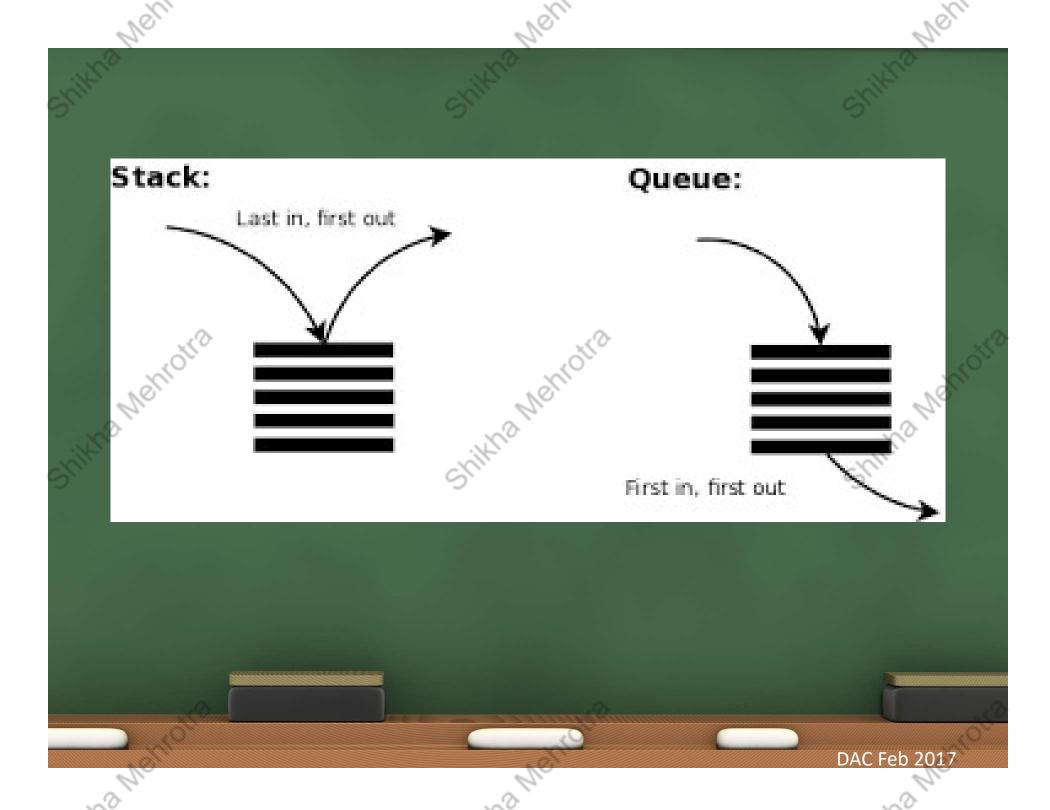
Data Structures and Algorithms Shikha Mehrotra

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Shillawehi The Queue **Shikha Mehrotra**



The Queue ADT

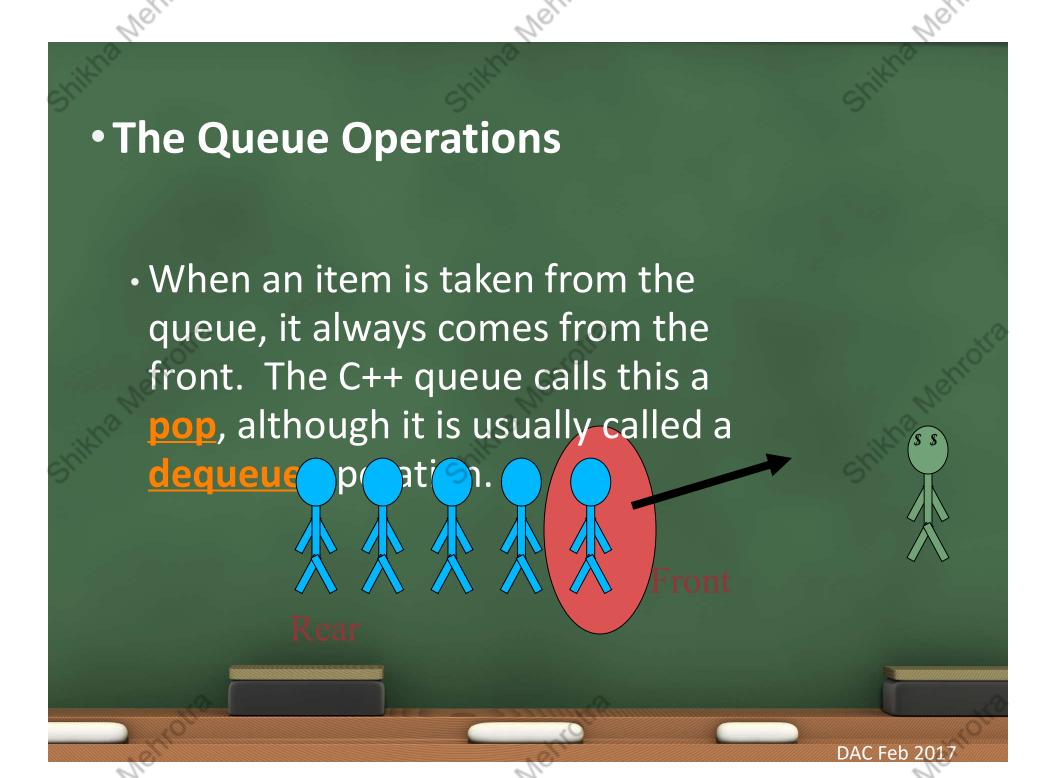
A list of collection with the restriction that insertion can be performed at one end (rear) and deletion can be performed at other end.

The Queue Operations A queue is like a line of people waiting for a bank teller. The queue has a front and a <u>rear</u>.

The Queue Operations

New people must enter the queue at the rear. The C++ queue class calls this a push, although it is usually called an enqueue operation.

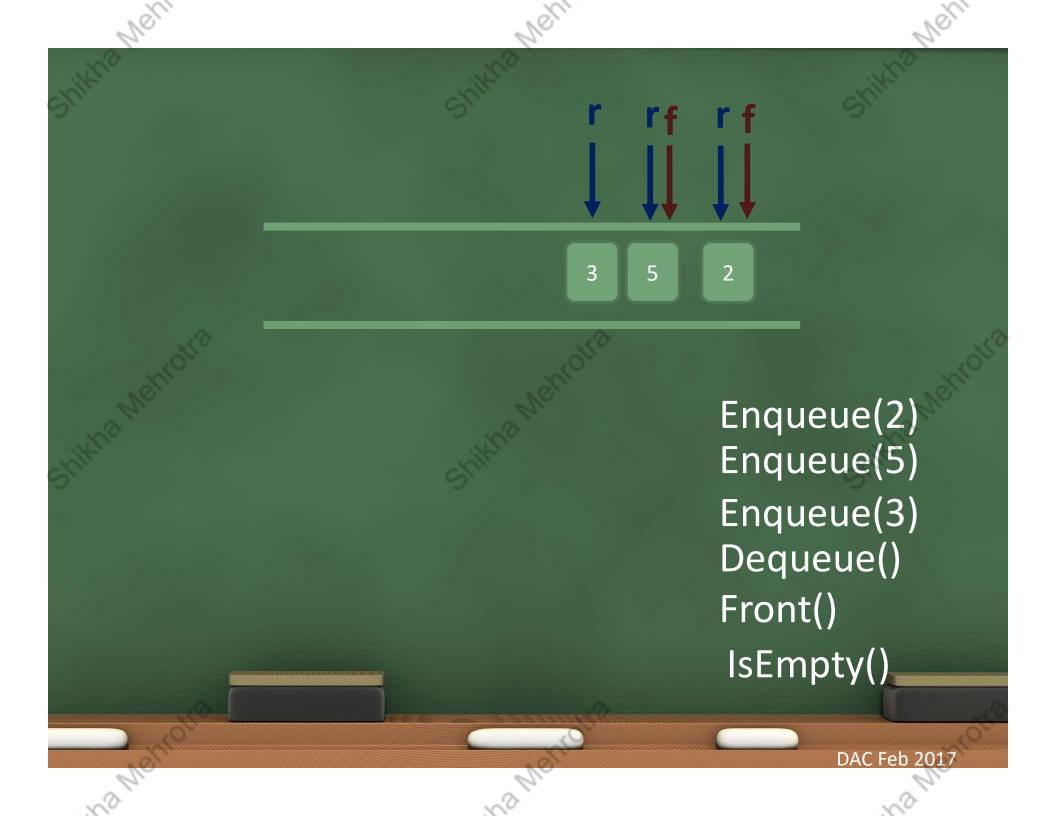




Queue ADT

```
AbstractDataType queue {
    instances
        ordered list of elements; one end is the front; the other is the rear;
    operations
        IsEmpty(): Return true if queue is empty, return false otherwise size(): Return the number of elements in the queue front(): Return the front element of queue
        dequeue(): Remove an element from the queue
        enqueue(x): Add element x to the queue
```

•Queue Enqueue Dequeue DAC Feb 2017



Array Implementation

• A queue can be implemented with an array, as shown here. For example, this queue contains the integers 4 (at the front), 8 and 6 (at the rear).

[0] [1] [2] [3] [4] [5] ... 4 8 6

An array of integers to implement a queue of integers

We don't care what's in this part of the array.

Array Implementation

• The easiest implementation also keeps track of the number of items in the queue and the index of the first element (at the front of the queue), the last element (at the rear).

3 size0 first2 last

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A Dequeue Operation

 When an element leaves the queue, size is decremented, and first changes, too. sizefirst

last

[0] [1] [2] [3] [4] [5] ... 8 6

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An Enqueue Operation When an element enters the size queue, size is incremented, and last changes, too. first last

At the End of the Array

• There is special behavior at the end of the array. For example, suppose we want to add a new element to this queue, where the last index is [5]:

3 size3 first

5 last

 [0]
 [1]
 [2]
 [3]
 [4]
 [5]

 2
 6
 1

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At the End of the Array The new element goes at the size front of the array (if that spot isn't already used): first last [5]

Array Implementation

- Easy to implement
- But it has a limited capacity with a fixed array
- Or you must use a dynamic array for an unbounded capacity
- Special behavior is needed when the rear reaches the end of the array.



3

size

0

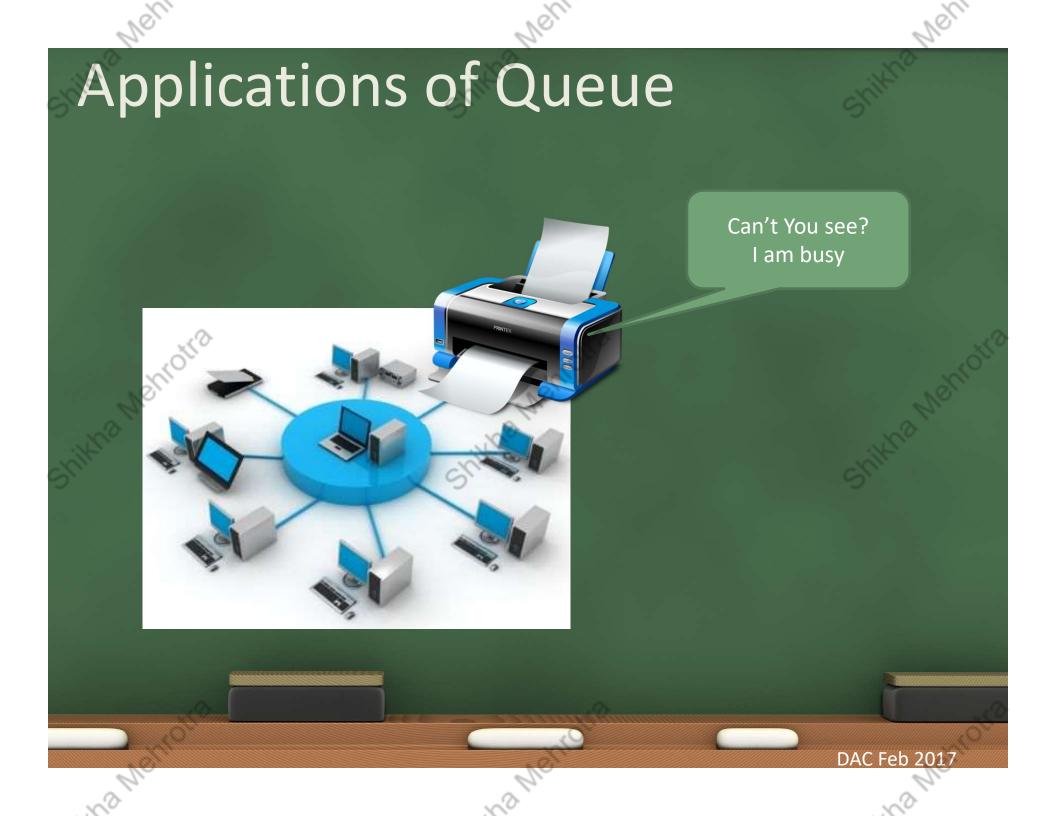
first

2

last

Applications of Queue





Applications of Queue



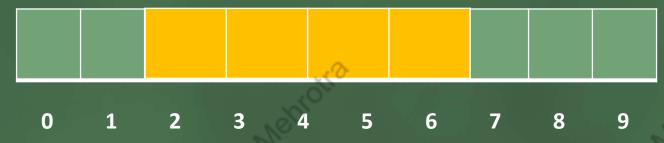


 Serving requests of a single shared resource (printer, disk, CPU),



Implementation of Queue

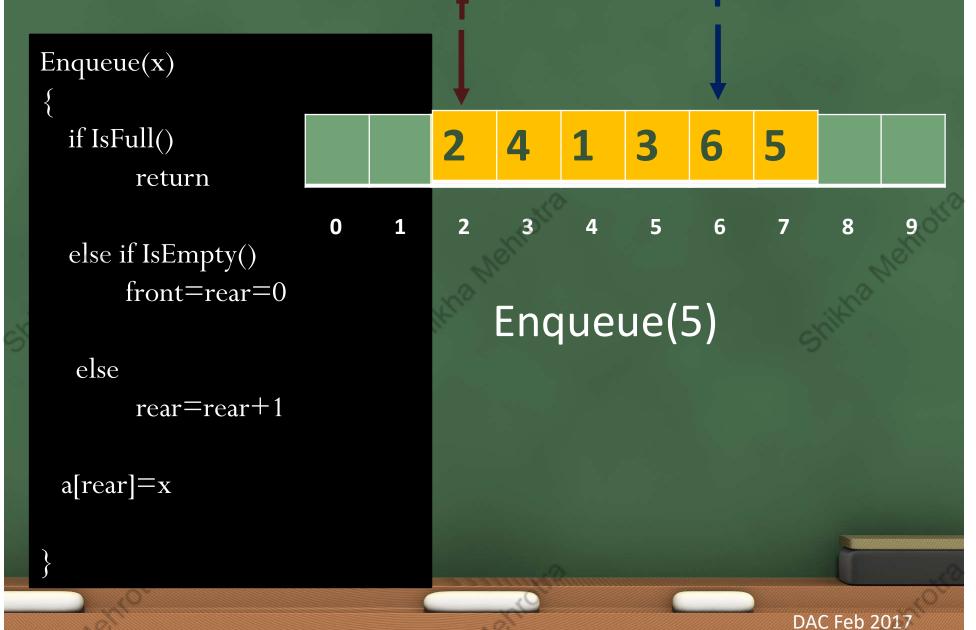


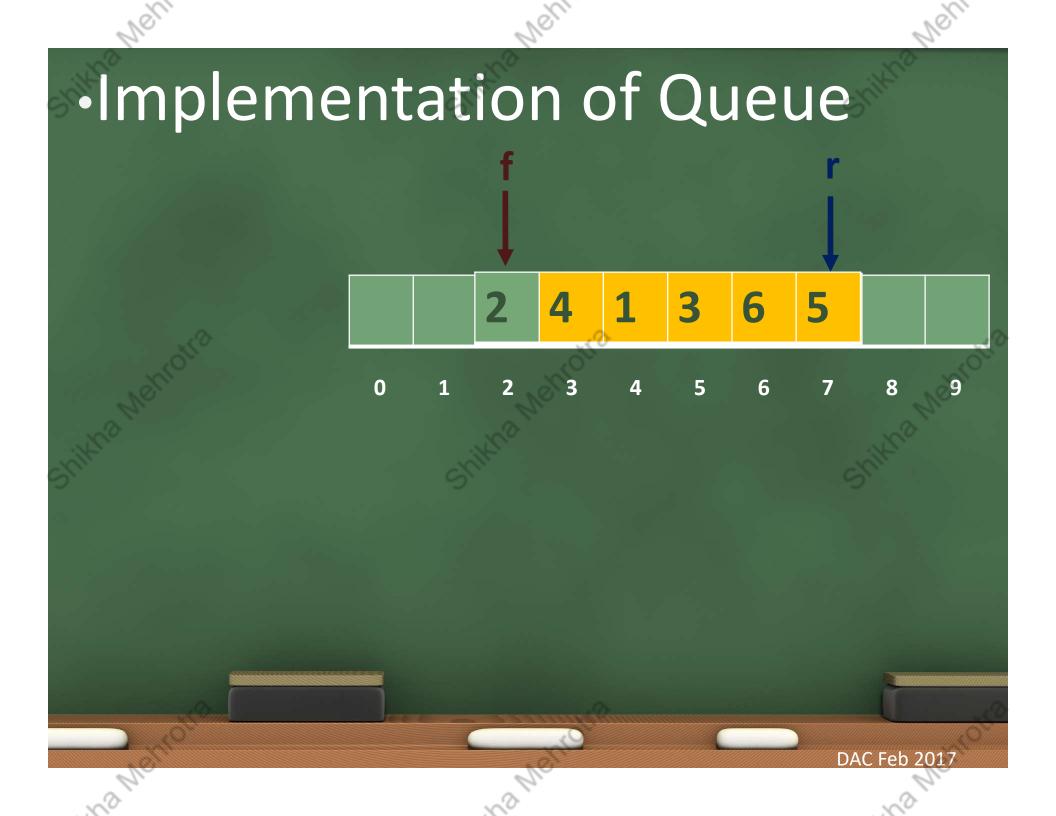


```
int A[10]
front=-1,
rear=-1
```

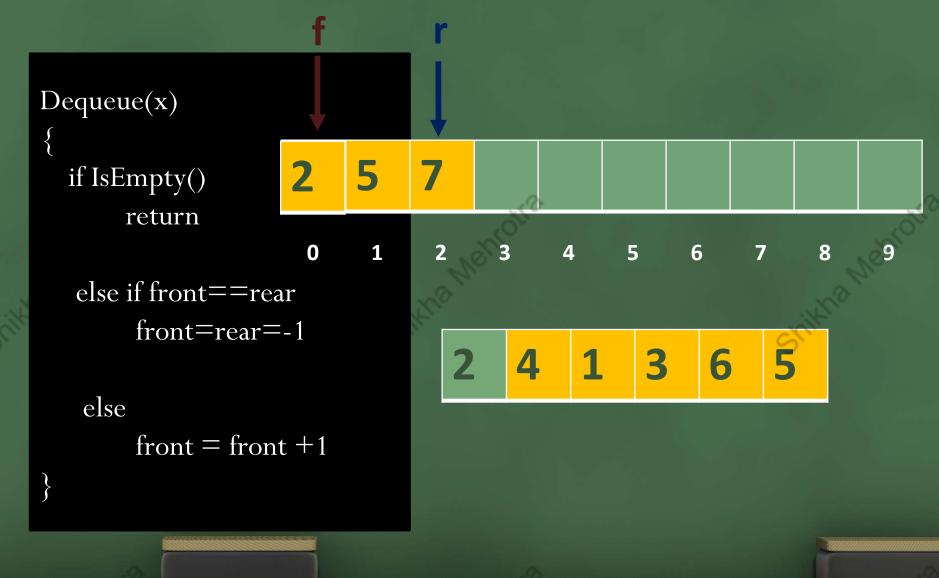
```
IsEmpty()
{
  if front ==-1 && rear ==-1
    return true
  else
    retrun false
}
```

Implementation of Queue





·Implementation of Queue

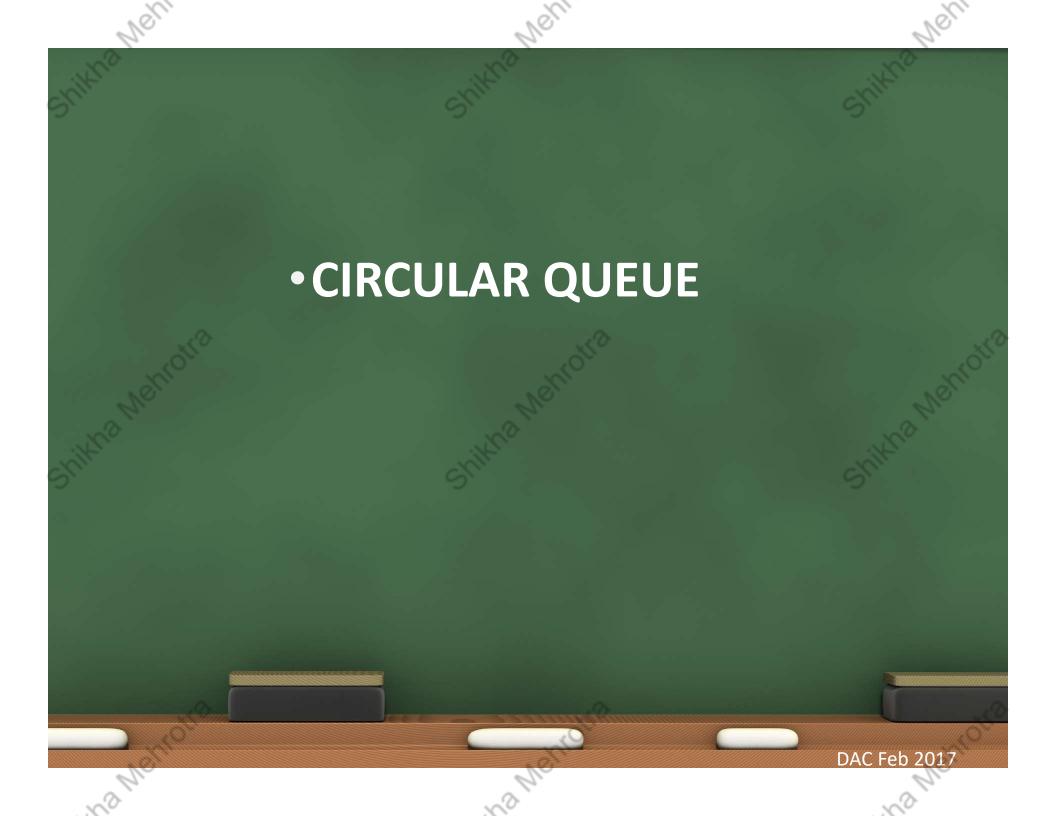


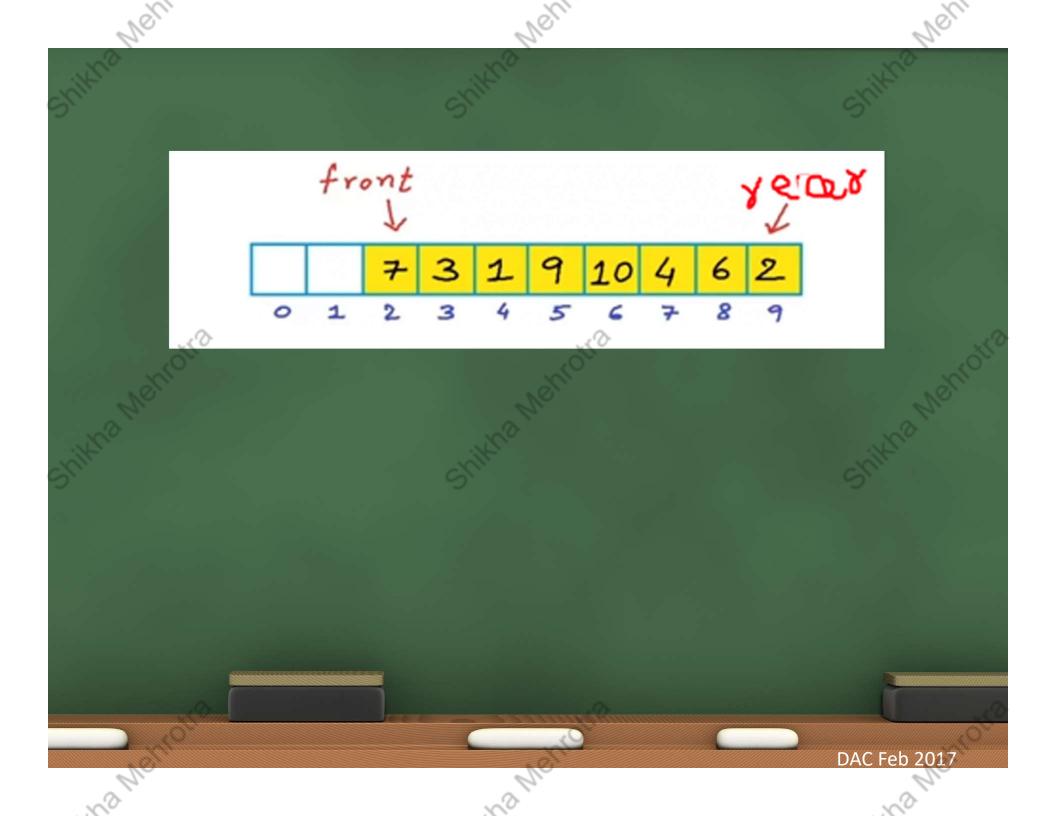
Implementation of Queue

```
Enqueue(x)
  if IsFull()
        return
  else if IsEmpty()
     front=rear=0
                                             3
  else
        rear=rear+1
 a[rear]=x
                               Enqueue(2)
                               Enqueue(5)
Dequeue(x)
                               Enqueue(7)
 if IsEmpty()
     return
                               Dequeue()
  else if front==rear
                               Enqueue(3)
        front=rear=-1
  else
                               Enqueue(1)
        front = front + 1
                                                                      DAC Feb 2017
```

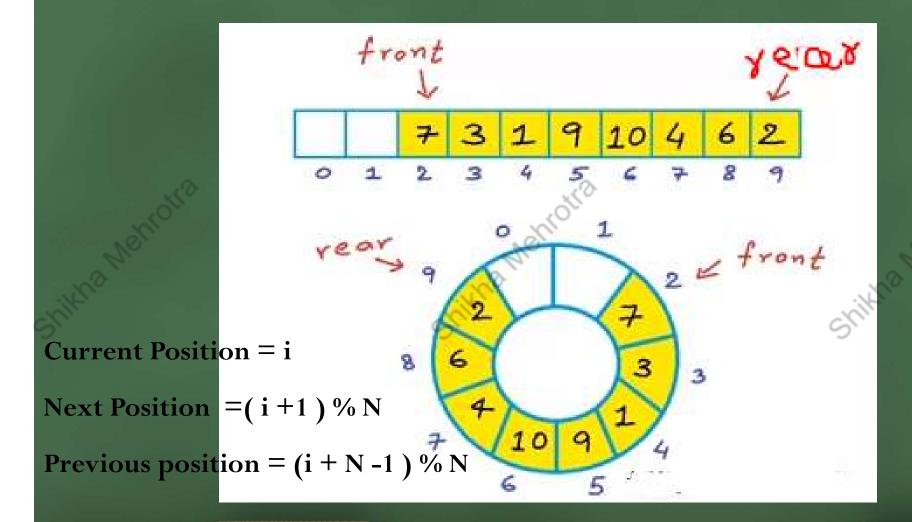
Implementation of Queue

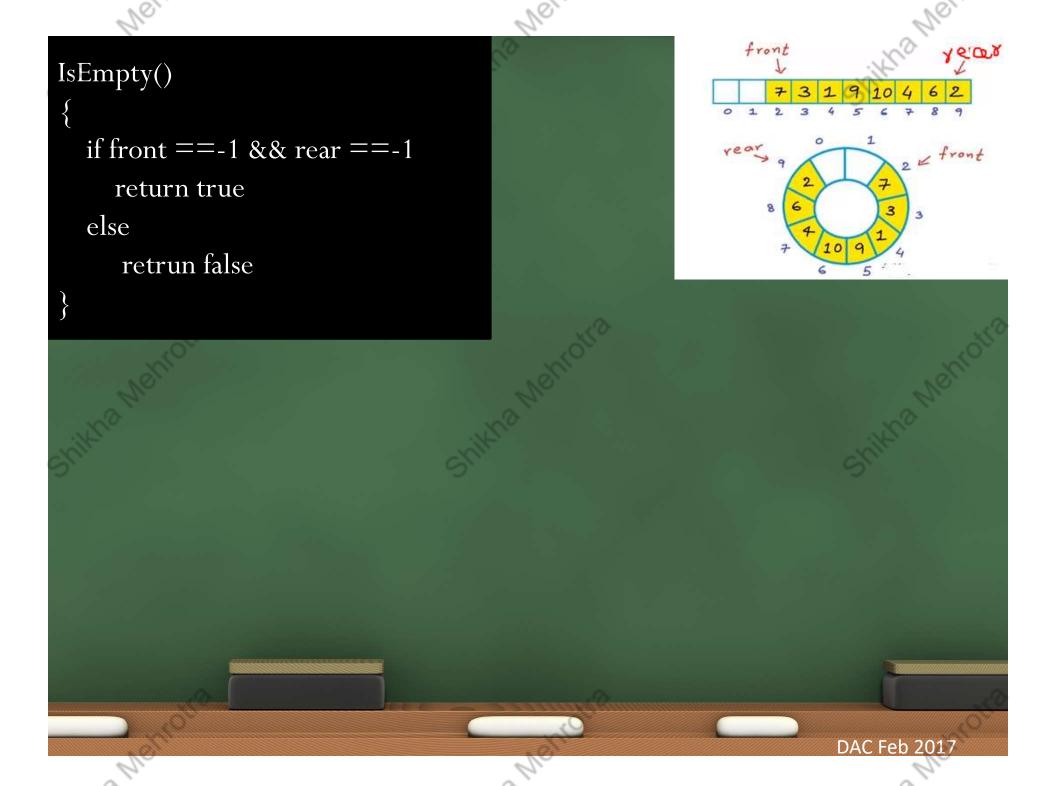
```
Enqueue(x)
 if IsFull()
       return
 else if IsEmpty()
     front=rear=0
                                                      10
                                       3
  else
       rear=rear+1
 a[rear]=x
                                                 Enqueue(9)
                           Enqueue(2)
                           Enqueue(5)
                                                 Enqueue(10)
Dequeue(x)
                                                 Enqueue(4)
                           Enqueue(7)
 if IsEmpty()
    return
                           Dequeue()
                                                 Enqueue(6)
 else if front==rear
                           Enqueue(3)
                                                 Dequeue()
       front=rear=-1
  else
                                                 Enqueue(2)
                           Enqueue(1)
       front = front + 1
                                                              DAC Feb 2017
```



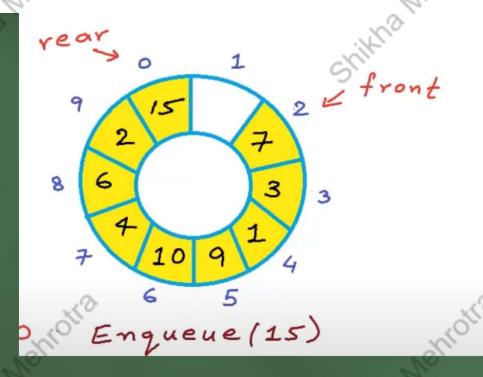


• CIRCULAR QUEUE

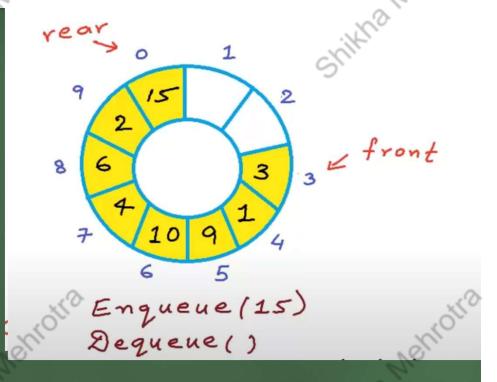




```
Enqueue(x)
 if (rear + 1) \% N == front
  else if IsEmpty()
      front=rear=0
  else
       rear= (rear + 1) \% N
 a[rear]=x
```



```
Dequeue(x)
  if IsEmpty()
      return
  else if front==rear
       front=rear=-1
   else
    front = (front + 1) \% N
```

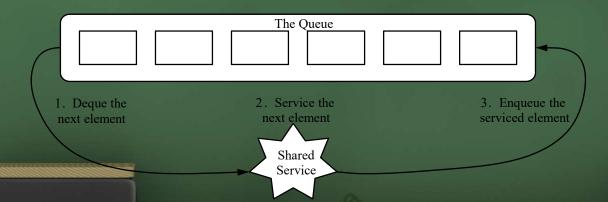


```
#include<stdio.h>
#define n 5
int main()
  int queue[n],ch=1,front=0,rear=0,i,j=1,x=n;
  printf("Queue using Array");
  printf("\n1.Insertion \n2.Deletion \n3.Display
\n4.Exit");
  while(ch)
    printf("\nEnter the Choice:");
    scanf("%d",&ch);
    switch(ch)
    case 1:
      if(rear==x)
         printf("\n Queue is Full");
      else
         printf("\n Enter no %d:",j++);
         scanf("%d",&queue[rear++]);
      break;
```

```
if(front==rear)
         printf("\n Queue is empty");
       else
         printf("\n Deleted Element is %d",queue[front++]);
         X++;
       break;
case 3:
       printf("\nQueue Elements are:\n ");
       if(front==rear)
         printf("\n Queue is Empty");
       else
         for(i=front; i<rear; i++)</pre>
            printf("%d",queue[i]);
           printf("\n");
         break;
       case 4:
         exit(0);
       default:
         printf("Wrong Choice: please see the options");
  return 0;
```

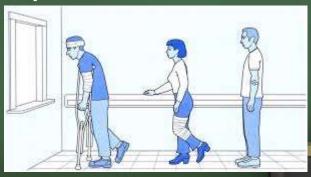
Application: Round RobinSchedulers

- We can implement a round robin scheduler using a queue, Q, by repeatedly performing the following steps:
 - 1. e = Q. dequeue()
 - 2. Service element *e*
 - з. **Q.enqueue(***e*)



Priority Queue

- In Priority queue items are ordered by key value so that item with the lowest value of key is at front and item with the highest value of key is at rear or vice versa.
- So we're assigned priority to item based on its key value.
- Lower the value, higher the priority.



Priority Queue Abstract Data Type

- Two fundamental methods:
 - enqueue
 - dequeue
- Supporting menthods
 - PeekMin
 - removeMin
 - isFull check if queue is full.
 - isEmpty check if queue is empty.

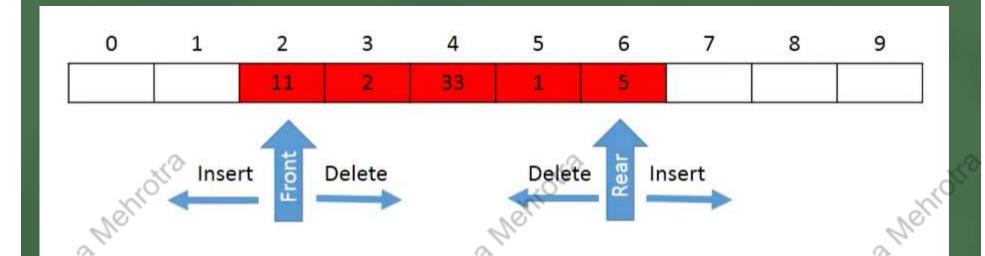
Double Ended Queue (Deque)

Double Ended Queue is also a Queue data structure in which the insertion and deletion operations are performed at both the ends (front and rear).



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Functions of Double Ended Queue (Deque)



Four functions possible:

- Insertion at front
- Deletion at front
- Insertion at Rear
- Deletion at Rear

Deque Abstract Data Type

- Deque: it creates a new deque that is empty. It needs no parameters and returns an empty deque.
- Fundamental Methods:
 - addFront(item) adds a new item to the front of the deque. It needs the item and returns nothing.
 - addRear(item) adds a new item to the rear of the deque. It needs the item and returns nothing.
 - removeFront() removes the front item from the deque. It needs no parameters and returns the item. The deque is modified.
 - removeRear() removes the rear item from the deque. It needs no parameters and returns the item. The deque is modified.
- Supporting Methods
 - isEmpty() tests to see whether the deque is empty. It needs no parameters and returns a boolean value.
 - size() returns the number of items in the deque. It needs no parameters and returns an integer.

Deque Operations

Queue Operation	Queue Contents	Return Value
q.isEmpty()		true
q.enqueue(4)	[4]	
q.enqueue(10)	[10][4]	
q.enqueue(5)	[5][10][4]	
q.size()	[5][10][4]	3
q.isEmpty()	[5][10][4]	false
q.enqueue(9)	[9][5][10][4]	
q.dequeue()	[9][5][10]	
q.dequeue()	[9][5]	
q.size()	[9][5]	2