

南京航空航天大学计算机科学与技术学院计算机组成原理

实验 Bomblab 二进制炸弹实验报告

实验信息

姓名：李元成
学号：161630216
班级：1616302
指导教师：李博涵
实验日期：2018.6
我的炸弹编号：bomb3
我的答案：

```
For NASA, space is still a high priority.  
8 9 26 61 148 357  
1 -458  
5 15 HappyHangHangXueZhang  
jdoefg  
3 2 6 4 5 1  
50
```

第一关

考察内容

寻找地址的所存的字符串

反汇编代码分析

```
08048ae0 <phase_1>:  
8048ae0: 55                push    %ebp  
8048ae1: 89 e5            mov     %esp, %ebp  
8048ae3: 83 ec 18        sub     $0x18, %esp  
8048ae6: c7 44 24 04 80 a1 04 movl    $0x804a180, 0x4(%esp) 答案  
8048aed: 08  
8048aee: 8b 45 08        mov     0x8(%ebp), %eax  
8048af1: 89 04 24        mov     %eax, (%esp)  
8048af4: e8 99 04 00 00  call    8048f92 <strings_not_equal>
```

8048af9: 85 c0	test %eax,%eax 比较，不同就炸
8048afb: 74 05	je 8048b02 <phase_1+0x22>
8048afd: e8 b3 06 00 00	call 80491b5 <explode_bomb>
8048b02: c9	leave
8048b03: c3	ret

操作过程或解题思路

```
(gdb) x/s 0x804a180
0x804a180: "For NASA, space is still a high priority."
```

第二关

考察内容

for 循环

反汇编代码分析

08048b04 <phase_2>:	
8048b04: 55	push %ebp
8048b05: 89 e5	mov %esp,%ebp
8048b07: 56	push %esi
8048b08: 53	push %ebx
8048b09: 83 ec 30	sub \$0x30,%esp
8048b0c: 8d 45 e0	lea -0x20(%ebp),%eax
8048b0f: 89 44 24 04	mov %eax,0x4(%esp)
8048b13: 8b 45 08	mov 0x8(%ebp),%eax
8048b16: 89 04 24	mov %eax,(%esp)
8048b19: e8 d9 06 00 00	call 80491f7 <read_six_numbers>
读数	
8048b1e: 8b 45 e0	mov -0x20(%ebp),%eax
8048b21: 83 f8 07	cmp \$0x7,%eax 第一个大于 7 就炸
8048b24: 7e 05	jle 8048b2b <phase_2+0x27>
8048b26: 3b 45 e4	cmp -0x1c(%ebp),%eax 第二个小于
第一个就炸	
8048b29: 7e 22	jle 8048b4d <phase_2+0x49>
8048b2b: e8 85 06 00 00	call 80491b5 <explode_bomb>

8048b30: eb 1b	jmp 8048b4d <phase_2+0x49>
8048b32: 8b 53 fc	mov -0x4(%ebx), %edx edx=ebx-4
8048b35: 8b 43 f8	mov -0x8(%ebx), %eax eax=ebx-8
8048b38: 8d 04 50	lea (%eax, %edx, 2), %eax
eax+=edx*2	
8048b3b: 39 03	cmp %eax, (%ebx) eax 和 ebx 不等
就炸	
8048b3d: 74 05	je 8048b44 <phase_2+0x40>
8048b3f: e8 71 06 00 00	call 80491b5 <explode_bomb>
8048b44: 83 c3 04	add \$0x4, %ebx ebx+=4
8048b47: 39 f3	cmp %esi, %ebx for 循环控制
8048b49: 75 e7	jne 8048b32 <phase_2+0x2e>
8048b4b: eb 08	jmp 8048b55 <phase_2+0x51>
8048b4d: 8d 5d e8	lea -0x18(%ebp), %ebx
8048b50: 8d 75 f8	lea -0x8(%ebp), %esi
8048b53: eb dd	jmp 8048b32 <phase_2+0x2e>
8048b55: 83 c4 30	add \$0x30, %esp
8048b58: 5b	pop %ebx
8048b59: 5e	pop %esi
8048b5a: 5d	pop %ebp
8048b5b: c3	ret

操作过程或解题思路

```
(gdb) u *0x08048b38
0x08048b38 in phase_2 ()
(gdb) p $edx
$2 = 9
(gdb) p $eax
$3 = 8
```

eax 放的是第一个数。edx 是第二个
 $eax = eax + edx * 2$

而 ebx 放的是第三个数

所以 $f(n) = f(n-1) * 2 + f(n-2)$

我前两个是 8 9

所以 8 9 26 61 148 357

第三关

考察内容

Switch 语句

反汇编代码分析

08048b5c <phase_3>:

8048b5c: 55	push %ebp
8048b5d: 89 e5	mov %esp, %ebp
8048b5f: 83 ec 28	sub \$0x28, %esp
8048b62: 8d 45 f0	lea -0x10(%ebp), %eax
8048b65: 89 44 24 0c	mov %eax, 0xc(%esp)
8048b69: 8d 45 f4	lea -0xc(%ebp), %eax
8048b6c: 89 44 24 08	mov %eax, 0x8(%esp)
8048b70: c7 44 24 04 dd a3 04	movl \$0x804a3dd, 0x4(%esp) %d %d
8048b77: 08	
8048b78: 8b 45 08	mov 0x8(%ebp), %eax
8048b7b: 89 04 24	mov %eax, (%esp)
8048b7e: e8 5d fc ff ff	call 80487e0
<__isoc99_sscanf@plt>	读数
8048b83: 83 f8 01	cmp \$0x1, %eax 输入个数小于 2 就炸
8048b86: 7f 05	jg 8048b8d <phase_3+0x31>
8048b88: e8 28 06 00 00	call 80491b5 <explode_bomb>
8048b8d: 83 7d f4 07	cmpl \$0x7, -0xc(%ebp) 如果第一个数大于 7 就炸
8048b91: 77 65	ja 8048bf8 <phase_3+0x9c>
8048b93: 8b 45 f4	mov -0xc(%ebp), %eax
8048b96: ff 24 85 dc a1 04 08	jmp *0x804a1dc(, %eax, 4)
8048b9d: b8 00 00 00 00	mov \$0x0, %eax
8048ba2: eb 05	jmp 8048ba9 <phase_3+0x4d>
8048ba4: b8 2b 01 00 00	mov \$0x12b, %eax
8048ba9: 2d 2b 02 00 00	sub \$0x22b, %eax
8048bae: eb 05	jmp 8048bb5 <phase_3+0x59>
8048bb0: b8 00 00 00 00	mov \$0x0, %eax
8048bb5: 05 c7 02 00 00	add \$0x2c7, %eax
8048bba: eb 05	jmp 8048bc1 <phase_3+0x65>
8048bbc: b8 00 00 00 00	mov \$0x0, %eax
8048bc1: 2d 66 02 00 00	sub \$0x266, %eax
8048bc6: eb 05	jmp 8048bcd <phase_3+0x71>
8048bc8: b8 00 00 00 00	mov \$0x0, %eax
8048bcd: 05 66 02 00 00	add \$0x266, %eax
8048bd2: eb 05	jmp 8048bd9 <phase_3+0x7d>
8048bd4: b8 00 00 00 00	mov \$0x0, %eax
8048bd9: 2d 66 02 00 00	sub \$0x266, %eax
8048bde: eb 05	jmp 8048be5 <phase_3+0x89>
8048be0: b8 00 00 00 00	mov \$0x0, %eax

8048be5: 05 66 02 00 00	add	\$0x266, %eax
8048bea: eb 05	jmp	8048bf1 <phase_3+0x95>
8048bec: b8 00 00 00 00	mov	\$0x0, %eax
8048bf1: 2d 66 02 00 00	sub	\$0x266, %eax
8048bf6: eb 0a	jmp	8048c02 <phase_3+0xa6>
8048bf8: e8 b8 05 00 00	call	80491b5 <explode_bomb>
8048bfd: b8 00 00 00 00	mov	\$0x0, %eax
8048c02: 83 7d f4 05	cmpl	\$0x5, -0xc(%ebp)
8048c06: 7f 05	jg	8048c0d <phase_3+0xb1>
8048c08: 3b 45 f0	cmp	-0x10(%ebp), %eax 第二个数和
各种操作后的第一个数相比, 不等就炸		
8048c0b: 74 05	je	8048c12 <phase_3+0xb6>
8048c0d: e8 a3 05 00 00	call	80491b5 <explode_bomb>
8048c12: c9	leave	
8048c13: c3	ret	

操作过程或解题思路

```
(gdb) x/d $ebp-0x10
0xffffd5e8: -458
```

答案是

1 -458

第四关

考察内容

递归

反汇编代码分析

08048c72 <phase_4>:		
8048c72: 55	push	%ebp
8048c73: 89 e5	mov	%esp, %ebp
8048c75: 83 ec 28	sub	\$0x28, %esp
8048c78: 8d 45 f0	lea	-0x10(%ebp), %eax
8048c7b: 89 44 24 0c	mov	%eax, 0xc(%esp)
8048c7f: 8d 45 f4	lea	-0xc(%ebp), %eax
8048c82: 89 44 24 08	mov	%eax, 0x8(%esp)
8048c86: c7 44 24 04 dd a3 04	movl	\$0x804a3dd, 0x4(%esp)

```

8048c8d: 08
8048c8e: 8b 45 08          mov     0x8(%ebp),%eax
8048c91: 89 04 24          mov     %eax, (%esp)
8048c94: e8 47 fb ff ff    call    80487e0
<__isoc99_sscanf@plt> 和 phase_3 一样的套路
8048c99: 83 f8 02          cmp     $0x2,%eax 个数不是 2 就炸
8048c9c: 75 06             jne     8048ca4 <phase_4+0x32>
8048c9e: 83 7d f4 0e       cmpl    $0xe,-0xc(%ebp) 第一个数小于 0xe, 不然就炸
8048ca2: 76 05             jbe     8048ca9 <phase_4+0x37>
8048ca4: e8 0c 05 00 00    call    80491b5 <explode_bomb>
8048ca9: c7 44 24 08 0e 00 00 movl    $0xe, 0x8(%esp) 准备参数
8048cb0: 00
8048cb1: c7 44 24 04 00 00 00 movl    $0x0, 0x4(%esp) 准备参数
8048cb8: 00
8048cb9: 8b 45 f4          mov     -0xc(%ebp),%eax
8048cbc: 89 04 24          mov     %eax, (%esp) 准备参数
8048cbf: e8 50 ff ff ff    call    8048c14 <func4>调用 func4
8048cc4: 83 f8 0f          cmp     $0xf,%eax 返回值不是 15 就炸
8048cc7: 75 06             jne     8048ccf <phase_4+0x5d>
8048cc9: 83 7d f0 0f       cmpl    $0xf,-0x10(%ebp) 第二个数不是 15 也要炸
8048ccd: 74 05             je      8048cd4 <phase_4+0x62>
8048ccf: e8 e1 04 00 00    call    80491b5 <explode_bomb>
8048cd4: c9               leave
8048cd5: c3               ret

```

08048c14 <func4>:

```

8048c14: 55               push    %ebp
8048c15: 89 e5            mov     %esp,%ebp
8048c17: 56               push    %esi
8048c18: 53               push    %ebx
8048c19: 83 ec 10         sub     $0x10,%esp
8048c1c: 8b 55 08         mov     0x8(%ebp),%edx 取出参数
8048c1f: 8b 45 0c         mov     0xc(%ebp),%eax 取出参数
8048c22: 8b 75 10         mov     0x10(%ebp),%esi 取出参数
8048c25: 89 f1            mov     %esi,%ecx
8048c27: 29 c1            sub     %eax,%ecx ecx-=eax
8048c29: 89 cb            mov     %ecx,%ebx
8048c2b: c1 eb 1f         shr     $0x1f,%ebx ebx>>31, 取零
8048c2e: 01 d9            add     %ebx,%ecx ecx+=ebx
8048c30: d1 f9            sar     %ecx ecx/=2
8048c32: 8d 1c 01         lea     (%ecx,%eax,1),%ebx

```

ebx=ecx+eax

8048c35: 39 d3	cmp %edx,%ebx 如果 edx<=ebx 就跳
8048c37: 7e 17	jle 8048c50 <func4+0x3c>
8048c39: 8d 4b ff	lea -0x1(%ebx),%ecx 准备参数
8048c3c: 89 4c 24 08	mov %ecx,0x8(%esp) ebx-1
8048c40: 89 44 24 04	mov %eax,0x4(%esp) eax
8048c44: 89 14 24	mov %edx, (%esp) edx
8048c47: e8 c8 ff ff ff	call 8048c14 <func4> 调用 func4
8048c4c: 01 d8	add %ebx,%eax ebx+=func4
8048c4e: eb 1b	jmp 8048c6b <func4+0x57> return
8048c50: 89 d8	mov %ebx,%eax else eax=ebx
8048c52: 39 d3	cmp %edx,%ebx if(edx<=ebx)就跳
8048c54: 7d 15	jge 8048c6b <func4+0x57> return

ebx

8048c56: 89 74 24 08	mov %esi,0x8(%esp) 准备参数 esi
8048c5a: 8d 43 01	lea 0x1(%ebx),%eax ebx++
8048c5d: 89 44 24 04	mov %eax,0x4(%esp) eax
8048c61: 89 14 24	mov %edx, (%esp) edx
8048c64: e8 ab ff ff ff	call 8048c14 <func4> 调用 func4
8048c69: 01 d8	add %ebx,%eax ebx+=func4
8048c6b: 83 c4 10	add \$0x10,%esp return ebx
8048c6e: 5b	pop %ebx
8048c6f: 5e	pop %esi
8048c70: 5d	pop %ebp
8048c71: c3	ret

操作过程或解题思路

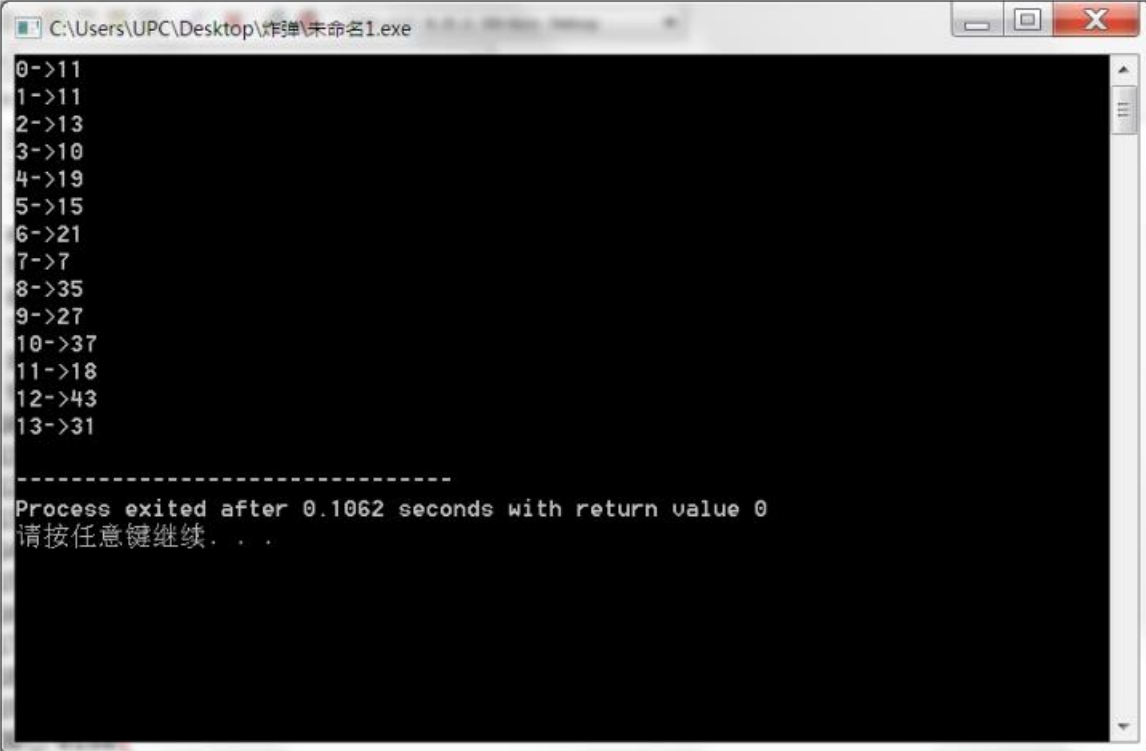
很明显，第二个数是 15

难点是第一个数输进去多少，写个 func4 程序轻松解决

```
int func(int x,int y,int z)
{int k;
int zz=z-y;
k=0;
zz/=2;
k=zz+y;
if(k>x)
{
y=k+func(x,y,k-1); //func(5,0,6)
return y;
}
else
{
```

```
y=k;
if(x<=k)
{
return
y; }
else{
y=k+func(x, k+1, z);

}
} }
```



```
C:\Users\UPC\Desktop\炸弹\未命名1.exe
0->11
1->11
2->13
3->10
4->19
5->15
6->21
7->7
8->35
9->27
10->37
11->18
12->43
13->31

-----
Process exited after 0.1062 seconds with return value 0
请按任意键继续...
```

第一个是 5

所以答案是 5 15

第五关

考察内容

数组

反汇编代码分析

08048cd6 <phase_5>:

8048cd6: 55	push %ebp
8048cd7: 89 e5	mov %esp, %ebp
8048cd9: 53	push %ebx
8048cda: 83 ec 24	sub \$0x24, %esp
8048cdd: 8b 5d 08	mov 0x8(%ebp), %ebx
8048ce0: 89 1c 24	mov %ebx, (%esp)
8048ce3: e8 88 02 00 00	call 8048f70 <string_length>
8048ce8: 83 f8 06	cmp \$0x6, %eax 字符串长度不是 6

就炸

8048ceb: 74 45	je 8048d32 <phase_5+0x5c>
8048ced: e8 c3 04 00 00	call 80491b5 <explode_bomb>
8048cf2: eb 3e	jmp 8048d32 <phase_5+0x5c>接着

跳

8048cf4: 0f b6 14 03	movzbl (%ebx, %eax, 1), %edx
----------------------	------------------------------

edx=ebx+eax

8048cf8: 83 e2 0f	and \$0xf, %edx
8048cfb: 0f b6 92 fc a1 04 08	movzbl 0x804a1fc(%edx), %edx
8048d02: 88 54 05 f1	mov %dl, -0xf(%ebp, %eax, 1)
8048d06: 83 c0 01	add \$0x1, %eax
8048d09: 83 f8 06	cmp \$0x6, %eax eax 控制循环
8048d0c: 75 e6	jne 8048cf4 <phase_5+0x1e> 循环

结束跳到这里

8048d0e: c6 45 f7 00	movb \$0x0, -0x9(%ebp)
8048d12: c7 44 24 04 d2 a1 04	movl \$0x804a1d2, 0x4(%esp) 答案
8048d19: 08	
8048d1a: 8d 45 f1	lea -0xf(%ebp), %eax 输入字符串

被转变后放在 ebp-0xf

8048d1d: 89 04 24	mov %eax, (%esp)
8048d20: e8 6d 02 00 00	call 8048f92 <strings_not_equal>
8048d25: 85 c0	test %eax, %eax 相比较, 不等就炸
8048d27: 74 10	je 8048d39 <phase_5+0x63>
8048d29: e8 87 04 00 00	call 80491b5 <explode_bomb>
8048d2e: 66 90	xchg %ax, %ax
8048d30: eb 07	jmp 8048d39 <phase_5+0x63>
8048d32: b8 00 00 00 00	mov \$0x0, %eax eax=0
8048d37: eb bb	jmp 8048cf4 <phase_5+0x1e>
8048d39: 83 c4 24	add \$0x24, %esp
8048d3c: 5b	pop %ebx
8048d3d: 5d	pop %ebp
8048d3e: c3	ret

操作过程或解题思路

```
(gdb) x/s 0x804a1d2
0x804a1d2: "oilers"
```

答案是 oilers

很明显，对于一个字母，它被转变后的字母是唯一的，挨个试就好了

```
phase_5 码表
e->e
o->l
i->f
l->v
r->d
q->a
s->u
d->i
f->r
g->s
oilers
jdoefg
```

第六关

考察内容

链表

反汇编代码分析

```
08048d3f <phase_6>:
8048d3f: 55                push    %ebp
8048d40: 89 e5             mov     %esp, %ebp
8048d42: 56                push    %esi
8048d43: 53                push    %ebx
8048d44: 83 ec 40          sub     $0x40, %esp
8048d47: 8d 45 e0           lea     -0x20(%ebp), %eax
8048d4a: 89 44 24 04        mov     %eax, 0x4(%esp)
8048d4e: 8b 45 08           mov     0x8(%ebp), %eax
8048d51: 89 04 24           mov     %eax, (%esp)
8048d54: e8 9e 04 00 00    call   80491f7 <read_six_numbers>
```

读 6 个数

8048d59: be 00 00 00 00	mov \$0x0,%esi esi 置 0
8048d5e: 8b 44 b5 e0	mov -0x20(%ebp,%esi,4),%eax
8048d62: 83 e8 01	sub \$0x1,%eax
8048d65: 83 f8 05	cmp \$0x5,%eax 大于 5
8048d68: 76 05	jbe 8048d6f <phase_6+0x30>
8048d6a: e8 46 04 00 00	call 80491b5 <explode_bomb>
8048d6f: 83 c6 01	add \$0x1,%esi
8048d72: 83 fe 06	cmp \$0x6,%esi
8048d75: 74 1b	je 8048d92 <phase_6+0x53>
8048d77: 89 f3	mov %esi,%ebx
8048d79: 8b 44 9d e0	mov -0x20(%ebp,%ebx,4),%eax
8048d7d: 39 44 b5 dc	cmp %eax,-0x24(%ebp,%esi,4)
8048d81: 75 05	jne 8048d88 <phase_6+0x49>
8048d83: e8 2d 04 00 00	call 80491b5 <explode_bomb>
8048d88: 83 c3 01	add \$0x1,%ebx
8048d8b: 83 fb 05	cmp \$0x5,%ebx
8048d8e: 7e e9	jle 8048d79 <phase_6+0x3a>
8048d90: eb cc	jmp 8048d5e <phase_6+0x1f>
8048d92: 8d 45 e0	lea -0x20(%ebp),%eax
8048d95: 8d 5d f8	lea -0x8(%ebp),%ebx
8048d98: b9 07 00 00 00	mov \$0x7,%ecx
8048d9d: 89 ca	mov %ecx,%edx
8048d9f: 2b 10	sub (%eax),%edx
8048da1: 89 10	mov %edx,(%eax)
8048da3: 83 c0 04	add \$0x4,%eax
8048da6: 39 d8	cmp %ebx,%eax
8048da8: 75 f3	jne 8048d9d <phase_6+0x5e>
8048daa: bb 00 00 00 00	mov \$0x0,%ebx
8048daf: eb 1d	jmp 8048dce <phase_6+0x8f>
8048db1: 8b 52 08	mov 0x8(%edx),%edx
8048db4: 83 c0 01	add \$0x1,%eax
8048db7: 39 c8	cmp %ecx,%eax
8048db9: 75 f6	jne 8048db1 <phase_6+0x72>
8048dbb: eb 05	jmp 8048dc2 <phase_6+0x83>
8048dbd: ba 54 c1 04 08	mov \$0x804c154,%edx
8048dc2: 89 54 b5 c8	mov %edx,-0x38(%ebp,%esi,4)
8048dc6: 83 c3 01	add \$0x1,%ebx
8048dc9: 83 fb 06	cmp \$0x6,%ebx
8048dcc: 74 17	je 8048de5 <phase_6+0xa6>
8048dce: 89 de	mov %ebx,%esi
8048dd0: 8b 4c 9d e0	mov -0x20(%ebp,%ebx,4),%ecx
8048dd4: 83 f9 01	cmp \$0x1,%ecx
8048dd7: 7e e4	jle 8048dbd <phase_6+0x7e>
8048dd9: b8 01 00 00 00	mov \$0x1,%eax

8048dde: ba 54 c1 04 08	mov	\$0x804c154, %edx	链表的头
8048de3: eb cc	jmp	8048db1 <phase_6+0x72>	
8048de5: 8b 5d c8	mov	-0x38(%ebp), %ebx	
8048de8: 8d 45 cc	lea	-0x34(%ebp), %eax	
8048deb: 8d 75 e0	lea	-0x20(%ebp), %esi	
8048dee: 89 d9	mov	%ebx, %ecx	
8048df0: 8b 10	mov	(%eax), %edx	
8048df2: 89 51 08	mov	%edx, 0x8(%ecx)	
8048df5: 83 c0 04	add	\$0x4, %eax	
8048df8: 39 f0	cmp	%esi, %eax	循环
8048dfa: 74 04	je	8048e00 <phase_6+0xc1>	
8048dfc: 89 d1	mov	%edx, %ecx	
8048dfe: eb f0	jmp	8048df0 <phase_6+0xb1>	
8048e00: c7 42 08 00 00 00 00	movl	\$0x0, 0x8(%edx)	
8048e07: be 05 00 00 00	mov	\$0x5, %esi	
8048e0c: 8b 43 08	mov	0x8(%ebx), %eax	
8048e0f: 8b 00	mov	(%eax), %eax	
8048e11: 39 03	cmp	%eax, (%ebx)	前后项对比，不是
后>前就炸			
8048e13: 7d 05	jge	8048e1a <phase_6+0xdb>	
8048e15: e8 9b 03 00 00	call	80491b5 <explode_bomb>	
8048e1a: 8b 5b 08	mov	0x8(%ebx), %ebx	
8048e1d: 83 ee 01	sub	\$0x1, %esi	
8048e20: 75 ea	jne	8048e0c <phase_6+0xcd>	
8048e22: 83 c4 40	add	\$0x40, %esp	
8048e25: 5b	pop	%ebx	
8048e26: 5e	pop	%esi	
8048e27: 5d	pop	%ebp	
8048e28: c3	ret		

操作过程或解题思路

其实中间大段的 for 循环我是没有看的，因为是没有必要的。重点是这个链表头

```
(gdb) x/d 0x804c154
0x804c154 <node1>:      624

(gdb) x/d 0x804c154+0x4
0x804c158 <node1+4>:      1
(gdb) x/d 0x804c154+0xc
0x804c160 <node2>:      369
(gdb) x/d 0x804c154+0xc+0x4
0x804c164 <node2+4>:      2
```

可以推断出这个链表

```

struct node{
int x;
int no;
node
*next; };

```

x 对于当前地址, no 对应+0x4, *next 对应 0x8
x 和 no 一一对应, 对于每个 no, 都有唯一的 x 相对应。
最后对 x 进行倒叙检查, 后面节点的 x 要小于前面的
现在的任务是排序, 先写出整个对应关系

1->624

2->369

3->472

4->977

5->656

6->359

排序结果为

节点的顺序也和输入数的顺序有关系

1 2 3 4 5 6 输入

4 5 6 1 2 3 映射结束

所以答案是

3 2 6 4 5 1

隐藏关

考察内容

二叉搜索树

进入本关的方法

```

08048e29 <fun7>:
8048e29: 55

```

(其实是看大佬日了 7 个才知道要进入隐藏关)

首先在 phase_defused

这个说明可能是第三和第四关再输一个特殊的字符串

```

(gdb) x/s 0x804c8f0
0x804c8f0 <input_strings+240>: "5 15

```

很明显, 这是第四关的数字

```
(gdb) x/s 0x804a440
0x804a440: "HappyHangHangXueZhang"
```

密码已给出，再四关输入个字符串即可

```
(gdb) x/s 0x804a437
0x804a437: "%d %d %s"
```

反汇编代码分析

08048e7c <secret_phase>:

8048e7c: 55	push %ebp
8048e7d: 89 e5	mov %esp, %ebp
8048e7f: 53	push %ebx
8048e80: 83 ec 14	sub \$0x14, %esp
8048e83: e8 be 03 00 00	call 8049246 <read_line> 读一个数
8048e88: c7 44 24 08 0a 00 00	movl \$0xa, 0x8(%esp)
8048e8f: 00	
8048e90: c7 44 24 04 00 00 00	movl \$0x0, 0x4(%esp)
8048e97: 00	
8048e98: 89 04 24	mov %eax, (%esp)
8048e9b: e8 a0 f9 ff ff	call 8048840 <strtol@plt>
8048ea0: 89 c3	mov %eax, %ebx eax 是我们输入的数
8048ea2: 8d 40 ff	lea -0x1(%eax), %eax eax--
8048ea5: 3d e8 03 00 00	cmp \$0x3e8, %eax eax 要比 0x3e8 小
8048eaa: 76 05	jbe 8048eb1 <secret_phase+0x35>
8048eac: e8 04 03 00 00	call 80491b5 <explode_bomb>
8048eb1: 89 5c 24 04	mov %ebx, 0x4(%esp) 准备参数
8048eb5: c7 04 24 a0 c0 04 08	movl \$0x804c0a0, (%esp) 这个地址放的是 36
8048ebc: e8 68 ff ff ff	call 8048e29 <fun7>
8048ec1: 83 f8 01	cmp \$0x1, %eax 返回值和 1 比较
8048ec4: 74 05	je 8048ecb <secret_phase+0x4f>
8048ec6: e8 ea 02 00 00	call 80491b5 <explode_bomb>
8048ecb: c7 04 24 ac a1 04 08	movl \$0x804a1ac, (%esp)
8048ed2: e8 a9 f8 ff ff	call 8048780 <puts@plt>
8048ed7: e8 a2 04 00 00	call 804937e <phase_defused>
8048edc: 83 c4 14	add \$0x14, %esp
8048edf: 5b	pop %ebx
8048ee0: 5d	pop %ebp
8048ee1: c3	ret
8048ee2: 66 90	xchg %ax, %ax

8048ee4: 66 90	xchg %ax,%ax
8048ee6: 66 90	xchg %ax,%ax
8048ee8: 66 90	xchg %ax,%ax
8048eea: 66 90	xchg %ax,%ax
8048eec: 66 90	xchg %ax,%ax
8048eee: 66 90	xchg %ax,%ax
08048e29 <fun7>:	
8048e29: 55	push %ebp
8048e2a: 89 e5	mov %esp,%ebp
8048e2c: 53	push %ebx
8048e2d: 83 ec 14	sub \$0x14,%esp
8048e30: 8b 55 08	mov 0x8(%ebp),%edx
8048e33: 8b 4d 0c	mov 0xc(%ebp),%ecx
8048e36: 85 d2	test %edx,%edx 放的\$符号, 自与,
zf=0, je 不会跳	
8048e38: 74 37	je 8048e71 <fun7+0x48>
8048e3a: 8b 1a	mov (%edx),%ebx \$码是 36
8048e3c: 39 cb	cmp %ecx,%ebx 我们输入的数和 36
比	
8048e3e: 7e 13	jle 8048e53 <fun7+0x2a>大于等于
就跳	
8048e40: 89 4c 24 04	mov %ecx,0x4(%esp)
8048e44: 8b 42 04	mov 0x4(%edx),%eax
8048e47: 89 04 24	mov %eax,(%esp)
8048e4a: e8 da ff ff ff	call 8048e29 <fun7>
8048e4f: 01 c0	add %eax,%eax
8048e51: eb 23	jmp 8048e76 <fun7+0x4d>
8048e53: b8 00 00 00 00	mov \$0x0,%eax eax 置零
8048e58: 39 cb	cmp %ecx,%ebx 等于就跳
8048e5a: 74 1a	je 8048e76 <fun7+0x4d>
8048e5c: 89 4c 24 04	mov %ecx,0x4(%esp) 准备参数 50
8048e60: 8b 42 08	mov 0x8(%edx),%eax
8048e63: 89 04 24	mov %eax,(%esp) “2” 码为 50
8048e66: e8 be ff ff ff	call 8048e29 <fun7>
8048e6b: 8d 44 00 01	lea 0x1(%eax,%eax,1),%eax
8048e6f: eb 05	jmp 8048e76 <fun7+0x4d>
8048e71: b8 ff ff ff ff	mov \$0xffffffff,%eax
8048e76: 83 c4 14	add \$0x14,%esp
8048e79: 5b	pop %ebx
8048e7a: 5d	pop %ebp
8048e7b: c3	ret

操作过程或解题思路

`fun7 (int x, char c)`

x 是要寻找的数, c 是 ascii 是左边或者右边的数

左边是 36, 右边是 50

返回值可能是 bool 0~1 表示找没找到

也有可能是寻找的次数

所以我先输了 36, 发现 `secret_phase` 的 `eax` 是 0

所以, 答案是 50

思考与体会

对汇编代码和 gdb 调试程序更加熟练。

遇到不会的可以跳过, 把它当成黑箱, 不用关心它的原理, 只关心它的输入和输出, 再通过输入输出寻找关系, 逆推实现原理。

`cmp` 附近有 `bomb` 函数, 突破点往往在这里。

回忆起大一上也做过这个实验, 自己运气好蒙对了 5 关。这次很顺畅的通完 7 关, 我变强了。

6	bomb3	Sun Jun 10 16:49	7	16	69	valid
---	-------	------------------	---	----	----	-------