



## 1.1 Graph

A collection of objects (i.e., nodes) with a set of interactions (i.e., edges)

- Proteins and biological interactions.
- Relationships between club members.

*simple graphs:*

- At most one edge between each pair of nodes
- No edges between a node and itself
- Edges are all undirected



## 1.1.1 Multi-relational graphs

- Undirected, directed, weighted edges
- Edges are of different types

1.  $\mathcal{R}$  is the set of relations
2. One  $\mathbf{A}_\tau$  per edge type  $\tau \in \mathcal{R}$ ,  $(u, \tau, v) \in \mathcal{E}$
3. Adjacency tensor  $\mathcal{A} \in \mathbb{R}^{|\mathcal{V}| \times |\mathcal{R}| \times |\mathcal{V}|}$

- e.g., *Heterogeneous* and *multiplex* graphs

## 1.1.1 Heterogeneous graphs

Nodes are in disjoint sets  $\mathcal{V} = \mathcal{V}_1 \cup \mathcal{V}_2 \cup \dots \mathcal{V}_k$   
where  $\mathcal{V}_i \cap \mathcal{V}_j = \emptyset, \forall i \neq j$ .

- Certain edges only connect nodes of certain types, (i.e.,  $(u, \tau_i, v) \in \mathcal{E} \rightarrow u \in \mathcal{V}_j, v \in \mathcal{V}_k$ ).
- Edges representing "treatments" only occur between drug nodes and disease nodes.
- Special case: *multipartite* graphs, (i.e.,  $(u, \tau_i, v) \in \mathcal{E} \rightarrow u \in \mathcal{V}_j, v \in \mathcal{V}_k \wedge j \neq k$ ).



## 1.1.1 Multiplex graphs

- Graph is decomposed in a set of  $k$  layers, (e.g., transportation network).
- Every node belong to every layer (city).
- Each layer corresponds to a unique relation (different mode of transportation).
- Inter-layer edges connect the same node across layers (possibility of switching modes of transportation within a city).

## 1.1.2 Feature information

Often node-level attributes:  $\mathbf{X} \in \mathbb{R}^{|V| \times m}$   
(e.g., a profile picture associated with a user in a social network)



# Graph or network?

The difference is a historical one.

- "Graph": machine learning community
- "Network": data mining and network science communities

In this book

- *Graph* for abstract data structure
- *Network* for specific, real-world instantiations of this data structure



## 1.2.1 Node classification

Nodes in a graph are not i.i.d.

- *Homophily*: adjacent nodes share attributes
- *Structural equivalence*: nodes with similar local neighborhood structures have similar labels
- *Heterophily*: nodes tend to be connected to nodes with different labels



# Supervised or semi-supervised?

## **Semi-supervised:**

- Have access to the full graph during training including unlabeled test nodes
- Can use neighborhood information about test nodes to improve model during training

## **Supervised:**

- Unlabeled nodes are completely unobserved during training



## 1.2.2 Relation prediction

Given nodes  $\mathcal{V}$  and an incomplete set of edges  $\mathcal{E}_{\text{train}} \in \mathcal{E}$ , infer missing edges.

1. Also known as link prediction, graph completion, and relational inference
2. Real-world application: recommendation system, predicting drug side-effects ...

## 1.2.3 Community detection

The graph analogue of unsupervised clustering

## 1.2.4 Graph classification, regression, and clustering

Graph are an i.i.d. data with labels.

- Challenge: define useful features upon relational structure
- E.g., predict molecule's toxicity or solubility given their structure



## 2.1 Graph statistics and kernel methods

Traditional machine learning: use node-level features based on heuristic functions or domain knowledge as input to a classifier.

- Popular prior to the advent of deep learning

## 2.1.1 Node-level statistics and features

- Node degree:  $d_u = \sum_{v \in \mathcal{V}} \mathbf{A}[u, v]$
- Node centrality

**Eigenvector centrality** for  $\forall u \in \mathcal{V}$   
 $e_u = \frac{1}{\lambda} \sum_{v \in \mathcal{V}} \mathbf{A}[u, v] e_v$ , constant  $\lambda$ .

Then  $\mathbf{e}$  is the eigenvector of  $\mathbf{A}$ , since this equation in vector notation is  $\lambda \mathbf{e} = \mathbf{A} \mathbf{e}$

Assume that centrality values are positive, by Perron-Frobenius Theorem, such  $\mathbf{e}$  corresponds to the largest eigenvalue of  $\mathbf{A}$ .

# Convergence of power iteration

For  $\mathbf{A} \in \mathbb{R}^{n \times n}$  with a dominant eigenvalue, there exists a nonzero vector  $\mathbf{x}$  such that the sequence  $\mathbf{Ax}$ ,  $\mathbf{A}^2\mathbf{x}$ , ...,  $\mathbf{A}^k\mathbf{x}$ , ... approach a multiple of  $\mathbf{e}$ .

$\mathbf{A}$  has  $n$  linearly independent eigenvectors  $\mathbf{x}_1, \dots, \mathbf{x}_n$  with  $\lambda_1 \leq \dots \leq \lambda_n$ . Define a  $\mathbf{x}$  with  $c_n \neq 0$  as  $\mathbf{x} = \sum_{i=1}^n c_i \mathbf{x}_i$ .  $\mathbf{A}^k \mathbf{x} = \sum_{i=1}^n c_i \lambda_i^k \mathbf{x}_i = \lambda_n^k \sum_{i=1}^n c_i \left( \frac{\lambda_i}{\lambda_n} \right)^k \mathbf{x}_i$

Since  $\frac{\lambda_i}{\lambda_n} < 1$  for  $i = 1:n-1$ ,  $\lim_{k \rightarrow \infty} \mathbf{A}^k \mathbf{x} = \lambda_n^k c_n \mathbf{x}_n$ .

■ Start off with  $[1, \dots, 1]^T$  to compute  $\mathbf{e}$



## 2.1.1 Node-level statistics and features

### **Degree centrality**

- All nodes are equivalent

### **Eigenvector centrality**

- A node is important if it is linked to by other important nodes
- Low degree node also becomes important by linking to a hub node

High degree  $\neq$  High eigenvector centrality

## 2.1.1 Node-level statistics and features

- Clustering coefficient: how tightly clustered a node neighborhood is

**local variant** for  $\forall u \in \mathcal{V}$

$$c_u = |(\nu_1, \nu_2) \in \mathcal{E} : \nu_1, \nu_2 \in \mathcal{N}(u)| / d_u^2$$

$c_u = 1$ : all of  $u$ 's neighbors are also neighbors of each other



## 2.1.2 Graph-level features

- Bag of nodes
- The Weisfieler-Lehman kernel
- Count the occurrence of different small subgraph structures (graphlets)
- Path-based methods: running walks to count occurrence of different degree sequences
  - e.g., random walk





## 2.2 Neighborhood overlap detection

Quantify the extent to which a pair of nodes are related, e.g., the simplest neighborhood overlap measure

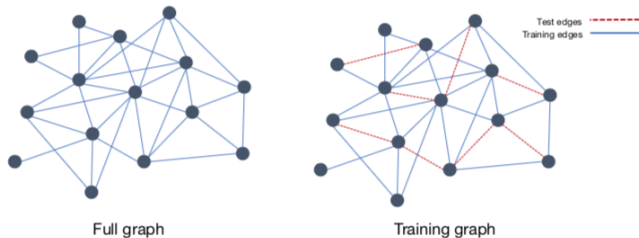
The similarity matrix  $\mathbf{S} \in \mathbb{R}^{|\mathcal{V}| \times |\mathcal{V}|}$

$$\mathbf{S}[u, v] = |\mathcal{N}(u) \cap \mathcal{N}(v)|$$

$$P(\mathbf{A}[u, v] = 1) \propto \mathbf{S}[u, v]$$

■ Relation prediction: set a threshold

## 2.2.1 Local overlap measures



Design functions of  $|\mathcal{N}(u) \cap \mathcal{N}(v)|$  to avoid bias towards predicting edges for nodes with large degrees

- $\mathbf{S}_{\text{Sorenson}}[u, v] = 2|\mathcal{N}(u) \cap \mathcal{N}(v)| / (d_u + d_v)$
- $\mathbf{S}_{\text{Jaccard}}[u, v] = |\mathcal{N}(u) \cap \mathcal{N}(v)| / |\mathcal{N}(u) \cup \mathcal{N}(v)|$

## 2.2.1 Local overlap measures

Consider the importance of common neighbors

$$\begin{aligned} \blacksquare \mathbf{S}_{\text{Resource Allocation}}[v_1, v_2] &= \sum_{u \in n(v_1) \cap n(v_2)} \frac{1}{d_u} \\ \blacksquare \mathbf{S}_{\text{Adamic-Adar}}[v_1, v_2] &= \sum_{u \in n(v_1) \cap n(v_2)} \frac{1}{\log(d_u)} \end{aligned}$$

- A shared low-degree neighbor is more informative than a shared high-degree one.

## 2.2.2 Global overlap measures

Two nodes could have no local overlap in their neighborhoods but still be members of the same community.

$$\textbf{Katz index } \mathbf{S}_{\text{Katz}}[u, v] = \sum_{i=1}^{\infty} \beta^i \mathbf{A}^i[u, v], \beta \in \mathbb{R}^+$$

- small  $\beta < 1$  down-weight the importance of long paths.
- If  $\lambda_{\max}$  of  $\mathbf{A}$  is  $< \frac{1}{\beta}$  and  $(\mathbf{I} - \beta\mathbf{A})$  is non-singular,  $\mathbf{S}_{\text{Katz}} = (\mathbf{I} - \beta\mathbf{A})^{-1} - \mathbf{I}$

## 2.2.2 Global overlap measures

- High-degree nodes are in more paths
- Katz index is biased by node degree
- $\Rightarrow$  LHN similarity considers  $\frac{\mathbf{A}^i}{E\{\mathbf{A}^i\}}$

### Configuration model

Draw a random graph with the same set of degrees as the given graph

- There are  $d_u$  edges leaving  $u$
- Each has a  $\frac{d_v}{2|\mathcal{E}|}$  chance of ending at  $v$
- $E\{\mathbf{A}[u, v]\} = \frac{d_u d_v}{2|\mathcal{E}|}$ .

## 2.2.2 Global overlap measures

For  $E\{\mathbf{A}^2[v_1, v_2]\}$ , consider path of length 2 with any intermediate vertex  $u$

$$\blacksquare E\{\mathbf{A}[v_1, u]\} = \frac{d_{v_1} d_u}{2|\mathcal{E}|}, \quad E\{\mathbf{A}[u, v_2]\} = \frac{d_{v_2}(d_u - 1)}{2|\mathcal{E}|}$$

$$\blacksquare \Rightarrow E\{\mathbf{A}^2[v_1, v_2]\} = \frac{d_{v_1} d_{v_2}}{(2|\mathcal{E}|)^2} \sum_{u \in \mathcal{V}} (d_u - 1) d_u$$

The analytical computation of expected node path becomes intractable as paths length  $> 3$ .

- Use  $\lambda_{\max}$  to approximate the growth in the number of paths

The expected number of paths of length  $k$  from  $i$  to  $j$  is the  $j^{\text{th}}$  element of  $\mathbf{p}_k$  where  $\mathbf{p}_k = \mathbf{A}^k \mathbf{e}_i$ . As  $k \rightarrow \infty$ ,  $\mathbf{p}_{k+1} = \lambda_{\max} \mathbf{p}_k$ .

The number of paths from  $i$  to  $j$  increase by a factor of  $\lambda_{\max}$  each time we add one extra step to the path length when  $k$  is large.

## 2.2.2 Global overlap measures

With approximation  $E\{\mathbf{A}^i[u, v]\} = \frac{d_u d_v \lambda_{\max}^{i-1}}{2|\mathcal{E}|}$ , the LNH index is defined as

$$\mathbf{S}_{\text{LNH}}[u, v] = \mathbf{I}[u, v] + \frac{2|\mathcal{E}|}{d_u d_v} \sum_{i=1}^{\infty} \beta^i \lambda_{\max}^{1-i} \mathbf{A}^i[u, v], \quad \beta \in \mathbb{R}^+$$

- small  $\beta < 1$  also down-weight long paths
- it only gives a high similarity score if two nodes occur on more paths than we expect



## 2.2.2 Global overlap measures

**Random walk methods:** considers random walks rather than exact counts of paths.

The walker starting from node  $u$  will move to a random neighbor with probability  $c$  and return to itself with probability  $1 - c$ .

Denote  $q_{uv}$  the probability this random walker locates at node  $v$  in the steady state, we have  $\mathbf{q}_u = c\mathbf{P}\mathbf{q}_u + (1 - c)\mathbf{e}_u$ , and  $\mathbf{q}_u = (1 - c)(\mathbf{I} - c\mathbf{P})^{-1}\mathbf{e}_u$

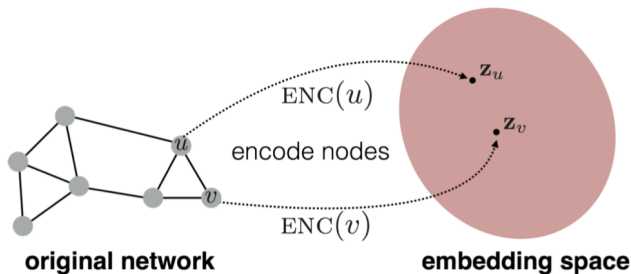
**Random walk similarity index**

$$S_{\text{RW}}[u, v] = \mathbf{q}_u[v] + \mathbf{q}_v[u]$$

# Introduction

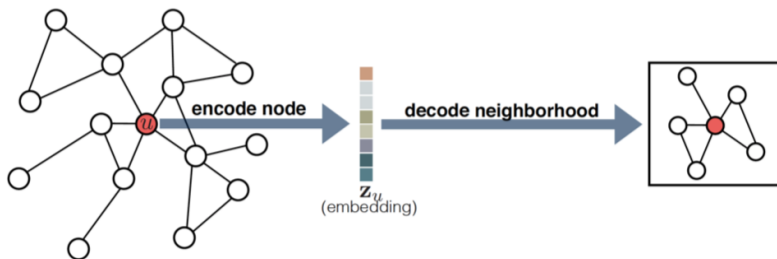
## Node embedding

- Project nodes into a latent space
- Geometric relations in this latent space correspond to edges
- No node feature or extra information is used



## 3.1 An encoder-decoder perspective

- **Encoder** model: maps each node into a low-dimensional embeddings
- **Decoder** model: use low-dimensional embeddings to reconstruct information about node's neighborhood





### 3.1.1 The encoder (ENC)

**Encoder**  $\text{ENC} : \mathcal{V} \rightarrow \mathbb{R}^d$

$\mathbf{z}_v \in \mathbb{R}^d$  is the embedding for node  $v$

Simplest encoding approach: encoder is just an embedding-lookup

**Shallow encoding**

$$\text{ENC}(v) = \mathbf{Z}[v], \quad \mathbf{Z} \in \mathbb{R}^{|\mathcal{V}| \times d}$$



### 3.1.2 The decoder (DEC)

**Pairwise decoder**  $\text{DEC} : \mathbb{R}^d \times \mathbb{R}^d \rightarrow \mathbb{R}^+$

predict similarity of pairs of nodes

**Reconstruction**  $\text{DEC}(\text{ENC}(u), \text{ENC}(v)) = \text{DEC}(\mathbf{z}_u, \mathbf{z}_v)$

**Similarity measure**  $\mathbf{S} \in \mathbb{R}^{|\mathcal{V}| \times |\mathcal{V}|}$ , e.g.,  $\mathbf{S} = \mathbf{A}$ ,  
neighborhood overlap

**Goal:** minimize reconstruction loss so  
that  $\text{DEC}(\mathbf{z}_u, \mathbf{z}_v) \approx \mathbf{S}[u, v]$



### 3.1.3 Optimizing the model

Given ENC, DEC and  $\mathbf{S}$ , the empirical reconstruction loss is defined as:

$$\mathcal{L} = \sum_{(u,v) \in \mathcal{D}} \ell(\text{DEC}(\mathbf{z}_u, \mathbf{z}_v), \mathbf{S}[u, v])$$

where  $\mathcal{D}$  is a set of training node pairs, and

$$\ell : \mathbb{R} \times \mathbb{R} \rightarrow \mathbb{R}$$

e.g., mean-squared error, classification loss

## 3.2 Factorization-based approaches

### Laplacian eigenmaps

- $\text{DEC}(\mathbf{z}_u, \mathbf{z}_v) = \|\mathbf{z}_u - \mathbf{z}_v\|_2^2$
- $\mathcal{L} = \sum_{(u,v) \in \mathcal{D}} \text{DEC}(\mathbf{z}_u, \mathbf{z}_v) \cdot \mathbf{S}[u, v]$

### Intuition

- Penalize the model when very similar nodes have embeddings that are far apart

## 3.2 Factorization-based approaches

### Inner-product methods

- $\text{DEC}(\mathbf{z}_u, \mathbf{z}_v) = \mathbf{z}_u^\top \mathbf{z}_v$
- $\mathcal{L} = \sum_{(u,v) \in \mathcal{D}} \|\text{DEC}(\mathbf{z}_u, \mathbf{z}_v) - \mathbf{S}[u, v]\|_2^2$  by  
stacking embedding into  $\mathbf{Z} \in \mathbb{R}^{|\mathcal{V}| \times d}$

### Intuition

- Often define  $\mathbf{S}$  as some polynomial function of the adjacency matrix
- Embeddings are optimized with  $\mathbf{z}_u^\top \mathbf{z}_v \approx \mathbf{S}[u, v]$





## 3.3 Random walk embeddings

**Idea:** Node embeddings are optimized so that two nodes with similar embeddings if they tend to co-occur on short random walks over the graph

## 3.3 Random walk embeddings

### DeepWalk and node2vec:

- We aim to find a shallow embedding, where  $\mathbf{z}_u$  is the embedding of node  $u$
- $N_R(u)$ , the neighborhood of node  $u$  is obtained by some random walk
- Denote  $P(v|\mathbf{z}_u)$  as the predicted probability of visiting node  $v$  on the random walk starting from node  $u$
- $\text{DEC}(\mathbf{z}_u, \mathbf{z}_v) = P(v|\mathbf{z}_u) = \exp(\mathbf{z}_u^T \mathbf{z}_v) / \sum_{k \in \mathcal{V}} \exp(\mathbf{z}_u^T \mathbf{z}_k)$
- $\mathcal{L} = -\sum_{u \in \mathcal{V}} \sum_{v \in N_R(u)} \log(P(v|\mathbf{z}_u))$
- Find embeddings that minimize  $\mathcal{L}$



## 3.3 Random walk embeddings

**Large-scale information network embeddings:**

- $\text{DEC}(\mathbf{z}_u, \mathbf{z}_v) = 1/(1 + \exp(-\mathbf{z}_u^\top \mathbf{z}_v))$
- Use adjacency matrix to measure similarity



## 3.4 Limitations of shallow embeddings

- No shared parameters between nodes in the encoder
- Do not leverage node features in the encoder
- Can only generate embeddings for nodes that were present during training - need additional optimizations for new nodes



# Knowledge graph completion

## **Knowledge graph** $\mathcal{G} = (\mathcal{V}, \mathcal{E})$

Edge  $(u, \tau, v)$  indicates the presence of a particular relation  $\tau$  between two nodes.

- E.g.,  $(u, \text{TREATS}, v)$  in a biomedical knowledge graph.
- Relation prediction.



## 4.1 Reconstructing multi-relational data

Given  $\mathbf{z}_u$  and  $\mathbf{z}_v$

- Predict the existence of an edge between node  $u$  and  $v$
- Determine the type of the predicted edge

### Multi-relational decoder

- $\text{DEC} : \mathbb{R}^d \times \mathcal{R} \times \mathbb{R}^d \rightarrow \mathbb{R}^+$
- interpret  $\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v)$  as the likelihood that the edge  $(u, \tau, v)$  exists



## Example: RESCAL

- $\mathbf{R}_\tau \in \mathbb{R}^{d \times d}$  is specific to relation  $\tau \in \mathcal{R}$
- $\mathbf{R}_\tau$  is learnable
- $\text{DEC}(u, \tau, v) = \mathbf{z}_u^\top \mathbf{R}_\tau \mathbf{z}_v$
- $\mathcal{A} \in \mathbb{R}^{|\mathcal{V}| \times |\mathcal{R}| \times |\mathcal{V}|}$
- $\mathcal{L} = \sum_{u \in \mathcal{V}} \sum_{v \in \mathcal{V}} \sum_{\tau \in \mathcal{R}} \|\mathbf{z}_u^\top \mathbf{R}_\tau \mathbf{z}_v - \mathcal{A}[u, \tau, v]\|^2$



# Loss functions, decoders, and similarity functions

Relation prediction reconstruct immediate neighbors

- Have various decoders and loss functions
- Define the similarity measure directly based on the adjacency tensor





## 4.2 Loss functions

- 1 Ideally, the computation of the loss function is  $O(|\mathcal{E}|)$  since many multi-relational graphs are sparse ( $|\mathcal{E}| \ll |\mathcal{V}|^2 |\mathcal{R}|$ ).
- 2 Use a loss for classification on edges to decode the adjacency tensor from the node embeddings.

# Cross-entropy with negative sampling

$$\mathcal{L} = \sum_{(u, \tau, v) \in \mathcal{E}} -\log(\sigma(\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v))) - \gamma \mathbb{E}_{v_n \sim P_{n,u}(\mathcal{V})} [\log(\sigma(-\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_{v_n})))]$$

- $\sigma$  denotes the logistic function
- $P_{n,u}(\mathcal{V})$  is a "negative sampling" distribution over the set of node  $\mathcal{V}$
- $P_{n,u}(\mathcal{V})$  might depend on  $u$
- $\gamma > 0$  is a hyperparameter
- $\log(\sigma(\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v)))$  is the log-likelihood for existence of edge  $(u, \tau, v)$
- $\mathbb{E}_{v_n \sim P_{n,u}(\mathcal{V})} [\log(\sigma(\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_{v_n})))]$  is the expected log-likelihood for non-existence of an edge
- Usually, we consider  $v_n \in \mathcal{P}_{n,v}$ , where  $\mathcal{P}_{n,v}$  is a small set sampled from  $P_{n,u}(\mathcal{V})$

# Max-margin loss

$$\mathcal{L} = \sum_{(u, \tau, v) \in \mathcal{E}} \sum_{v_n \in \mathcal{P}_{n, u}} \max(0, -\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v) + \text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_{v_n}) + \Delta).$$

**Contrastive estimation:** compare the decoded score for a true pair with a negative sample

The loss will equal 0 if the difference in scores  $\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v)$  and  $\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_{v_n})$  is at least  $\Delta$



## 4.3 Multi-relational decoders

The decoder for RESCAL  $\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v) = \mathbf{z}_u^T \mathbf{R}_\tau \mathbf{z}_v$  has  $O(d^2)$  parameters for each relation type.

- Ideally use  $O(d)$  parameters

# Translational decoders

$$\blacksquare \mathbf{r}_\tau \in \mathbb{R}^d$$

**TransE**

$$\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v) = -\|\mathbf{z}_u + \mathbf{r}_\tau - \mathbf{z}_v\|$$

- Represents relations as translations in the embedding space

**TransX**

$$\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v) = -\|g_{1,\tau}(\mathbf{z}_u) + \mathbf{r}_\tau - g_{2,\tau}(\mathbf{z}_v)\|$$

**TransH**

$$\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v) = -\|(\mathbf{z}_u - \mathbf{w}_r^T \mathbf{z}_u \mathbf{w}_r) + \mathbf{r}_\tau - (\mathbf{z}_v - \mathbf{w}_r^T \mathbf{z}_v \mathbf{w}_r)\|$$

# Multi-linear dot products

**DistMult**

$$\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v) = \langle \mathbf{z}_u, \tau, \mathbf{z}_v \rangle = \sum_{i=1}^d \mathbf{z}_u[i] \times \mathbf{r}_\tau[i] \times \mathbf{z}_v[i]$$

- Take a straightforward generalization of the dot product
- Can only encode symmetric relations

$$\begin{aligned}\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v) &= \langle \mathbf{z}_u, \mathbf{r}_\tau, \mathbf{z}_v \rangle \\ &= \langle \mathbf{z}_v, \mathbf{r}_\tau, \mathbf{z}_u \rangle \\ &= \text{DEC}(\mathbf{z}_v, \tau, \mathbf{z}_u)\end{aligned}$$

# Complex decoders

Many relation types in multi-relational graphs are directed and asymmetric

## Complex

$$\begin{aligned}\text{DEC}(\mathbf{z}_u, \tau, \mathbf{z}_v) &= \text{Re}(\langle \mathbf{z}_u, \tau, \bar{\mathbf{z}}_v \rangle) \\ &= \text{Re}\left(\sum_{i=1}^d \mathbf{z}_u[i] \times \mathbf{r}_\tau[i] \times \bar{\mathbf{z}}_v[i]\right)\end{aligned}$$

- $\mathbf{z}_u, \mathbf{z}_v, \mathbf{r}_\tau \in \mathbb{C}^d$  are complex valued embedding
- $\text{Re}$  takes the real component
- $\bar{\mathbf{z}}_v$  is the complex conjugate



## 4.3.1 Representational abilities

**Symmetry**  $(u, \tau, v) \in \mathcal{E} \Leftrightarrow (v, \tau, u) \in \mathcal{E}$

**Anti-symmetry**  $(u, \tau, v) \in \mathcal{E} \Rightarrow (v, \tau, u) \notin \mathcal{E}$

**Inversion**  $(u, \tau_1, v) \in \mathcal{E} \Leftrightarrow (v, \tau_2, u) \in \mathcal{E}$

**Compositionality**

$(u, \tau_1, y) \in \mathcal{E} \wedge (y, \tau_2, v) \in \mathcal{E} \Rightarrow (u, \tau_3, v) \in \mathcal{E}$