Homework 1

Due Date: Sunday, April 08, 2018

1. **(3 pts) For each of the following pairs of functions, either f(n) is O(g(n)), f(n) is Ω(g(n)), or f(n) is Θ (g(n)). Determine which relationship is correct and explain.**

**Answers:**

|  |  |  |  |  |
| --- | --- | --- | --- | --- |
| Problem |  |  | Relationship | Explanation |
| a. |  |  | is  BIG OH |  |
| b. |  |  | is  THETA |  |
| c. |  |  | is  BIG OH |  |
| d. |  |  | is  BIG OH |  |
| e. |  |  | is  THETA |  |
| f. |  |  | is  BIG OH |  |

1. **(3 pts) Use mathematical induction to show that when n is an exact power of 2, the solution of the recurrence**

**is**

**Answer:**

**Base Case:**

If ,

Then and both must equal 2 as well

Therefore,

**Inductive Hypothesis:**

If for integers where then we can assume that is true.

**Apply the Axiom of Induction:**

If , then

Now substitute in for ,

Thus, proven through mathematical induction as for both for and

1. **(4 pts) Let and be asymptotically positive non-decreasing functions. Prove or disprove each of the following conjectures. To disprove give counter example.**
   1. **If and then**
   2. **If and then**

**Answers:**

1. Disproved If and then through counter example.

Therefore,

Thus, disproved through counter example

1. Prove If and then

If

Then constants and also exist such that…

**Step 1:** for all and since

Therefore constants and also exist such that…

**Step 2:**  for all

Let

**Step 3:**  for

Combining the two items from step 3

for Therefore,

and

or

1. **(10 pts) Merge Sort and Insertion Sort Programs  
   Implement merge sort and insertion sort to sort an array/vector of integers. You may implement the algorithms in the language of your choice, name one program “mergesort” and the other “insertsort”. Your programs should be able to read inputs from a file called “data.txt” where the first value of each line is the number of integers that need to be sorted, followed by the integers.  
     
   Example values for data.txt:**

**4 19 2 5 11**

**8 1 2 3 4 5 6 1 2**

**The output will be written to files called “merge.out” and “insert.out”.**

**For the above examples the output would be:**

**2 5 11 19**

**1 1 2 2 3 4 5 6**

***Submit a copy of all your code files and a README file that explains how to compile and run your code in a ZIP file to TEACH. We will test execution with an input file named data.txt.***

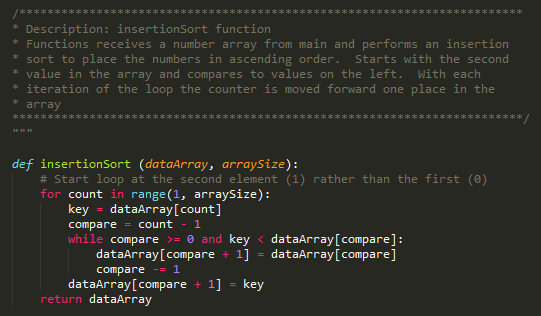
**Answer:**

My implementations for Merge Sort and Insertion Sort were submitted to TEACH in a zipped folder under the filenames mergesort.py and insertsort.py, respectively. They take arrays from a file called data.txt, sort them, and write the sorted results to files called merge.out and insert.out. I used python as my development language, it executes slower than C/C++ but abstracts away a lot of the lines of code necessary for implementing these algorithms. The algorithms were tested on my desktop and the university FLIP servers and confirmed to work correctly. Below are screen shots from my code of the functions that perform each version of sort. Please see my zip file TEACH submission for full versions of the code.

**Merge Sort Algorithm (Python)**



**Insertion Sort Algorithm (Python)**



1. **(10 pts) Merge Sort vs Insertion Sort Running time analysis  
     
   The goal of this problem is to compare the experimental running times of the two sorting algorithms.** 
   1. **Now that you have proven that your code runs correctly using the data.txt input file, you can modify the code to collect running time data. Instead of reading arrays from a file to sort, you will now generate arrays of size n containing random integer values from 0 to 10,000 and then time how long it takes to sort the arrays. We will not be executing the code that generates the running time data so it does not have to be submitted to TEACH or even execute on flip. Include a “text” copy of the modified code in the written HW submitted in Canvas.**

**Answer:**

**Insertion Sort Timing Code (Screen Shot)**



**Text Copy of Insertion Sort Time Code**

import time # Import the time library to measure time

import random # Import the random number generator library

"""

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\* Description: insertionSort function

\* Functions receives a number array from main and performs an insertion

\* sort to place the numbers in ascending order. Starts with the second

\* value in the array and compares to values on the left. With each

\* iteration of the loop the counter is moved forward one place in the

\* array

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"""

def insertionSort (dataArray, arraySize):

# Start loop at the second element (1) rather than the first (0)

for count in range(1, arraySize): # Increment through the array one element at a time using the array's length

key = dataArray[count] # Set the current index to the loop counter (counter starts at 1)

compare = count - 1 # Adjusted the comparative number one element left in the array

while compare >= 0 and key < dataArray[compare]: # While compare number is in range and key is less than the compare number

dataArray[compare + 1] = dataArray[compare] # Swap the compare number forward (right) in the array

compare -= 1 # Move comparitive number to the next element in the array (going left toward 0)

dataArray[compare + 1] = key # No swap occurs and the key remains in place

return dataArray # Return the sorted Array to main function

"""

/\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*

\* Description: main function

\* Program starts here and calls functions as necessary

\* Calls the mergesort function that completes most of the program's

\* functionality.

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"""

if \_\_name\_\_ == "\_\_main\_\_":

n = 1000 # Set number of variables that need to be generated

for count in range(0, 13):

starttime = time.time() # Collect the start time at the initial execution of the program

if count == 1:

n = 2000

if count == 2:

n = 5000

if count == 3:

n = 10000

unsortedArray = [] # Establish a new array to store the random integers

for x in range (n): # Run through a loop for the desired number of random elements

unsortedArray.append(random.randint (0, 10000)) # Append the random elements to the unsorted Array

insertionSort(unsortedArray, n) # Send the array to the sort function to be sorted

print("Array Size: {}".format(n)) # Output the Array size to the screen

print("Time to Complete: {} seconds".format(time.time() - starttime)) # Output the time to complete the program.

n += 10000

**Merge Sort Timing Code (Screen Shot)**



**Text Copy of Merge Sort Time Code**

import time # Import the time library to measure time

import random # Import the random number generator library

"""

/\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*

\* Description: mergeSort function

\* Functions receives a number array from main and performs a merge sort

\* to place the numbers in ascending order. The function takes and array

\* and recursively devides until until it reaches array sizes of one.

\* It passes data to a merge function that combines the elements in

\* ascending order into the original array

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"""

def mergeSort (dataArray, arraySize):

if arraySize < 1: # If Array Size is less than 1, this is invalid

print("Array Size {} is invalid".format(arraySize)) # print an error to the console

elif arraySize == 1: # If the array size is one it's already sorted

return dataArray # Return the array as is without performing any work

else: # If the array is greater than size one (less than one and already sorted)

first = dataArray[:(arraySize // 2)] # Split the array and assign the left (bottom) half to variable (Use Floor Division)

last = dataArray[(arraySize // 2):] # Split the array and assign the right (top) half to variable (Use Floor Division)

firstLength = len(first) # Get left array length and assign to variable

lastLength = len(last) # Get left array length and assign to variable

firstArray = mergeSort(first, firstLength) # Recursively call the mergeSort function and send the first half of the array

secondArray =mergeSort(last, lastLength) # Recursively call the mergeSort function and send the second half of the array

return merge (dataArray, firstArray, secondArray)

"""

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\* Description: merge function

\* Called by the mergeSort function to evaluate and merge the split array

\* elements back into a single sorted array. Receives the original array

\* the split right and left arrays and their respective lengths. Returns

\* a sorted integer array to mergerSort

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"""

def merge (dataArray, first, second):

firstcount = secondcount = originalcount = 0 # create three counter variables and initialize to zero

while firstcount < len(first) and secondcount < len(second): # while both the left and right counters are less than their array lengths

if first[firstcount] < second[secondcount]: # If the left array is less than the right

dataArray[originalcount]=first[firstcount] # Assign it to the current element in the original array

firstcount += 1 # Increment left array counter by one (hate that python can't just do ++)

else: # otherwise if the right array element is smaller

dataArray[originalcount]=second[secondcount] # Assign it to the current element in the original array

secondcount += 1 # Increment Right array counter by one

originalcount += 1 # Increment original array counter by one

while firstcount < len(first): # If the first count is still less than the left array length

dataArray[originalcount]=first[firstcount] # assign the current element from the left array to the current element in the original array

firstcount += 1 # Increment counter by one

originalcount += 1 # Increment counter by one

while secondcount < len(second): # While the left count is still less than the right array length

dataArray[originalcount]=second[secondcount] # Assign it's value to the current element in the original array

secondcount += 1 # Increment counter by one

originalcount += 1 # Increment counter by one

return dataArray # Return the sorted array to the main function for writing.

"""

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\* Description: main function

\* Program starts here and calls functions as necessary

\* Calls the mergesort function that completes most of the program's

\* functionality.

\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*\*/

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if \_\_name\_\_ == "\_\_main\_\_":

n = 1000 # Set number of variables that need to be generated

for count in range(0, 13):

starttime = time.time() # Collect the start time at the initial execution of the program

if count == 1:

n = 2000

if count == 2:

n = 5000

if count == 3:

n = 10000

unsortedArray = [] # Establish a new array to store the random integers

for x in range (n): # Run through a loop for the desired number of random elements

unsortedArray.append(random.randint (0, 10000)) # Append the random elements to the unsorted Array

mergeSort(unsortedArray, n) # Send the array to the sort function to be sorted

print("Array Size: {}".format(n)) # Output the Array size to the screen

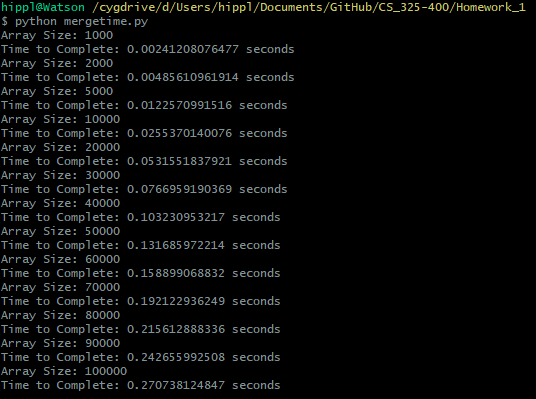
print("Time to Complete: {} seconds".format(time.time() - starttime)) # Output the time to complete the program.

n += 10000

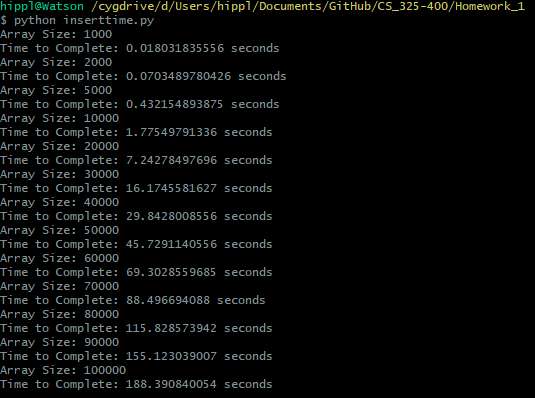
* 1. **Use the system clock to record the running times of each algorithm for n = 1000, 2000, 5000, 10,000, … You may need to modify the values of n if an algorithm runs too fast or too slow to collect the running time data. If you program in C your algorithm will run faster than if you use python. You will need at least seven values of t (time) greater than 0. If there is a variability in the times between runs of the same algorithm you may want to take the average time of several runs for each value of n.**

**Answer:**

**Merge Sort Average Run Times**



**Insertion Sort Average Run Times**



* 1. **For each algorithm plot the running time data you collected on an individual graph with n on the x-axis and time on the y-axis. You may use Excel, Matlab, R or any other software. Also plot the data from both algorithms together on a combined graph. Which graphs represent the data best?**

**Answer:**

Below are the individual runtime graphs generated from the runtime output on varying sized arrays with randomized integer values in both insertion sort and merge sort algorithms (actual results for each array size are presented as screenshots in part b). The charts were created in Excel rather than Matlab. I found that the Scatter Plot graph in Excel best represented the runtime data. I also felt that the graphs were best represented individually rather than in combination. The scales were far too extreme between the merge sort and the insertion sort algorithms. When the two graphs are presented together in comparison, the scaling is relatively poor (even the professor mentioned in the HW 5 Help video that she would not combine the two graphs because of this). That being said, also included below is a graph comparing the runtimes of both the insertion and merge sorting algorithms. The graphs combined does give a good visual of how much faster the insertion sort grows in time compared to merge sort.

**Insertion Sort Run Times (Graph)**

**Merge Sort Run Times (Graph)**

**Insertion Sort vs Merge Sort (Graph)**

* 1. **What type of curve best fits each data set? Again, you can use Excel, Matlab, any software or a graphing calculator to calculate a regression equation. Give the equation of the curve that best “fits” the data and draw that curve on the graphs of created in part c).**

**Answer:**

When individually graphed the 2nd degree polynomial (exponential) trend line matched insertion sort best with an equation of (gives Regression of ). However, a linear trend line was better suited for merge sort with an equation of (gives Regression of ). The graph screen shots above include the trend line curves necessary for this question. Just to be safe, I included the combination graph a second time below.

* 1. **How do your experimental running times compare to the theoretical running times of the algorithms? Remember, the experimental running times were “average case” since the input arrays contained random integers.**

**Answer:**

The experimental running times were in line with the theoretical running times expected from both algorithms. The insertion sort was almost exactly on point with the theoretical O( performance associated with the algorithm and we can certainly see that the runtime increases exponential as the array size gets larger. The running time of merge sort deviated slightly as it looked more linear O( rather than the expected O(. As mentioned in the lecture videos, this is likely because the effect of the logarithm was negligible with the size of our randomized test arrays. The “Why does my nlogn graph look linear” document states that the variation is so small when compared to the values of n that for all practical purposes it is a constant. In theory the algorithm may be but for some narrow ranges of n it may appear almost linear.

**EXTRA CREDIT: A Tale of Two Algorithms: *It was the best of times, it was the worst of times…***

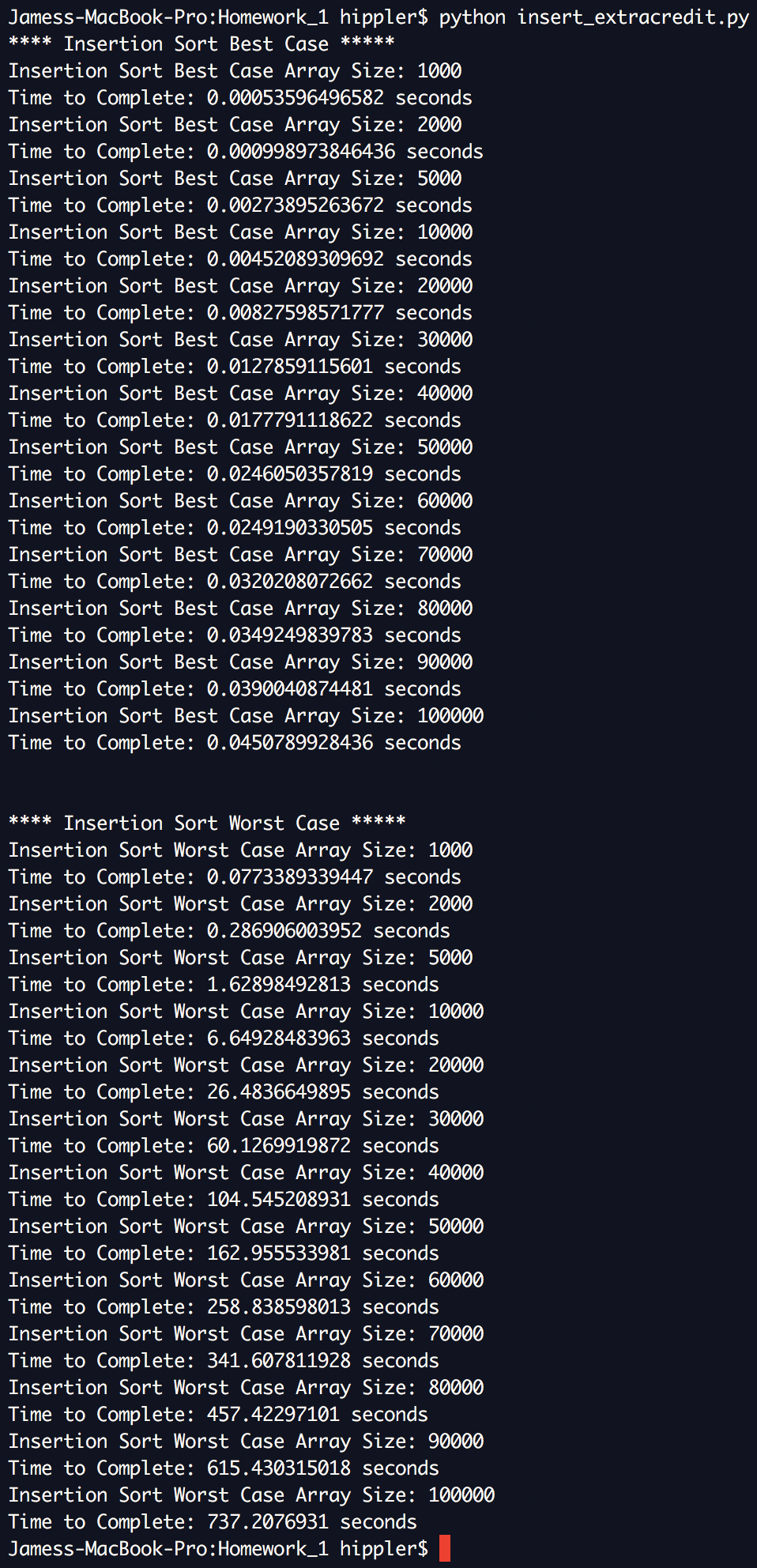
**Generate best case and worst-case inputs for the two algorithms and repeat the analysis in parts b) to e) above. To receive credit, you must discuss how you generated your inputs and your results. Submit your code to TEACH in a zip file with the code from Problem 4.**

**Answer:**

Code was for each sorting algorithm has been included in my submission to TEACH.

**Insertion Sort Best vs Worst Case Scenario**

The best-case scenario for Insertion sort would be having the array already sorted (ascending order). This would allow the program to function at linear time which is even faster than the output of merge sort. The worst-case scenario for Insertion sort would be when the array items are inserted in the exact opposite order of what would be considered sorted (descending order). This would cause the program to perform at a speed of approximately exponential where the algorithm needs to make a comparison with every preceding element for each element that it is attempting to sort. To accomplish the creation of these arrays, I had a function that generated an array of increasing length that inserted elements in ascending order to replicate best case scenarios and then descending order to replicate worst-case scenarios. I then used the system clock to record the time it took insertion sort to sort each back into ascending order. From the results, we can see that the function performs significantly slower than when random integers were assigned to those same array lengths. For example, an array of 100,000 elements already sorted is processed by insertion sort in less than 1 second (0.0450789928436 seconds) whereas the same sized array with integers inserted in the opposite order takes over 12 minutes (737.2076931 seconds) to complete. Code and Results provided below.





**Merge Sort Best vs Worst Case Scenario**

Merge Sort tends to operate at approximately the speed of as is has to recursively split the array and then remerge the elements linearly regardless of the order of the array elements. This means that merge sort will generally operate less efficiently than even insertion sort when processing arrays that are sorted or nearly sorted. For the best-case scenario, I inserted the integers in ascending order to match what would be considered sorted after merge sort completes. To emulate a worst-case scenario, I followed the recommendation provided in the Geeks for Geeks “Find a permutation that causes worst case of merge sort” (<https://www.geeksforgeeks.org/find-a-permutation-that-causes-worst-case-of-merge-sort/>) designed to maximize the number of required comparisons. To accomplish this, the integers need to be inserted into the array such that “the left and right sub-array involved in merge operation should store alternate elements of sorted array” resulting in each element being compared at least once. I managed this by inserting odd numbers in the left side of the array and even numbers on the right, both in ascending order. When processing the unsorted, worst-case array we can see that merge sort is slower than the best-case scenario but only marginally. The algorithm is only saving time when comparing ( but is still being dominated by the linear ( need to split and merge all elements of the array. Code and Results provided below.

