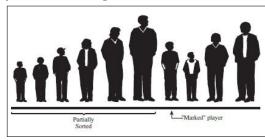
Data Structures - Simple Sorting 3

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Description

- ▶ Let us assume that the baseball team is partially sorted.
- ▶ Put an imaginary marker such that the players to the left of this marker are partially sorted.
- ► This means that they are sorted among themselves; but not necessarily in their final positions.



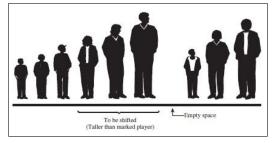
▶ Partial sorting did not take place in the bubble sort and selection sort.

Insertion Sort

- ▶ Best of the elementary sorts.
- ▶ Still $O(N^2)$ time.
- ► Twice as fast as bubble sort.
- ► Somewhat faster than the selection sort in normal situations.
- ▶ Slightly complex than bubble and selection sorts.
- ▶ Used as the final stage of more sophisticated sorts, such as quicksort.

Description (Contd.)

- ▶ Next insert the marked player in the appropriate place in the (partially) sorted group.
- ▶ To provide a space for this shift, first take the marked player out of line (stored in a temporary variable).



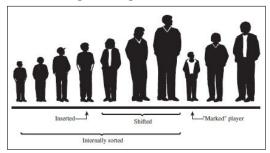
▶ Then shift the sorted players to make room (tallest \rightarrow marked player's spot, next tallest \rightarrow tallest's spot, and so on).

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Description (Contd.)

- ► This shifting process stops when the first player that is shorter than the marked player is met.
- ▶ In this instance, the marked player is inserted into the space the last shift opened up.



- ▶ Next move the marked player one position to the right.
- ► This process is repeated until all the unsorted players have been inserted into the right place in the partially sorted group.

Efficiency of the Insertion Sort

If there N number of elements in the array,

Avg. no of comparisons on the first pass = $1 = \frac{2}{2}$

Avg. no of comparisons on the second pass $=\frac{1+2}{2}=\frac{3}{2}$

Avg. no of comparisons on the third pass = $\frac{1+2+3}{2} = \frac{6}{2} = 2 = \frac{4}{2}$

Avg. no of comparisons on the fourth pass = $\frac{1+2+3+4}{4} = \frac{10}{4} = \frac{5}{2}$

:

Avg. no of comparisons on the last pass

$$\frac{1+2+3+\dots+(N-1)}{N-1} = \frac{N(N-1)}{2(N-1)} = \frac{N}{2}$$

InsertionSortArray class

Efficiency (Contd.)

: Average no of comparisons

$$= \frac{2}{2} + \frac{3}{2} + \frac{4}{2} + \frac{5}{2} + \dots + \frac{N}{2}$$

$$= (\frac{1}{2} + \frac{2}{2} + \frac{3}{2} + \frac{4}{2} + \frac{5}{2} + \dots + \frac{N}{2}) - \frac{1}{2}$$

$$= \frac{N(N+1)}{4} - \frac{1}{2}$$

$$\approx \frac{N(N+1)}{4}$$

Efficiency (Contd.)

- ► The number of copies is approximately the same as the number of comparisons.
- ▶ Copying does not take time as swapping.
- ▶ For random data, this algorithm runs twice as fast as the bubble sort and faster than the selection sort.
- ▶ The insertion sort runs in $O(N^2)$ time for random data.
- ▶ For data that is almost sorted, the algorithm runs in almost O(N) time.
- ► For data arranged in inverse sorted order, the algorithm runs no faster than the bubble sort.

Memory

- ▶ All three algorithms, operations carried out inside the array.
- ▶ Very little extra memory is required an extra variable to store an item temporarily while it's being swapped.

Comparing the simple sorts

- ▶ Bubble sort is suitable if the amount of data is small.
- ▶ The selection sort minimizes the number of swaps, but the number of comparisons is still high.
- ▶ Useful when amount of data is small and swapping data is very time consuming.
- ▶ The insertion sort is the most versatile bet in most cases, assuming amount of data is small or the data is almost sorted.