Introduction, or Why Bother With This Stuff, Anyway?

- Increased capacity to express ideas.
- Improved background for choosing languages.
- Increased ability to learn languages.
- Understanding of significance of implementation.
- Ability to design new languages.
- Overall advancement of computing.

Readability & Writability

- Simplicity
 - small # of basic components
 (subsets a poor solution!)
 - one syntax: one meaning
 Counter-example in C: four ways to increment x:

```
X++;
x=x+1;
x+=1;
++x;
In FORTRAN, two meanings for:
Y=SUM(I,J)
```

-- array reference

Orthogonality

- => Any composition of basic primitives is allowed
- need small set of primitives, ways to combine them
- · Pascal not very orthogonal.

Functions can't return structured types.

Type of formal parameter must be stated in function/ procedure heading <u>unless</u> parameter is a function or procedure.

Enumerated types can't be read or written. etc...

- Non-orthogonality is often to simplify implementation.
- LISP is much more orthogonal than Pascal.

Control Statements/Constructs

Importance and desirability of various control mechanisms with the language.

Data Types

Rich set of data types makes programs much easier to writunderstand. Provides abstraction.

Syntax

Matters more than you think!

Identifier length, reserved words, layout, etc.

Abstraction

Must be able to hide details, or complexity is too great.

process abstraction

data abstraction

Reliability

Definition: performs to specifications under all conditions

Impact from:

```
type checking (or lack thereof)

compile-time best
runtime good

exception handling

Special language features to help intercept and handle
unusual situations. No magic.

Somewhat controversial.

aliasing

Two or more names for same memory cell.

var p,q: ^int;

begin

new(p);
q:=p;
```

Cost

More than just runtime!

time to train programers

*program development time
compile time
runtime

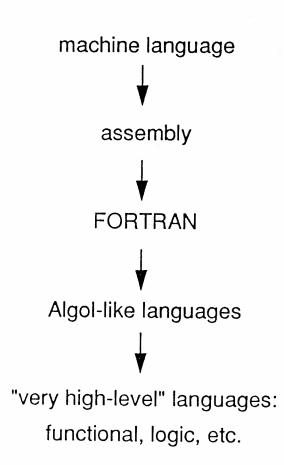
*maintenance time (>50%!)

* functions of writability and readability

=> most important

Influences on Language Design

- (1) Computer architecture
- (2) Programming methodologies
- Historically more of (1), moving toward (2). Getting higher and higher-level:



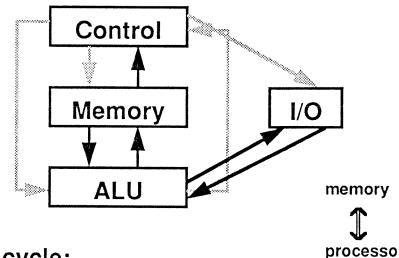
The von Neumann Architecture

- John von Neumann at Princeton
 - late 1940's

- has influenced the design of most programming

languages

control
data



- Fetch-execute cycle:
 - Control fetches next instr from memory
 - Control decodes instr
 - Execution

data from memory to ALU, or data from ALU to memory, or I/O

The von Neumann Architecture

- The von Neumann architecture is reflected in traditional programming languages in two ways:
 - sequential, step-by-step execution of instructions
 - modifiable variables -- "cubbyholes" in memory
- These languages became popular and drove further architectural designs. Vicious circle... other language designs didn't have much chance until recently.

Translation and Interpretation

 We could build a special machine to execute each language directly, but this is impractical. So how to get a program in a high-level language down to machine code?

Interpretation

An interpreter takes statements of a program one at a time and executes them directly as follows:

Get next statement

Determine actions

Perform actions

Repeat

(look familiar?)

Notes on Interpretation:

- data is provided to the interpreter as required
- one high-level instruction (HLI) => one sequence of machine-level instructions
- redetermine actions each time HLI is encountered
- highly dynamic

 Prog

 Data ► Interpreter ► Answer

Translation

Translation

 A translator takes a program in language A and produces an equivalent program in language B. If B is "closer" to machine code than A, it's called a compiler.

· Notes on translation:

high-level program --> machine-level prog

```
(not HLI-->MLI)
```

- parsing + code generation
- decode each statement once

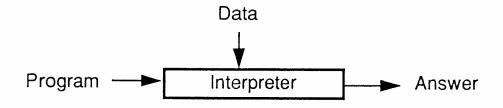
```
=> saves time
```

store expanded version of program

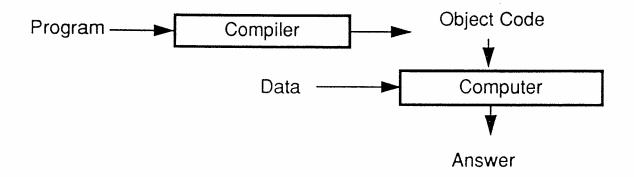
=> costs space

Interpretation vs. Translation

Interpretation



Translation



 Combine translation and interpretation by substituting Interpreter for Computer above.

- Wegner, "Prog. Langs-The First 25 Years", IEEE Transactions on Computers, 12/76, 1207-25
- 1950's "Discovery and description"
 - assembly
 - FORTRAN, ALGOL60, COBOL, LISP
 - basic implementation techniques

symbol tables
stack evaluation of arithmetic
activation records
garbage collection

- languages as tools
- late 1950's: first compilers (Hopper, etc.)
 grammars and automata (Chomsky and Miller)

- 1960's "Elaboration and analysis"
 - theories of programming languages
 - more formal development

formal languages

automata

formal semantics

verification

- bigger, more complex languages
- PL/I, Simula, ALGOL68
- late 1960's: theoretical work on compilers, program optimization

- <u>1970's</u> "Technology"
 - practical issues
 - applications of computer science
 - hardware cheaper, faster
 - software complexity increased
 - programming methodologies
 structured programming
 program verification
 - Pascal, C, Modula, Clu

• <u>1980's</u>

- parallel hardware => parallel language
- very high-level languages
 functional
- logic
- (object-oriented)

- 1944: EDVAC (Electronic Discrete Variable Automatic Calculator) Report (von Neumann)
 - first description of a stored-program computer
- 1950: First Assemblers
- 1954-57: FORTRAN ("FORmula TRANslating system")
 - Backus et al @ IBM
 - Goals:

efficiency -- less than twice as slow as assembler solve economic problem -- design, coding, debugging too expensive in assembly

- elegance of design secondary
- Versions I, II, III, IV
- introduced separate compilation with II because programs were getting too large to compile without hardware errors
- "An existence proof for higher-level languages..."

- 1958-60: ALGOL 60 ("ALGOrithmic Language")
 - by committee, including Backus
 - Goals:

elegant, universal language (FORTRAN was for IBM)

standard mathematical notation

major contributions:

BNF

block structure

recursion

call-by-value/name

stack model of evaluation

semi-dynamic arrays

but no formatted I/O -- too machine-dependent

- 1956-62: LISP ("LISt Processing")
 - McCarthy @ MIT
 - for symbolic computation in Al
 - free of von Neumann concepts
 - (roughly) based on lambda-calculus

- 1956-62: APL ("A Programming Language")
 - Iverson @ Harvard
 - array processing
 - functional flavor, fairly non-von Neumann
 - didn't catch on until 1970's
- 1960: COBOL ("COmmon Business Oriented Language")
 - at U Penn by representatives of computer manufacturers
 - alienated from CS community
 - developed by commercial community; didn't ask CS'ers no interest in scientific or research implications no BNF definition no good books
 - commercial applications thought trivial by CS'ers
 - main contribution: file/record structure
 - syntax wordy, English-like
 - very slow at first, but survived because use m
 - ref: Schneiderman, Annals of the History of Computing, 10/85

- 1960's: BASIC
 - Kemeny and Kurtz @ Dartmouth
 - for teaching
 - access through terminals
 - novel idea: user time more important than machine time!
 - commercial success a surprise -- intended for their students
 - no real contributions
- 1962-67: SNOBOL4 ("StriNg Oriented symBOlic Language")
 - Griswold @ Bell Labs
 - string processing
 - introduced pattern-matching
- 1964-69: PL/I ("Programming Language I")
 - by committee @ IBM
 - tried to unify commercial and scientific features
 - very large; programmers learn a subset

- 1963-68: ALGOL68
 - by committee
 - small number of orthogonal constructs
 - hard to learn -- too general and too flexible
 - poor implementations/documentation
- 1967-71: Pascal
 - _ Wirth
 - small, simple -- for teaching
 - structured programming, fairly rich data structures
- ~1973: C
 - Kernighan and Ritchie @ Bell Labs
 - low level, for systems programming
 - fairly small, fast
 - hard to read and maintain

- mid 1970's: Modula-2
 - Niklaus Wirth
 - Pascal and modules
 - better for systems programming and large projects
- mid 1970's: PROLOG ("PROgramming in LOGic")
 - Kowalski and Colmerauer @ Edinburgh and Marseilles
 - non-von Neumann, based on first-order logic (but impure)
 - most applications in Al
 - Japanese 5th generation computing project chose it
- mid 1970's: SMALLTALK
 - Xerox
 - object-oriented: shift in focus
 - not just a language; a whole system

- mid 1970's 80: Ada (after Ada Augusta, daughte of Lord Byron, associate of Babbage -- "the first programmer")
 - DoD
 - requirements developed slowly:

Strawman

Woodman

Tinman

Ironman

Steelman

- design contract won by CII-Honeywell Bull (Jean Ichbiah)
- based on Pascal
- large, complex
- features:

packages

tasks

real-time capabilities

exception handling

- 1980's: C++
 - Bjorne Stroustrup
 - C + classes
 - OOP in a popular language
- 1980's: Hope, Miranda, LML, Haskell
 - purely functional
 - based on lambda-calculus
 - higher-order functions, pattern matching, type inferencing
 - good for parallel machines?