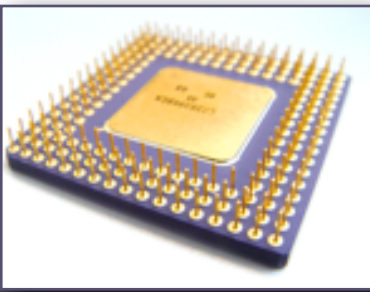




# Procedures (Part 4)

## §8.2



## Activity 12



# ESP Alignment & Locals



- ▶ ESP should always be aligned on a doubleword boundary, i.e., it must contain a memory address that's divisible by 4
  - ▶ Failure to do this may cause page faults, degraded performance
- ▶ **Round up local variable storage to a multiple of 4 bytes**
  - ▶ Need a 30-byte local array? Reserve 32 bytes.

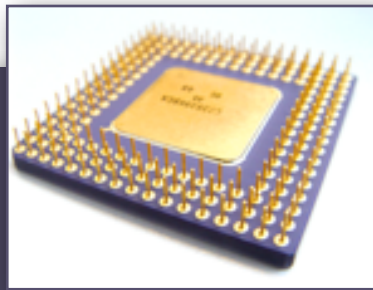
```
.code
sample PROC
    push ebp
    mov ebp, esp
    sub esp, 32    ; Reserve 32 bytes, even though we only use 30
    ...
    add esp, 32
    pop ebp
    ret
sample ENDP
```

# 8-, 16-, 64-bit Arguments



- ▶ ESP should always be aligned on a doubleword boundary, i.e., it must contain a memory address that's divisible by 4
  - ▶ Failure to do this may cause page faults, degraded performance
- ▶ Always push 32-bit values, including stack arguments
  - ▶ Expand 8-, 16-bit values to 32 bits

# 8-, 16-, 64-bit Arguments



- ▶ ESP should always be aligned on a doubleword boundary, i.e., it must contain a memory address that's divisible by 4
  - ▶ Failure to do this may cause page faults, degraded performance
- ▶ **Always push 32-bit values, including stack arguments**
  - ▶ Expand 8-, 16-bit values to 32 bits (MOVZX/MOVSX)
  - ▶ Pass multiword arguments in little endian order

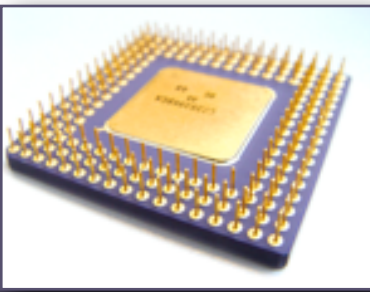
```
.data
q QWORD 1234567800ABCDEFh    ; In memory: EF CD AB 00 78 56 34 12
.code
                                ;      ^q      ^q+4
push DWORD PTR [q + 4]
push DWORD PTR q             ; Now the 8 bytes on the stack are in the same order as q
call WriteHex64                ; (little endian order in memory)
```

# ENTER and LEAVE



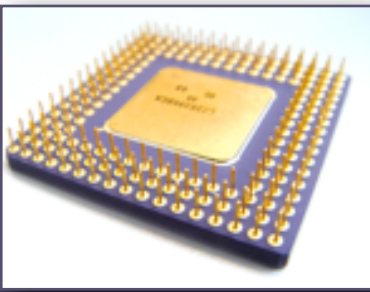
- ▶ The ENTER and LEAVE instructions create and terminate stack frames
- ▶ Simplify prologue/epilogue code
- ▶ ENTER *numbytes*, 0 is equivalent to
  - push ebp
  - mov ebp, esp
  - sub esp, *numbytes*
- ▶ LEAVE is equivalent to
  - mov esp, ebp
  - pop ebp

# LEA – Load Effective Address



- ▶ `.data`  
`array DWORD 10, 20, 30, 40, 50`
- ▶ Suppose we read an integer 0–4 into EAX, and we want to display the *memory address* of the element at that index
- ▶ We could retrieve that element using the indexed operand `[array + eax*4]`
- ▶ The *load effective address* (LEA) instruction determines the address of a memory operand and stores it in a register
- ▶ `call ReadDec` ; Read integer 0–4 into EAX  
`lea eax, [array+eax*4]` ; Store address in EAX  
`call WriteHex` ; Display address of that element

# LEA – Load Effective Address



- ▶ LEA is useful for creating an indirect operand from an indexed operand

*; Receives three 32-bit unsigned integers stack parameters*

```
AddThree PROC
```

```
    push ebp
```

```
    mov ebp, esp
```

*; Display each of the three stack parameters*

```
    lea esi, [ebp+8]           ; Point ESI at the first parameter
```

```
    mov ecx, 3
```

```
top:
```

```
    mov eax, [esi]
```

```
    call WriteDec
```

```
    add esi, SIZEOF DWORD ; Point ESI at the next parameter
```

```
    loop top
```

```
    ...
```



# Topics Covered in Notes:



- ▶ ENTER instruction
- ▶ LEAVE instruction
- ▶ LEA instruction

# Call by Value vs. Call by Reference



- ▶ *Value parameters* contain a value (e.g., integer)
  - ▶ For example, `min(int n, int m)`
  - ▶ This is what we've done so far
- ▶ *Reference parameters* contain a memory address
  - ▶ Passing an array “by value” would mean pushing every value in the array onto the stack – expensive!
  - ▶ Instead, pass the *address* (offset) of the array
- ▶ In Java, primitives (int, float, etc.) are passed by value; objects (arrays, strings, etc.) by reference

# Example: Reference Parameters



```
INCLUDE Irvine32.inc
```

```
.data
```

```
aWord          WORD ?
```

```
anotherWord WORD ?
```

```
.code
```

```
main PROC
```

```
    ; Set the value of aWord to 5
```

```
    push OFFSET aWord
```

```
    call SetToFive
```

```
    ; Set the value of anotherWord to 5
```

```
    push OFFSET anotherWord
```

```
    call SetToFive
```

```
    exit
```

```
main ENDP
```

```
    ; Sets a WORD variable to 5. (STDCALL)
```

```
    ; Receives: [ebp+8] Address of variable
```

```
    ; Returns: None
```

```
SetToFive PROC
```

```
    enter 0, 0
```

```
    push edi
```

```
    ; Copy the variable's address into EDI
```

```
    mov edi, [ebp+8]
```

```
    ; Set the variable's value to 5
```

```
    mov WORD PTR [edi], 5
```

```
    pop edi
```

```
    leave
```

```
    ret 4
```

```
SetToFive ENDP
```

```
end main
```

# Example: Reference Parameters



*In C++, reference parameters are denoted by an ampersand (&).*

*The following compiles to essentially the same code as the previous slide:*

```
unsigned short aWord, anotherWord;
```

```
void set_to_five(unsigned short &variable) {  
    variable = 5;  
}
```

```
void main() {  
    set_to_five(aWord);  
    set_to_five(anotherWord);  
}
```



# Example: Value Parameters



*The following code does **not** set the value of aWord or anotherWord. Why?*

*It uses a value parameter. The next slide shows what it essentially compiles to...*

```
unsigned short aWord, anotherWord;

void not_set_to_five(unsigned short variable) {
    variable = 5;
}

void main() {
    not_set_to_five(aWord);
    not_set_to_five(anotherWord);
}
```

# Example: Value Parameters



```
INCLUDE Irvine32.inc

.data
aWord          WORD ?
anotherWord WORD ?

.code
main PROC
    ; Push the (uninitialized) value of aWord
    push aWord
    call NotSetToFive

    ; Push the value of anotherWord
    push anotherWord
    call NotSetToFive

    exit
main ENDP
```

```
NotSetToFive PROC
    enter 0, 0

    ; Change the value of the parameter
    ; (on the stack) to 5
    mov WORD PTR [ebp+8], 5

    leave
    ret 4
    ; After RET, the parameter is gone
    ; since we destroyed the stack frame
NotSetToFive ENDP

end main
```