11. Searching and Sorting (Ch10)

- · Objectives when we have completed this set of notes, you should be familiar with:
 - Linear Search
 - Binary Search
 - Selection Sort
 - Insertion Sort

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Searching

- Finding a specific element in a group of items (search pool)
 - A student with a certain name in an array of Student objects
 - A bank account with a certain SS# in an online database
- Consideration:
 - What if the target is not present? (i.e., will not be found)
- Which search will be the most efficient?
- If class implements the Comparable interface, then two objects can be compared using the compareTo method (also a class must implement the Comparable interface in order for Collections.sort to work on objects of the class)

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Linear Search

- You have likely used linear searches in COMP 1210 for your projects
- Examines each element in a group starting at the first
- Elements do not have to be in a specific order
 - Must consider how you will find the specific item
- · Search ends when...
 - The element is found
 - The end of the list is reached
- NumberLinearSearch.java

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Linear Search

- Easy to implement
- The search pool does not have to be sorted
- Why not always use linear search?
- If the target is not present, you must go to the end of the search pool
 - Not good if the size of the search pool is large
- What if your target is near the end of the list?
 - Each comparison takes a certain amount of time, which can add up to minutes, days, weeks, or longer on a large database

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Binary Search

- An efficient alternative to linear search in some cases
- The search pool must be sorted first
- Starts at the middle of the candidate pool and then eliminates one half of the pool each time
- For 2ⁿ items, at most n compares are required
 - $16 = 2^4$ so at most 4 compares
 - 4 billion is approx = 2³² so at most 32 compares

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Binary Search

- Suppose that you are searching for the number 59 in the following array...
 - 2 4 9 23 27 29 30 34 45 59 67 76 89 92 97
- Binary search starts at the middle.



- 59 > 34, so the entire left side can be eliminated (cuts the candidate pool in half)
 - 2 4 9 23 27 29 30 34 45 59 67 76 89 92 97

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- Start again with the second half
 - 2 4 9 23 27 29 30 34 45 59 67 76 89 92 97
- 59 < 76, so the right half of the new candidate pool can be eliminated
 - 2 4 9 23 27 29 30 34 45 59 67 76 89 92 97
- 59 is the middle of the new candidate pool and is returned
 - 2 4 9 23 27 29 30 34 45 59 67 76 89 92 97



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Binary Search

- Binary search would have taken 3 comparisons to find the element
 - 2 4 9 23 27 29 30 <mark>34</mark> 45 <mark>59</mark> 67 <mark>76</mark> 89 92 97
- Linear search would have taken 10
 - 2 4 9 23 27 29 30 34 45 59 67 76 89 92 97
- Imagine if there were 4,000,000,000 elements in the search pool ... at most 32 compares
- NumberBinarySearchArray.java



Binary Search

- Why not always use binary search?
- Recall that it requires a sorted candidate pool
- Great if objects are pre-sorted, but sorting can also take a large number of comparisons
- You'll learn about many more sorting algorithms later in the CSSE program

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Sorting

- We've said that sorting also takes comparisons
- When you invoke Collections.sort or Arrays.sort the code has been written for you, but it will still take time to perform the sort
- How are elements sorted?
- Which sort is 'better' given certain conditions?
- Let's look at two sorting algorithms:
 - Selection Sort
 - Insertion Sort

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Selection Sort

- The approach of Selection Sort:
 - select a value and put it in its final place into the list
 - repeat for all other values
- In more detail:
 - find the smallest value in the list
 - switch it with the value in the first position
 - find the next smallest value in the list
 - switch it with the value in the second position
 - repeat until all values are in their proper places

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Selection Sort

→ Indicates that the two elements should be swapped • An example:

```
3 \leftarrow 9 \rightarrow 1
original:
smallest is 1:
                         9 ← 6 → 3 → 2
smallest is 2:
                         2
                              6 ←→3
                                       9
smallest is 3:
smallest is 6:
```

· Each time, the smallest remaining value is found and exchanged with the element in the "next" position to be filled

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Swapping

- The processing of the selection sort algorithm includes the *swapping* of two values
- Swapping requires three assignment statements and a temporary storage location:

```
temp = first;
first = second;
second = temp;
```

- See <u>NumberSelectionSort.java</u>
- See <u>NumberSelectionSortExample.java</u>

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Insertion Sort

- The approach of Insertion Sort:
 - pick any item and insert it into its proper place in a sorted sublist
 - repeat until all items have been inserted
- In more detail:
 - consider the first item to be a sorted sublist (of one item)
 - insert the second item into the sorted sublist, shifting the first item as needed to make room to insert the new addition
 - insert the third item into the sorted sublist (of two items), shifting items as necessary
 - repeat until all values are inserted into their proper positions

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Insertion Sort

• An example: Indicates where the next element should be inserted

original: 3 9 6 1 2 insert 9: 3 9 6 1 2 insert 6: 3 6 9 1 2 insert 1: 1 3 6 9 2 insert 2: 1 2 3 6 9

• See NumberInsertionSort.java

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Comparing Sorts

- The Selection and Insertion sort algorithms are similar in efficiency
- They both have outer loops that scan all elements, and inner loops that compare the value of the outer loop with almost all values in the list
- Approximately n² number of comparisons are made to sort a list of size n -> order n²
- Other sorts can be more efficient: order n log₂ n

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