

# CSCI-H 335: Homework 2: Number representation

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## 1. Answer to question 1:

Assume the largest number for  $N$  digits is  $L(N)$ .

Base case ( $N = 0$ ):

For binary number,  $L(0) = 0 = 1 - 1 = r^0 - 1 = r^N - 1$ .

Induction case :

Assume  $L(N) = r^N - 1$  for some  $r, N \in \mathbb{N}$ .

$$\begin{aligned} L(N+1) &= r^{N+1-1} \cdot (r-1) + L(N) \\ &= r^N - 1 + r^N \cdot (r-1)N \\ &= r^N - 1 + r^{N+1} - r^N \\ &= r^{N+1} - 1 \end{aligned}$$

Add largest number in r-radix to highest digit.  
By IH

## 2. Answer to question 2

Because  $B$  is magnitude less than  $2^{N2}$ ,  $-B$  can be represent in  $N$  bits. We define the twos complement of a binary number by inverting its bits and adding 1. Thus  $-(b_{N-1}..b_0)_2 = (\overline{b_{N-1}..b_0})_2 + 1$ .

## 3. Answer to question 3

Assume  $I$  is  $k$  digits binary number,  $m > k$ , and  $a_1$  is 0 or 1.

Positive case:

$$I = \sum_0^k a_i \cdot 2^i = 0 + \sum_0^k a_i \cdot 2^i = \sum_{k+1}^m 0 \cdot 2^i + \sum_0^k a_i \cdot 2^i \quad (1)$$

Thus, it's value prevente for positive value because the highest digit of positive number always be 0.

Negative case:

For all two's complement number, highest digit is always 1 if the given number is in range. Inverse 1 we get 0, and we proved in positive case. Thus extension will always keep value same.

## 4. Answer to question 4

Left-shift:

Assume  $b_{-1} = 0$

$$(b_{N-1}..b_0)_2 \cdot 2 = (-b_{N-1} \cdot 2^{N-1} + \sum_{i=0}^{N-2} b_i \cdot 2^i) \cdot 2 = -b_{N-1} \cdot 2^{(N+1)-1} + \sum_{i=-1}^{(N+1)-2} b_{i+1} \cdot 2^i = (b_{N-1}..b_0, b_{-1})_2 = (b_{N-1}..b_0, 0)_2 \quad (2)$$

Arithmetic right-shift:

$$(b_{N-1}..b_0)_2/2 = (-b_{N-1} \cdot 2^{N-1} + \sum_{i=0}^{N-2} b_i \cdot 2^i)/2 = -b_{N-1} \cdot 2^{N-3} + \sum_{i=1}^{N-3} b_i \cdot 2^i = (b_{N-1}..b_1)_2 = (b_{N-1}, b_{N-1}..b_1)_2 \quad (3)$$

5. Answer to question 5

Repeated Division (mod 2)															
(a)	Quotients	273	136	68	34	17	8	4	2	1					
	Remainders	1	0	0	0	1	0	0	0	1					
$(273)_{10} = (100010001)_2$															
Repeated Division (mod 2)															
(b)	Quotients	1327	663	331	165	82	41	20	10	5	2	1			
	Remainders	1	1	1	1	0	1	0	0	1	0	1			
$(1327)_{10} = (10101001111)_2$															
Repeated Division (mod 2)															
(c)	Quotients	31673	15836	7918	3959	1979	989	494	247	123	61	30	15	7	3
	Remainders	1	0	0	1	1	1	0	1	1	1	0	1	1	1
$(31673)_{10} = (111101110111001)_2$															
Repeated Division (mod 2)															
(d)	Quotients	107	53	26	13	6	3	1							
	Remainders	1	1	0	1	0	1	1							
$(107)_{10} = (1101011)_2$															

6. Answer to question 6

(a)	<div>0</div>	<div>0</div>	<div>0</div>	
		2	3	3
	-		2	1
		2	1	2
$233 - 21 = 212$				
(b)	<div>0</div>	<div>1</div>	<div>1</div>	<div>1</div>
		1	0	0
	-		9	0
			9	9
$1001 - 902 = 99$				
(c)	<div>0</div>	<div>0</div>	<div>0</div>	
		1	7	3
	-	1	2	3
			5	0
$173 - 123 = 50$				

(d)

7.

(a)

$$(111)_{10} = (01101111)_2$$

(b)

$$(59)_{10} = (00111011)_2$$

(c)

$$(105)_{10} = (01101001)_2$$

(d)

$$(31)_{10} = (00011111)_2$$

(e)

$$(127)_{10} = (01111111)_2$$

8. Answer to question 8:

Base case ( $i = 0$ ):

$$V((c_{i+1}, s_i)_2) = V((c_1, s_0)_2) = V((a_0)_2) + V((b_0)_2) + V((c_0)_2) = V((a_i)_2) + V((b_i)_2) + V((c_i)_2)$$

Induction case :

$$\text{Assume } V((c_N, s_{N-1}, \dots, s_0)_2) = V((a_{N-1}, \dots, a_0)_2) + V((b_{N-1}, \dots, b_0)_2) + V((c_0)_2).$$

$$\begin{aligned} V((c_{N+2}, s_{N+1}, \dots, s_0)_2) &= V((a_N, 0, \dots, 0)_2) + V((b_N, \dots, b_0)_2) + V((c_N, 0, \dots, 0)_2) + V((c_N, s_{N-1}, \dots, s_0)_2) \\ &= V((a_N, 0, \dots, 0)_2) + V((b_N, \dots, b_0)_2) + V((c_N, 0, \dots, 0)_2) + V((a_{N-1}, \dots, a_0)_2) + V((b_{N-1}, \dots, b_0)_2) + \\ &V((c_0)_2) \text{ By IH} \\ &= V((a_N, \dots, a_0)_2) + V((b_N, \dots, b_0)_2) + V((c_0)_2) \end{aligned}$$