



## Polymorphism in ML

Polymorphic functions and types, type inference, ...

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## Polymorphism Introduction

### Polymorphism

It permits to handle values of different data types by using a uniform interface.

- A function that can evaluate to or be applied to values of different types is known as a polymorphic function.
- A data type that can appear to be of a generalized type is designated as a polymorphic data type.

OCaML/ML natively supports polymorphism

```
let compose f g x = f (g x);;
```

```
[15:34]cazzola@surtur:~/lp/ml>ocaml
# #use "compose.ml" ;;
val compose : ('a -> 'b) -> ('c -> 'a) -> 'c -> 'b = <fun>
# compose char_of_int int_of_char ;;
- : char -> char = <fun>
# compose (not) (not) ;;
- : bool -> bool = <fun>
# compose (fun x -> x+1) int_of_char ;;
- : char -> int = <fun>
```



## Polymorphism Polymorphism Taxonomy

### Ad Hoc Polymorphism

- the function/method denotes different implementations depending on a range of types and their combination;
- it is supported in many languages by overloading.

### Parametric Polymorphism

- all the code is written without mention of any specific type and thus can be used transparently with any number of new types;
- it is widely supported in statically typed functional programming languages or in object-orientation by generics or templates.

### Subtype Polymorphism

- the code employs the idea of subtypes to restrict the range of types that can be used in a particular case of parametric polymorphism;
- in OO languages is realized by inheritance and sub-classing.



## Polymorphism Parametric Polymorphism in ML

OCaML supports parametric polymorphism.

- compose implements  $\text{fog}$  without any type binding;
- its (polymorphic) type is

$$(\alpha \rightarrow \beta) * (\gamma \rightarrow \alpha) * \gamma \rightarrow \beta$$

$\alpha, \beta$  and  $\gamma$  are type variables denoted by 'a', 'b' and 'c' respectively;

- the type is inferred from time to time; in compose' the possible values for  $\alpha$  and  $\beta$  are restricted to char and int

```
[17:13]cazzola@surtur:~/lp/ml>ocaml
# let compose f g x = f (g x);;
val compose : ('a -> 'b) -> ('c -> 'a) -> 'c -> 'b = <fun>
# let compose' = compose (fun c -> int_of_char c) ;;
val compose' : ('_a -> char) -> '_a -> int = <fun>
```

compose' is weak-typed ('\_a).





# Polymorphism

## Weak Typed

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Polymorphism  
Introduction  
Taxonomy

Polymorphism  
in ML  
parametric  
weak typed  
Type Inference

@ Work  
datatype

References

Nothing that is the result of the application of a function to an argument can be polymorphic

- if we don't know yet exactly what is its type, then it's a weak type.

The type 'a -> 'a means:

- for all type 'a, this is the type 'a -> 'a.

Whereas, the type '\_a -> '\_a means:

- there exist one and only one type '\_a such that this is the type '\_a -> '\_a.

Shall we say that what is potentially polymorphic turns to monomorphic in practice when the compiler deals with its polymorphic form.

```
# let a = ref [];;
val a : 'a list ref = {contents = []}
# let b = 1::!a ;;
val b : int list = [1]
# a;;
- : int list ref = {contents = []}
```

Slide 5 of 11



# Polymorphism

## Type Inference

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Introduction  
Taxonomy

Polymorphism  
in ML  
parametric  
weak typed  
Type Inference

@ Work  
datatype

References

```
let rec map f l = match l with
| h::l1 -> f h::map f l1
| - -> [];
```

Let us calculate the type of map

1.  $l = []$ ,  $[]$  is a zeroary function  $[] : \alpha \text{ list} \forall \alpha$ ;
2.  $h::l1$ ,  $::$  is a binary operator  $:: : \alpha \times \alpha \text{ list} \rightarrow \alpha \text{ list}$  so the type of  $h$  is  $\alpha$  and the type of  $l1$  is  $\alpha \text{ list}$ ;
3. the type of  $f$  is a function whose input has type  $\alpha$  nothing can be said on the return type (denoted by  $\beta$ );
4. so the second occurrence of  $::$  should be  $\beta \times \beta \text{ list} \rightarrow \beta \text{ list}$  due to the type of  $f$ ; that means
5.  $\text{map } f \text{ } l1$  should have type  $\beta \text{ list}$

and this is possible only if

6. the type of map is  $(\alpha \rightarrow \beta) \times \alpha \text{ list} \rightarrow \beta \text{ list}$

```
# use "map.ml" ;;
val map : ('a -> 'b) -> 'a list -> 'b list = <fun>
```

Slide 6 of 11



# Polymorphism @ Work

## Polymorphic ADT: Stack

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Introduction  
Taxonomy

Polymorphism  
in ML  
parametric  
weak typed  
Type Inference

@ Work  
datatype

References

```
module Stack = struct
type 'a stack = { mutable c : 'a list }
exception EmptyStackException
let empty () = { c = [] }
let push s x = s.c <- x :: s.c
let pop s =
match s.c with
| hd::tl -> s.c <- tl
| [] -> raise EmptyStackException
end;;
```

```
[22:40]cazzola@surtur:~/lp/ml>ocaml
# use "adtstack.ml";;
# let s = Stack.empty();;
val s : 'a Stack.stack = {c = []}
# Stack.push s 7;;
- : unit = ()
# Stack.push s 25;;
- : unit = ()
# s ;;
- : int Stack.stack = {c = [25; 7]}
# let s1 = Stack.empty();;
val s1 : 'a Stack.stack = {c = []}
# Stack.push s1 "Hello";;
- : unit = ()
# Stack.push s1 "World";;
- : unit = ()
# s1;;
- : string Stack.stack = {c = ["World"; "Hello"]}
```

Slide 7 of 11



# Polymorphism @ Work

## Iterating on Collections

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Polymorphism  
Introduction  
Taxonomy

Polymorphism  
in ML  
parametric  
weak typed  
Type Inference

@ Work  
datatype

References

Count the occurrences

```
let rec count ?(tot=0) x = function
| [] -> tot | h::l1 -> if (h=x) then count ~tot:(tot+1) x l1 else count ~tot:tot x l1
```

```
val count : ?tot:int -> 'a -> 'a list -> int = <fun>
# let il = [1;2;3;4;2;2;1;3;4;5;7;3;2;1] ;;
# let cl=['a';'b';'c';'a'];;
# count 'a' cl;;
- : int = 2
# count 3 il;;
- : int = 3
```

Reducing a List

```
let rec remove x = function
| [] -> [] | h::l1 -> if (h = x) then (remove x l1) else (h::(remove x l1))
```

```
val remove : 'a -> 'a list -> 'a list = <fun>
# remove 3 il;;
- : int list = [1; 2; 4; 2; 2; 1; 4; 5; 7; 2; 1]
# remove 'a' cl;;
- : char list = ['b'; 'c']
```

Iterating on strings

```
let rec iter f ?(k = 0) s =
if k < String.length s then ( f s.[k] ; iter f ~k:(k + 1) s ) ;;
```

```
val iter : (char -> 'a) -> ?k:int -> string -> unit = <fun>
```

Slide 8 of 11



## Polymorphism @ Work Sorting (Quicksort)

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Polymorphism  
Introduction  
Taxonomy

Polymorphism  
in ML  
parametric  
weak typed  
Type Inference

@ Work  
datatype

References

```
let qsort (>) l =
  let rec qsort = function
    [] -> []
  | h::tl -> (qsort (List.filter (fun x -> (x >: h)) tl) )
    @ [h] @
    (qsort (List.filter (fun x -> (h >: x)) tl) )
  in qsort l
```

```
[14:58]cazzola@surtur:~/lp/ml>ocaml
# #use "qsort.ml" ;;
val qsort : ('a -> 'a -> bool) -> 'a list -> 'a list = <fun>
# let l=[11; 4; 123; 7; -8; 0; 15; 11; -7; 77; 99; 100; 1; 2; 4; -77] ;;
val l : int list = [11; 4; 123; 7; -8; 0; 15; 11; -7; 77; 99; 100; 1; 2; 4; -77]
# let l'=['a'; 'z'; 'w'; 'b'; 'f'; 'a'; 'x'] ;;
val l' : char list = ['a'; 'z'; 'w'; 'b'; 'f'; 'a'; 'x']
# qsort (>) l ;;
- : int list = [123; 100; 99; 77; 15; 11; 7; 4; 2; 1; 0; -7; -8; -77]
# qsort (<) l ;;
- : int list = [-77; -8; -7; 0; 1; 2; 4; 7; 11; 15; 77; 99; 100; 123]
# qsort (<) l' ;;
- : char list = ['a'; 'b'; 'f'; 'w'; 'x'; 'z']
```

### Note

- (>) represents a binary operator, you can use any sort of symbol.

Slide 9 of 11



## Polymorphism @ Work Sorting (Selection Sort)

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Polymorphism  
Introduction  
Taxonomy

Polymorphism  
in ML  
parametric  
weak typed  
Type Inference

@ Work  
datatype

References

```
let lmin (<) l =
  let rec lmin m = function
    [] -> m
  | h::tl -> lmin (if (m <: h) then m else h) tl
  in lmin (List.hd l) (List.tl l)

let filter_out x l =
  let rec filter_out acc x = function
    [] -> List.rev acc
  | h::tl when h=x -> List.rev_append tl acc
  | h::tl -> filter_out (h::acc) x tl
  in filter_out [] x l

let selection (<) l =
  let rec selection acc = function
    [] -> List.rev acc
  | l' -> let m = (lmin (<) l') in selection (m::acc) (filter_out m l')
  in selection [] l
```

```
[10:56]cazzola@surtur:~/lp/ml> ocaml
# let l1 = [-7;1;25;-3;0;15;77;-7] ;;
val l1 : int list = [-7; 1; 25; -3; 0; 15; 77; -7]
# #use "selection.ml";;
val lmin : ('a -> 'a -> bool) -> 'a list -> 'a list = <fun>
val filter_out : 'a -> 'a list -> 'a list = <fun>
val selection : ('a -> 'a -> bool) -> 'a list -> 'a list = <fun>
# selection (<) l1 ;;
- : int list = [-7; -7; -3; 0; 1; 15; 25; 77]
# selection (>) l1 ;;
- : int list = [77; 25; 15; 1; 0; -3; -7; -7]
```

Slide 10 of 11



## References

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Polymorphism  
Introduction  
Taxonomy

Polymorphism  
in ML  
parametric  
weak typed  
Type Inference

@ Work  
datatype

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Slide 11 of 11