

University of Waterloo
CS 462 — Formal Languages and Parsing
Winter 2013
Problem Set 3

Distributed Tuesday, January 22 2013.

Due Tuesday, January 29 2013, in class.

All answers should be accompanied by proofs.

1. [10 marks] Must every primitive word of length > 1 have at least two unbordered conjugates? Prove or disprove.

Yes they do.

Proof by Contradiction. \models

Let $w \in \Sigma^*$ and $\alpha \in \Sigma$ s.t. $\neg \exists v^i = w, |x| > 0, |x|_\alpha > 0$

$\text{shuffle}(\dots, w) := \bigcup_{W \in \text{shuffle}(\dots)} \text{shuffle}(W, w)$

□

2. [10 marks] Consider the equation in words $x \text{III} y = z^2$, where III is the perfect shuffle. Describe all solutions to this equation. Hint: there are separate cases depending on whether $|x|, |y|$ are both even or both odd.

(By “describe all solutions” we mean come up with a characterization of the solutions that is somewhat analogous to the descriptions in Theorems 2.3.2 and 2.3.3 of the course text, giving necessary and sufficient conditions for x, y, z .)

if $|x| = |y| = 2 * i \models$

Let $x = x_a x_b$ and $y = y_a y_b$ where $|x_a| = |x_b| = |y_a| = |y_b| = i$

$x \text{III} y = x_a \text{III} y_a x_b \text{III} y_b = z z$

$x_a \text{III} y_a = z = x_b \text{III} y_b$

$x_a = x_b$ and $y_a = y_b$

if $|x| = |y| = 2 * i + 1 \models$

Let $x = x_a x_c x_b$ and $y = y_a y_c y_b$ where $|x_a| = |x_b| = |y_a| = |y_b| = i$ and $|x_c| = |y_c| = 1$

$x \text{III} y = x_a \text{III} y_a x_c y_c x_b \text{III} y_b = z z$

$x_a x_c = y_c y_b$ and $y_a = x_b$

3. [10 marks] Call a word w *uneven* if every nonempty subword has the property that at least one letter appears an odd number of times. For example, **abac** is uneven.

- (a) [5 marks] Show that if w is an uneven word over an alphabet with k letters, then $|w| < 2^k$.
- (b) [5 marks] Prove that the bound in (a) is sharp, by exhibiting an uneven word of length $2^k - 1$ over every alphabet of size $k \geq 1$.

For (a)

Given $W \in \Sigma^*$ and $|\Sigma| = k$

Assume $|W| \geq 2^k$

Let $\text{prefix}(w, j) = w'$ where $w'w'' = w$ and $|w'| = j$

Let $\bigcup_{i \in Z_{|\Sigma|}} \{\alpha_i\} = \Sigma$

Let $\text{par}(w) = v$ where $v_i = |w|_{\alpha_i} \bmod 2$

Let $V = \{\text{par}(\text{prefix}(w, j)) \mid j \in Z_{|W|+1} \setminus \{0\}\}$

if $\exists i$ where $V_i = \text{par}(\text{prefix}(W, j)) = 0$ then $\text{prefix}(W, j)$ is an even word

otherwise $\exists i, j$ where $m \leq n$, $V_i = \text{par}(\text{prefix}(W, m)) = V_j = \text{par}(\text{prefix}(W, n))$ because there $2^k - 1$ unique non-zero vector in the space and $|V| = |W| \geq 2^k$

Let $\text{prefix}(W, m)z = \text{prefix}(W, n)$

Then $\text{par}(\text{prefix}(W, m)) \text{par}(z) = \text{par}(\text{prefix}(W, n))$

Then $\text{par}(z) = \text{par}(\text{prefix}(W, n)) - \text{par}(\text{prefix}(W, m)) = 0$

Then z is an even word

In both cases, W end up contains an even word. Thus $|W| \not\geq 2^k$.

For (b)

Given $|\Sigma| = k$

Let $\bigcup_{i \in Z_{|\Sigma|}} \{\alpha_i\} = \Sigma$

Let $w_0 = \alpha_0$

Let $w_{n+1} = w_n \amalg \alpha_{n+1}^{|w_n|} \alpha_{n+1}$ for $n \geq 0$

By observation, $|w_0| = 2^1 - 1$ and w_0 is uneven

Given $|w_n| = 2^n - 1$ and w_n is uneven

$|w_{n+1}| = |w_n| * 2 + 1 = 2^{n+1} - 1$

Let w be a subword of w_{n+1}

if $|w| = 1$, then w is uneven otherwise $|w| \geq 2$, then $w \in \text{perm}(w' \alpha_{n+1}^i)$ where w' is a subword of w_n

Then $\text{par}(w) = \text{par}(w') + \text{par}(\alpha_{n+1}^i)$

Then $\text{par}(w') \neq 0$ because w' is uneven

Then $\text{par}(w') + \text{par}(\alpha_{n+1}^i) \neq 0$ because $|w|_{\alpha_{n+1}} = 0$

Then w is uneven because $\text{par}(w) \neq 0$ In both cases, w is uneven. Therefore w_{n+1} is uneven

By induction, $|w_k| = 2^k - 1$ and w_k is uneven.