### **Error Control and Detection**

#### **Fundamentos de Redes**

Mestrado Integrado em Engenharia de Computadores e Telemática DETI-UA



### Error Control in a Communication Channel

#### Causes:

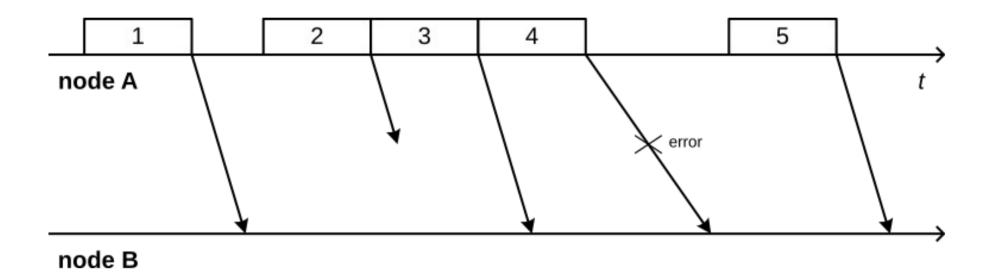
- Corrupted packets (data received with errors);
- Lost packets;
- Packets received out of order.

#### Solutions:

- Sender and receiver must be able to coordinate between them to retransmit a lost or corrupted packet.
- Retransmit protocols: ARQ (Automatic Repeat reQuest).

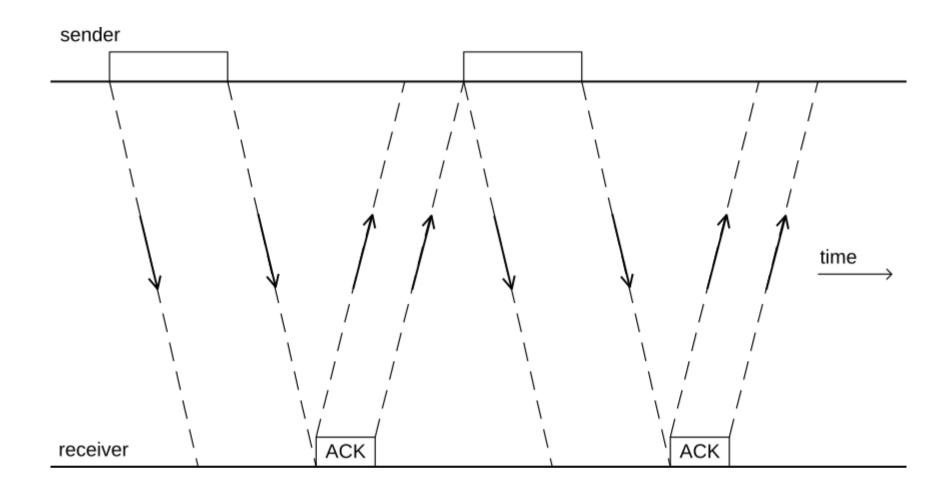
### **Error Control Assumptions**

- Errors are always detected;
- Frames/packets may have variable (limited) delays:
- Some frames/packets may be lost.



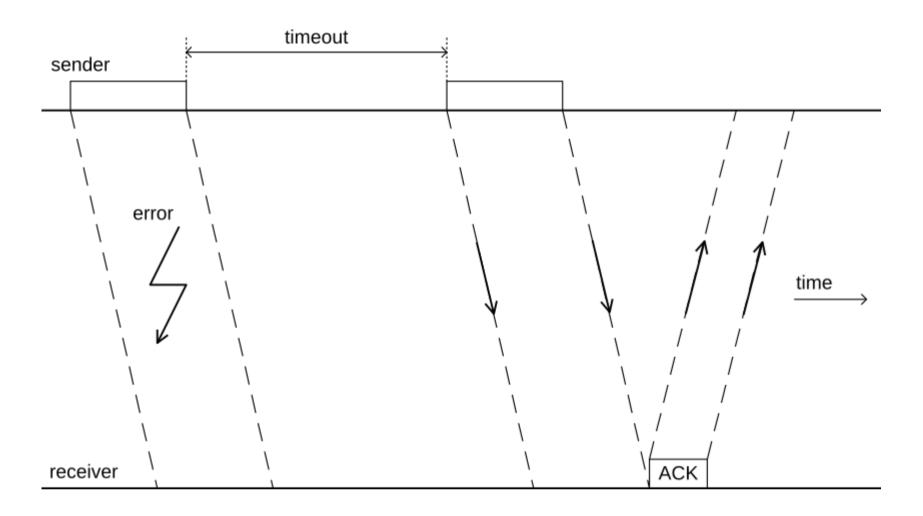
# Stop-and-Wait (SW)

#### Operation without errors:



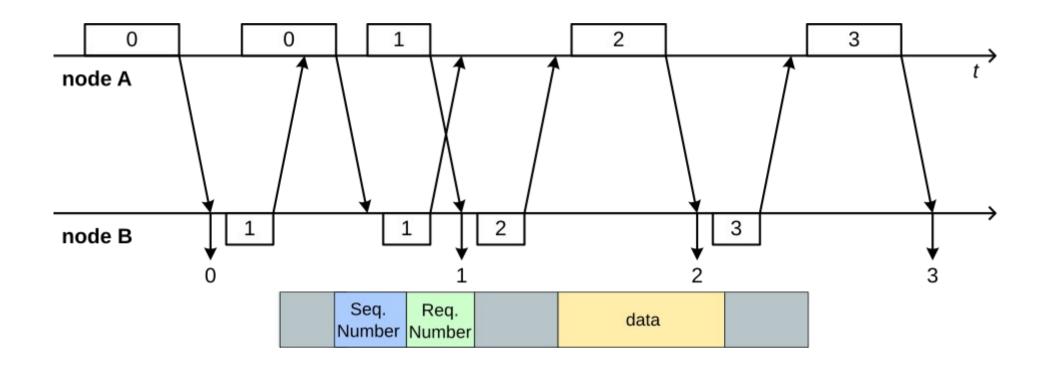
# Stop-and-Wait (SW)

Operation with error and recover:



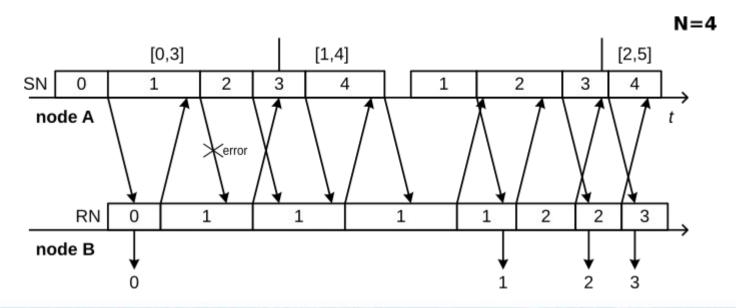
# Stop-and-Wait (SW)

- Messages have a sequence number and a request number.
  - The sequence number identifies the messages sent.
  - The request number allow the receiver to notify the sender about the received message and next message expected.



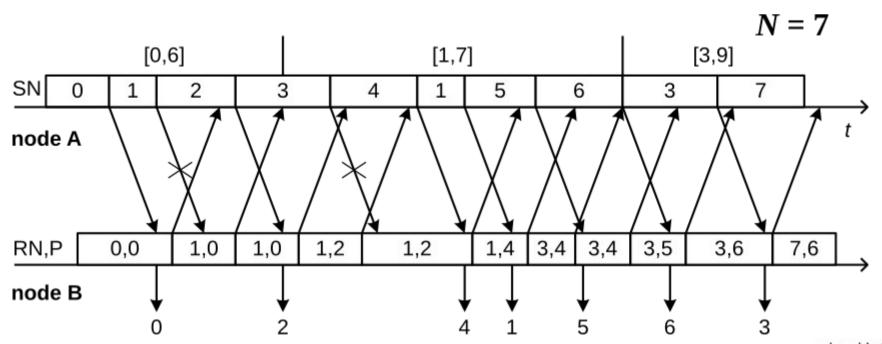
# Go-Back-N (GB-N)

- The sender is allowed to send more than one message before receiving the confirmation of reception of the previous messages.
  - Window with size N N defines the number of messages that may be sent without confirmation of reception of the first.
- •The sender, after a timeout, resends the first message without confirmation of reception and all of the following.
- •The receiver, after receiving all messages until a sequence number equal to n (SN=n), accepts only the message with SN=n+1.
  - Drops all the others.
- •The receiver, responds with a request number (RN) to identified the next message not received.
  - Implicitly, indicates that all messages up to RN were correctly received.



# Selective Repeat (SR)

- The sender is allowed to send more than one message before receiving the confirmation of reception of the previous messages.
  - Window with size N N defines the number of messages that may be sent without confirmation of reception of the first.
- •The sender, after a timeout, resends a message without confirmation of reception.
- •The receiver, after receiving all messages until a sequence number equal to n (SN=n), accepts to receive all messages with SN>n and SN  $\leq$  N+n.
- •The receiver, responds with a request number (RN) to identified the next message not received, and the higher payload/message number correctly received (P).



#### **Error Detection**

#### Performed at multiple Layers.

- Nowadays the (new protocols) tendency is to be performed only at the physical/link layer and application layer.
  - Has a performance impact.
  - E.g., IPv4 supports it, IPv6 does not.

#### Methods

- CRC (Cyclic-Redundancy Check)
  - Based on the theory of cyclic error-correcting codes.
  - Adds a fixed-length check value to messages.
  - Requires of a generator polynomial.
  - The binary message with the CRC field is divided by the generator polynomial.
    - No errors imply remainder equal to zero.
- Checksums and hash/digest values
  - Generated by a predefined function that receives as input all message bits.
  - Value is appended to message.