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主要内容

本次实验的基础版本是使用libfuse实现一个只读的FAT16文件系统。

实验步骤

助教已经给出了一个基础版本的代码,只需要填充9个部分的内容即可实现这个文件系统。因为接口已经事先给好,所以这个实验总的来说还是比较简单的。在搭建好实验环境后,按照需求编写代码就可以完成实验(而且debug的体验比lab3好很多)。

核心代码解析

1. path_split():

```
1
    char **path split(char *pathInput, int *pathDepth ret)
 2
 3
      int i,j;
 4
      int pathDepth = 0;
 5
      for (i=0;pathInput[i]!='\0';++i)
 7
        if (pathInput[i]=='/')
8
9
          pathDepth++;
        }
10
11
      char **paths = (char **)malloc(pathDepth * sizeof(char *));
12
      for (i=0,j=1;i<pathDepth;++i)</pre>
13
14
        paths[i]=(char *)malloc(12 * sizeof(char));
15
        int k1=-1;
16
        int k;
17
18
        for (k=j;pathInput[k]!='\0' && pathInput[k]!='/';++k)
19
20
          if (pathInput[k]=='.')
21
          {
            k1=k;
22
23
        }//找点号所在的位置, 若没有则为-1
24
        if (k1==-1)
25
2.6
        {
27
          int 1;
28
          for (1=0;1<12;++1)
```

```
29
30
             if (1<8)
31
             {
               if (j<k)
32
33
34
                 paths[i][l]=upper(pathInput[j]);
35
                 ++j;
36
               }
37
               else
                 paths[i][l]=' ';
38
39
40
             else if (1>=8 && 1<11)
41
42
               paths[i][1]=' ';
43
44
             else
45
               paths[i][1]='\0';
46
47
             }
48
           }
49
         }
50
        else
51
        {
52
           int 1;
53
           for (1=0;1<12;++1)
54
             if (1<8)
55
56
               if (j<k1)
57
58
59
                 paths[i][l]=upper(pathInput[j]);
60
                 ++j;
61
               }
62
               else
                 paths[i][l]=' ';
63
64
65
             else if (1==8)
66
67
               j=k1+1;
68
               if (j<k)
69
70
                 paths[i][l]=upper(pathInput[j]);
71
72
               }
73
               else
                 paths[i][1]=' ';
74
75
             }
             else if (1>8 && 1<11)
76
77
             {
```

```
78
               if (j < k)
79
80
                 paths[i][l]=upper(pathInput[j]);
81
                 j++;
82
              else
83
84
                 paths[i][1]=' ';
85
86
            else
              paths[i][1]='\0';
87
88
         }
89
       }
90
        j=k+1;
91
92
      *pathDepth ret=pathDepth;
93
      return paths;
94
   }
```

对于我来讲,其实整个实验部分最难的两个代码就是前两个对字符串处理的函数(菜的真实)。要 实现对路径的分割,首先需要统计路径中出现的所有的 / ,这个就代表目录的深度。其次,再遍历 每个字路径中所有的自负,若名称中存在扩展名,则将名称中点号之后的字符复制到格式化名称的 后三位,否则只需要复制前八位即可。如此就可以实现路径分割。

2. path decode()

```
BYTE *path_decode(BYTE *path)
 2
 3
 4
      BYTE *pathDecoded = malloc(MAX_SHORT_NAME_LEN * sizeof(BYTE));
 5
      int i,j;
      for (i=0; i<8; ++i)
 6
 7
       if (path[i]!=' ')
8
9
10
          pathDecoded[i]=lower((char)path[i]);
11
        }
12
        else
13
         break;
14
      if (path[8]==' ')//无扩展名
15
16
17
       pathDecoded[i]='\0';
18
      }
19
      else {
        pathDecoded[i]='.';
20
21
        ++i;
22
       for (j=8; j<11; ++j)
23
         if (path[j]!=' ')
24
```

```
25
26
            pathDecoded[i]=lower((char)path[j]);
27
            ++i;
28
          }
29
        }
        pathDecoded[i]='\0';
30
31
32
      return pathDecoded;
33
    }
```

将路径解码的过程要相对简单的多,因为编码是是规范的,所以只需要将编码名称多前八位进行复制,再根据是否有扩展名来确定扩展名部分的内容是否需要复制即可。

3. pre init fat16()

```
FAT16 *pre init fat16(void)
 2
 3
      /* Opening the FAT16 image file */
 4
      FILE *fd;
      FAT16 *fat16 ins;
 5
 6
 7
      fd = fopen(FAT FILE NAME, "rb");
8
9
      if (fd == NULL)
10
      {
        fprintf(stderr, "Missing FAT16 image file!\n");
11
        exit(EXIT_FAILURE);
12
13
      }
14
15
      fat16 ins = malloc(sizeof(FAT16));
      fat16_ins->fd = fd;
16
17
      /** TODO:
18
      * 初始化fat16_ins的其余成员变量
19
      * Hint: root directory的大小与Bpb.BPB RootEntCnt有关、并且是扇区对齐的
20
      **/
21
22
      char buffer[BYTES_PER_SECTOR];
23
      sector_read(fd, 0, (void *)buffer);
24
     int i;
25
      for (i=0; i<3; ++i){
26
        fat16_ins->Bpb.BS_jmpBoot[i]=buffer[i];
27
      }
28
      for (i=0; i<8; ++i) {
29
        fat16_ins->Bpb.BS_OEMName[i]=buffer[0x3+i];
30
      }
      fat16 ins->Bpb.BPB BytsPerSec = *((WORD *)(&buffer[0xb]));
31
32
      fat16_ins->Bpb.BPB_SecPerClus = buffer[0xd];
      fat16_ins->Bpb.BPB_RsvdSecCnt = *((WORD *)(&buffer[0xe]));
33
      fat16 ins->Bpb.BPB NumFATS = buffer[0x10];
34
```

```
35
      fat16 ins->Bpb.BPB RootEntCnt = *((WORD *)(&buffer[0x11]));
36
      fat16 ins->Bpb.BPB TotSec16 = *((WORD *)(&buffer[0x13]));
37
      fat16 ins->Bpb.BPB Media = buffer[0x15];
      fat16_ins->Bpb.BPB_FATSz16 = *((WORD *)(&buffer[0x16]));
38
      fat16 ins->Bpb.BPB SecPerTrk = *((WORD *)(&buffer[0x18]));
39
      fat16 ins->Bpb.BPB NumHeads = *((WORD *)(&buffer[0x1a]));
40
41
      fat16 ins->Bpb.BPB HiddSec = *((DWORD *)(&buffer[0x1c]));
42
      fat16 ins->Bpb.BPB TotSec32 = *((DWORD *)(&buffer[0x20]));
      fat16_ins->Bpb.BS_DrvNum = buffer[0x24];
43
      fat16 ins->Bpb.BS Reserved1 = buffer[0x25];
44
45
     fat16_ins->Bpb.BS_BootSig = buffer[0x26];
      fat16 ins->Bpb.BS VollID = *((DWORD *)(&buffer[0x27]));
46
      for (i=0;i<11;++i){
47
        fat16_ins->Bpb.BS_VollLab[i] = buffer[0x2b+i];
48
49
      for (i=0;i<8;++i){
5.0
51
        fat16 ins->Bpb.BS FilSysType[i] = buffer[0x36+i];
52
      }
      for (i=0; i<448; ++i){
53
54
        fat16 ins->Bpb.Reserved2[i] = buffer[i+0x3e];
55
56
      fat16 ins->Bpb.Signature word = *((WORD *)(&buffer[0x1fe]));
57
      fat16 ins->FirstRootDirSecNum = fat16 ins->Bpb.BPB RsvdSecCnt+
5.8
    (fat16 ins->Bpb.BPB NumFATS)*(fat16 ins->Bpb.BPB FATSz16);
      fat16_ins->FirstDataSector = fat16_ins->FirstRootDirSecNum+32*
59
    (fat16 ins->Bpb.BPB RootEntCnt)/(fat16 ins->Bpb.BPB BytsPerSec);
60
61
      return fat16_ins;
62
    }
```

这部分的内容同样十分简单。在我一开始编写的时候,我不是很清楚结构体中每个变量代表的内容,但是对照首个扇区的内容后我发现,结构体声明的顺序与首个扇区的里面填充的内容顺序相同,只需要将第一个扇区中的抄到结构体中即可。

4. fat entry by cluster()

```
WORD fat entry by cluster(FAT16 *fat16 ins, WORD ClusterN)
2
     BYTE sector buffer[BYTES PER SECTOR];
3
4
     DWORD Fat1 Start = fat16 ins->Bpb.BPB RsvdSecCnt;
     WORD Fat1_Large = fat16_ins->Bpb.BPB_FATSz16;//可以在检查是否越界的时候
5
   使用, 以增加安全性
     WORD Sec_Large = fat16_ins->Bpb.BPB_BytsPerSec;
6
7
     //需要寻找对应的扇区来找到对应簇的后继,首先要确定ClusterN簇号存储在哪一个扇
8
   区,再确定偏移了多少;
     //因为使用两个字节来存储一个簇的簇号所以...
9
     int Sec Shift = (ClusterN*2)/Sec_Large;
10
```

```
int Byte_Shift = (ClusterN*2)%Sec_Large;
sector_read(fat16_ins->fd, (Fat1_Start+Sec_Shift), sector_buffer);
WORD result = *((WORD *)(&sector_buffer[Byte_Shift]));

return result;
}
```

FAT表中每个项占两个字节,所以要想知道簇N对应的下一个块,就需要先根据要查找的簇号算出扇区的偏移,读取对应扇区的内容,再算出扇区内字节的偏移,就得到了要查找的簇号和后继。

5. find root()

```
1
    for (i = 0; i < fat16_ins->Bpb.BPB_RootEntCnt; i++)
 2
        //先判断是否超出当前扇区;
 3
        if (i*32 >= (RootDirCnt+1)*BYTES PER SECTOR){
 4
 5
          RootDirCnt++;
 6
          sector_read(fat16_ins->fd, (fat16_ins-
    >FirstRootDirSecNum) +RootDirCnt, buffer);
7
8
        char Name Buffer[12];
9
        int Start Read=(i*32)%BYTES PER SECTOR;
10
        for (j=0; j<11; ++j){
         Name Buffer[j] = buffer[Start Read+j];
11
12
        }
        Name_Buffer[12]='\0';
13
        if (strncmp(Name Buffer, paths[0],11)==0){
14
15
          for (j=0; j<11; ++j) {
            Dir->DIR Name[j] = Name Buffer[j];
16
17
          }
18
          Dir->DIR Attr = buffer[Start Read+0x0b];
19
          Dir->DIR NTRes = buffer[Start Read+0x0c];
          Dir->DIR_CrtTimeTenth = buffer[Start_Read+0x0d];
20
21
          Dir->DIR CrtTime = *((WORD *)(&buffer[Start Read+0x0e]));
          Dir->DIR CrtDate = *((WORD *)(&buffer[Start Read+0x10]));
2.2
          Dir->DIR_LstAccDate = *((WORD *)(&buffer[Start_Read+0x12]));
23
          Dir->DIR_FstClusHI = *((WORD *)(&buffer[Start_Read+0x14]));
24
25
          Dir->DIR WrtTime = *((WORD *)(&buffer[Start Read+0x16]));
26
          Dir->DIR WrtDate = *((WORD *)(&buffer[Start Read+0x18]));
27
          Dir->DIR_FstClusLO = *((WORD *)(&buffer[Start_Read+0x1a]));
28
          Dir->DIR FileSize = *((DWORD *)(&buffer[Start Read+0x1c]));
          if (pathDepth==1) {
2.9
            return 0;
30
31
          }
32
          else{
33
            return find_subdir(fat16_ins, Dir, paths, pathDepth, 1);
          }
34
35
        }
```

这个函数的核心代码如上。扇区的根目录中每项有32个byte,共有512项,所以只要按照每个项查找匹配的名称,查找到后,再根据提供的目录的深度来判断是否需要调用寻找子目录查找的函数,如果不需要就直接返回文件的信息,需要的话就进行下一步查找。

6. find subdir

```
int find_subdir(FAT16 *fat16_ins, DIR_ENTRY *Dir, char **paths, int
    pathDepth, int curDepth)
 2.
      BYTE buffer[BYTES PER SECTOR];
 3
 4
      if (curDepth<pathDepth) {</pre>
 5
        int i, j, k;//i簇内扇区偏移,j为扇区内字节偏移
        WORD ClusterN = Dir->DIR FstClusLO;
 6
 7
        WORD FatClusEntryVal, FirstSectorofCluster;
        //第一个变量是当前簇的后继簇号,第二个是该簇的第一个扇区号。
 8
 9
     first_sector_by_cluster(fat16_ins,ClusterN,&FatClusEntryVal,&FirstSec
    torofCluster, buffer);
        //开始进行查找
10
        while (ClusterN >= 0x0002 && ClusterN <= 0xffef){
11
12
     first_sector_by_cluster(fat16_ins,ClusterN,&FatClusEntryVal,&FirstSec
    torofCluster, buffer);
          for (i=0;i<fat16_ins->Bpb.BPB_SecPerClus;++i){
13
            sector read(fat16 ins->fd, FirstSectorofCluster+i, buffer);
14
            for (j=0; j<BYTES PER SECTOR; j+=32) {</pre>
15
16
              char Name Buffer[12];
              for (k=0; k<11; ++k) {
17
                Name_Buffer[k] = buffer[j+k];
18
19
              }
20
              Name Buffer[12]='\0';
21
              if (strncmp(Name_Buffer, paths[curDepth],11)==0){
                for (k=0; k<11; ++k){
2.2
23
                  Dir->DIR Name[k] = Name Buffer[k];
24
                int Start_Read=j;
2.5
26
                Dir->DIR Attr = buffer[Start Read+0x0b];
                Dir->DIR NTRes = buffer[Start Read+0x0c];
27
                Dir->DIR CrtTimeTenth = buffer[Start Read+0x0d];
28
2.9
                Dir->DIR_CrtTime = *((WORD *)(&buffer[Start_Read+0x0e]));
30
                Dir->DIR CrtDate = *((WORD *)(&buffer[Start Read+0x10]));
                Dir->DIR_LstAccDate = *((WORD *)
31
    (&buffer[Start Read+0x12]));
                Dir->DIR FstClusHI = *((WORD *)
32
    (&buffer[Start_Read+0x14]));
33
                Dir->DIR_WrtTime = *((WORD *)(&buffer[Start_Read+0x16]));
                Dir->DIR WrtDate = *((WORD *)(&buffer[Start Read+0x18]));
34
```

```
35
                Dir->DIR_FstClusLO = *((WORD *)
    (&buffer[Start_Read+0x1a]));
36
                Dir->DIR FileSize = *((DWORD *)
    (&buffer[Start_Read+0x1c]));
37
                return find_subdir(fat16_ins, Dir, paths, pathDepth,
    curDepth+1);
38
39
              if (Name_Buffer[0]==0 && Name_Buffer[1]==0){
                return 1;
40
41
42
            }
43
          }
          ClusterN = FatClusEntryVal;
44
45
46
       }
47
      else if (curDepth == pathDepth){
48
49
      return 0;
50
      }
51
      return 1;
52
    }
```

这个函数也是实现中比较简单的一个函数。其思路是按照目录的深度进行递归查找,只需通过判断 当前查找的目录深度来判断是否需要进行下一步递归或者返回值。需要注意的是,这个与 find_root不是很相同,原因在于。对于findroot函数,其文件中存储的内容在物理上是连续的,所 以只需要按照顺序检查每个扇区内是否有符合要求的文件名。但是对于寻找子目录这个函数,由于 子目录没有特定的大小限制,所以是有可能跨簇寻找匹配项的,所以还需要考虑到是否读取到文件 结束,最后的代码如上。

7. fat16 readdir()

```
int fat16 readdir(const char *path, void *buffer, fuse fill dir t
    filler,
 2
                      off t offset, struct fuse file info *fi)
 3
 4
     FAT16 *fat16_ins;
     BYTE sector_buffer[BYTES_PER_SECTOR];
 5
 6
     struct fuse_context *context;
 8
     context = fuse_get_context();
9
     fat16 ins = (FAT16 *)context->private data;
10
11
      sector_read(fat16_ins->fd, fat16_ins->FirstRootDirSecNum,
    sector_buffer);
12
13
      if (strcmp(path, "/") == 0)
14
15
        DIR ENTRY Root;
16
        int i,j;
```

```
17
        int RootDirCnt = 0;
18
19
        /** TODO:
         * 将root directory下的文件或目录通过filler填充到buffer中
20
         * 注意不需要遍历子目录
21
        **/
22
2.3
24
        for (i = 0; i < fat16 ins->Bpb.BPB RootEntCnt; i++)
25
          if (i*32 >= (RootDirCnt+1)*BYTES PER SECTOR){
2.6
27
            RootDirCnt++;
            sector read(fat16 ins->fd, (fat16 ins-
28
    >FirstRootDirSecNum) + RootDirCnt, sector buffer);
29
30
          BYTE Name Buffer[12];
          int Start_Read=(i*32)%BYTES_PER_SECTOR;
31
          for (j=0; j<11; ++j) {
32
            Name_Buffer[j] = sector_buffer[Start_Read+j];
33
          }
34
35
          Name Buffer[12]='\0';
36
          if ((BYTE)Name_Buffer[0]==0 && (BYTE)Name_Buffer[1]==0){
37
            break;
38
          }//根目录遍历结束
39
40
          for (j=0;j<11;++j){
41
42
            Root.DIR Name[j] = Name Buffer[j];
          }
43
          Root.DIR_Attr = sector_buffer[Start_Read+0x0b];
44
          Root.DIR NTRes = sector buffer[Start Read+0x0c];
45
          Root.DIR CrtTimeTenth = sector buffer[Start Read+0x0d];
46
          Root.DIR_CrtTime = *((WORD *)
47
    (&sector buffer[Start Read+0x0e]));
          Root.DIR CrtDate = *((WORD *)
48
    (&sector buffer[Start Read+0x10]));
          Root.DIR_LstAccDate = *((WORD *)
49
    (&sector buffer[Start Read+0x12]));
          Root.DIR FstClusHI = *((WORD *)
50
    (&sector_buffer[Start_Read+0x14]));
51
          Root.DIR WrtTime = *((WORD *)
    (&sector buffer[Start Read+0x16]));
52
          Root.DIR_WrtDate = *((WORD *)
    (&sector buffer[Start Read+0x18]));
          Root.DIR FstClusLO = *((WORD *)
53
    (&sector buffer[Start Read+0x1a]));
54
          Root.DIR FileSize = *((DWORD *)
    (&sector_buffer[Start_Read+0x1c]));
55
```

```
if ((BYTE)Name_Buffer[0]==0x00 | (BYTE)Name_Buffer[0]==0xe5 |
    (BYTE)Root.DIR Attr==0x0f){
57
           continue;
58
          }
          //printf("%x\n", Name_Buffer[0]);
59
          const char *filename = (const char
60
    *)path decode(Root.DIR Name);
61
          filler(buffer, filename, NULL, 0);
62
           * const char *filename = (const char
6.3
    *)path_decode(Root.DIR_Name);
          * filler(buffer, filename, NULL, 0);
64
          **/
65
66
67
        }
68
      }
69
      else
70
      {
71
       DIR_ENTRY Dir;
72
       int i,j,k;
73
       WORD ClusterN, FatClusEntryVal, FirstSectorofCluster;
74
75
       /** TODO:
76
        * 通过find root获取path对应的目录的目录项,
77
         * 然后访问该目录,将其下的文件或目录通过filler填充到buffer中,
        * 同样注意不需要遍历子目录
78
         * Hint: 需要考虑目录大小,可能跨扇区,跨簇
79
        **/
80
        find_root(fat16_ins, &Dir, path);
81
        ClusterN = Dir.DIR FstClusLO;
82
83
     first_sector_by_cluster(fat16_ins,ClusterN,&FatClusEntryVal,&FirstSe
    ctorofCluster, sector buffer);
84
       int end flag=0;
85
86
        while (ClusterN >= 0x0002 && ClusterN <= 0xffef){
87
     first sector by cluster(fat16 ins,ClusterN,&FatClusEntryVal,&FirstSe
    ctorofCluster, sector buffer);
88
          for (i=0;i<fat16 ins->Bpb.BPB SecPerClus;++i){
            sector read(fat16 ins-
89
    >fd,FirstSectorofCluster+i,sector_buffer);
            for (j=0;j<BYTES_PER_SECTOR;j+=32){</pre>
90
              char Name Buffer[12];
91
              for (k=0; k<11; ++k) {
92
93
                Name_Buffer[k] = sector_buffer[j+k];
94
              }
              Name Buffer[12]='\0';
95
96
              if ((BYTE)Name_Buffer[0]==0 && (BYTE)Name_Buffer[1]==0){
```

```
97
                 end_flag=1;
 98
                 break;
               }//当前目录遍历结束
99
100
               for (k=0; k<11; ++k) {
101
                   Dir.DIR_Name[k] = Name_Buffer[k];
102
103
               int Start Read=j;
104
               Dir.DIR Attr = sector buffer[Start Read+0x0b];
105
               Dir.DIR_NTRes = sector_buffer[Start_Read+0x0c];
               Dir.DIR CrtTimeTenth = sector buffer[Start Read+0x0d];
106
               Dir.DIR_CrtTime = *((WORD *)
107
     (&sector buffer[Start Read+0x0e]));
108
               Dir.DIR CrtDate = *((WORD *)
     (&sector_buffer[Start_Read+0x10]));
               Dir.DIR LstAccDate = *((WORD *)
109
     (&sector buffer[Start Read+0x12]));
               Dir.DIR FstClusHI = *((WORD *)
110
     (&sector_buffer[Start_Read+0x14]));
               Dir.DIR_WrtTime = *((WORD *)
111
     (&sector buffer[Start Read+0x16]));
112
               Dir.DIR_WrtDate = *((WORD *)
     (&sector buffer[Start Read+0x18]));
113
               Dir.DIR FstClusLO = *((WORD *)
     (&sector buffer[Start Read+0x1a]));
               Dir.DIR FileSize = *((DWORD *)
114
     (&sector_buffer[Start_Read+0x1c]));
115
116
               if ((BYTE)Name_Buffer[0]==0x00 |
     (BYTE)Name_Buffer[0]==0xe5 | (BYTE)Dir.DIR_Attr==0x0f){
117
                 continue:
118
               }
119
               const char *filename = (const char
     *)path decode(Dir.DIR Name);
               filler(buffer, filename, NULL, 0);
120
121
             }
122
             if (end_flag==1){
123
               break;
124
             }
125
126
           if (end flag==1){
             break;
127
128
129
           ClusterN = FatClusEntryVal;
130
         }
131
       }
132
       return 0;
133 }
```

这个函数的实现更是非常简单,两个部分中所有的逻辑与 find_root 和 find_subdir 完全一致,只不过是把查找特定的文件或者目录改成将所有的文件和目录的名字录入。但是,也不是完全相同的,因为助教给出的镜像中是包括长文件名的镜像,所以就需要判断当前在读的项是LFN项还是文件项。在上课的时候讲到,LFN的文件属性是特定的 0x0F 所以可以根据这个来识别是否是LFN。同样还需要判断这个文件是不是被删除了,所以需要判断文件名的第一个字符,判断是不是无效。

8. fat16_read()

```
int fat16 read(const char *path, char *buffer, size t size, off t
    offset,
 2
                   struct fuse file info *fi)
 3
    {
 4
     FAT16 *fat16_ins;
 5
      struct fuse context *context;
 6
      context = fuse get context();
7
     fat16_ins = (FAT16 *)context->private_data;
 8
      DIR ENTRY File;
9
      BYTE sector buffer[BYTES PER SECTOR];
      WORD ClusterN, FatClusEntryVal, FirstSectorofCluster;
10
11
     find_root(fat16_ins, &File, path);
      ClusterN = File.DIR FstClusLO;
12
13
     DWORD File Size = File.DIR FileSize;
     if ((DWORD)offset >= File Size){
14
15
      return 0;
     }
16
17
      first sector by cluster(fat16 ins,ClusterN,&FatClusEntryVal,&FirstSe
    ctorofCluster, sector buffer);
      int i;
18
19
      int Cluster Shift = (offset)/((fat16 ins->Bpb.BPB SecPerClus)*
    (fat16 ins->Bpb.BPB BytsPerSec));
2.0
      for (i=0;i<Cluster Shift;++i){</pre>
        ClusterN = FatClusEntryVal;
2.1
22
     first_sector_by_cluster(fat16_ins,ClusterN,&FatClusEntryVal,&FirstSec
    torofCluster, sector buffer);
23
      }
24
      DWORD Real Read Size = ((DWORD)offset+(DWORD)size)<=File Size ?</pre>
    (DWORD)size : (DWORD)File_Size-(DWORD)offset;
      DWORD Start Sector = (offset - Cluster Shift*(fat16 ins-
26
    >Bpb.BPB_SecPerClus)*(fat16_ins->Bpb.BPB_BytsPerSec))/fat16_ins-
    >Bpb.BPB_BytsPerSec;
      DWORD Start Byte = offset - Cluster Shift*(fat16 ins-
2.7
    >Bpb.BPB_SecPerClus)*(fat16_ins->Bpb.BPB_BytsPerSec) - Start_Sector*
    (fat16 ins->Bpb.BPB BytsPerSec);
28
29
30
      int CurSector = FirstSectorofCluster + Start Sector;
      sector read(fat16 ins->fd, CurSector, sector buffer);
31
```

```
32
33
34
      for (i=0;i<Real Read Size;++i,++Start Byte){</pre>
        if ((Start_Byte) >= BYTES_PER_SECTOR){
35
36
          Start_Byte = 0;
          if (CurSector+1-FirstSectorofCluster>=fat16 ins-
37
    >Bpb.BPB SecPerClus) {
38
            ClusterN = FatClusEntryVal;
39
     first sector by cluster(fat16 ins,ClusterN,&FatClusEntryVal,&FirstSec
    torofCluster,sector_buffer);
            CurSector = FirstSectorofCluster;
40
41
          }
42
          else {
43
            CurSector++;
44
            sector_read(fat16_ins->fd, CurSector, sector_buffer);
45
          }
46
        }
47
        buffer[i] = sector_buffer[Start_Byte];
48
49
      return (int)Real_Read_Size;
50
```

这个函数实现的细节在于首先需要根据offset算出要读的位置的簇号、簇内扇区的偏移和扇区内字 节偏移。然后将这个之后的real_size内容复制到buffer中。

实验截图

1. 实验测试截图

```
🛊 拼 🗊 12:58 🔆
 🙆 🖨 🗊 root@pc: ~/OSLab_Ubuntu1404/oslab/lab4/lab4-code
success in test_path_decode
#3 running test_pre_init_fat16 success in test_pre_init_fat16
#4 running test_fat_entry_by_cluster
test case 1: OK
test case 2: OK
test case 3: OK
success in test_fat_entry_by_cluster
#5 running test_find_root
test case 1: OK
test case 2: OK
test case 3: OK
success in test_find_root
#6 running test_find_subdir
test case 1: OK
test case 2: OK
test case 3: OK
success in test_find_subdir
root@pc:~/OSLab_Ubuntu1404/oslab/lab4/lab4-code#
```

2. 读取pdf截图



实验要点

本次实验的基础部分说句实话不是很难,只需要时刻注意当前读取的数据到底是在哪个簇、哪个扇区,哪个字节即可,当读取的指针变动的时候,也要同时进行这些的更新,才能保证不会出错。

实验总结

- 通过这次实验,我对FAT这一类的文件系统有了更深刻的了解,对其基本的结构和实现有了更加深刻的认识。
- 发现了多线程对于实验结果的干扰,帮助助教找出了实验指导书中的一个错误。