

CS172: COMPUTER SYSTEMS II

Lecture 5

Propositional Logic

Expressiveness

James Power, Kevin Casey



**Maynooth
University**
National University
of Ireland Maynooth

Expressiveness: “The” five operators?

LiA §2.9

Looking at the five propositional connectives \neg , \wedge , \vee , \rightarrow and \leftrightarrow we can ask

- *Q1: Do we need all we have?*

Do we need all five connectives? Could we “make do” with less?

- *Q2: Do we have all we need?*

Are the five connectives we’ve defined the *only* connectives?

Could we get extra expressive power by using ternary connectives (3 operands) or more?

Q1: Do we need all we have?

- Note the following equivalences:
 1. $\phi \leftrightarrow \psi$ and $(\phi \rightarrow \psi) \wedge (\psi \rightarrow \phi)$.
 2. $\phi \rightarrow \psi$ and $\neg\phi \vee \psi$.
 - 3a. $\phi \wedge \psi$ and $\neg(\neg\phi \vee \neg\psi)$.
 - 3b. $\phi \vee \psi$ and $\neg(\neg\phi \wedge \neg\psi)$.

Q1: Do we need all we have?

- Note the following equivalences:
 1. $\phi \leftrightarrow \psi$ and $(\phi \rightarrow \psi) \wedge (\psi \rightarrow \phi)$.
 2. $\phi \rightarrow \psi$ and $\neg\phi \vee \psi$.
 - 3a. $\phi \wedge \psi$ and $\neg(\neg\phi \vee \neg\psi)$.
 - 3b. $\phi \vee \psi$ and $\neg(\neg\phi \wedge \neg\psi)$.
- Alternative:
 - Rewrite and/or/iff in terms of implies and not
 - Let \perp represent the formula that is always false
 - Rewrite $\neg\phi$ to $\phi \rightarrow \perp$

A minimal set of operators

- In fact we can get away with a single operator that represents a combination of *not* with *and* (sometimes called NAND)
- This was discovered multiple times, most notably by Henry Sheffer in 1913

In mathematical logic it is usually called the [Sheffer stroke](#):

p	q	$p \mid q$
F	F	T
F	T	T
T	F	T
T	T	F

- Actually, *not* with *or* (NOR) will do equally well

Q2: Do we have all that we need?

What other **unary operators** could there be?

p	$f_1(p)$	$f_2(p)$	$f_3(p)$	$f_4(p)$
0				
1				

Q2: Do we have all that we need?

What other **unary operators** could there be?

p	$f_1(p)$	$f_2(p)$	$f_3(p)$	$f_4(p)$
0	0	0	1	1
1	0	1	0	1

Q2: Do we have all that we need?

What other **binary operators** could there be?

p	q	
0	0	
0	1	
1	0	
1	1	

Q2: Do we have all that we need?

What other **binary operators** could there be?

p	q	f_1	f_2	f_3	f_4	f_5	f_6	f_7	f_8	f_9	f_{10}	f_{11}	f_{12}	f_{13}	f_{14}	f_{15}	f_{16}
0	0	1	1	1	1	1	1	1	1	0	0	0	0	0	0	0	0
0	1	1	1	1	1	0	0	0	0	1	1	1	1	0	0	0	0
1	0	1	1	0	0	1	1	0	0	1	1	0	0	1	1	0	0
1	1	1	0	1	0	1	0	1	0	1	0	1	0	1	0	1	0

Q2: Do we have all that we need?

What **ternary operators** could there be?

p	q	r	\dots	f_i	\dots
0	0	0	\dots	0	\dots
0	0	1	\dots	0	\dots
0	1	0	\dots	0	\dots
0	1	1	\dots	0	\dots
1	0	0	\dots	1	\dots
1	0	1	\dots	1	\dots
1	1	0	\dots	0	\dots
1	1	1	\dots	1	\dots

- Can we define f_i as a function of p, q, r , using only our existing connectives?