

# CS222/CS122C: Principles of Data Management

UCI, Fall 2019  
Lecture #02

## Storing Data: Record/Page Formats

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# *Files of Records*

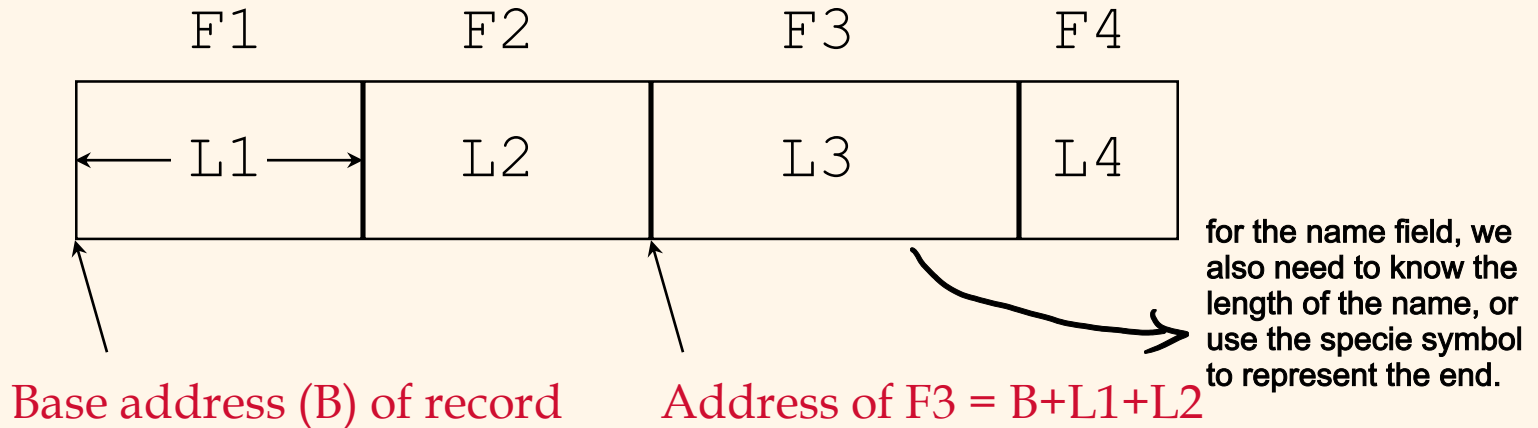
- ❖ Page or block is OK when doing I/O, but higher levels of DBMS operate on *records*, and thus want *files of records*.
- ❖ FILE: A collection of pages, each containing a collection of records. Must support:
  - Insert (append)/delete/modify record
  - Read a particular record (specified using *record id*)
  - Scan all records (possibly with some conditions on the records to be retrieved)

# *Example*

```
CREATE TABLE Emp(id INT,  
gender CHAR(1),  
name VARCHAR(30),  
Salary float  
);
```

# Record Formats: Fixed Length

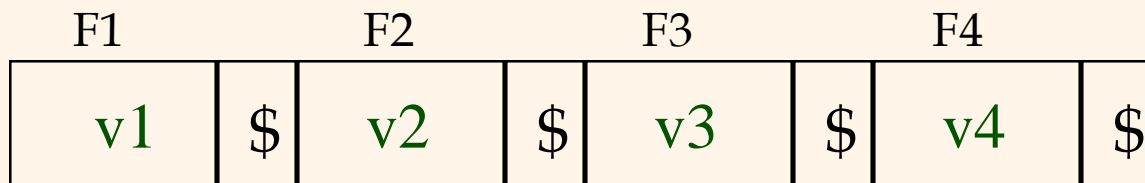
Each field is fixed length



- ❖ Information about field types is the same for all records in file; it is stored in the *system catalogs*.  
(Note: Record field info in Project 1 passed in “from above” ...!)
- ❖ Finding the *i*'th field of a record does *not* require scanning the record.

# Record Formats: Variable Length

## ❖ Several alternative formats (# fields is fixed):

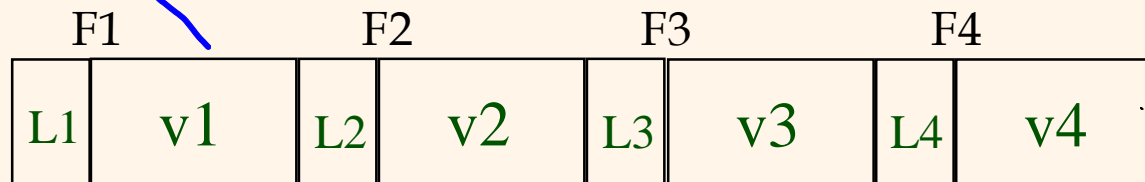


Fields Delimited by Special Symbols



Consider is the data contains the special symbol, then it doesn't work.

If V1 is null, put some special to L1, like (-1) to represent V1 is null



Fields Preceded by Field Lengths



Use one byte -> L to shown the length of the field data.  
If we know V1 is integer, V2 is float, we don't need to use L1 and L2.

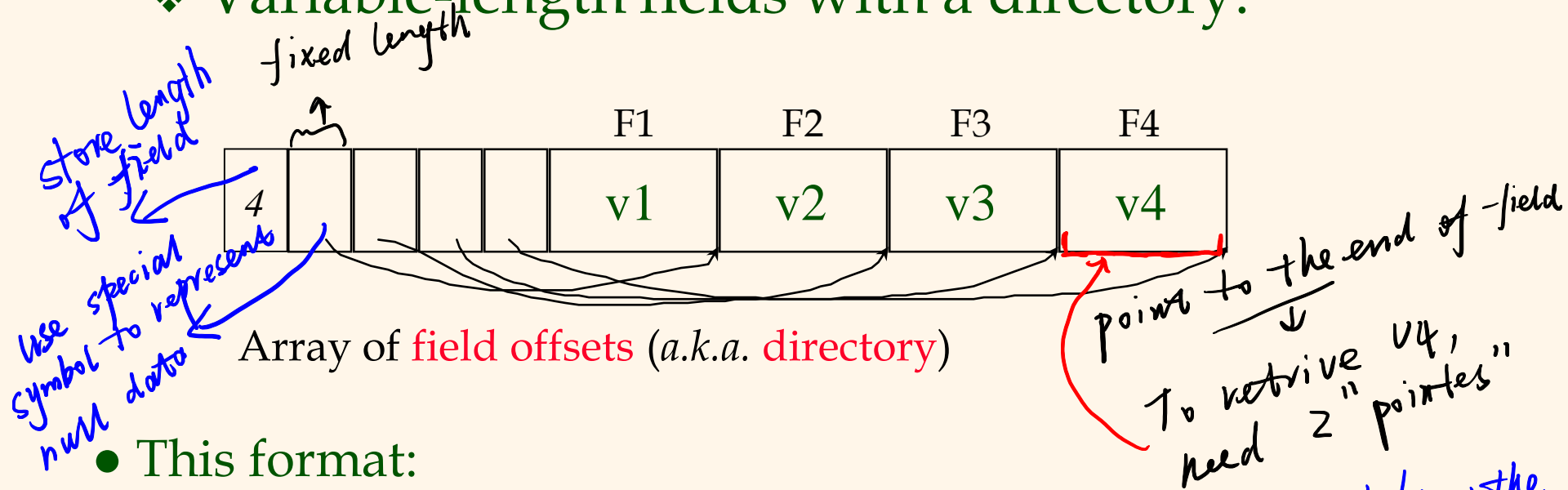
## ● Some thought questions for you:

Because the length is variable, if you want query a special field data, you may need to scan the entire record, and this is expensive.

- (1) What's true of the second format but not the first?
- (2) What annoying disadvantage do both formats share?
- (3) And, how do we know the field count in each case?

# Record Formats: Variable Length (continued)

## ❖ Variable-length fields with a directory:



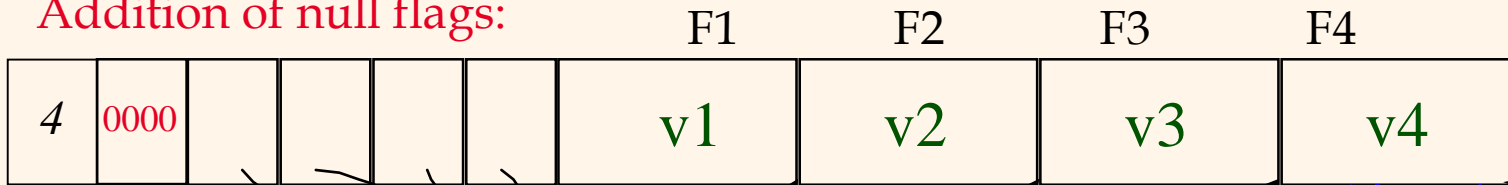
### ● This format:

- (1) Offers direct access to the  $i$ 'th field.
- (2) Helps support efficient storage of null values. (Q: How?)  
→ means I don't know the value
- (3) Just requires a small directory overhead.
- (4) Can even help with ALTER TABLE ADD COLUMN! (Q: How?)

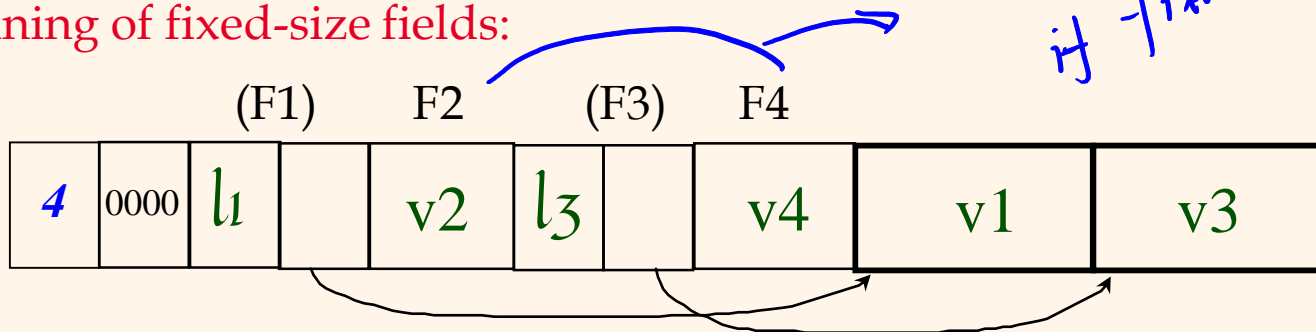
# Record Formats: Variable Length

## ❖ More variations on a theme...

Addition of null flags:



Inlining of fixed-size fields:

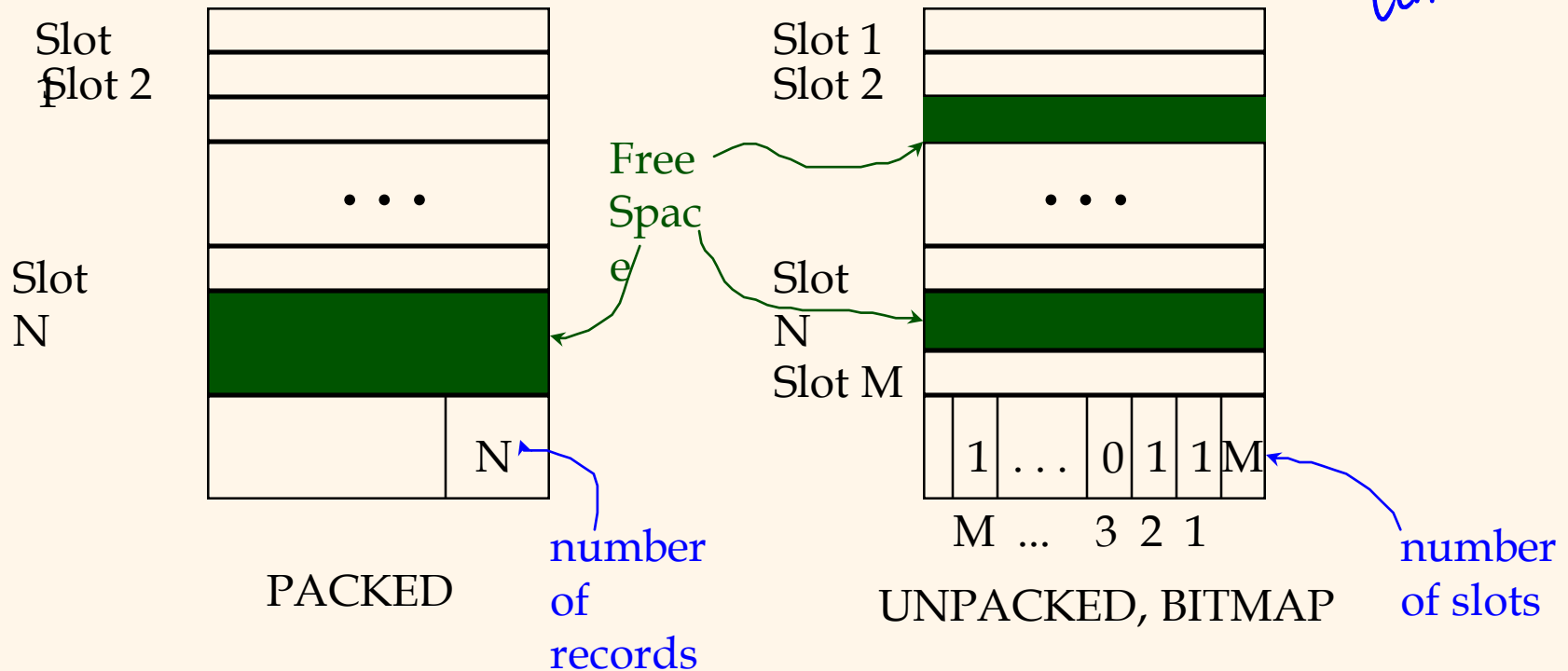


*v2 and v4  
if fixed length for int  
and float...*

*Next: Page formats*



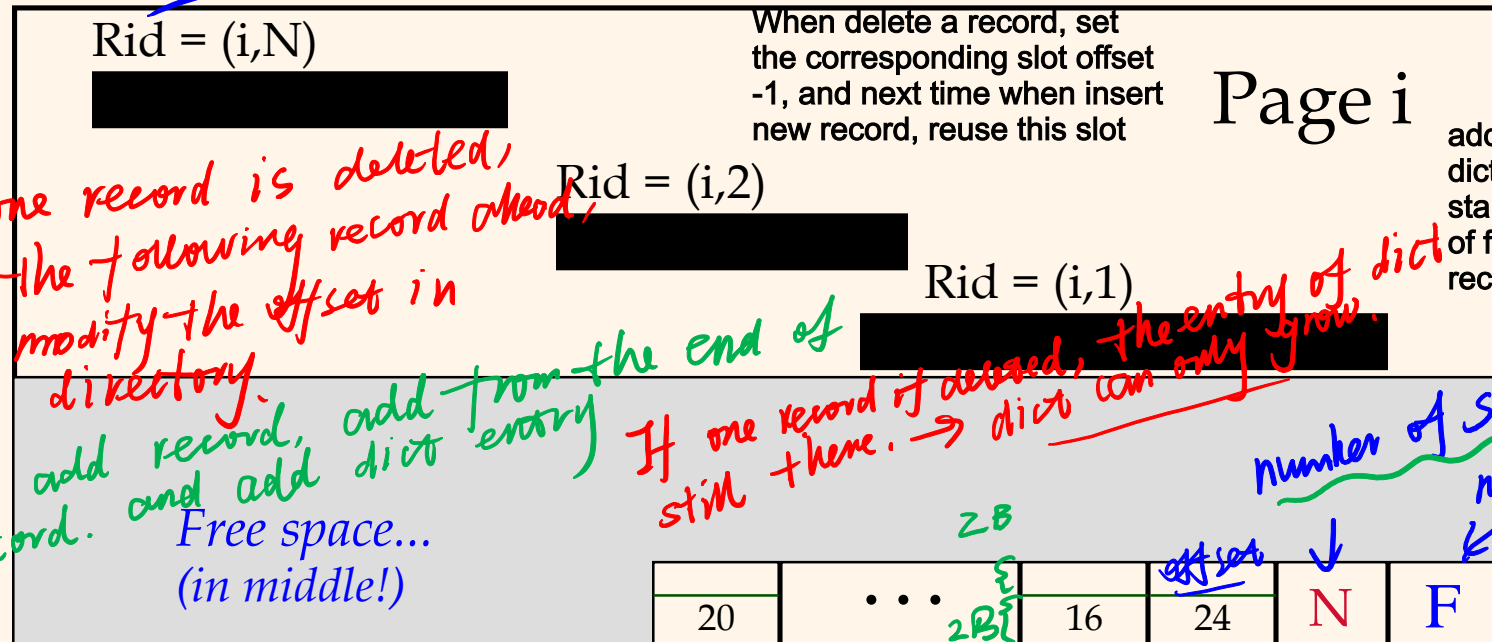
# Page Formats: Fixed Length Records



- Record id =  $\langle \text{page id, slot \#} \rangle$ . In the first (packed) alternative, records will move around for free space management: Rids change  $\square$  may be unacceptable!

location needs 2 Byte  $\rightarrow \log_2^{4K} = 12 \text{ bit} \rightarrow 2 \text{ Bytes}$

## Page Formats: Variable Length Records

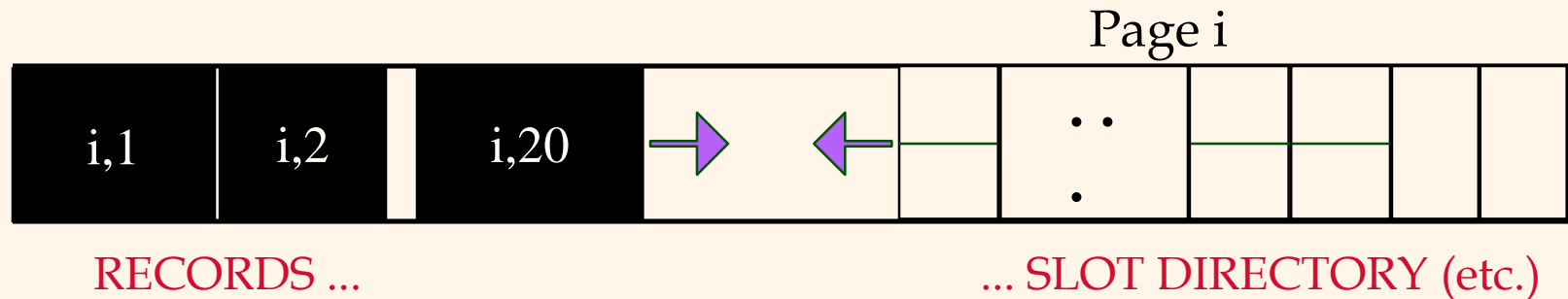


RID { page #  
dir. id: → points to directory entry.  
Can move records within page w/o changing RIDs; not as

← dict grow SLOT DIRECTORY (offset, length) record 1

- Can move records within page w/o changing RIDs; not so unattractive for fixed-length records as a result.
- Record movement? (1) Tombstones, or (2) PKeys (vs. RIDs)

## ... *Variable Length Records (cont.)*



- *Two variable-sized areas growing towards to each other (living within a one-page space budget!)*
- Other variations on these formats are possible as well
  - Could track free space holes with an offset-based list structure
  - Could use a different *record* format (e.g., PAX, which clusters values by field in page rather than by record and then field)
  - ....

## *PAX format*

# Traditional Format

PAGE HEADER					RH1		0962	
Jane		30	RH2		7658		John	
45	RH3	3589		Jim	20	RH4		
5523		Susan		52				
				•	•	•	•	

# PAX Format

<b>PAGE HEADER</b>				0962	7658
3859	5523				
Jane	John	Jim	Susan		
					• • • •
30	52	45	20		

- ❖ PAX partitions each page into minipages based on fields
  - Good caching behaviors for “select fields from ...”;
  - Compression
  - [www.pdl.cmu.edu/PDL-FTP/Database/pax.pdf](http://www.pdl.cmu.edu/PDL-FTP/Database/pax.pdf)
- ❖ Column store (e.g., Vertica)

# Project 1: RecordBasedFileManager

insertRecord

readRecord

printRecord

start of free space:  $PAGE\_SIZE - 2 - 2 - 4 \times 4 - F$   
↑  
freespace  
PageDirectory SlotDirectory

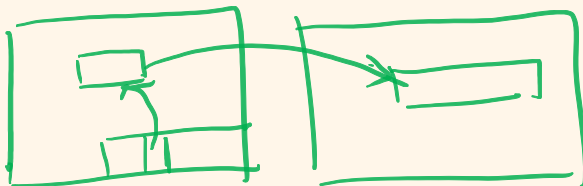
deleteRecord: set the slot special symbol, shift the following records, change the offset of the following record. When insert a new record, if there is a slot unused (previously deleted), first use this slot.

updateRecord:

case 1: update record needs less space: shift the following record from the start to end and change the offset and freespace.

case2: update record needs more space: If the page is available to contain this updated record, shift the following record from the end to start. We can't change the RID, and rewrite the record corresponding to this record as a pointer, which points to the location of another page where the record is actually stored. Remember to change the free space of the new page.

Update  
①



update  
②

