

CS172: COMPUTER SYSTEMS II

Lecture 3

Propositional Logic

Truth Tables

James Power



Truth Tables

Topics on truth tables:

- Definition of the propositional operators
- Valuations, evaluating complex formulas
- Special kinds of formulas: satisfiable, unsatisfiable, tautology
- Using truth tables to decide if an inference is valid

Truth Tables: defining and, or, not

Truth tables **define the semantics** of the propositional operators

ϕ	$\neg\phi$
<i>F</i>	<i>T</i>
<i>T</i>	<i>F</i>

ϕ	ψ	$\phi \wedge \psi$
<i>F</i>	<i>F</i>	<i>F</i>
<i>F</i>	<i>T</i>	<i>F</i>
<i>T</i>	<i>F</i>	<i>F</i>
<i>T</i>	<i>T</i>	<i>T</i>

ϕ	ψ	$\phi \vee \psi$
<i>F</i>	<i>F</i>	<i>F</i>
<i>F</i>	<i>T</i>	<i>T</i>
<i>T</i>	<i>F</i>	<i>T</i>
<i>T</i>	<i>T</i>	<i>T</i>

- Invented by:
C.S. Peirce (1893), Ludwig Wittgenstein (1912), Emil Post (1920)
- If you prefer, use 0 for “false” and 1 for “true”.

Truth Tables: compound propositions

Truth tables can also be used to *calculate the values* for compound propositions.

Example: write out the truth table for $(p \vee q) \wedge \neg (p \wedge q)$

p	q	$(p \vee q) \wedge \neg (p \wedge q)$
0	0	
0	1	
1	0	
1	1	

Truth Tables: implication and equivalence

You may not have seen these ones before:

ϕ	ψ	$\phi \rightarrow \psi$
<i>F</i>	<i>F</i>	<i>T</i>
<i>F</i>	<i>T</i>	<i>T</i>
<i>T</i>	<i>F</i>	<i>F</i>
<i>T</i>	<i>T</i>	<i>T</i>

ϕ	ψ	$\phi \leftrightarrow \psi$
<i>F</i>	<i>F</i>	<i>T</i>
<i>F</i>	<i>T</i>	<i>F</i>
<i>T</i>	<i>F</i>	<i>F</i>
<i>T</i>	<i>T</i>	<i>T</i>

Truth Tables: all five propositional operators

ϕ	$\neg\phi$	ϕ	ψ	$\phi \wedge \psi$	$\phi \vee \psi$	$\phi \rightarrow \psi$	$\phi \leftrightarrow \psi$
0	1	0	0	0	0	1	1
1	0	0	1	0	1	1	0
		1	0	0	1	0	0
		1	1	1	1	1	1

One way of remembering these is to treat ϕ and ψ as numbers:

<i>Logic</i> :	$\neg\phi$	$\phi \wedge \psi$	$\phi \vee \psi$	$\phi \rightarrow \psi$	$\phi \leftrightarrow \psi$
<i>Arithmetic</i> :	$1 - \phi$	$\min(\phi, \psi)$	$\max(\phi, \psi)$	$\phi \leq \psi$	$\phi = \psi$

Truth Tables: more complex propositions (again)

Example: write out the truth table for $(\neg(p \vee q) \rightarrow r)$

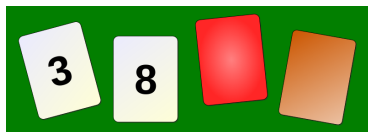
p	q	r	$(\neg(p \vee q) \rightarrow r)$
0	0	0	
0	0	1	
0	1	0	
0	1	1	
1	0	0	
1	0	1	
1	1	0	
1	1	1	

Wason selection task

A test to see if you understand **implication**:

- Four cards
- Each card has a number on one side, and a patch of color on the other.

Which card(s) must be turned over to test the idea that if a card shows an even number on one face, then its opposite face is red?



Devised by Peter Cathcart Wason (1966).
See: [Wikipedia](#), or textbook §2.12

Valuations

- In propositional logic, a *valuation* corresponds to a single row in the truth table.
- Notation:
 - $V \models \varphi$ means "the proposition φ is *true* under the valuation V "
 - $V \not\models \varphi$ means "the proposition φ is *false* under the valuation V "

Valuations

- In propositional logic, a *valuation* corresponds to a single row in the truth table.
- Notation:
 - $V \models \varphi$ means "the proposition φ is *true* under the valuation V "
 - $V \not\models \varphi$ means "the proposition φ is *false* under the valuation V "
- Example: Suppose we let $V(p) = 0$ and $V(q) = 1$
Decide which of the following are correct **using their syntax trees**:
 - $V \models p \wedge q \vee p$
 - $V \models p \rightarrow (\neg p \rightarrow q)$
 - $V \models p \rightarrow (q \rightarrow p)$
 - $V \models p \vee \neg p$
 - $V \models p \leftrightarrow (p \vee q)$