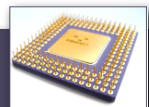




Procedures (Part 2)

§5.5

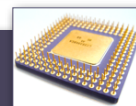
One Way to Implement a Stack



- `array DWORD 256 DUP(?)`
- `top_address DWORD (OFFSET array + SIZEOF array)`
- *Push (push 32-bit value in EAX onto stack):*
 - `sub top_address, 4 ; Stack grows downward in memory!`
 - `mov esi, top_address`
 - `mov [esi], eax`
- *Pop (remove 32-bit top element, return in EAX):*
 - `mov esi, top_address`
 - `mov eax, [esi]`
 - `add top_address, 4 ; Omit this to implement Top`

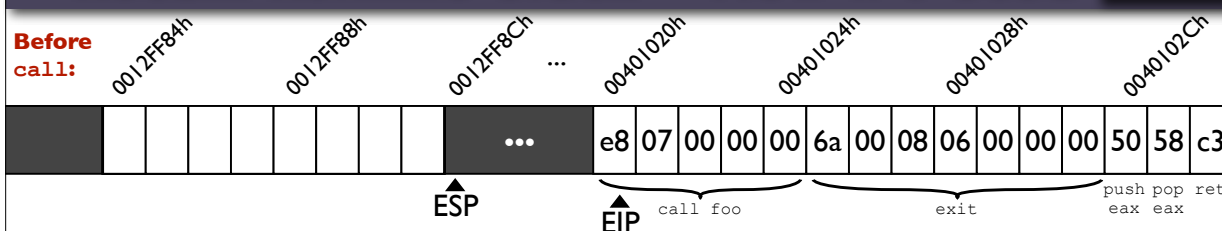
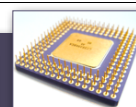
This is essentially how the runtime stack works
(but the top address is stored in ESP)

Topics Covered in Notes:



- ▶ PUSH instruction
- ▶ POP instruction
- ▶ CALL instruction
- ▶ RET instruction

Recall: Runtime Stack – How It's Used



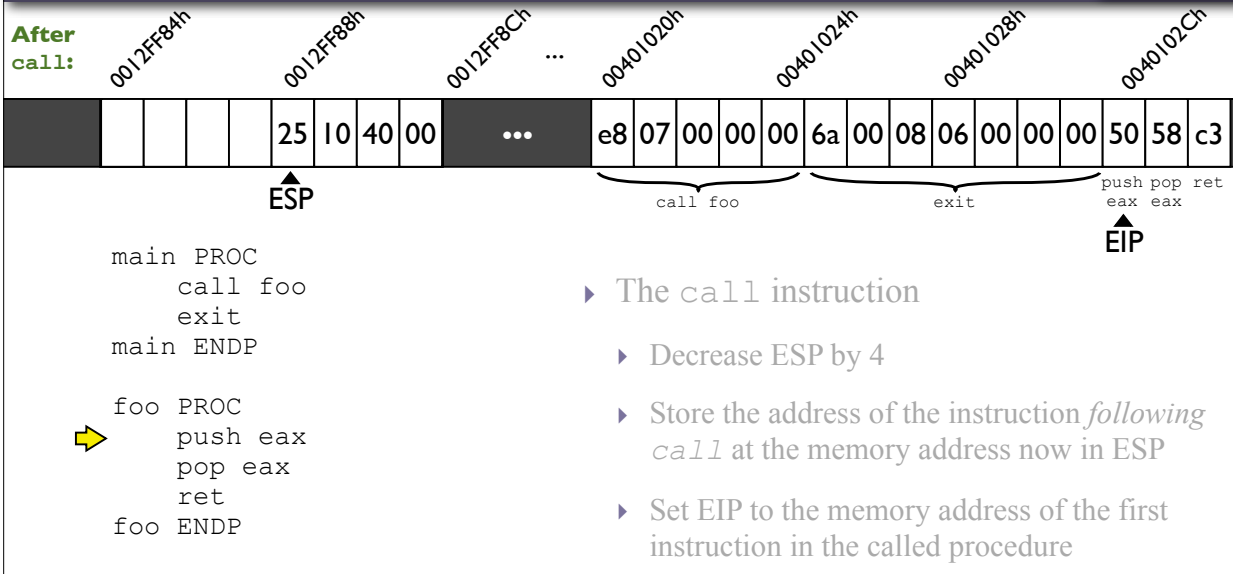
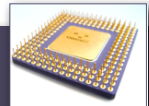
```

main PROC
    call foo
    exit
main ENDP

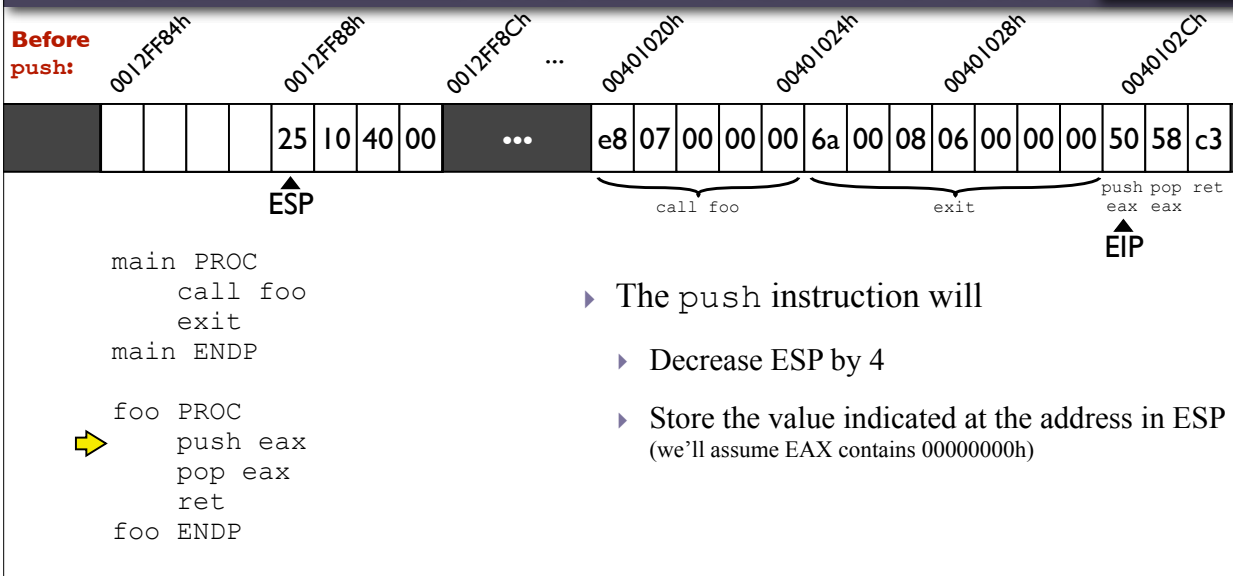
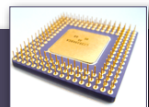
foo PROC
    push eax
    pop eax
    ret
foo ENDP
    
```

- ▶ The `call` instruction will
 - ▶ Decrease ESP by 4
 - ▶ Store the address of the instruction *following* `call` at the memory address now in ESP
 - ▶ Set EIP to the memory address of the first instruction in the called procedure

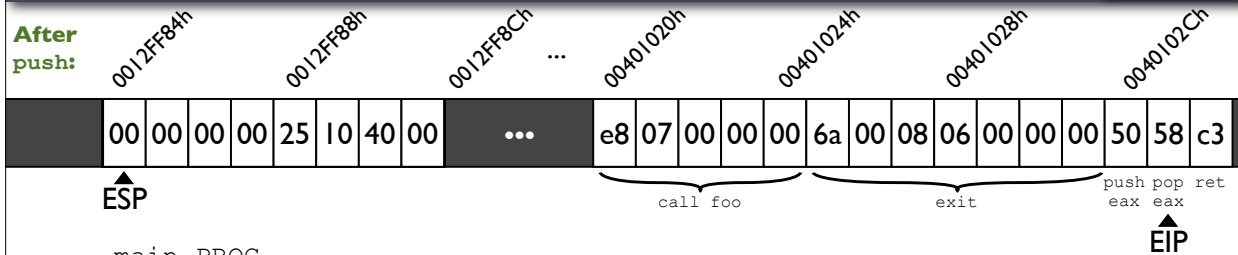
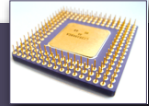
Recall: Runtime Stack – How It's Used



Recall: Runtime Stack – How It's Used



Recall: Runtime Stack – How It's Used

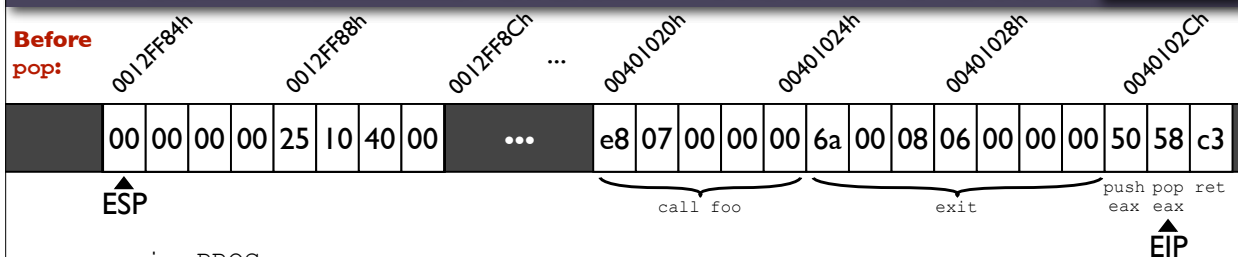
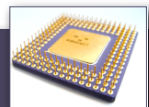


```
main PROC
    call foo
    exit
main ENDP
```

```
foo PROC
    push eax
    pop eax
    ret
foo ENDP
```

- ▶ The push instruction will
 - ▶ Decrease ESP by 4
 - ▶ Store the value indicated at the address in ESP (we'll assume EAX contains 00000000h)

Recall: Runtime Stack – How It's Used

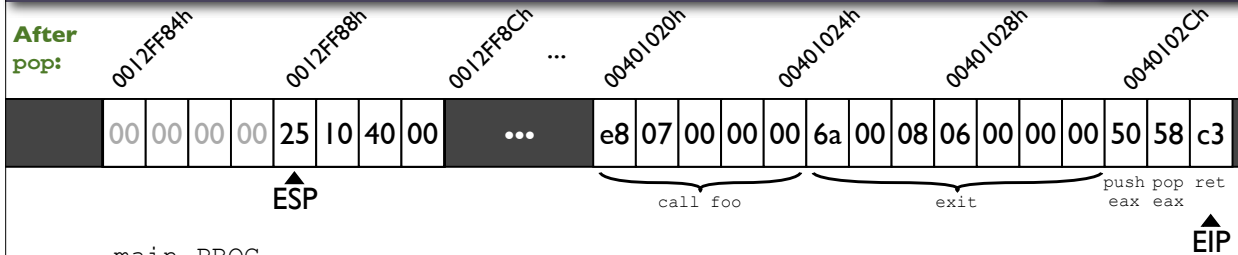
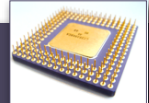


```
main PROC
    call foo
    exit
main ENDP
```

```
foo PROC
    push eax
    pop eax
    ret
foo ENDP
```

- ▶ The pop instruction will
 - ▶ Load the value from the address given by ESP, copying it into the given register
 - ▶ Increase ESP by 4

Recall: Runtime Stack – How It's Used



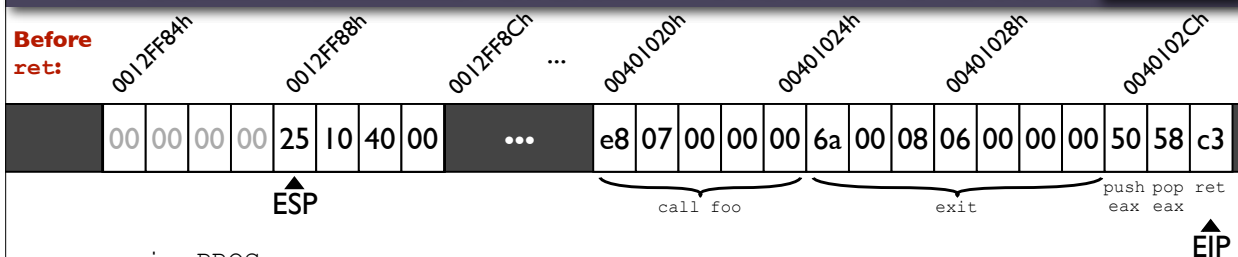
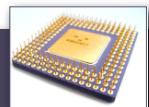
```
main PROC
    call foo
    exit
main ENDP
```

```
foo PROC
    push eax
    pop eax
    ret
foo ENDP
```



- ▶ The pop instruction will
 - ▶ Load the value from the address given by ESP, copying it into the given register
 - ▶ Increase ESP by 4

Recall: Runtime Stack – How It's Used



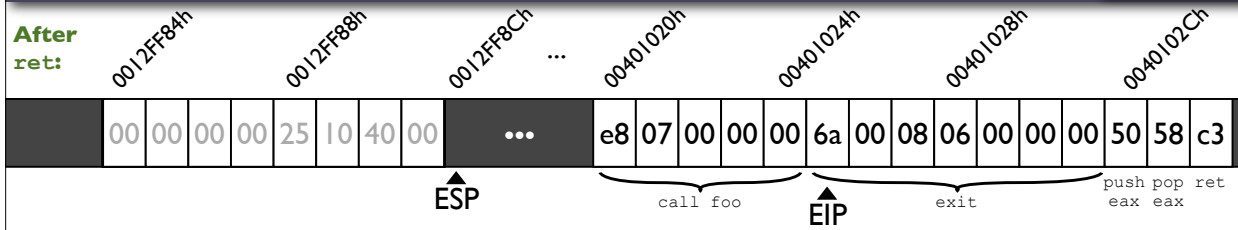
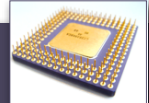
```
main PROC
    call foo
    exit
main ENDP
```

```
foo PROC
    push eax
    pop eax
    ret
foo ENDP
```



- ▶ The ret instruction will
 - ▶ Read the 32-bit value at ESP (in this example, 00401025h)
 - ▶ Increase ESP by 4
 - ▶ Set EIP to the value it just read (00401025h)

Recall: Runtime Stack – How It's Used



```

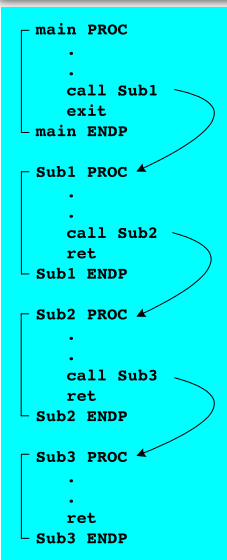
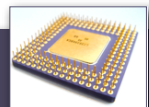
main PROC
    call foo
    exit
main ENDP
    
```

```

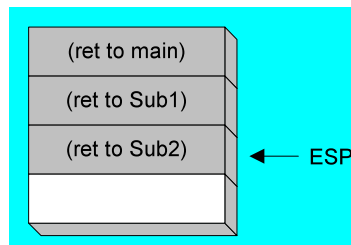
foo PROC
    push eax
    pop eax
    ret
foo ENDP
    
```

- ▶ The ret instruction will
 - ▶ Read the 32-bit value at ESP (in this example, 00401025h)
 - ▶ Increase ESP by 4
 - ▶ Set EIP to the value it just read (00401025h)

Nested Procedure Calls



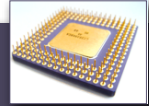
- ▶ *Nested procedure calls:*
you call a procedure, and it calls other procedures before returning to you
- ▶ main calls Sub1
Sub1 calls Sub2
Sub2 calls Sub3



- ▶ Recall: Stacks are last-in-first-out (LIFO) data structures
- ▶ You always want to return to the *last* procedure that CALLED
- ▶ The last procedure that CALLED will be the first address POPPED

Irvine, Kip R. *Assembly Language for x86 Processors* 6/e, 2010.

Nested Procedure Calls



```
main PROC
    call A
    exit
main ENDP
```

```
A PROC
    push eax
    push ebx
    call B
    pop ebx
    pop eax
    ret
```

```
A ENDP
```

```
B PROC
    push eax    at this point, the stack contains:
    pop eax
    ret
```

```
B ENDP
```

Return address for main
Saved value of EAX from A
Saved value of EBX from A
Return address for A
Saved value of EAX from B

<bottom of stack

<top of stack