David Utyuganov's File System

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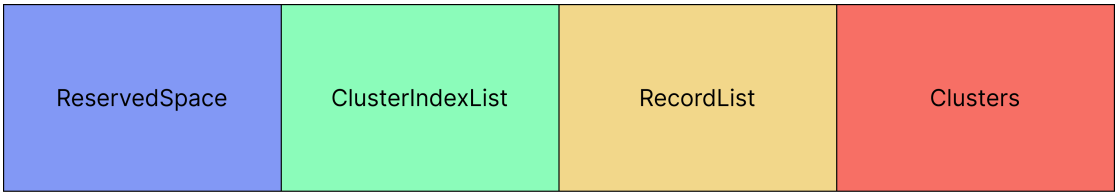
DUFS Architecture Overview



General DUFS architecture contains 4 blocks. First 3 of them are static and generated in the process of volume's mount, while the last used as storage of the content.
Both directory and file structures are linear lists and this is the key disadvantage of the filesystem at this moment. It could be remade on b-trees to significantly reduce disk operations.
Currently it does not support logging/journaling or any other features for safe use. If there is met any defect during work of filesystem, most probably, the volume will become corrupted and some part of it's data will be lost forever.

ReservedSpace

ReservedSpace is a block that contains specific service information for fast and correct work of the filesystem. Total size of this block is 60 bytes and the fields are the following:

ReservedSpace		
Field	Bytes	Description
DUFSNoseSignature	4	Constant value equal to 0x44554653 (DUFS). Used as key-signature to detect if the file is DUFS volume
VolumeName	16	8 UTF-16 symbols
ClusterSize	4	Cluster size value. Must be divisible by 4
VolumeSize	8	Brutto size of the volume
ReservedClusters	4	Number of netto reserved clusters
CreateDate	2	Date of volume creation
CreateTime	2	Time of volume creation
LastDefragmentationDate	2	Date of volume's last defragmentation
LastDefragmentationTime	2	Time of volume's last defragmentation
NextClusterIndex	4	Index of next free cluster index
FreeClusters	4	Number of free clusters in the volume
NextRecordIndex	4	Index of next free record index
DUFSTailSignature	4	Constant value equal to 0x4A455432 (JETB). Same as DUFSNoseSignature

ClusterIndexList

ClusterIndexList is a linear list that keeps information about relationships between clusters. It is similar to FAT.
Each element of this list is ClusterIndexElement that takes 12 bytes.
Length of ClusterIndexList is equal to the number of reserved clusters and it cannot be greater than $2^{32} - 1$. This limit could be increased up to $2^{64} - 1$.
ClusterIndexElement structure:

ClusterIndexElement		
Field	Bytes	Description
NextClusterIndex	4	Index of next cluster in the chain
PrevClusterIndex	4	Index of previous cluster in the chain. Uses to perform fast defragmentation
RecordIndex	4	Index of record that contains this cluster. Uses to perform fast removal of record from parent directory's cluster

RecordList

RecordList is a linear list that keeps information about records in the volume.
Record is an entity to generalize files and directories in one term. It is useful because files and directories mostly contains topologically identical meta data.
Length of RecordList is equal to the number of reserved clusters.
One Record takes 93 bytes of data:

Record		
Field	Bytes	Description
Name	64	32 UTF-16 symbols
CreateDate	2	Date of record creation
CreateTime	2	Time of record creation
FirstClusterIndex	4	Index of first cluster in the chain
LastEditDate	2	Date of last edition of the record
LastEditTime	2	Time of last edition of the record
Size	8	Size of record in bytes
ParentDirectoryIndex	4	Index of record of parent directory
ParentDirectoryIndexOrderNumber	4	Order number of this record in parent directory's cluster. Used to find/delete/remove records from parent directory's cluster by O(1)
IsFile	1	Byte that says if this record is file or directory

Clusters

This block uses as a storage of the content for files and as a storage of nesting records for directories.

Comparison With Other File Systems

File system	Max. record name	Max. file size	Max. number of files	Bake
NTFS	255	16TiB	$2^{32}-1$	Yes
FAT32	8.3	4GiB	?	No
ext4	255	16TiB	$2^{32}-1$	No
DUFS	32	16TiB (?)	$2^{32}-1$	Yes

Performance And Scalability Analysis

RAM

During filesystem usage RAM is always very low. Only 60 bytes of additional data keeps in the RAM all the time, everything else is imperceptible. Read/write/append uses bytes in RAM equal to cluster size through buffered read/write.

CPU

During filesystem usage CPU load is also imperceptible. Read/write/append operations overhead grows linearly to the length of the content. Find operation overhead grows linearly to the number of records in the directory. Low-level overhead (CPU L-cache/registers) grows probably linearly to the number of records. Defragmentation overhead depends on the quality of the initial approximation – if the volume is almost defragmented, the overhead of the defragmentation on top of it will be very low.

DISK

DISK complexity is arguable. Generally it is ok from the algorithmic point of view (besides linear search complexity in cause of linear structure of record list), but it could be optimized from the software point of view to reduce the "constant" complexity (memorization/caching of the results of DISK read operations).
Algorithmically it could be significantly reduced if linear structure of record list will be replaced with the b-tree structure so the complexity of find operation will be reduced from linear to logarithmical

Scalability

DUFS is opened to select any cluster size that is divisible by 4 and any volume size that is less than 16EiB (2^{64} bytes), but the number of clusters $= \left\lceil \frac{\text{volume size}}{\text{cluster size}} \right\rceil$ must be less than 2^{32}. This value theoretically could be increased up to 2^{64}.

DUFS API

DUFS is implemented as java class with the following public interface:

```
public class Dufs {
    public RandomAccessFile getVolume();
    public void mountVolume(String path, int clusterSize, long nettoVolumeSize);
    public void attachVolume(String path);
    public void closeVolume();
    public void createRecord(String path, String name, byte isFile);
    public void writeFile(String path, File file);
    public void appendFile(String path, File file);
    public void readFile(String path, File file);
    public void deleteRecord(String path, byte isFile);
    public void renameRecord(String path, String newName, byte isFile);
    public void moveRecord(String path, String newPath, byte isFile);
    public void printDirectoryContent(String path);
    public void printVolumeInfo();
    public void printVolumeRecords();
    public void printDirectoryTree();
    public void defragmentation();
    public void bake();
    public void unbake();
}
```

This class simulates file system on top of volume that could be mounted or attached through methods. It supports basic I/O operations like read/write/append, move/delete operations, rename, defragmentation, bake/unbake and several printing methods to represent nested data via console. DUFS volume itself is implemented on top of `java.io.RandomAccessFile`, I/O operations are implemented on top of `java.io.File`.

`void`	`mountVolume(String path, int clusterSize, long nettoVolumeSize)`
Mounts (creates) a volume by the following `path` with the following `clusterSize` and the following `nettoVolumeSize`. Size of volume will be bigger after the creation of file since `nettoVolumeSize` is the size that indendent only for clusters.	
Throws `DufsException` if such volume already exists by this `path`; Throws `DufsException` if given volume name length is beyond 8 symbols; Throws `DufsException` if given volume name length contains prohibited symbols; Throws `DufsException` if given `nettoVolumeSize` is greater than $1.1 \cdot 10^{12}$ *(1 TiB or 1024 GiB);* Throws `DufsException` if there is not enough space on disk to mount the volume; Throws `DufsException` if `clusterSize` is not divisible by 4.	

`void`	`attachVolume(String path)`
Attaches volume by the following `path`.	
Throws `DufsException` if there is no volume with such name by this `path`; Throws `DufsException` if volume signatures doesn't match with the expected values.	

`void`	`closeVolume()`
Closes the volume.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation).*	

void	createRecord(String path, String name, byte isFile)
Creates the record with the following `name` by the following `path` in the volume. Record could be either file (`isFile` == 1), either directory (`isFile` == 0).	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*; Throws `DufsException` if given record name length is beyond 32 symbols; Throws `DufsException` if given record contains prohibited symbols; Throws `DufsException` if record with such name and type already contained in the `path`; Throws `DufsException` if there is not enough space on the volume to create new record.	

void	writeFile(String path, File file)
Writes the content from `file` into the file *(record)* in the volume by the following `path`. Overwrites any content that is already contained in the file *(record)*	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*; Throws `DufsException` if there is not enough space on the volume to write given content; Throws `DufsException` if file *(record)* does not exist.	

void	appendFile(String path, File file)
Appends the content from `file` to the file *(record)* in the volume by the following `path`.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*; Throws `DufsException` if there is not enough space on the volume to append given content; Throws `DufsException` if file *(record)* does not exist.	

void	readFile(String path, File file)
Reads the content from file *(record)* in the volume by the following `path` and writes it into the `file`.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*; Throws `DufsException` if file *(record)* does not exist.	

void	deleteRecord(String path, byte isFile)
Deletes record by the following `path` from the volume.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*; Throws `DufsException` if `isFile` == 0 and the directory *(record)* is not empty; Throws `DufsException` if file *(record)* does not exist.	

void	renameRecord(String path, String newName, byte isFile)
Sets record's name by the following `path` in the volume as `newName`.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*; Throws `DufsException` if given `newName` length is beyond 32 symbols; Throws `DufsException` if given `newName` contains prohibited symbols; Throws `DufsException` if record with `newName` and type already contained in the `path`; Throws `DufsException` if record by the following `path` does not exist.	

void	moveRecord(String path, String newPath, byte isFile)
Moves record by the following `path` into the `newPath` in the volume.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*; Throws `DufsException` if record with such name and type already contained in the `newPath`; Throws `DufsException` if record does not exist.	

void	printDirectoryContent(String path)
Prints records that contained by the following `path` in the volume. Used for debug purposes, so does not catch any errors. **Unsafe to use**.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*.	

void	printVolumeInfo()
Prints volume info, such as date/time of creation, reserved/free size etc..	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*.	

void	printVolumeRecords()
Prints all records that contained in the volume. Used for debug purposes, so the data is specific, but it prints full cluster chain for each record.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*.	

void	printDirectoryTree()
Prints volume's directory tree.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*.	

void	defragmentation()
Performs defragmentation of the volume. Order of clusters is defined by the order of indexes in the record list.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*.	

void	bake()
Bakes the volume. calls `defragmentation()` inside and then truncates the volume to the minimal possible size without losing any data. After baking the volume becomes read-only. Non-read operations's behaviour on the baked volume is undefined.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*.	

void	unbake()
Unbakes the volume to the initial size. After unbaking the volume is ready to be used for any operation.	
Throws `DufsException` if volume is `null` *(was not mounted or attached before this operation)*.	