

University of Central Florida
Department of Computer Science
COP 3402: System Software
Fall 2020

Homework #3 (Parser- Code Generator)

Due Sunday, November 8th, 2020 by 11:59 p.m.

This is a team project (Same team)

REQUIRMENT:

All assignments must compile and run on the Eustis server. Please see course website for details concerning use of Eustis.

Objective:

In this assignment, you must implement a Recursive Descent Parser and an Intermediate Code Generator for tiny PL/0. In addition, you must create a compiler driver to combine all of the compiler parts into one single program.

Your compiler driver must support the following compiler directives:

- l : print the list of lexemes/tokens (scanner output) to the screen
- a : print the generated assembly code (parser/codegen output) to the screen
- v : print virtual machine execution trace (virtual machine output) to the screen
- <filename>.txt : input file name, for e.g. input.txt

Example commands:

- | | |
|------------------------------|--|
| ./compile input.txt -l -a -v | Print all types of output to the console |
| ./compile input.txt -v | Print only the VM execution trace to the console |
| ./compile input.txt | Print nothing to the console except for “in” and “out” |

Note: You may want to also print all forms of output to a single output file in order to avoid losing points on the other requirements if your implementation of compiler directives does not work.

Example of a program written in PL/0:

```
var x, w;  
begin  
  x:= 4;  
  read w;  
  if w > x then  
    w:= w + 1  
  write w;  
end
```

Component Descriptions:

The **compiler driver** is a program that manages the parts of the compiler. It must handle the input, output, and execution of the Scanner (HW2), the Parser (HW3), the Intermediate Code Generator (HW3) and the Virtual Machine (HW1).

The **Parser** is a program that reads in the output of the Scanner (HW2) and parses the lexemes (tokens). It must be capable of reading in the tokens produced by your Scanner (HW2) and produce, as output, a message that states whether the PL/0 program is well-formed (syntactically correct) if it follows the grammar rules. (see Grammar in Appendix B). Otherwise, if the program does not follow the grammar, a message indicating the type of error present must be printed (**reminder: if the scanner detects an error the compilation process must stop and the error must be indicated**), the compilation process must stop). A list of the some errors that might be considered can be found in Appendix C. In addition, the Parser must fill out the Symbol Table, which contains all of the variables, procedure and constants names within the PL/0 program. See Appendix E for more information regarding the Symbol Table. If the program is syntactically correct and the Symbol Table is created without error, the execution of the compiler driver continues with intermediate code generation.

The **Intermediate Code Generator** is a program that takes, as input, the output from the Parser, i.e. the Symbol Table and parsed code. As output, it produces the assembly language for your Virtual Machine (HW1). This functionality may be interleaved with the Parser functionality (i.e. generate assembly code as you parse the token list). Once the code has been generated for your Virtual Machine, the execution of the compiler driver continues by executing the generated assembly code on your Virtual Machine

Submission Instructions:

1.- Submit via WebCourses:

1. Source code of the tiny- PL/0 compiler.
2. A text file with instructions on how to use your program entitled readme.txt.
3. A text file composed of the input file to your Scanner and the output of your Parser/codegen to demonstrate a correctly formed tiny- PL/0 program. The Parser/codegen output should indicate the program is syntactically correct. Following the statement that the program is syntactically correct, the text file should contain the generated code from your intermediate code generator and the stack output from your Virtual Machine running your code.
4. A text file composed of the input file to your Scanner and the output of your Parser/codegen to demonstrate all possible errors. This may require many runs and the Parser output should indicate which error is being identified.
5. All files should be compressed into a single .zip format.
6. Late assignments will not be accepted.

Appendix A:

Traces of Execution:

Example 1, if the input is:

```
var x, y;  
begin  
  x := y + 56;  
end.
```

The output should look like:

1.- A print out of the token (internal representation) file:

```
29 2 x 17 2 y 18 21 2 x 20 2 y 4 3 56 18 22 19
```

And its symbolic representation:

```
varsym identsym x commasym identsym y semicolon beginsym identsym x  
becomessym identsym y plussym numbersym 56 semicolonsym endsym  
periodsym
```

2.- Print out the message “ No errors, program is syntactically correct”

3.- Print out the generated code. For example:

```
LOD 0 0 5  
LIT 1 0 56  
ADD 0 0 1  
STO 0 0 4  
SIO 0 0 3
```

4.- Run the program on the virtual machine (HW1)

Example 2, if the input is:

```
var x, y;  
begin  
  x := y + 56;  
end           ← ( notice period expected after the “end” reserved word)
```

The output should look like:

1.- A print out of the token (internal representation) file:

```
29 2 x 17 2 y 18 21 2 x 20 2 y 4 3 56 18 22
```

And its symbolic representation:

```
intsym identsym x commasym identsym y semicolonsym beginsym identsym x  
becomessym identsym y plussym numbersym 56 semicolonsym endsym
```

2.- Print the message “Error number xxx, period expected”

```
var x, y;  
begin  
  x := y + 56;  
end  
  ***** Error number xxx, period expected
```

Appendix B:

EBNF of tiny PL/0:

```
program ::= block "." .
block ::= const-declaration var-declaration statement.
constdeclaration ::= [ "const" ident "=" number { "," ident "=" number } ";" ].
var-declaration ::= [ "var" ident { "," ident } ";" ].
statement ::= [ ident ":" expression
                | "begin" statement { ";" statement } "end"
                | "if" condition "then" statement
                | "while" condition "do" statement
                | "read" ident
                | "write" ident
                | e ] .
condition ::= "odd" expression
            | expression rel-op expression.
rel-op ::= "=" | "<" | "<=" | ">" | ">=" .
expression ::= [ "+" | "-" ] term { ("+" | "-") term } .
term ::= factor { ("*" | "/" ) factor } .
factor ::= ident | number | "(" expression ")".
number ::= digit { digit } .
ident ::= letter { letter | digit } .
digit ::= "0" | "1" | "2" | "3" | "4" | "5" | "6" | "7" | "8" | "9".
letter ::= "a" | "b" | ... | "y" | "z" | "A" | "B" | ... | "Y" | "Z".
```

Based on Wirth's definition for EBNF we have the following rule:

[] means an optional item.

{ } means repeat 0 or more times.

Terminal symbols are enclosed in quote marks.

A period is used to indicate the end of the definition of a syntactic class.

Appendix C:

Error messages for the tiny PL/0 Parser:

1. Use = instead of :=.
2. = must be followed by a number.
3. Identifier must be followed by =.
4. **const**, **var**, **procedure** must be followed by identifier.
5. Semicolon or comma missing.
6. Incorrect symbol after procedure declaration.
7. Statement expected.
8. Incorrect symbol after statement part in block.
9. Period expected.
10. Semicolon between statements missing.
11. Undeclared identifier.
12. Assignment to constant or procedure is not allowed.
13. Assignment operator expected.
14. **call** must be followed by an identifier.
15. Call of a constant or variable is meaningless.
16. **then** expected.
17. Semicolon or } expected.
18. **do** expected.
19. Incorrect symbol following statement.
20. Relational operator expected.
21. Expression must not contain a procedure identifier.
22. Right parenthesis missing.
23. The preceding factor cannot begin with this symbol.
24. An expression cannot begin with this symbol.
25. This number is too large.

Note: Not all of these error messages may be used, and you may choose to create some error messages of your own to more accurately represent certain situations.

Appendix D:

Recursive Descent Parser for a PL/0 like programming language in pseudo code:

As follows you will find the pseudo code for a PL/0 like parser. This pseudo code will help you out to develop your parser and intermediate code generator for tiny PL/0:

```
procedure PROGRAM;
begin
  GET(TOKEN);
  BLOCK;
  if TOKEN != "periodsym" then ERROR
end;

procedure BLOCK;
begin
  if TOKEN = "constsym" then begin
    repeat
      GET(TOKEN);
      if TOKEN != "identsym" then ERROR;
      GET(TOKEN);
      if TOKEN != "eqsym" then ERROR;
      GET(TOKEN);
      if TOKEN != NUMBER then ERROR;
      GET(TOKEN)
    until TOKEN != "commasym";
    if TOKEN != "semicolonsym" then ERROR;
    GET(TOKEN)
  end;
  if TOKEN = "intsym" then begin
    repeat
      GET(TOKEN);
      if TOKEN != "identsym" then ERROR;
      GET(TOKEN)
    until TOKEN != "commasym";
    if TOKEN != "semicolonsym" then ERROR;
    GET(TOKEN)
  end;
  while TOKEN = "procsym" do begin
    GET(TOKEN);
    if TOKEN != "identsym" then ERROR;
    GET(TOKEN);
    if TOKEN != "semicolonsym" then ERROR;
    GET(TOKEN);
```



```

        BLOCK;
        if TOKEN != "semicolon" then ERROR;
        GET(TOKEN)
    end;
    STATEMENT
end;

```

```

procedure STATEMENT;
begin
    if TOKEN = "ident" then begin
        GET(TOKEN);
        if TOKEN != "becomes" then ERROR;
        GET(TOKEN);
        EXPRESSION
    end
    else if TOKEN = "call" then begin
        GET(TOKEN);
        if TOKEN != "ident" then ERROR;
        GET(TOKEN)
    end
    else if TOKEN = "begin" then begin
        GET TOKEN;
        STATEMENT;
        while TOKEN = "semicolon" do begin
            GET(TOKEN);
            STATEMENT
        end;
        if TOKEN != "end" then ERROR;

        GET(TOKEN)
    end
    else if TOKEN = "if" then begin
        GET(TOKEN);
        CONDITION;
        if TOKEN != "then" then ERROR;
        GET(TOKEN);
        STATEMENT
    end
    else if TOKEN = "while" then begin
        GET(TOKEN);
        CONDITION;
        if TOKEN != "do" then ERROR;
        GET(TOKEN);
        STATEMENT
    end
end;

```

```

procedure CONDITION;

```

```

begin
  if TOKEN = "oddsym" then begin
    GET(TOKEN);
    EXPRESSION
  else begin
    EXPRESSION;
    if TOKEN != RELATION then ERROR;
    GET(TOKEN);
    EXPRESSION
  end
end;

procedure EXPRESSION;
begin
  if TOKEN = "plussym" or "minussym" then GET(TOKEN);
  TERM;
  while TOKEN = "plussym" or "minussym" do begin
    GET(TOKEN);
    TERM
  end
end;

procedure TERM;
begin
  FACTOR;
  while TOKEN = "multsym" or "slashsym" do begin
    GET(TOKEN);
    FACTOR
  end
end;

procedure FACTOR;
begin
  if TOKEN = "identsym" then
    GET(TOKEN)
  else if TOKEN = NUMBER then
    GET(TOKEN)
  else if TOKEN = "(" then begin
    GET(TOKEN);
    EXPRESSION;
    if TOKEN != ")" then ERROR;
    GET(TOKEN)
  end
  else ERROR
end;

```

Appendix E:

Symbol Table

Recommended data structure for the symbol.

```
typedef struct
{
    int kind;           // const = 1, var = 2, proc = 3
    char name[10];      // name up to 11 chars
    int val;            // number (ASCII value)
    int level;          // L level
    int addr;           // M address
} symbol;
```

```
symbol_table[MAX_SYMBOL_TABLE_SIZE];
```

For constants, you must store kind, name and value.

For variables, you must store kind, name, L and M.

For procedures, you must store kind, name, L and M.