System Programming

Lecture 1

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About me

- 조영필 (Yeongpil Cho)
 - A system security researcher
 - Designing new SW/HW techniques for better security
 - Kernels & Firmware
 - Applications
 - (co-)Processors

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Objectives

- Computer systems consists of various hardware and software components.
- Most system implementations share some fundamental design principles.
 - e.g., Android Phone, iPhone, Windows desktop
- We are going to generally explore those principles
 - to understand how programs run on the system
 - to understand what affect programs
 - to, ultimately, develop better programs

Materials

■ Textbook

- Computer Systems: A Programmer's Perspective 3rd Edition
 - Randal E Bryant, David R O'Hallaron

Reference Course

Carnegie Mellon University's Intro to Computer Systems

Evaluation

Exam	60%
Midterm	30%
Final	30%
Lab	30%
Attendance	10%

■ 3 tardies → 1 absence

Topics

Overview of the system programming course

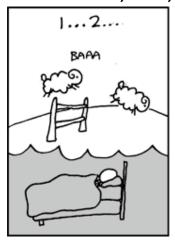
Abstraction Is Good But Don't Forget Reality

- Most CS and CE courses emphasize abstraction
 - Abstract Data Type
 - List, Tree, ...
 - Asymptotic analysis
 - Big-O/Θ/Ω, ...
- These abstractions are good for understanding, but we need to know reality
 - for better performance optimization
 - for better secure implementation

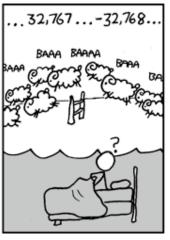
Great Reality #1: Ints are not Integers, Floats are not Reals

- **■** Example 1: Is $x^2 \ge 0$?
 - Float's: Yes!
 - Int's:
 - 40,000 * 40,000 → 1,600,000,000
 - $50,000 * 50,000 \rightarrow 2,500,000,000 ?? \rightarrow No! \rightarrow -1,794,967,296$
 - Range of Integer value

» -2,147,483,648 ~ 2,147,483,647









Great Reality #1: Ints are not Integers, Floats are not Reals

- **Example 2:** Is (x + y) + z = x + (y + z)?
 - Unsigned & Signed Int's: Yes!
 - Float's:
 - $(1e20 + -1e20) + 3.14 \rightarrow 3.14$
 - $1e20 + (-1e20 + 3.14) \rightarrow 3.14 ?? \rightarrow No! \rightarrow 0$
 - Why?
 - Addition of float numbers
 - » (1) equalize exponents to the larger one
 - » (2) adjust mantissa according to the modified exponent
 - $-1e20 \rightarrow 0b0 11000001 01011010111100011101100$
 - $-1e20 \rightarrow 0b1 11000001 01011010111100011101100$
 - $-3.14 \rightarrow 0b0 \ 10000000 \ 10010001111010111000011$

Computer Arithmetic

Does not generate random values

Arithmetic operations have important mathematical properties

Cannot assume all "usual" mathematical properties

- Due to finiteness of representations
- Integer operations satisfy "ring" properties
 - Commutativity, associativity, distributivity
- Floating point operations satisfy "ordering" properties
 - Monotonicity, values of signs

Observation

- Need to understand which abstractions apply in which contexts
- Important issues for compiler writers and serious application programmers

Great Reality #2: You've Got to Know Assembly

- Chances are, you'll never write programs in assembly
 - Compilers are much better & more patient than you are
- But: Understanding assembly is key to machine-level execution model
 - Behavior of programs in presence of bugs
 - High-level language models break down
 - Tuning program performance
 - Understand optimizations done / not done by the compiler
 - Understanding sources of program inefficiency
 - Implementing system software
 - Compiler has machine code as target
 - Operating systems must manage process state
 - Creating / fighting malware
 - x86 assembly is the language of choice!

Assembly Code Example

■ Time Stamp Counter

- a 64-bit register on x86
- increment on every CPU clock
- readable using rdtsc instruction

Elapsed time of a procedure.

```
unsinged long long t1, t2;
t1 = read_counter();
P();
t2 = read_counter();
printf("execution time of P():%d clock cycles\n", t2-t1);
```

Assembly Code Example

- Implementing through inline assembly code
- Or, linking an object file of an assembly code

Great Reality #3: Memory MattersRandom Access Memory Is an Unphysical Abstraction

Memory is not unbounded

- It must be allocated and managed
- Many applications are memory dominated

Memory referencing bugs especially pernicious

Effects are distant in both time and space

Memory performance is not uniform

- Cache and virtual memory effects can greatly affect program performance
- Adapting program to characteristics of memory system can lead to major speed improvements

Memory Referencing Bug Example

```
typedef struct {
  int a[2];
  double d;
} struct_t;

double fun(int i) {
  volatile struct_t s;
  s.d = 3.14;
  s.a[i] = 1073741824; /* Possibly out of bounds */
  return s.d;
}
```

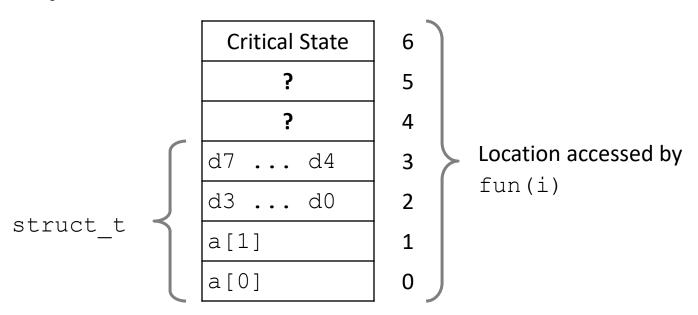
```
fun(0) → 3.14
fun(1) → 3.14
fun(2) → 3.1399998664856
fun(3) → 2.00000061035156
fun(4) → 3.14
fun(6) → Segmentation fault
```

Memory Referencing Bug Example

```
typedef struct {
  int a[2];
  double d;
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```

```
fun(0) → 3.14
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fun(6) → Segmentation fault
```

Explanation:



Memory Referencing Errors

C and C++ do not provide any memory protection

- Out of bounds array references
- Invalid pointer values
- Abuses of malloc/free

Can lead to nasty bugs

- Whether or not bug has any effect depends on system and compiler
- Action at a distance
 - Corrupted object logically unrelated to one being accessed
 - Effect of bug may be first observed long after it is generated

How can I deal with this?

- Program in Java, Ruby, Python, ML, ...
- Understand what possible interactions may occur
- Use or develop tools to detect referencing errors (e.g. Valgrind)

Great Reality #4: There's more to performance than asymptotic complexity

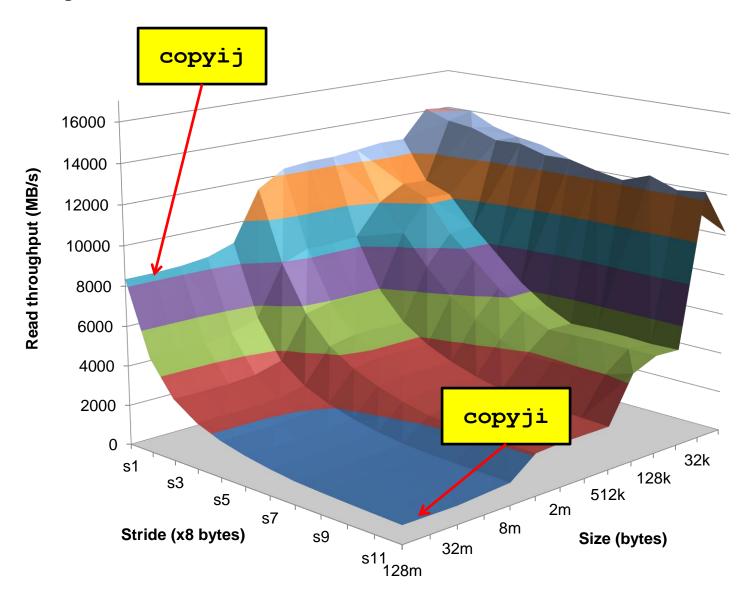
- Constant factors matter too!
- And even exact op count does not predict performance
 - Easily see 10:1 performance range depending on how code written
 - Must optimize at multiple levels: algorithm, data representations, procedures, and loops
- Must understand system to optimize performance
 - How programs compiled and executed
 - How to measure program performance and identify bottlenecks
 - How to improve performance without destroying code modularity and generality

Memory System Performance Example

4.3ms 2.0 GHz Intel Core i7 Haswell 81.8ms

- Hierarchical memory organization
- Performance depends on access patterns
 - Including how step through multi-dimensional array

Why The Performance Differs



Great Reality #5: Computers do more than execute programs

- They need to get data in and out
 - I/O system critical to program reliability and performance

■ They communicate with each other over networks

- Many system-level issues arise in presence of network
 - Concurrent operations by autonomous processes
 - Coping with unreliable media
 - Cross platform compatibility
 - Complex performance issues

Amdahl's Law

- When we improve a part of the system, how does it affect the entire performance?
 - T_{old} = Original execution time
 - T_{new} = New execution time

- α = Ratio of the improved part
- k = Degree of improvement
- $T_{new} = (1 \alpha)T_{old} + (\alpha T_{old})/k$
- (Speed up) $S = T_{old}/T_{new} = \frac{1}{(1-\alpha) + \alpha/k}$
- ex1) improve 60% of the system (α = 0.6) by three times (k = 3)
 - S = 1/[(1-0.6) + 0.6/3] = 1.67
- ex2) improve 30% of the system (α = 0.3) by twenty times (k = 20)
 - S = 1/[(1-0.3) + 0.3/3] = 1.4
- Conclusion: Try to improve a major part of the system as possible.