls are identical except the colors. Consider the following stochastic process: each round we pick a ball uniformly at random and observe its color; then we return the ball, and add an additional ball of the same color into the urn. We repeat the process, and our goal is to study the sequence of colors of the selected balls.

W.l.o.g. assume that we start from only one white ball and one black ball in the urn, and the index of rounds starts from 2. Then after round n , there are exactly n balls in the urn. Let X_n be the number of black balls after round n , and $Z_n = \frac{X_n}{n}$ is the ratio of black balls after round n . Clearly $Z_2 = \frac{1}{2}$. Then we have

$$\begin{aligned} \mathbb{E}[Z_{n+1} | \bar{X}_{2,n}] &= \frac{1}{n+1} \mathbb{E}[X_{n+1} | \bar{X}_{2,n}] \rightarrow \text{分析随机过程, 有 } \frac{1}{n+1} \text{ 概率为 } X_{n+1} \text{ 个黑球, 有 } 1 - \frac{1}{n+1} \text{ 概率为 } X_n \text{ 个黑球} \\ &= \frac{1}{n+1} (Z_n(X_n + 1) + (1 - Z_n)X_n) = \frac{Z_n + X_n}{n+1} = Z_n. \end{aligned}$$

(因为上一轮摸黑我们就加一个黑球)
然后化简即可.

That is, $\{Z_n\}_{n \geq 2}$ is a martingale w.r.t. $\{X_n\}_{n \geq 2}$.

Example 5 shows that X_n does not have to be i.i.d..

References