

# Data Structures

## Arrays

99% arrays problem

- \* If array is sorted use binary search

$\text{mid} = \text{start} + (\text{end} - \text{start})/2;$  → In the loop  
while ( $\text{start} < \text{end}$ )

## Maximum tip calculator

- find difference (absolute difference) →  $c = \text{abs}(a - b)$   
and sort array  $a, c$ , and continue

## Find if string is palindrome.

→ input  $\text{fgets(string, sizeof(string), stdin);}$   
If you're taking another input before enter might  
be taken so call one extra fgets)

Also here if you enter simran, length is 8

so  $\text{length} = \text{strlen(string)} - 2;$  \*

↳ index of last character

- \* To ignore all special characters and take only alphanumeric

$\text{if} (i >= 65 \& i <= 90 \text{ || } i >= 97 \& i <= 122 \text{ || } i >= 48 \& i <= 57)$   
 $i++; \text{continue}; \text{(ignore)}$

To convert to lowercase

$\text{if} (i >= 65 \& i <= 90)$   
 $s[i] = s[i] + 32;$

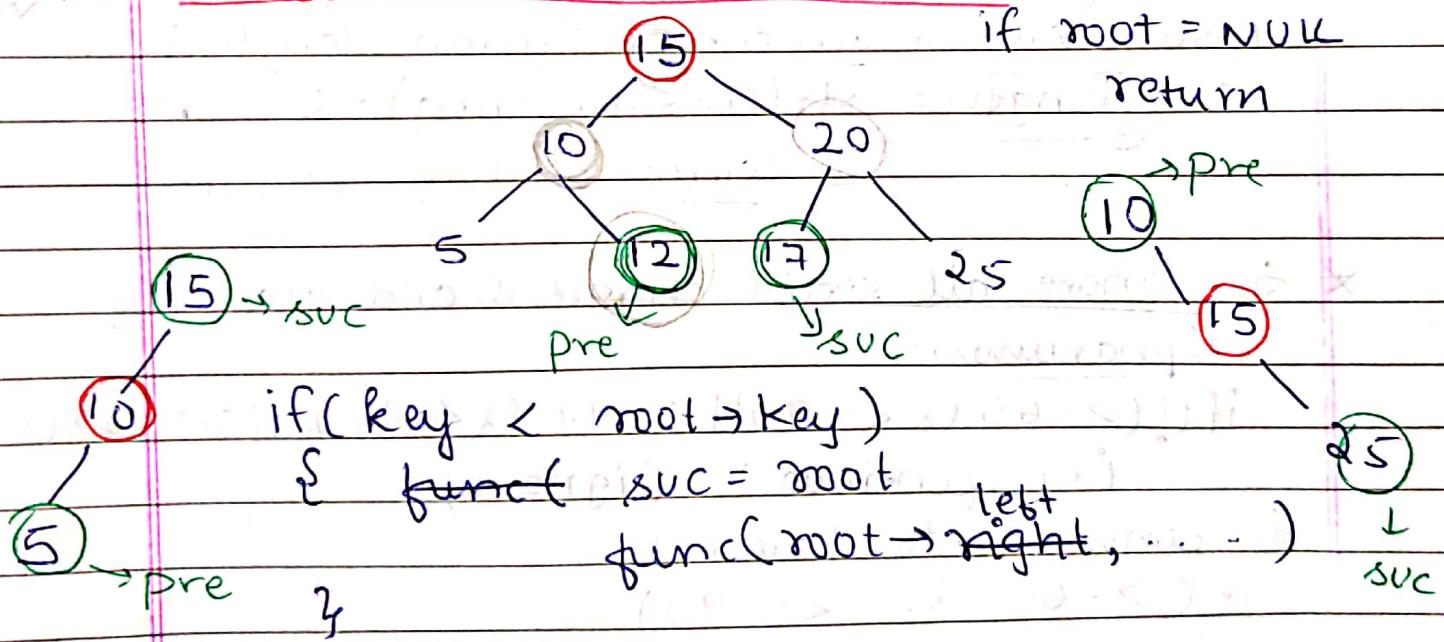
If copying one string to other do character wise  
and at last  $\text{string2}[i] = '\backslash 0'$

↳ last char →  $\text{strlen(string2)} - 1$

## Binary Search Tree

- Inserting ✓ O(n)  
if (node == NULL) create a new node (root)
- Inorder traversal of BST always produces sorted output
- So can always get inorder traversal by sorting any given traversal.
- Deleting ✓ O(n)

### Predecessor and Successor



→ If a given Binary tree is BST

do <sup>inorder</sup> preorder traversal store in array, check if array sorted XD

## lowest common Ancestor in a BST

```
while (root == null)
```

```
{
```

```
if (n1 < root->data && n2 < root->data)
```

```
{ root = root->left;
```

```
else if (n2 > root->data && n1 > root->data)
```

```
root = root->right;
```

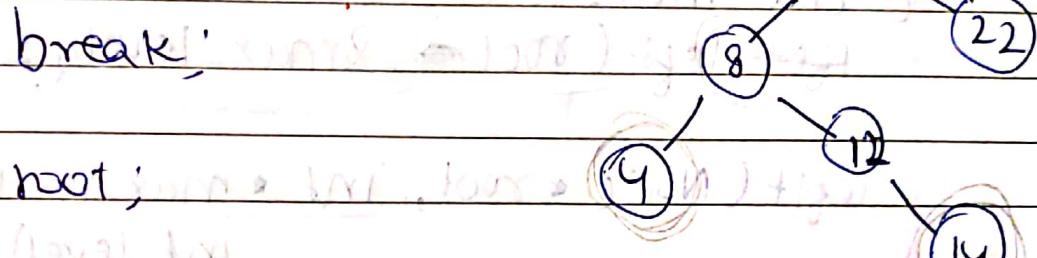
```
else
```

```
break;
```

```
}
```

```
return root;
```

```
}
```



The value always lies b/w  $n_1$  and  $n_2$  and is the first such value even if its  $n_1$  or  $n_2$ .

\* Distance b/w  $n_1$  and  $n_2$  = distance from root to  $n_1$  + distance from root to  $n_2$  -  $2 * \text{distance from root to LCA}$

Unsorted array, find max of three elements

sort array return  $\max(a[0], a[1], a[n-1])$  OR  
 $a[n-1], a[n-2], a[n-3]$

Left view of a binary tree.

left view(Node \* root)

{ int max\_level = 0;

left(left(root, &max\_level, 1))

left(left(root, int \* max\_level,  
int level))

{ if (root == NULL) return int;

if (\*max\_level < level)

root->data = print(root->data);

Print(root->data);

root->max\_level = level; }

left(left(root, &max\_level, level+1));

left(left(root->right, max\_level, level+1));

}

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## Middle of a Linked List (Slow, fast pointer)

while (fast != NULL && fast->next != NULL)

### Aptitude

Prime no.  $\Rightarrow$  up to  $p^2$  if <sup>not</sup> divisible

- \* Find smallest prime number  $n$  such that  
 $n^2 \geq p$ , check if  $p$  divisible by a prime no.  $\leq n$

### Divisibility test

- By 3 → sum divisible by 3 + digit
- By 4 → last two digits divisible by 4
- By 6 → if divisible by 2 or 3
- By 7 → remove last digit, subtract double of last digit from remaining no, if left 0, 7 ✓
- By 8 → last 3 digits divisible by 8
- By 9 → sum of digits
- By 11 → sum of digits at even position - odd pos. = multiple of 11 or 0

$$\frac{x^n - a^n}{x - a} \text{ divisible } \times \text{ all } \frac{x^n - a^n}{x + a} \text{ even}$$

$$\frac{x^n + a^n}{x + a} \checkmark \text{ for } x \in \text{odd}$$

## LCM and HCF.

- For 2 numbers a and b  $\text{LCM} \times \text{HCF} = a \times b$
- Find the greatest number which on dividing 70 and 50 leaves remainders 1 & 4 respectively  
 → Number exactly divides 69 and 46  
 $\text{HCF}(69, 46) = 23$
- Find the largest number which divides 64, 136 and 238 to leave the same remainder in each case  
 $\rightarrow \text{HCF}(\frac{136-64}{}, (238-136), (238-64))$   
 $= 6$
- Find the least number which when divided by 5, 7, 9 and 12 leaves remainder 3 in each case  
 $\rightarrow \text{LCM}(5, 7, 9, 12) = 1260$   
 $= 1260 + 3 = 1263$
- Find the largest four digit number exactly divisible by 15, 21 and 28  
 $\rightarrow \text{LCM}(15, 21, 28) = 420$   
 largest 4 digit number = 9999  
 $\rightarrow 9999 \div 420 \rightarrow \text{remainder } 339$   
 $\rightarrow 9999 - 339 = 9660 //$

→ The policeman at 3 different places on a ground blow a whistle after every 42 sec, 60 sec and 78 sec respectively. If they all blow whistle at 9: 30 : 00, when will they blow whistle again together?

$$\rightarrow \text{LCM}(42, 60, 78) = 56460 = 1 \text{ hours} 31 \text{ min}$$

at 11:01

→ Find the least no. which when divided by 6, 7, 8 leaves remainder 3, but when divided by 9

$$\rightarrow \text{LCM}(6, 7, 8) = 168$$

$$= 168m + 3$$

→ 3 rectangular fields area  $60\text{m}^2$ ,  $84\text{m}^2$ ,  $108\text{m}^2$  are to be divided into identical rectangular flower beds, each having length 6m, breadth = ?

$$\rightarrow \text{HCF}(60, 84, 108) = 12$$

$$6 \times x = 12 \quad x = 2$$

## WORK & WAGES

Comparison of work and efficiency

$$\frac{M_1 D_1 H_1 E_1}{W_1} = \frac{M_2 D_2 H_2 E_2}{W_2}$$

$M_1$  = # workers

$D_1$  = days

$H_1$  = hours

$E$  = efficiency

$W$  = units of work

Q. A  $\rightarrow$  10 days to complete work, B  $\rightarrow$  15 days to complete work together, 9 days?

$$\rightarrow \text{LCM}(10, 15) = 30$$

A  $\rightarrow$  3 units/day, B  $\rightarrow$  2 units/day  $\rightarrow$  5 units per day

$$B \rightarrow 2 \text{ units/day}$$

$$30/5 = 6 \text{ days}$$

Q. Two friends A & B working together complete work in 4 days, A does in 12 days, B = ?

$\rightarrow$  Let's say work = 12 units

$$A \& B = 3 \text{ units/day}$$

$$A \rightarrow 12/12 \rightarrow 1 \text{ unit/day}$$

$$B \rightarrow 2 \text{ units}, \frac{12}{2} = 6 \text{ days}$$

$$A \& B = 18, B \& C = 24, A \& C = 36$$

Find A, B, C? (together & individual)

$$\text{LCM}(A \& B, B \& C, A \& C) = (18, 24, 36) = 72$$

$$A \& B = 4 \text{ units/day}$$

$$B + C = 3$$

$$A + C = 2$$

$$2(A + B + C) = 9 \quad A + B + C = 4.5 \text{ /day}$$

$$72/4.5 = 16 \text{ days}$$

## Overview of data structures

### Array

$O(n)$  insertion [at the beginning]

$O(n)$  deletion [at the beginning]

[ $head \rightarrow head + 1$  for  $next \leftarrow next + 1$ ]

### Linked List ( $JUN = 1, 2, 3, \dots, N$ )

Circular linked list (any node can be made starting element)

Access time  $O(n)$

Search  $O(n)$ :  $q.mot = \text{target}$

Insertion, deletion  $O(1)$  [depending on implementation]

### Stack

Access time  $O(n)$

Search  $O(n)$

### Queue same.

Circular queue (if implemented using array) space efficient as can add elements at  $tail$ ,  $O(1)$  if  $tail + present < N$

## Q Print all elements of linked list in reverse order

```
void func1( void * head )
```

```
2 if (head == NULL) return;
```

```
func1( head -> next );
```

```
3 print( "r.d ", head -> data );
```

```
between 67702 and (spn)0
```

```
2008-09-14 11:00
```

return; you have to do it yourself  $\leftarrow *$

return from stack [printing is ended]

2008-09-14

Write code to reverse doubly linked list.

void fun( Struct node \*\* head )

{

    Struct node \* temp=NULL;

    Struct node \* curr = \*head;

    while (curr != NULL)

    {

        temp = curr -> prev;

        curr->prev = curr->next;

        curr->next = temp;

        curr = curr->prev;

    if (temp.p != NULL)

        \*head = temp->prev;

}

Sort a linked list?

Merge sort and insertion sort

$O(n \log n)$

$O(n^2)$

Sorting algorithms

Selection sort

$O(n^2)$  time

sorted    unsorted

$O(1)$  auxiliary space

\* → Never takes more than  $O(n)$  swaps, use  
when memory write is costly operation  
in place

## Bubble Sort

(compares adjacent elements, swaps last element in place.  $O(n^2)$ )

WORST CASE

array is reverse sorted.

BEST CASE

$O(n) : 1 + 2 + \dots + m = 10$

$$m = 8 = 60$$

In an almost sorted array, (swap of just two elements) can be fixed in linear time  $O(2n)$

## Merge Sort

Divide and Conquer

$O(n \log n) \rightarrow$  always

\* Sort linked list in  $O(n \log n)$

merge operation can be implemented without extra space.

Merge sort access data sequentially so can be applied on linked list, quick sort needs random access.

```
void mergesort(arr[], l, r)
{
    if(l < r)
    {
        m = (l+r)/2;
        mergesort(arr, l, m);
        mergesort(arr, m+1, r);
        merge(arr, l, m, r);
    }
}
```

$$T(n) = 2T(n/2)$$

$$T(n) = 2T(n/2) + cn$$

$$T(n) = cn + 2T(n/2)$$

Teacher's Signature.....

void merge(int arr[], int l, int m, int r)

{

    int i, j, k;

    n1 = m - l + 1;

    n2 = r - m;

    int L[n1], R[n2];

    for (i=0; i < n1; i++) L[i] = arr[l+i];

    for (j=0; j < n2; j++) R[j] = arr[m+j+n1];

    i = 0; j = 0; k = l;

    while (i < n1 && j < n2)

    {

        if (L[i] <= R[j])

            arr[k] = L[i];

        else

            arr[k] = R[j];

        k++;

    while (i < n1)

        arr[k++] = L[i++];

    while (j < n2)

        arr[k++] = R[j++];

}

Used to find inversion count

inv-c = ms

inv-c + = ms

inv-c + = m

## Prefix of strings

char arr[n][100];  
 scanf(" %s", arr[i]); → to get input  
 scanf("%[^\n]", arr[i]); → on newline

## New1.c

### Level order Tree Traversal

\* Breadth first traversal

#### Find the height of a tree

```
int height(struct node *node)
```

```
{ if(node == NULL)
```

```
    return 0;
```

```
    lheight = height(node->left);
```

```
    rheight = height(node->right);
```

```
    if(lheight > rheight)
```

```
        return (1 + lheight);
```

```
    else return (1 + rheight);
```

```
}
```

```
void print_level(struct node *root)
```

```
{ int h = height(root);
```

```
for(i=1; i<=h; i++)
```

```
    printlevelnew(root, i); }
```

```
void printlevelnew(root, int level)
```

```
{ if(root == NULL) return;
```

```
if(level == 1) print(root->data);
```

```
else if(level > 1)
```

```
    printlevelnew(root->left, level-1);
```

```
    printlevelnew(root->right, level-1);
```

```
}
```

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## Largest Sum Contiguous Subarray

```
int max( int arr[], size)
```

```
    { int max = INT_MIN;
```

```
        int max_here = 0;
```

```
        for( i=0; i < n; i++ )
```

```
            { max_here += arr[i];
```

```
                if( max < max_here )
```

```
                    max = max_here;
```

```
                (if( maxhere < 0) maxhere = 0;
```

```
                maxhere = 0;
```

```
}
```

```
        return max;
```

```
}
```

## Largest Product contiguous Subarray

```
int max( int arr[], size)
```

```
{
```

```
    int max = arr[0];
```

```
    int min = arr[0];
```

```
    int maxP = arr[0];
```

```
    for( i=1; i < n; i++ )
```

```
        { if( arr[i] < 0)
```

```
            { (i, true), swap(max, min);
```

```
                max = max(arr[i], arr[i]*max);
```

```
                min = (min(arr[i]), arr[i]*min);
```

```
            }
```

```
            maxP = max(maxP, max);
```

```
}
```

```
return maxP;
```

```
}
```

Teacher's Signature.....

## Find the Missing Number

Method 1 for sum =  $n(n+1)/2$  subtract.

array a, n      sum =  $(n+1)(n+2)/2$

Method 2

$x_1 = arr[0]$

$x_2 = -1; R = > i; L = i$

for ( $i=1; i < n; i++$ )

$i \& (x_1 = x_1 \wedge arr[i])$

for ( $i=2; i < n+R; i++$ )

$x_2 = x_2 \wedge i;$

return  $x_1 \wedge x_2$

## Reverse words in a given string

input hey what you doing  
gets()

int t; scanf(&t);

getchar();  $\rightarrow$  after scan

int main()

{ char s[100];

gets(s); }

PROGRAM GO THROUGH GFG

## Sort an array containing 0, 1, 2

1 1 1 0 0 0

0 0 0 1 1 1

0 0 0 1 1 1

0 0 0 1 1 1

0 0 0 1 1 1

## Find permutation of a string

$l=0 \ r=n-1$

void perm(char \*s, int l, int r)

{

if ( $l == r$ )

printf("%s", s);

else

{

for (i = l; i <= r; i++)

{ swap(a[i], a[r]);

perm(s, l + 1, r);

swap(a[i], a[r]);

}

## Count no of set bits.

count = 0

while (n)

{

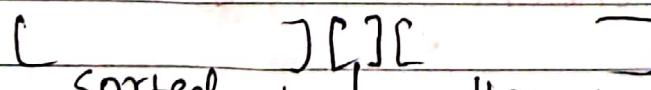
count = count + (n & 1),

n = n >> 1

}

## Insertion sort

Can be used for linked list



insert in sorted array

Check from last and find where to insert  
meanwhile keep shifting the array elements



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## Subarray with a given sum.

sum → given sum  $u = [a[0], \dots]$ , start = 0

for ( $i=1$ ;  $1 \leq n$ ;  $i++$ )  
 {

if (while ( $\text{sum} \geq u \& \& \text{start} < i-1$ ))

{ for ( $j=0$ ;  $\text{sum}_j = \text{sum} - a[\text{start}]$ ;

if ( $\text{sum}_j = u$ )  $\text{print}(start + 1, i - 1)$

if ( $i < n$ )  $\text{print}(sum - i, a[i])$

}

## Trapping Rain Water

arr = [0, 1, 0, 2, 1, 0, 0, 3, 2, 2, 1, 2, 1, 0, 0, 0]

find left\_max, right\_max works backward

0 1 1 0 2 2 3 2 3 3 3 3 3 3 3 3

3 3 3 3 3 3 3 3 3 2 2 2 2 2 2 2

sum = sum + min(leftmax[i], rightmax[i]) - arr[i]

## Pythagorean triplet in an array

3 1 4 6 5

rain 0.8

find sq. 9 1 16 36 25 for (int i=n-1; i>=2; i--)

sort 1 9 16 25 36 { l=0, r=i-1;

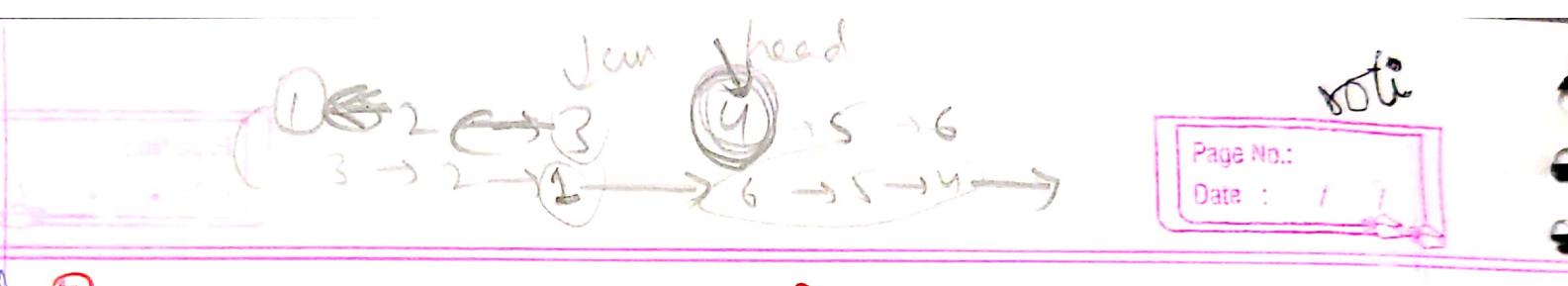
ans = ans + while (l < r)

if (arr[l] + arr[r] == arr[i])

return true.

ans = ans + (arr[l] + arr[r] > arr[i]) ? r-- : l++;

return false.



Reverse a linked list in groups of given size.

## Permutation and Combination

Permutation: arrangement.  $(x A, B) \rightarrow AB, BA$

$$Pr = n! / (n-r)! \quad G =$$

Combination:  $(A, B) \rightarrow AB$

$${}^n C_r = {}^n C_{n-r}$$

$$n! / (n-r)! r!$$

Words from DRIVE RIVE → all novels together.

D, R, I, E, V, O, R, I) down ↴

$$= 5 \times 4 \times 3 \times 2 \times 1 \times 2$$

↓ board = book ↴ book

UN = fresh ↴ UN = ! book ↴ (R is repeated)  
2 novels can  
(stob ← board => stob ← book) ↴ be interchanged

In how many ways can 20 boys and 18  
girls make a queue such that no 2  
girls are together?

$$20! \times {}^{21} P_{18} \quad \{$$

say

There are 5 floating stones on a river. A  
man wants to cross the river. He can move  
1 or 2 steps at a time. Find the number  
of ways in which he can cross the river?

Total board steps

[ 6 unit steps ]

$${}^5 C_4$$

[ 4 unit, 1 double ]

$${}^5 C_2$$

[ 3 unit, 2 double ]

[ 3 double ]

3 ( 1 UN = 1 board ) ↴ 13 + very ↴

• Answer = 13 + very ↴

Out of 7 boys and 4 girls, how many queues  
of 3 boys and 2 girls can be formed?

$$A = \frac{7!}{3!} \times \frac{4!}{2!} = 210 \times 120 = 25200$$

$$15!/(5!10!) = 12A = (A, A)$$

Merge two sorted = linked lists

if ( $\text{head}_2 \rightarrow \text{data} < \text{head}_1 \rightarrow \text{data}$ )

{ swap ( $\text{head}_1$ ,  $\text{head}_2$ ), }

Node <sup>15</sup>  $\text{head} = \text{head}_1$ ;

while ( $\text{head}_1 \neq \text{NULL}$  &  $\text{head}_2 \neq \text{NULL}$ )

if ( $\text{head}_1 \rightarrow \text{data} \leq \text{head}_2 \rightarrow \text{data}$ )

{

$\text{prev} = \text{head}_1$

$\text{head}_1 = \text{head}_1 \rightarrow \text{next}$

}

else

{

$\text{prev} \rightarrow \text{next} = \text{head}_2$ ;

$\text{prev} = \text{head}_2$ ;

$\text{temp} = \text{head}_2 \rightarrow \text{next}$ ;

$\text{head}_2 \rightarrow \text{next} = \text{head}_1$  [toot]

$\text{head}_2 = \text{temp}$ ; + in u  $\Delta$  7

}

{

[lab 1 + in u  $\Delta$  7]

[lab 6 + in u  $\Delta$  6]

if ( $\text{head}_2 = \text{NULL}$ ) { lab 2 }

{  $\text{prev} \rightarrow \text{while}(\text{head}_2 \neq \text{NULL}) \{$

~~$\text{head}_2 \rightarrow \text{prev} \rightarrow \text{next} = \text{head}_2$~~ ;

$\text{prev} = \text{prev} \rightarrow \text{next}$ ;

$\text{head}_2 = \text{head}_2 \rightarrow \text{next}$ ; }

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Learn Speed

2 pipes together fill bucket in 12 min  
 First pipe 10 min faster than second pipe  
 Second pipe alone?

time taken by first pipe  $x = 2t$

second =  $t + 10$

$$\frac{1}{t} + \frac{1}{t+10} = \frac{1}{12}$$

$$2t + 10 = \frac{1}{12}$$

$$\frac{b}{2} = \frac{2t + 10}{12} = \frac{1}{6} = 5mt$$

$$24t + 120 = t^2 + 10t$$

$$t = 20$$

work B takes  $\frac{(20+10)}{2} = 30$  mins

TSD (Aptitude)  $\Rightarrow d = b$

Km/hour  $\rightarrow \times 5/18$  to m/s

m/s  $\rightarrow \times 18/5$  to km/h

Average speed =  $\frac{n}{\sum (1/s_i)}$

A person travels 10km 3 times at speed

4 km/h, ~~5 km/h~~, 6 km/h.

Avg speed:-  $\frac{3}{(1/4 + 1/5 + 1/6)} \text{ m/h}$

$\rightarrow$  Same direction  $\Rightarrow$  relative speed  $(x-y)$  km/h

$\rightarrow$  Opp direction  $\Rightarrow$   $x+y$  km/h

$x-y, x+y$ : - basically the distance b/w them

after 1 hour  $+ MA II \Leftarrow$

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Q.  total time = 2 hour 54 min  
 total distance = ? (one side)

$$d_1 = d_2$$

$$25 \times t_{in} = 4 \times (2.9 + t) \text{ km}$$

$$t_{in} = \text{brake}$$

Q. Walking at speed  $5 \text{ km/hr}$  + 7 min late  
 if  $1 \text{ km/hr}$  faster + 5 min early

$$O = \left( \frac{4}{5} + \frac{12}{6} \right) (25 - t)$$

difference b/w time = true time

$$= \frac{40}{6} + 20 - t$$

$$= \frac{20}{3} + 20 - t$$

$$= \frac{80}{3} - t$$

$$= 26\frac{2}{3} - t$$

$$= 26.67 - t$$

$$= 26.67 - 12$$

$$= 14.67$$

$\therefore$  time taken by  $O$  =  $14.67$  hours

Q.

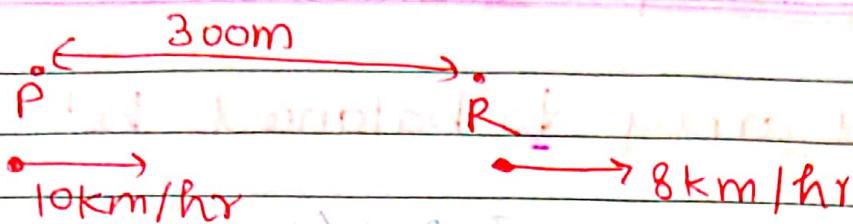
$$\text{A} \xrightarrow[11\text{ AM}]{400} \text{B} \xleftarrow[\text{11 AM}]{\text{E}} \text{C}$$

$65 \text{ km/hr} \left( \frac{1}{2} + \frac{1}{2} + \frac{1}{2} \right) 35 \text{ km/h}$

$$\text{Time} = \frac{\text{Distance}}{\text{Speed}} = \frac{400 \text{ km}}{100 \text{ km/hr}} = 4 \text{ hours}$$

Arrival time = 11 AM + 4 hours = 3 PM

Q.



From the diagram (two people) to edge ribbon I think

(top) How long will R run before being caught?  
 $\rightarrow$  relative speed = 2 km/hr.

$$\text{time} = \frac{0.3}{2} = 0.15 \text{ hours}$$

$$(2 \text{ km}) \times 0.15 \text{ hours} = \text{distance} = 0.3 \text{ km}$$

From 0.15 hours R travels

$$\text{distance} = 0.15 \times 8 = 1.2 \text{ km}$$

OR

both run the same time.

$$\text{time}_1 = \text{time}_2$$

$$\Rightarrow \frac{x}{8} = \frac{0.3+x}{10}$$

$$(x+7) = 1.2, \text{ so } x = 0.2 \text{ km}$$

$$25 \text{ m/min} = 0.25 \text{ km/hr}$$

$$5x(0.25) = 6x(0.25)$$

$$75 = 6x$$

$$d = 5x(6x(0.25))$$

$$= 6 \times 0.25 \text{ km/hr}$$

$$(75 + 6x) \text{ km/hr}$$

$$11 \text{ km/hr} = 100 \text{ m/min}$$

$$11 \text{ km/hr} = 100 \text{ m/min}$$

$$(100 - 6x) \text{ m/min}$$

$$100 - 6x = 25$$

$$x = 12.5$$

## Binary Tree problems

### Sorted array to balanced bst

```

struct TNoder* sToBst (int arr[], int start,
{                                , int end)
    if (start > end) {
        return NULL;
    }
    int mid = (start + end) / 2;
    struct TNoder* root = new Node(arr[mid]);
    root->left = sToBst(arr, start, mid - 1);
    root->right = sToBst(arr, mid + 1, end);
    return root;
}
    
```

Root to leaf path sum =  $\frac{x}{8} + \dots$

```

bool hasPathSum (Node* node, int sum)
{
    if (node == NULL)
        return (sum == 0);
    else
    {
        bool dsum = 0;
        int dsum = sum - node->data;
        if (node->left)
            dsum = dsum || hasPathSum(node->left,
                                         ans);
        if (node->right)
            dsum = dsum || hasPathSum(node->right,
                                         dsum);
        return dsum;
    }
}
    
```

$j + \text{col}[k] \geq 0 \& k \leq \text{MAX}$

while(n)

{  
  n = n & (n - 1)

  Count++

}

1110	110	101	1011
100	100	011	1111
011	011	110	10000
000	110	110	11
	2	2	

DFS → When find connected components

(from \* components → (BFS also) island)

①

7

111

$n = 3$

110

110

\* \* `memset(array, 0, sizeof(array) + M)` ② 101

Find number of islands (components)

int visited[MAX][MAX];

int find(int A[MAX][MAX], int N, int M) 0

{

: (0, i, toor) + nira

  int count = 0;

  for(i = 0; i < MAX; i++)

    for(j = 0; j < MAX; j++)

      if (A[i][j] && !visited[i][j])

        DFS(A, i, j);

        Count++;

    }

  }

  memset(visited, 0, sizeof(visited));

  return count;

  int bDFS(int row, int col, int toor) {

    visited[r][c] = 1;

    static int rowr = {-1, 0, 1, -1, 1, 0, 1, 0};

    static int colr = {-1, 0, 1, 0, -1, 1, 0, 1};

    for(int k = 0; k < 8; k++) {

      if (isValid(r + rowr[k], c + colr[k]))

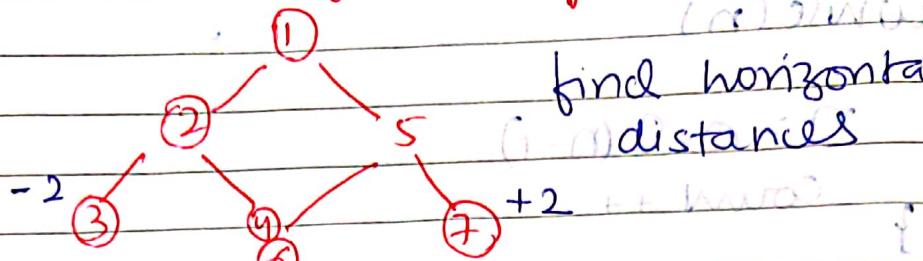
        if (!visited[r + rowr[k]][c + colr[k]])

          DFS(A, r + rowr[k], c + colr[k]);

}

Teacher's Signature.....

## Vertical order of a binary tree



Output 3 2 1 4 6 5 7

```
void verticalOrder( Node *root )
```

```
int min=0, int max=0  
findMin( root, &min, &max, 0 );  
for( i=min; i<=max; i++ )  
    print( root, i, 0 );
```

```
}
```

```
void find( Node *root, int *min, int *max,
```

```
int level, int hd ) {
```

```
if( root == NULL ) return;
```

```
if( hd < *min )
```

```
*min = hd
```

```
if( hd > *max ) *max = hd;
```

```
find( root->left, min, max, hd-1 );
```

```
find( root->right, min, max, hd+1 );
```

```
void print( Node *root, int level, int hd )
```

```
if( root == NULL ) return
```

```
if( level != hd ) printf( "%d ", root->data );
```

```
print( root->left, level, hd-1 );
```

```
print( root->right, level, hd+1 );
```

```
(3) 100 + 3 [K] work A 270
```

Teacher's Signature.....

## Find next larger element (Using stack)

```
void printNGE( int arr[], int n )  
{  
    int i = 0  
    push(&s, arr[0]); // stack s:  
    for( i = 1; i < n; i++ )  
    {  
        char next = arr[i]; // store current element  
        if( isEmpty(s) == false )  
        {  
            ele = pop(&s); // number  
            while( ele < next ) // when  
            {  
                cout << ele << " -> " << next; // print  
                if( isEmpty(s) == true )  
                    break; // break  
                ele = pop(&s); // previous  
            }  
            if( ele > next ) // when  
            {  
                push(&s, ele); // store number  
                push(&s, next); // store next  
            }  
        }  
        else // when stack is empty  
        {  
            cout << "-1" << " -> " << next; // print  
            ele = pop(&s);  
            nge = -1;  
            pf(" -. d --> %. d", ele, next);  
        }  
    }  
}
```

## Add 2 numbers represented by linked lists

$5 \rightarrow 6 \rightarrow 3$

$8 \rightarrow 4 \rightarrow 2$

$1 \rightarrow 4 \rightarrow 0 \rightarrow 5$

If same size following

```
node * addSameSize(Node * head1, Node * head2, int carry)
{
    if (head1 == NULL)
        return NULL;
    Node * result = (Node *) malloc(sizeof(Node));
    result->data = addSameSize(head1->next, head2->next, carry);
}
```

Sum = head1->data + head2->data + carry

carry = sum/10;

sum = sum%10;

result->data = sum;

return result;

$3 \rightarrow 5 \rightarrow 6 \rightarrow 3$  come to this pointer

$8 \rightarrow 6 \rightarrow 9$  call addSame

then pass  $3 \rightarrow 5$  and carry and  
head1 and add

and insert node at front of list.

Teacher's Signature.....

$$\begin{array}{l} 0 \rightarrow g \\ L \\ 0 \rightarrow g \end{array} \quad \begin{array}{l} A \rightarrow f \\ \downarrow \\ 10 \rightarrow 15 \end{array}$$

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## Intersection and union of LL.

Use merge sort first

TODO : Merge Sort LL ( $18 > 1 = 11$ )

(a) Union : (1)  $\text{frag}_1(18 \times)$

## BIT Manipulation

### Find $n^{\text{th}}$ Magic number

→ either a power of 5, or sum of  $(c.s^k, b.r^k)$  using powers of it.

5, 25, 125 ( $25 + 5$ ) ...

$7^{\text{th}}$  number → 11111111  $5^{\text{th}} \rightarrow 101$

b value ( $5^2 + 5^1 + 5^0$ ) mid value ( $5^2 + 5^1 + 5^0$ )

Swap bits ( $101110001$ )  $\rightarrow$  (01110010)

Get all even bits  $\rightarrow$  unsigned int even =

$(b - c5) >> n) | (b << m) \& 0xAAAAAAAA$

odd =  $b \& 0x55555555$

A ~~ans~~ = even = even  $>> 1$

odd = odd  $<< 1$

ans = even  $\#$  odd.

### Find the element that appears once.

(3 3 3 3 5 5 5 1 1 2) ~~13032~~

for(i=0; i < INT\_SIZE; i++)

{ sum = 0

l = 1 << i;

for(j=0; j < arr[i]; j++)

if(arr[j] & l)

count++;

if(sum % 3),

result = result | ~~second X'~~

return result;

Teacher's Signature.....

Find binary representation of a number.

```
for(i = 1 << 31; i > 0; i = i/2)
    if(x & i) ? printf("1") : printf("0");
```

OR

```
bin(n)
```

```
{ if(n >= 1)
    bin(n/2)
    printf('.%d', n % 2); }
```

Rotate bits in a number

left rotation ( $n, d$ ) by  $d$

```
return (n << d) | (n >> (32 - d));
```

right rotation ( $n, d$ )

```
return (n >> d) | (n << (32 - d));
```

Count no. of bits to be flipped to convert  $A \rightarrow B$

do XOR then count set bits

Count triplets with sum smaller than a given element.

sort array [see gfg xD]

16

Find  $a^b \% m$

~~\*y in O(logn) time.~~

(while ( $y > 0$ ))

```

if (y & 1) res = res * x
y = y >> 1
x = x * (x)gog = 10v
(2)gog = 610v

```

[if gog not true]

Convert infix to postfix

10v + 610v, 8) debug : +

```

for (i = 0 ; k = -1 ; exp[i] ; i++)
{
    if (isOperand(exp[i])) push(exp[i])
    else if (exp[i] == ')') while (!isEmpty(stack) && peek(stack) != '(')
        exp[++k] = pop(stack)
    else
        if (peek(stack) <= prec(exp[i])) push(exp[i])
        else
            exp[++k] = pop(stack)
}
```

\*

else

```

if (!isEmpty(s) && prec(exp[i]) <= prec(peek(s)))
    exp[++k] = pop(s)
    push(s, exp[i]);
```

{ } w3073 w3073

while (!isEmpty(s))

exp[++k] = pop(s);

exp[++k] = '0'

printf("%s", exp);

9018

9018

9018

9018

9018

9018

9018

## Evaluate postfix expression

evaluate C exp char [] exp )

```
for(i=0; i < strlen(exp); i++)  
    {
```

```
        if(isdigit(exp[i]))  
            push(s, exp[i] - '0')  
        else
```

```
            val1 = pop(s)
```

```
            val2 = pop(s)
```

```
            switch(exp[i])
```

```
                {
```

```
                    '+': push(&, val2 + val1);
```

```
(++i, exp[i] : + - * / % == )
```

```
)
```

```
printf(pop(s))
```

No of digits in a number

$\lceil \log_{10}(n) \rceil$

$\lceil \log_{10}(n) \rceil = \lceil \log_{10}(n+1) - 1 \rceil$

Permutation in lexicographically inc order.

↓ first      ↓ second

A G S B

find first (so that  $a[first] < a[first+1]$ )

find second = [greater than  $a[first]$ ]

but the (smallest)

swap

sort after  $a[first]$

A S G

B (2nd)

A S B

(2nd)

(2nd)

DFS → inorder, pre, post

BFS → level order.

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## Print min level (Leaves sum)

### 1) Find the minimum level

```
int minLevel (Node *root, level) → 1.
```

```
{ if (root == NULL)  
    return 0;
```

```
if (!root->left && !root->right)
```

```
    return level; → int left, right = INT_MAX;
```

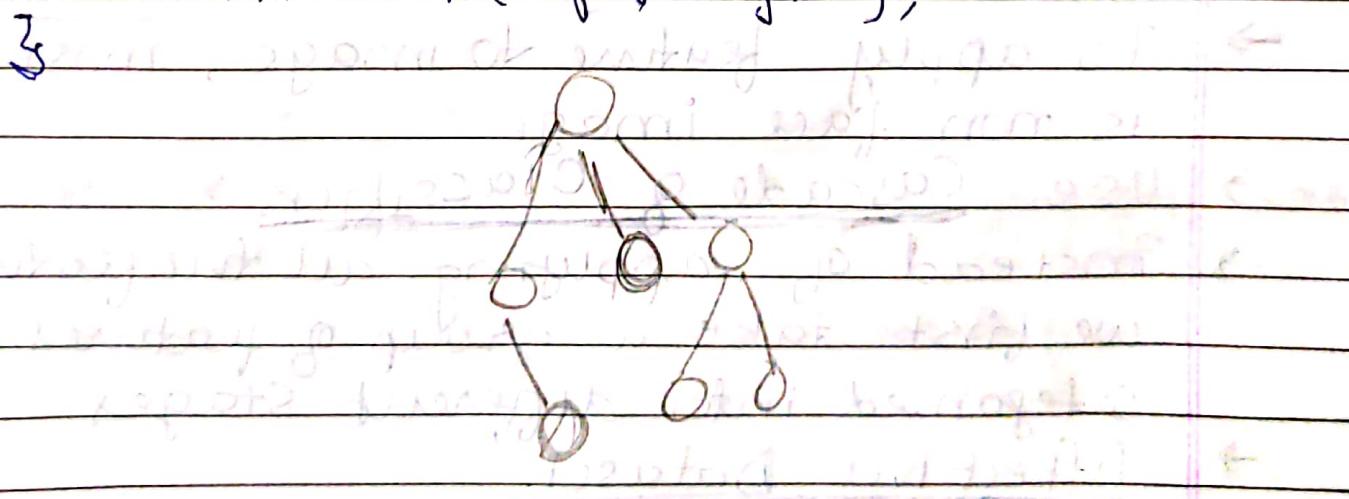
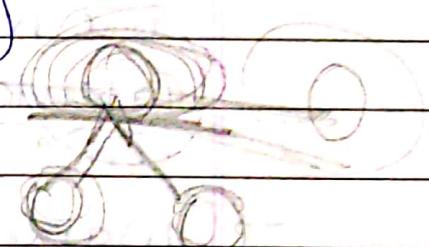
```
if (root->left != NULL)
```

```
left = minLevel (root->left, level + 1)
```

```
if (root->right != NULL)
```

```
right = minLevel (root->right, level + 1)
```

```
return min (left, right);
```



CIP  
Ex. (3)