Lenstra Elliptic Curve Factorization

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- ECM is third-fastest known factoring algorithm and the best algorithm for finding divisors not exceeding 50-60 digits.
- The largest factor found using ECM has 83 digits.

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- Observe f and g are inverses.

$$f(g(h)) = f(ah^{-1}) = (ah^{-1})^{-1}a = h$$

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- Therefore, |G| = n|H|.



A different perspective

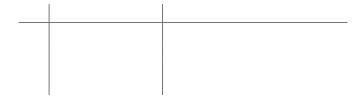
Lemma 2.2.5 Suppose that $m, n \in \mathbb{N}$ and gcd(m, n) = 1. Then the map

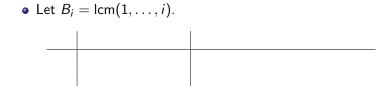
$$\psi: (\mathbb{Z}/\mathsf{mn}\mathbb{Z})^* \to (\mathbb{Z}/\mathsf{m}\mathbb{Z})^* \times (\mathbb{Z}/\mathsf{n}\mathbb{Z})^*$$

defined by

$$\psi(c) = (c \pmod{m}, c \pmod{n})$$

is a bijection.





• Let $B_i = \operatorname{lcm}(1, \ldots, i)$.

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6	570	(37, 11)
60	575	(1, 16)

• We compute gcd(574, 1763) = 41

• Let E be an elliptic curve over $\mathbb{Z}/N\mathbb{Z}$ of the form

$$y^2 = x^3 + ax + 1$$

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• Definition 6.3.1 (Power Smooth). Let B be a positive integer. If n is a positive integer with prime factorization

$$n=\prod p_i^{e_i},$$

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• Example $30 = 2 \cdot 3 \cdot 5$ is B power smooth for $B \ge 5$, but $150 = 2 \cdot 3 \cdot 5^2$ is not 5-power smooth.

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- The idea of ECM is to replace modular exponentiation on $(\mathbb{Z}/N\mathbb{Z})^*$ by repeated addition of points on $E((\mathbb{Z}/N\mathbb{Z})^*)$
- Recall, by the Hasse-Weil bound we can reduce the size of our group by $2 \cdot \sqrt{p}$.





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	1
	1

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Table: Let E be an elliptic curve, and m = lcm(1, 2, ..., B) for some B

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- If Pollard p-1 fails, we have no choice but to increase B.
- However, ECM has a second option. We can choose another random elliptic curve.



Why ECM "Works"

We can consider an analogous mapping

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$$g: E(\mathbb{Z}/N\mathbb{Z}) \to \prod_{p|N} E(\mathbb{Z}/p\mathbb{Z})$$
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where p are prime divisors of N.

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• Note the quotations. There is a subtly in the difference between $E(\mathbb{Z}/N\mathbb{Z})$ and $\mathbb{Z}/N\mathbb{Z}$.

Implementation

- Generate a random elliptic curve $E \pmod{N}$ and let P = (0,1).
- Compute m = lcm(1, 2, ..., B).
- Compute mP (don't be naive!).
- If the calculation fails, you have found a non-trivial factor of N.
- Otherwise, just generate a new Elliptic curve and try again.

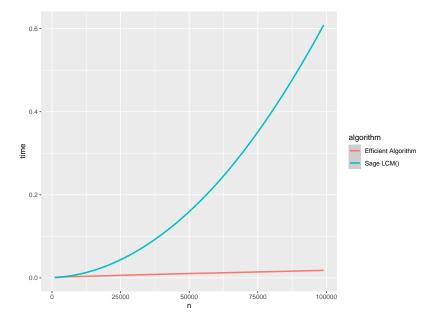
Computing lcm(1,2,...,B)

Recall,

$$lcm(1,2,...B) = \prod_{p \in P} p^r$$

where $r = \max\{r \in \mathbb{Z} \mid p^r \leq B\}$.

$$p^r \le B$$
 $r \log(p) \le \log(B)$
 $r \le \log_p(B)$
 $r = \lfloor \log_p(B) \rfloor$



Computing mP

$$mP = \overbrace{P + P + P \dots P}^{m \text{ times}}$$

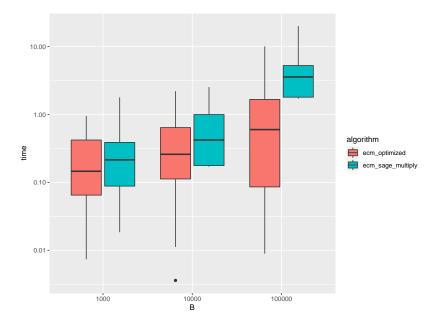
A very bad way to compute mP.

There are many algorithms for computing general elliptic curve point multiplication efficiently, but given the very specific make-up of m, we can save time by being thoughtful here. Consider,

$$m_n=q_1^{r_1}\cdot q_2^{r_2}\ldots q_n^{r_n}$$

then

$$m_n P = q_n^{r_n} \cdot m_{n-1} P$$



Coded Example

```
1 def ecm(n, B=10^4, trials=100):
      R = Zmod(n)
2
      primes = list(prime_range(B+1))
4
      for _ in range(trials):
5
           while True:
6
               a = R.random_element()
7
               if gcd(4 * Integer(a)^3 + 27, n) == 1:
8
                   break
9
10
          E = EllipticCurve([a, 1])
          P = E([0,1])
13
          try:
14
               for p in primes:
15
                   P = P * p^floor(math.log(B,p))
16
           except ZeroDivisionError as e:
18
               return gcd(Integer(str(e).split()[2]), n)
19
20
21
      return -1
```