

Bitwise Operations (Part 2)

§6.2, §7.2

Administrivia



- ▶ **Exam 2** Wednesday, November 5
 - ▶ Make-up exams must be scheduled **before** the exam is given in class; no make-ups afterward
- ▶ **Homework 5** out today
- ▶ Reading on procedures (no reading questions):

	<u>6/e</u>	<u>7/e</u>	<u>Title</u>
▶	§5.4	§5.1	Stack Operations
▶	§5.5	§5.2	Defining and Using Procedures
▶	§8.2	§8.2	Stack Frames
▶	§8.3	§8.3	Recursion – not covered in lecture

Topics Covered in Notes:



- ▶ TEST instruction

AND Masks – Clearing Bits



- ▶ Notice what happens to each bit when you apply an AND mask...

- ▶ $\begin{array}{r} 1010 \\ \& \underline{0011} \\ 0010 \end{array}$ $\begin{array}{l} < \textit{Original number} \\ < \textit{Bit mask} \end{array}$

- ▶ If there is a 0 bit in the mask, the corresponding bit is **cleared**
 - ▶ If there is a 1 bit in the mask, the corresponding bit is **retained**

- ▶ Bitwise AND can be used to **clear** particular bits

- ▶ $\begin{array}{r} 101101011110011001110110 \\ \& \underline{111111110000000011111111} \\ 101101010000000001110110 \end{array}$ $\begin{array}{l} < \textit{Original number} \\ < \textit{Bit mask} \end{array}$

OR Masks – Setting Bits



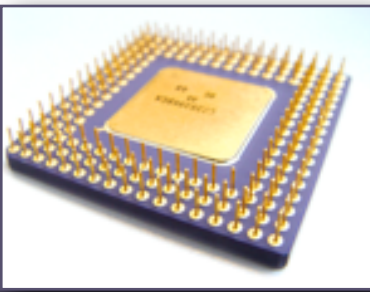
- ▶ A bit mask can be used with OR to **set** particular bits

- ▶
|
1010 < *Original number*
0011 < *Bit mask*
1011

- ▶ If there is a 0 bit in the mask, the corresponding bit is **retained**
- ▶ If there is a 1 bit in the mask, the corresponding bit is **set**

- ▶
|
101101011110011001110110 < *Original number*
000000001111111100000000 < *Bit mask*
101101011111111101110110

XOR Masks – Flipping Bits



- ▶ A bit mask can be used with **xOR** to **flip** particular bits

- ▶ $1010 < \textit{Original number}$

- $\oplus \underline{0011} < \textit{Bit mask}$

1001

- ▶ If there is a 0 bit in the mask, the corresponding bit is **retained**

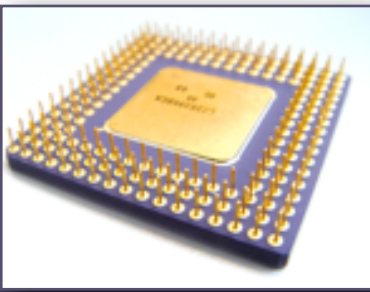
- ▶ If there is a 1 bit in the mask, the corresponding bit is **flipped**

- ▶ $101101011110011001110110 < \textit{Original number}$

- $\oplus \underline{000000001111111100000000} < \textit{Bit mask}$

101101010001100101110110

Expressions with AND, OR, XOR



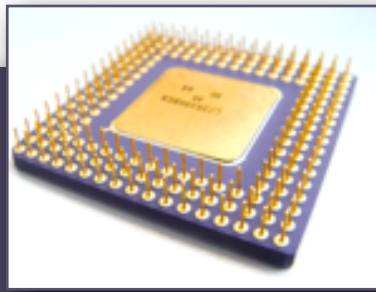
- ▶ Sometimes we will find it helpful to write mathematical expressions using bitwise operators
- ▶ Example: Given a 4-bit integer n , give an expression that is equal to n with bit 0 cleared and bit 3 set
 - ▶ n with bit 0 cleared is: $n \& 1110$
 - ▶ n' with bit 3 set is: $n' | 1000$
 - ▶ So, n with bit 0 cleared and bit 3 set is given by both of these:

$$(n \& 1110) | 1000$$

$$(n | 1000) \& 1110$$

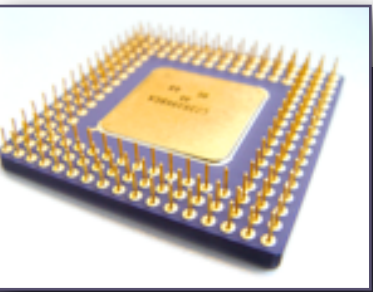
- ▶ Implementation in assembly language is straightforward:
; Suppose n is in AL
and al, 1110b
or al, 1000b
; Result is in AL

Expressions with AND, OR, XOR



- ▶ Sometimes we will find it helpful to write mathematical expressions using bitwise operators
- ▶ Example: Given an 8-bit unsigned integer, give an expression that evaluates to the largest even number not greater than n
 - ▶ I.e., $0 \mapsto 0, 1 \mapsto 0, 2 \mapsto 2, 3 \mapsto 2, 4 \mapsto 4, 5 \mapsto 4, 6 \mapsto 6, 7 \mapsto 6, 8 \mapsto 8, \dots$
 - ▶

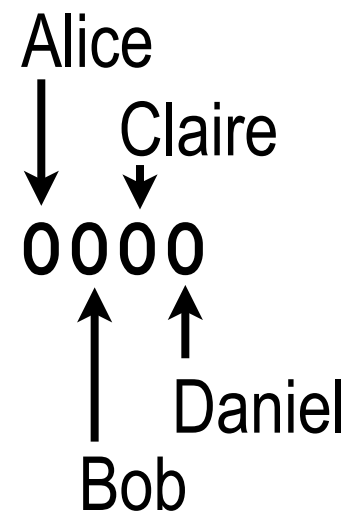
$0 = 00000000_2$	$1 = 00000001_2$
$2 = 00000010_2$	$3 = 00000011_2$
$4 = 00000100_2$	$5 = 00000101_2$
$6 = 00000110_2$	$7 = 00000111_2$
$8 = 00001000_2$	$9 = 00001001_2$
 - ▶ Even numbers: bit 0 clear Odd numbers: bit 0 set
 - ▶ Every odd number is equal to the previous even number + 1
 - ▶ So, the solution is to clear bit 0: the expression is $n \& 11111110$



Application: Bit Sets



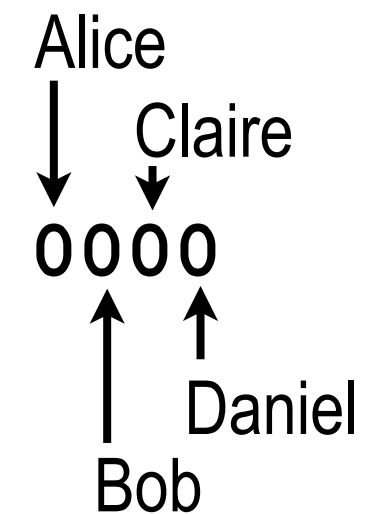
- ▶ Bit strings can represent sets
- ▶ Each bit represents one element
- ▶ The bit is 1 if that element is present
- ▶ The bit is 0 if that element is absent



Application: Bit Sets



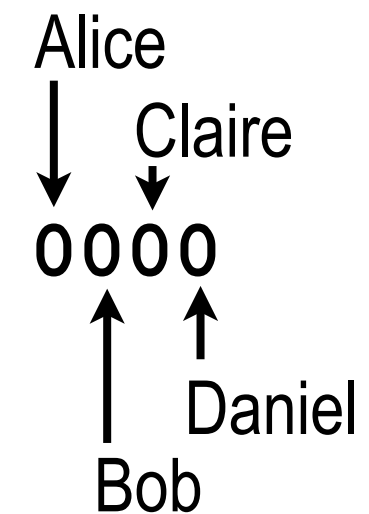
- ▶ **Combine the elements in two sets using $|$ (*set union*)**
 - ▶ 1010 denotes { Alice, Claire }; 0110 denotes { Bob, Claire }
 - ▶ $1010 | 0110 = 1110$, which denotes { Alice, Bob, Claire }
 - ▶ In set theory notation:
 $\{ \text{Alice, Claire} \} \cup \{ \text{Bob, Claire} \} = \{ \text{Alice, Bob, Claire} \}$



Application: Bit Sets



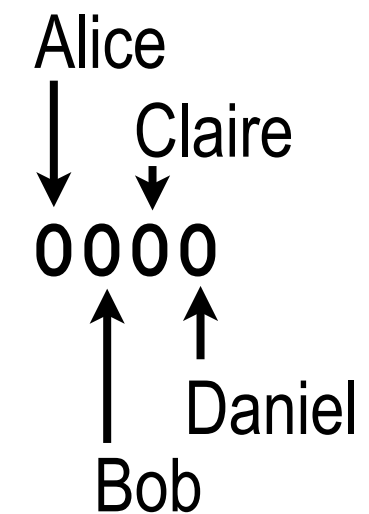
- ▶ **Find common elements using $\&$ (*intersection*)**
 - ▶ 1010 denotes { Alice, Claire }; 0110 denotes { Bob, Claire }
 - ▶ $1010 \& 0110 = 0010$, which denotes { Claire }
 - ▶ In set theory notation,
 $\{ \text{Alice, Claire} \} \cap \{ \text{Bob, Claire} \} = \{ \text{Claire} \}$



Application: Bit Sets



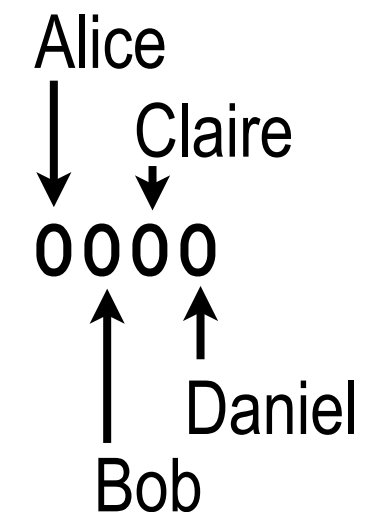
- ▶ **A set is empty iff the bits are all zeroes**
 - ▶ 0000 denotes $\{ \}$
- ▶ **Do two sets have any elements in common?**
 - ▶ Test whether their intersection is nonempty
 - ▶ In set theory notation, determine whether $S \cap T \neq \emptyset$
 - ▶ Use a bitwise AND, then determine if the result is nonzero
 - ▶ 1010 denotes $\{ \text{Alice, Claire} \}$; 0110 denotes $\{ \text{Bob, Claire} \}$
 - ▶ $1010 \& 0110 = 0010$, which denotes $\{ \text{Claire} \}$
 - ▶ Since 0010 is nonzero, the two sets have elements in common



Application: Bit Sets



- ▶ **Test for membership using &**
 - ▶ Bob is represented by bit 2
 - ▶ So, test if bit 2 is set
 - ▶ 1110 denotes { Alice, Bob, Claire }
 - ▶ $1110 \& 0100 = 0100$, which is nonzero
 - ▶ So, Bob *is* in { Alice, Bob, Claire }
 - ▶ `test al, 0100b ; Bit set in AL – is Bob in it?`
`jnz some_label ; Jump if Bob was in the set`
- ▶ Note that the expression $(1110 \& 0100)$ above is also used to compute the intersection of { Alice, Bob, Claire } and { Bob }
- ▶ So, testing whether Bob is a member of { Alice, Bob, Claire } is the same as testing if $\{ \text{Alice, Bob, Claire} \} \cap \{ \text{Bob} \} \neq \emptyset$



Application: Bit Sets



- ▶ Bit sets are good for representing sets where
 - ▶ there is a **finite** universe of elements
 - ▶ i.e., every element that could be in the set is known ahead of time
 - ▶ and the sets are **small** (very few possible elements) or **dense** (many of the possible elements will be in the set)
 - ▶ E.g., there are 24 students in this class;
can represent any set of these students with 24 bits!
- ▶ To represent sets from an infinite universe, or sets that are large and sparse, do not use bit sets: use, e.g., hash tables or binary trees
 - ▶ E.g., there are 7.076 billion people in the world – more are being added all the time (so the universe of elements is not fixed), and anyway a bit set that large would occupy 844 MB of memory
- ▶ In the Java library: BitSet, HashSet, TreeSet

Bit Shifting

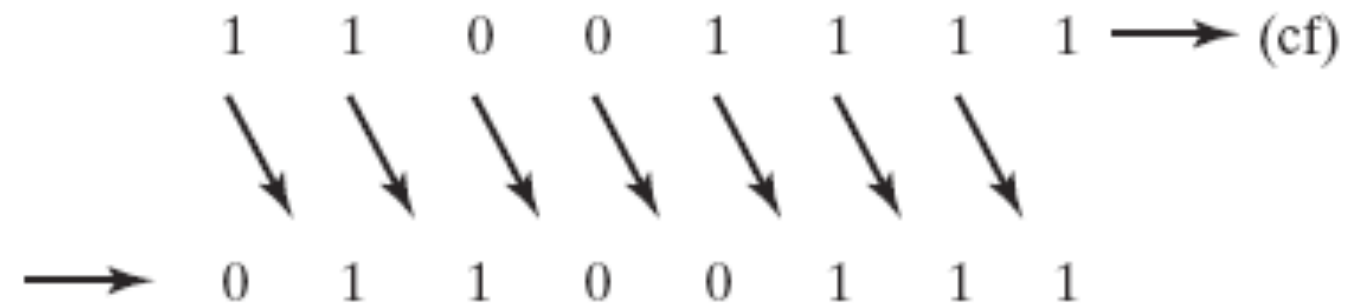
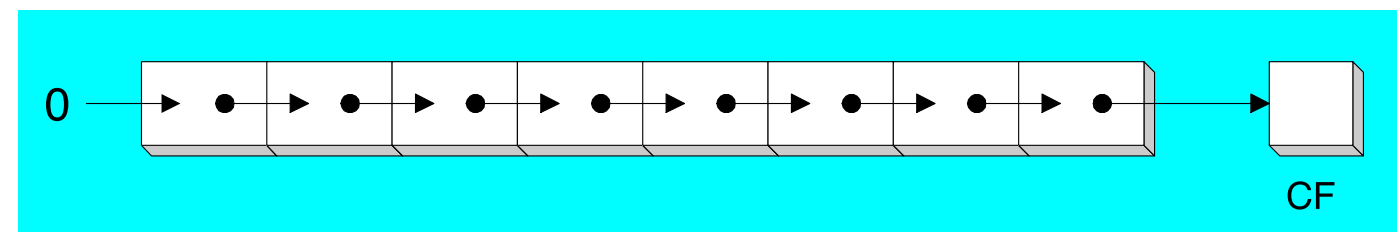


- ▶ The bits in a number can be *shifted* left or right
- ▶ When writing formulas, we will denote this by
 - ▶ $a \ll n$ Left shift by n bits
 - ▶ $a \gg^u n$ Right shift by n bits with 0-fill (logical shift)
 - ▶ $a \gg^s n$ Right shift by n bits with sign-fill (arithmetic shift)
- ▶ Examples (8-bit numbers):
 - ▶ $10111001 \ll 3 = 11001\mathbf{000}$
 - ▶ $10111001 \gg^u 2 = \mathbf{00}101110$
 - ▶ $10111001 \gg^s 2 = \mathbf{11}101110$
 - ▶ $00111001 \gg^s 2 = \mathbf{00}001110$

Logical Shift



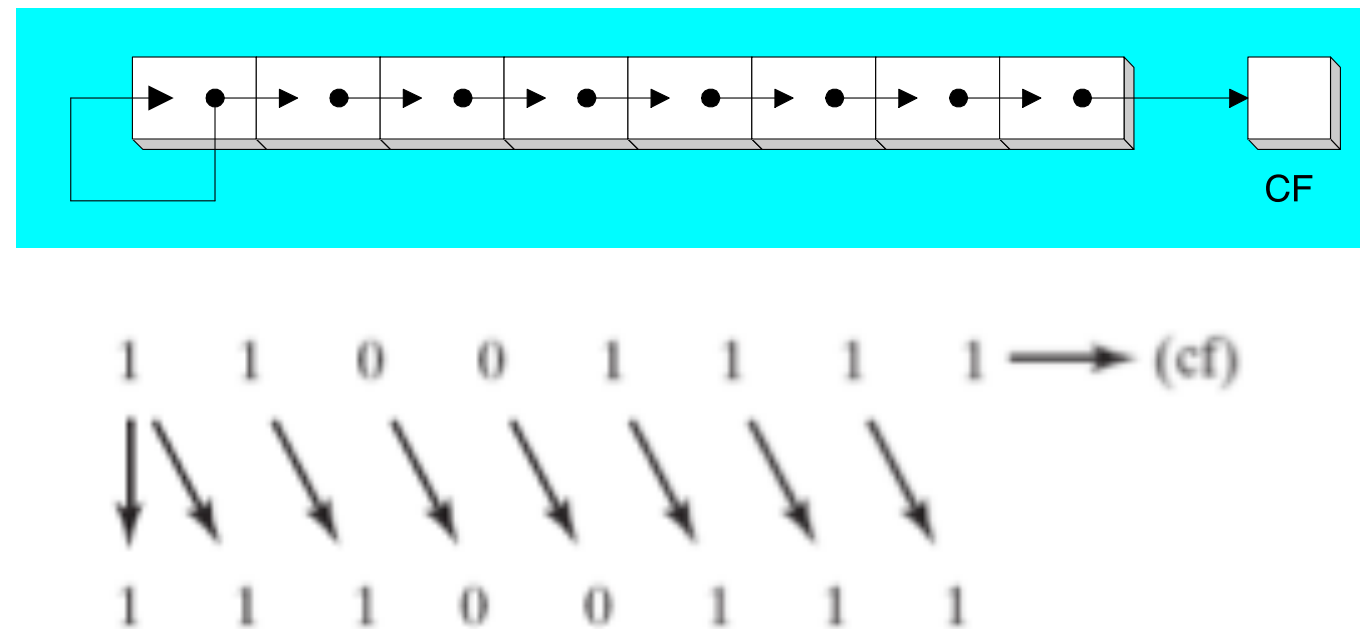
- ▶ A *logical* shift fills the newly created bit position with zero:



Arithmetic Shift



- ▶ An *arithmetic* right shift fills the newly created bit position with a copy of the number's sign bit:



- ▶ An arithmetic left shift is the same as a logical left shift!
(nothing special about the sign bit when shifting left)

SHL Instruction



- ▶ The SHL (shift left) instruction performs a logical left shift on the destination operand, filling the lowest bit with 0.



Operand types for SHL:

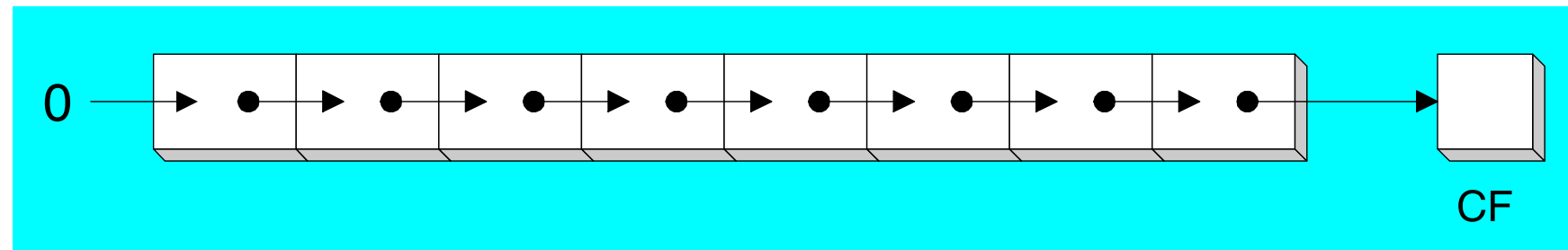
```
SHL reg,imm8  
SHL mem,imm8  
SHL reg,CL  
SHL mem,CL
```

(Same for all shift and rotate instructions)

SHR Instruction



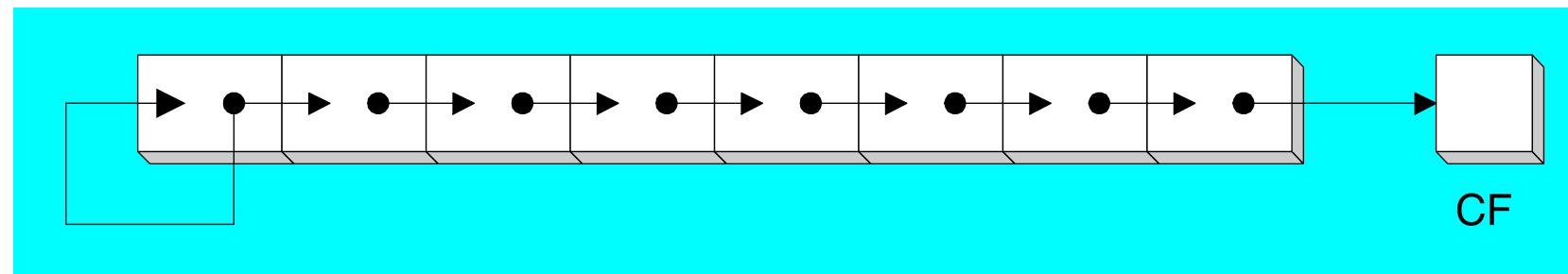
- ▶ The SHR (shift right) instruction performs a logical right shift on the destination operand. The highest bit position is filled with a zero.



SAL and SAR Instructions



- ▶ SAL (shift arithmetic left) is identical to SHL.
- ▶ SAR (shift arithmetic right) performs a right arithmetic shift on the destination operand.



An arithmetic right shift preserves the number's sign.

```
mov dl,-80
sar dl,1      ; DL = -40
sar dl,2      ; DL = -10
```

Application: Powers of 2



- ▶ What happens when you left shift the number 1?
- ▶ $1 \ll n$ has a 1 in bit position n and zeros elsewhere
- ▶ $1 \ll 0 = 00000001_2 = 1$
- ▶ $1 \ll 1 = 00000010_2 = 2$
- ▶ $1 \ll 2 = 00000100_2 = 4$
- ▶ $1 \ll 3 = 00001000_2 = 8$
- ▶ $1 \ll 4 = 00010000_2 = 16$
- ▶ $1 \ll 5 = 00100000_2 = 32$
- ▶ $1 \ll 6 = 01000000_2 = 64$
- ▶ $1 \ll 7 = 10000000_2 = 128$
- ▶ So, to compute powers of 2: $1 \ll n = 2^n$

Application: Shifting the Sign Bit



- ▶ Example: Given an SDWORD n , give an expression that is equal to 1 if n is negative and 0 otherwise
 - ▶ n is negative iff the sign bit is set
 - ▶ So, shift the sign bit into the ones' position
 - ▶ The expression is: $n \gg^u 31$
- ▶ Example: Given an SDWORD n , give an expression that is equal to -1 if n is negative and 0 otherwise
 - ▶ Recall that -1 is represented by all one bits
 - ▶ The expression is: $n \gg^s 31$