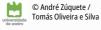
Digital signatures

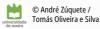


Applied Cryptography

1

Digital signatures: goals

- > Authenticate the contents of a document
 - Ensure its integrity
- > Authenticate its author
 - Ensure the identity of the creator/originator
- ⊳ Non-repudiation
 - Prevent signing repudiation



Applied Cryptography

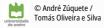
Digital signatures: fundamental approach

> Signature generation

- Encrypt with a private key
- Signer (or signatory) is the private key owner

> Signature validation

- Decrypt with the public key
- Anyone can verify
 - · Since public keys can be universally known
- Signature can be linked to the public key owner



Applied Cryptography

3

Signature schemes

- The message is fully recovered upon a signature validation
- Signature validation is mandatory prior to message observation

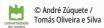
- The signature is detached from the message
- The message can be observed anytime



Applied Cryptography

Key elements of a digital signature

- > The message (or document)
 - It only makes sense with the signed object
- > The signature date
 - Because is usually required
 - Because key pairs have validity periods
- > The identity of the signatory
 - Otherwise it would not mean anything

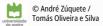


Applied Cryptography

5

The document to sign

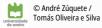
- ▷ It may accommodate digital signatures as appendixes
 - PDF, XML
 - DOCX (archive of XML components)
- Other formats may group document and signature
 - S/MIME (mail)
 - JOSE (JSON Object Signing and Encryption)



Applied Cryptography

The signature date

- ▷ It may be given by the signatory machine
 - Does not protect against time forgery attacks by the signatory
- It may be given by a Time Stamping Authority (TSA)
 - Does not protect against the future discovery of the private keys used



Applied Cryptography

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The identity of the signatory

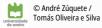
- □ Usually provided by a X.509 public key certificate
 - It provides several attributes of the identity
 - It provides the public key for signature validation
 - It provides the acceptable signing time frame
 - · Together with the respective CRL



Applied Cryptography

Optional elements of a digital signature

- > Attributes that can help to interpret it
 - Location
 - · Where it was signed
 - Reason
 - · Why it was signed
 - Appearance
 - Handwritten signature (usually without legal value)
 - Name of the signatory
 - · Date of signature
 - · Some kind of logo



Applied Cryptography

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Digital signatures' algorithms

- Message recovery scheme
 - Asymmetric encryption and decryption
 - Only for RSA
- Verification info→K_x
 D(K_x, A_x(doc))

Check integrity of doc

- Message appendix scheme
 - Digest functions
 - Asymmetric signature and validation
 - RSA, ElGamal (DSA), EC
- ⊳ Signing

```
A_x(doc) = info + E(K_x^{-1}, h(doc+info))

A_x(doc) = info + S(K_x^{-1}, h(doc+info))
```

Verification
 info→K_x

 $D(K_{x'} A_x(doc)) \equiv h(doc + info)$ $V(K_{x'} A_x(doc), h(doc + info)) = True$

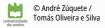


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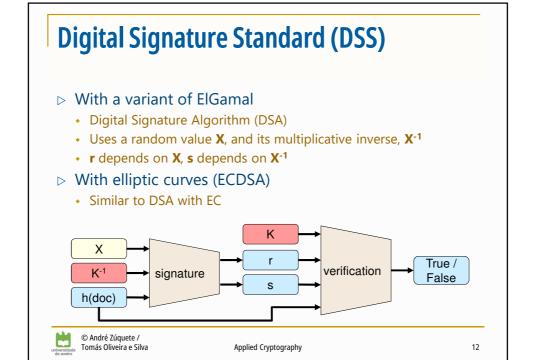
Applied Cryptography

RSA signatures

- Creation with private key
 - Validation with the corresponding public key
- Special padding for Signature Scheme w/ Appendix
 - RSASSA-PKCS#1 (v1.5)
 - Deterministic
 - RSASSA-PSS (Probabilistic Signature Scheme)
 - · Randomized (EMSA-PSS)
- Hash function prefixing
 - ASN.1 algorithm OID

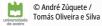


Applied Cryptography



Blind signatures

- ▷ Signatures made by a "blinded" signer
 - · Signer cannot observe the contents it signs
 - Similar to a handwritten signature on an envelope containing a document and a carbon-copy sheet
- □ Useful for ensuring anonymity of the signed information holder, while the signed information provides some extra functionality
 - Signer X knows who requires a signature (Y)
 - X signs T₁, but Y afterwards transforms it into a signature over T₂
 - Not any T₂, a specific one linked to T₁
 - Requester Y can present T₂ signed by X
 - \cdot But it cannot change T_2
 - X cannot link T₂ to the T₁ that it observed when signing



Applied Cryptography

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Chaum Blind Signatures

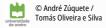
- - Blinding
 - Random blinding factor K
 - $\cdot \mathbf{k} \times \mathbf{k}^{-1} \equiv 1 \pmod{N}$
 - $\cdot m' = k^e \times m \mod N$
 - Ordinary signature (encryption w/ private key)
 - $\cdot A_x (m') = (m')^d \mod N$
 - Unblinding
 - $\cdot A_x (m) = k^{-1} \times A_x (m') \mod$



Applied Cryptography

Qualified electronic signature

- - Regulation No 910/2014
- - · Over long periods of time

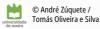


Applied Cryptography

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Qualified electronic signature

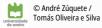
- > Three main requirements:
 - The signatory must be linked and uniquely identified to the signature
 - The data used to create the signature must be under the sole control of the signatory
 - Must have the ability to identify if the data that accompanies the signature has been tampered with since the signing of the message



Applied Cryptography

Qualified electronic signature

- Must be created using a qualified signature creation device
 - This device uses specific hardware and software that ensures that the signatory only has control of their private key
- > A qualified trust service provider manages the signature creation data that is produced
 - But the signature creation data must remain unique, confidential and protected from forgery



Applied Cryptography

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Signature devices

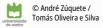
- - Smartcards
 - Cartão de Cidadão
- - Mainly for mobile devices
 - Chave Móvel Digital



Applied Cryptography

PKCS #11

- > Crypto tokens' standard interface
 - Cryptoki
- - Developed for a specific set of crypto tokens
- > Specification in C
 - There are interfaces for other languages



Applied Cryptography

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Microsoft Cryptographic API (CAPI)

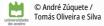
- - Applications use the abstractions it provides
- - Target-specific software module under the CAPI
 - It enables a particular functionality
 - Signature capabilities can be added with CSPs
 - For local crypto tokens
 - · For remote, cloud-based HSMs



Applied Cryptography

Long-Term Validation (LTV)

- A document signature may become invalid upon an initial verification
 - Due to a late certification revocation
- > Signature algorithms may become vulnerable
 - · Allowing signatures with old credentials to be forged
- > LTV attempts to handle both issues
 - With successive signature layers
 - Performed by signed documents' holders

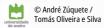


Applied Cryptography

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LTV Advanced Electronic Signatures (AdES)

- > PAdES
 - PDF Advanced Electronic Signature
- CAdES
 - Cryptographic Message Syntax Advanced Electronic Signatures
- > XAdES
 - XML Advanced Electronic Signatures



Applied Cryptography