

# CSE373 Assignment 1

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## 1 Optimal Network Transmission Combination

Lets assume that  $Request[i]$  contains 3 variables which are startime, endtime and size.

Lets assume that the list is sorted by the start time of each element.

At each position in the array, the maximum non overlapping sum is the maximum value out of the max overlapping sum till previous index or the max overlapping sum of current index where the max sum upto index  $j$  such that stop time doesn't overlap with current element.

$$T[i] = \begin{cases} Request[i].size & \text{if } (i == 0) \\ \max(T[i-1], T[j] + Request[i].time) & \text{s.t. } j = \max_{j \in [0, i-1]} Request[j] \text{ does not overlap } Request[i] \end{cases}$$

The asymptotic running time is  $O(n \log n)$

The asymptotic space requirement is  $O(n)$

## 2 Mining For Gold

At each position we can either go up or right. So the maximum gold value that can be found is the maximum value out of the max gold value obtained by going up and the max gold value obtained by going right plus the gold value of the current spot. If we can't go up then the max value is obtained by going right and if we can't go right then max value is from going up. At the top right corner the max value is the value of the position itself. At the end we return the gold value of  $T[0,0]$  where

$$T[i, j] = \begin{cases} Gold[i, j] & \text{if } (i == M - 1 \wedge j == N - 1) \\ Gold[i, j] + T[i, j + 1] & \text{if } (i == M - 1) \\ Gold[i, j] + T[i + 1, j] & \text{if } (j == N - 1) \\ Gold[i, j] + \max(T[i, j + 1], T[i + 1, j]) & \text{otherwise} \end{cases}$$

The asymptotic running time is  $O(mn)$

The asymptotic space requirement is  $O(mn)$

## 3 Maximum Product

For an array  $A$ , to find the subseries with the maximum product, we fill the dynamic programming table using the rules :

$$T[i, j] = \begin{cases} A[i] & \text{if } i == j \\ A[i] * T[i - 1, j] & \text{if } i > j \\ A[i] * T[i + 1, j] & \text{if } i < j \end{cases}$$

where  $i$  and  $j$  are indices in the array

The asymptotic running time is  $O(n^2)$

The asymptotic space requirement is  $O(n)$

## 4 String Decomposition

Lets say a and b are the 2 strings and m and n are the length of Our recursive function takes in all 3 strings: a, b and c. If a and b are both of length 0 then we return false. If the first lett of c is equal to the first letter of a then we we return the value returned by the recursive call containing a[1:n],b and c[1:m+n] . If not then we compare the first letter of b and c and if equal then the make the recursive call with a, b[1:n], c[1:m+n]. If that is also unequal then we return false.

$$T[i, j] \begin{cases} \text{false} & \text{if } m \text{ and } n \text{ are both } 0 \\ T[i-1, j] & \text{if } a[0] == c[0] \\ T[i, j-1] & \text{if } b[0] == c[0] \\ \text{false} & \text{otherwise} \end{cases}$$

The asymptotic running time is  $O(mn)$

The asymptotice space requirement is  $O(mn)$

Psuedocode:

```
def isStringDecomposition(a, b, c, i, j):
    if (len(a) == 0 && len(b) == 0):
        T[0,0] = False
        return false
    if (a[0] == c[0]):
        if (T[i-1, j] != None):
            return T[i-1, j]
        else:
            success = isStringDecomposition(a[1:i], b, c[1:i+j], i-1, j)
            T[i-1, j] = success
            return success
    elif (b[0] == c[0]):
        if (T[i, j-1] != None):
            return T[i, j-1]
        else:
            success = isStringDecomposition(a, b[1:j], c[1:i+j], i, j-1)
            T[i, j-1] = success
            return success
    T[i, j] = False
    return False
```

## 5 Two salesmen

We get the following recursive structure to this question.

Let  $C = \text{distance}(A)/2$

Let  $T[i]$  be the cost of splitting the list of cities  $A[0] \dots [i-1]$  into 2 arrays the we have

$$T[i] = \min_{t \in [0, i]} \sum_{l=t}^{i-1} |\text{distance}(l) - C|$$

The time complexity of this problem is  $O(n^2)$  The space complexity is  $O(n)$