From decimal expansions to reduced fractions

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In [1]: from math import gcd

A real number is rational if and only if a pattern eventually appears in its decimal expansion that repeats forever. So π , being irrational, is such that no finite sequence of consecutive digits in 3.14159265358979... eventually repeats forever. On the other hand,

- $\frac{25}{12} = 2.08333...3...$
- $\frac{97}{21}$ = 4.619047619047619047...619047...
- $\frac{11941}{49950} = 0.23905905905...905...$

The decimal expansion is unique except for fractions that in reduced form, have a power of 10 as denominator: those fractions have two decimal expansions, one that ends in 0 repeating forever, another one that ends in 9 repeating forever. For instance, $\frac{1234567}{1000} = 1234.567000...0... = 1234.566999...9...$

We want to, given two nonempty strings of digits σ and τ (that we treat as strings or numbers depending on the context), find out the unique natural numbers p and q such that the decimal expansion of $\frac{p}{q}$ reads as $0.\sigma\tau\tau\tau...\tau...$ and

- either p=0 and q=1 (case where σ and τ consist of nothing but 0's), or
- p and q are coprime, so $\frac{p}{q}$ is in reduced form (including the case where p=1 and q=1 because σ and τ consist of nothing but 9's).

For instance, if $\sigma=23$ and $\tau=905$, then p=11941 and q=49950. Writing $|\sigma|$ for the length (number of digits) in a string of digits σ , we compute:

$$\begin{split} 0.\sigma\tau\tau\tau\dots\tau\dots &= \sigma 10^{-|\sigma|} + \tau (10^{-|\sigma|-|\tau|} + 10^{-|\sigma|-2|\tau|} + 10^{-|\sigma|-3|\tau|} + \dots) \\ &= \sigma 10^{-|\sigma|} + \frac{\tau 10^{-|\sigma|-|\tau|}}{(1-10^{-|\tau|})} \\ &= \sigma 10^{-|\sigma|} + \frac{\tau 10^{-|\sigma|}}{(10^{|\tau|}-1)} \\ &= \frac{\sigma 10^{-|\sigma|} (10^{|\tau|}-1) + \tau 10^{-|\sigma|}}{(10^{|\tau|}-1)} \\ &= \frac{\sigma (10^{|\tau|}-1) + \tau}{(10^{|\tau|}-1)10^{|\sigma|}} \end{split}$$

Reducing the last fraction if needed provides the desired answer.

The result of the previous computation immediately translates to the function that follows:

To reduce a fraction, it suffices to divide its numerator and its denominator by their gcd (greatest common divisor). The math module has a gcd function:

```
In [3]: gcd(1234541976, 99999900000)
Out[3]: 24
```

Let us implement the gcd function ourselves, following Euclid's algorithm, which is based on the following reasoning. Let a and b be two natural numbers with b > 0. Since $a = \lfloor \frac{a}{b} \rfloor b + a \mod b$:

- if n divides both a and b then it divides both a and $\lfloor \frac{a}{b} \rfloor b$, hence it divides $a \lfloor \frac{a}{b} \rfloor b$, hence it divides $a \mod b$;
- conversely, if n divides both b and a mod b then it divides $\left| \frac{a}{b} \right| b + a \mod b$, hence it divides a.

Hence n divides both a and b iff n divides both b and a mod b. So $\gcd(a,b)=\gcd(b,a \bmod b)$. Since a mod b

<math>b, we get a sequence of equalities of the form: $\gcd(a,b)=\gcd(a_1,b_1)=\gcd(a_2,b_2)=\cdots=\gcd(a_{k-1},b_{k-1})=\gcd(a_k,0)$ with $k\geq 1$ and $b>b_1>b_2>\cdots>b_{k-1}>0$; as $\gcd(a_k,0)=a_k$, a_k is the \gcd of a and b.

To compute $\left|\frac{a}{b}\right|$, Python offers the // operator; to compute $a \mod b$, the % operator:

```
Out[4]: (4, 4.0, 4.0, 4.0)
Out[4]: (1, 1.0, 1.0, 1.0)
```

If a and b are arbitrary numbers (not necessarily integers) with $b \neq 0$, then the equality a = qb + r together with the conditions

- q is an integer
- |r| < |b|
- $r \neq 0 \rightarrow (r > 0 \leftrightarrow b > 0)$

determine q and r uniquely; // and % operate accordingly:

```
In [5]: 5 // 2, 5 % 2

-5 // 2, -5 % 2

5 // -2, 5 % -2

-5 // -2, -5 % -2

print()

7.5 // 2, 7.5 % 2

-7.5 // 2, -7.5 % 2

7.5 // -2, 7.5 % -2

-7.5 // -2, -7.5 % -2

Out[5]: (2, 1)

Out[5]: (-3, 1)

Out[5]: (-3, -1)

Out[5]: (3.0, 1.5)

Out[5]: (-4.0, 0.5)
```

The divmod() function offers an alternative to the previous combined use of // and %:

Out[5]: (-4.0, -0.5)

Out[5]: (3.0, -1.5)

```
Out[6]: (2, 1)
Out[6]: (-3, 1)
Out[6]: (-3, -1)
Out[6]: (2, -1)
Out[6]: (3.0, 1.5)
Out[6]: (-4.0, 0.5)
Out[6]: (-4.0, -0.5)
Out[6]: (3.0, -1.5)
```

In [7]: a = 30

b = 18

Let us get back to Euclid's algorithm, so assume again that a and b are two natural numbers with b>0. To implement the algorithm and compute $\gcd(a,b)$, it suffices to have two variables, say $\mathbf a$ and $\mathbf b$, initialised to a and b, and then change the value of $\mathbf a$ to b and change the value of $\mathbf b$ to a mod b, and do that again and again until $\mathbf b$ gets the value 0. To change the value of $\mathbf a$ to a mod b and change the value of $\mathbf b$ to a, it seems necessary to introduce a third variable:

```
c = a % b
a = b
b = c
a, b

Out[7]: (18, 12)

But Python makes it easier:

In [8]: a = 30
b = 18
# Evaluate the expression on the right hand side;
# the result is the tuple (18, 12).
# Then assign that result to the tuple on the left,
# component by component.
a, b = b, a % b
a, b

Out[8]: (18, 12)
```

Note that when the value of a is strictly smaller than the value of b, then a, b = b, a % b exchanges the values of a and b:

On the other hand, if the value of a is at least equal to the value of b, then this holds too after a, b = b, a % b has been executed. Let us trace all stages in the execution of Euclid's algorithm. The code makes use of a while statement whose condition is not a boolean expression. Applying bool() to an expression reveals which one of True or False the expression evaluates to in contexts where one or the other is expected:

```
In [10]: bool(None)
         bool(0), bool(5), bool(-3)
         bool(0.0), bool(0.1), bool(-3.14)
         bool([]), bool([0]), bool([[]])
         bool({}), bool({0: 0}), bool({0: None, 1: None})
         bool(''), bool(''), bool('0000')
Out[10]: False
Out[10]: (False, True, True)
In [11]: def trace_our_gcd(a, b):
             while b:
                 a, b = b, a % b
                 print(a, b)
         for a, b in (1233, 1233), (1233, 990), (990, 1233):
             print(f'\nTracing the computation of gcd of {a} and {b}:')
             trace_our_gcd(a, b)
Tracing the computation of gcd of 1233 and 1233:
1233 0
Tracing the computation of gcd of 1233 and 990:
990 243
243 18
18 9
9 0
```

```
Tracing the computation of gcd of 990 and 1233: 1233 990 990 243 243 18 18 9 9 0
```

The gcd is the value of a when exiting the while loop:

compute_fraction() returns the numerator and denominator of a fraction that another function, say reduce(), can easily reduce thanks to $our_gcd()$. It is natural to let reduce() take two arguments, the numerator and the denominator of the fraction to simplify, respectively. But $compute_fraction()$ returns those as the first and second elements of a tuple; a function always returns a single value. Between the parentheses that surround the arguments of a function f(), one can insert an expression that evaluates to a tuple and precede it with the *symbol, which "unpacks" the members of the tuple and make them the arguments of f():

```
Out[14]: (2, 6)
Out[14]: (4, 12)
Out[14]: (8, 24)
```

The * symbol can also be used in the definition of a function and precede the name of a parameter. It then has the opposite effect, namely, it makes a tuple out of all arguments that are provided to the function:

```
In [15]: \# x is the tuple of all arguments passed to f().
         def f(*x):
             return x * 2
         f()
         f(0)
         f(f(0))
         f(*f(0))
         f(f(f(0)))
         f(f(*f(0)))
         f(*f(*f(0)))
Out[15]: ()
Out[15]: (0, 0)
Out[15]: ((0, 0), (0, 0))
Out[15]: (0, 0, 0, 0)
Out[15]: (((0, 0), (0, 0)), ((0, 0), (0, 0)))
Out[15]: ((0, 0, 0, 0), (0, 0, 0, 0))
Out[15]: (0, 0, 0, 0, 0, 0, 0, 0)
```

Thanks to this syntax, it is possible to let reduce() as well as another function output() take two arguments numerator and denominator, and "pipe" compute_fraction(), reduce() and output() together so that the unpacked returned value of one function becomes the arguments of the function that follows: