# Overloading

An operator or function is overloaded if its meaning depends on the number or types of its arguments

- + (real, integer)
- (unary, binary)

Also called ad-hoc polymorphism

Some languages permit overloading of user-defined functions.

Good idea??

### Coercion

Implicit, or automatic, type conversion.

```
var x:integer;
    y,z:real;
    ...
y := x + z;
```

- → "mixed-mode" expression
- Coercion is different from explicit type conversion

# Coercion Extremes (Algol 68)

```
int i;
real r;
[1:1] int rowi;
ref int refi;
union (int,real) ir;
proc int p;

r := i/r; -- widening (of i)
ir := i; -- uniting
i := refi; -- dereferencing
i := p; -- deproceduring
rowi := 5 -- rowing
```

# Relational and Boolean Expressions

Relational operators compare two operands and return a boolean

Lower precedence than arithmetic operators

$$a + b < c + d \equiv$$
  
 $(a + b) < (c + d)$ 

- Boolean values: true, false
- Boolean operators: and, or, xor, not, =
  - True boolean values are helpful.
    - → In C, use integers:

 $\rightarrow$  a > b > c is legal!

### **Short-Circuit Evaluation**

Get a result without evaluating entire expression.

```
x and y ≡ if x then y else false
x or y ≡ if x then true else y
(x and y are arbitrary boolean expressions)
```

- Same as regular evaluation in the absence of sideeffects.
- May have a choice, e.g. Ada:

x or else y

x and then y

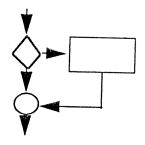
Otherwise, need to know if it's used.

### Statement-Level Control

- Three types:
  - → sequencing (not much to be said about this)
  - → selection
  - → repetition

### Selection

- Issues:
  - → How is selection controlled?
  - → How many choices? Can none of the choices be selected
- Single-way selectors

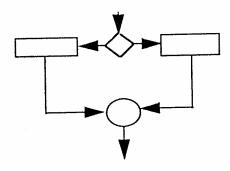


if <bool> then <stmt>

→ Without compound statements, need GOTO

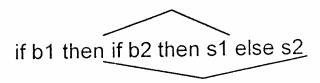
### Two-Way Selection

Two-way selectors



if <bool> then <stmt> else <stmt>

Recall potential ambiguity with nested 2-way selectors:



→ Several possible remedies...

### **Avoiding Ambiguity**

#### Disallow nested conditionals

if b1 then

begin if b2 then s1 else s2 end

if b1 then

begin if b2 then s1 end else s2

# Multiple-Way Selection

# Multiple-way Selection

→ New issues:

How to handle unrepresented selector values?

Should selectable segments be followed by implicit branches of construct?

ightarrow Lots of forms of multiple selectors. . .

# Multiple-Way Selection -- Early Approaches

#### Three-way selectors (Fortran)

IF (<arith\_exp>) L1, L2, L3

→ Must jump out of segment:

10

**GO TO 40** 

20

GO TO 40

30 ...

40 ..

### Computed GO TO (Fortran)

- → goes to first label if exp=1, second if exp=2, etc.
- → if exp < 1 or exp > n, no effect
- → must still jump out of segment

# Multiple-Way Selection -- Modern Approache

#### Pascal example

- → <exp> of ordinal type
- → implicit branch to end after each segment
- → "otherwise" option
- → selector value unrepresented => error

## More Modern Approaches

#### C example

```
switch (<exp>)
{
  case <const_exp> : <stmt>
  ...
  case <const_exp> : <stmt>
[default : <stmt>]
}
```

- → <exp> and <const\_exp> yield integers
- → No implicit branches -- use "break" to leave switch Is this a good idea?

#### Elsif (Ada)

- → Like LISP cond
- → Cuts down on indenting
- → More flexible than case statement

# Repetition

- Iteration (statement-level)
- Recursion (unit-level) -- later
- Issues:
  - → How is iteration controlled?
  - → Where in loop is control mechanism?

top

bottom

anywhere

# Counter-Controlled Loops

#### Issues:

- → Type and scope of loop variable
- → Value of loop variable at loop term
- → Change loop parameterss in body? Effect?
- → Branch into loop?
- → Test at top or bottom?
- → Loop parameters evaluated when?

### FORTRAN loops

- DO <label> <var> = <init>,<term>
  - → Can't change loop variable/parameters inside loop
  - → Loop variable undefined after normal termination; last value if jump out
  - → Can jump out of, back into loop

100 CONTINUE

→ Note!

DO 100 I = 1.10

# More Loop Examples

- ALGOL 60
  - → Very complicated! see book.
- ALGOL 68

for 
$$(i = 0; i \le 10; i++)$$
  
 $sum = sum + a[i];$ 

- → No special loop variables.
- → Can have multiple statements for each <exp>

# Logically Controlled Loops

- Issues:
  - → Pretest or posttest?
  - → Jump into loop?
- Pascal constructs

```
while <exp> do <stmt> repeat <stmt> until <exp>
```

Other exits, e.g. Modula 2, Ada:

```
loop
...
if <exp> then EXIT
...
end
```

CONTINUE, BREAK

# More Statement-Level Control Constructs

### Explicit loop exits

→ Basic model: (Modula)

```
loop
...
if . . . then EXIT
...
end
```

- Multi-level exits (Ada)
  - → named loops
  - → conditional EXIT

exit [<loop name>] [when <condition>]

# Ada Example

```
Outer_loop:

for i in 1. .10 loop

Inner_loop:

for j in 1. .20 loop;

...

EXIT Outer_loop when <cond>;

...

end loop Inner_loop;

...

end loop Outer_loop;
```

### Iterators (Clu)

 Lets user specify range of values over which a loop iterates.

```
for atom:node in list(x) do
...

list = iter(z: linked_list) yields (node)
...

yield (n)
...
end list
```

→ list produces elements of type node one at a time

# **Guarded Commands**

```
if <bool> -> <stmt>
W <bool> -> <stmt>
...
W <bool> -> <stmt>
fi
```

#### Semantics

- → Evaluate all <bool>s
- → If ≥ 1 is true, choose one nondeterministically and execute its <stmt>
- → If none true, error

# Guarded Commands (continued)

#### Semantics

- → Evaluate all <bool>s
- → If ≥ 1 is true, choose one nondeterministically and execute its <stmt>
- → Repeat until none true, then normal termination