A TUTORIAL FOR NEW ENGINEERS

How to Think About Systems

We start with one engineer's Tuesday. We design Instagram, then run the whole method again on a parking app you could ship this month.

Where this goes.

01 A Tuesday — one engineer, wake to sleep

02 The words — naming what you just watched

03 Protagonist and antagonist — where a design starts

04 Solution types — which answer is wanted

05 Six kits — data, compute, network, scale, reliability, security

06 Two designs — Instagram, then a parking app from nothing

PART ZERO

Before anything is a system, it is a Tuesday.

THE RULES OF THIS PART

Maya has been an engineer for seven months. Today she wants one thing.

She wants pull request 482 merged before the release window closes at four o'clock.

Watch her Tuesday. This part uses no systems-design words on purpose. Everything Part one names happens here first, so you see each idea before you have a word for it.

MORNING

Five things happen to Maya before she reaches her desk.

06:55

The shower

The water runs too hot. Maya turns it down and it runs too cold. She corrects it twice more.

08:05

The plan

Maya reads her plan for the day. It has the build, the review and the four o'clock window in it.

08:20

A signal failure

A signal fails and the train sits nine minutes. Nobody aboard can change it. Maya opens Instagram.

08:50

At her desk

Maya uses the wifi, the VPN and the package registry. She thinks about none of them.

09:05

The registry goes down

Three teams cannot build. Maya thinks about the registry now.

09:15 · STANDUP

Four people met twice this week, and only one meeting worked.

On Monday each person reported to the manager in turn. That meeting ran twenty minutes and settled nothing. Today they talked to each other instead, and Maya learned in one sentence that Priya had already read the code she was about to start.

The four people did not change. Only who talked to whom changed.

Get 482 merged before the window closes.

MAYA WROTE THIS ON A STICKY NOTE AT 09:30. NOTHING ELSE IS WRITTEN ON IT.

09:45 · IN THE WAY

Four limits shape Maya's day, and she chose none of them.

The build takes twelve minutes

Every push costs twelve minutes. Maya cannot shorten that or skip it.

The team is three people

Nobody on the team is free to review her code this morning.

The one reviewer is asleep

He is six time zones away and will not read anything before three o'clock.

Production data stays in production

She may not copy it to her laptop, even to reproduce the bug.

MIDDAY

By lunchtime, other people decide what Maya works on.

10:15

A favor

Another team asks Maya to fix a flaky test in their repository. She says no.

10:40

A page

The checkout service is failing. Maya does not own it, and she is second on call.

11:00

The loop

Maya pushes, waits twelve minutes, finds her own mistake and pushes again. She does this twice.

12:30

A branch left behind

Her second attempt sits half-finished on a branch, and she forgets it.

13:30

The queue

Five pull requests are waiting for one reviewer, and he is asleep.

AFTERNOON

Maya spends the afternoon answering questions about the work she is not doing.

15:30

The spiral

Nobody has reviewed 482. Someone asks for a status update, so Maya stops work to answer.

15:50

She stops answering

Maya turns down the fourth request and saves every reply until four o'clock.

16:00

The window closes

The reviewer wakes at three. One hour is not enough, and 482 does not merge.

16:30

One thing held

The main branch stayed ready to deploy every minute of the day.

17:00

A small realization

Maya checks the history. Nothing merges on Thursdays either, and nobody planned that.

22:40 · THE REVEAL

You just watched a system run for sixteen hours.

Her day had a purpose it did not meet, and one hour that decided it. Two things ran in circles: the shower, and the questions about 482. Something she never thinks about held the day up until the registry stopped at nine. And main stayed deployable throughout — the one promise that held.

The reviewer woke at three and the window shut at four. Everything Maya did that day reached the outcome only through that hour. She spent twenty of its sixty minutes answering questions about 482, instead of waiting ready to act on whatever came back.

Ten ideas ran through that day. Part one gives you the words for them.

01

PART ONE

Now we give each of those things its name.

WORD ONE · SYSTEM

A system is parts, connected, doing something together.

Maya's morning is one. The shower, the train, the laptop and the registry between them do what none of them does alone: put Maya at a desk, able to work, at nine.

IN MAYA'S DAY

Nothing on that list is a system. The order she does them in is.

LISTING THE PARTS IS EASY

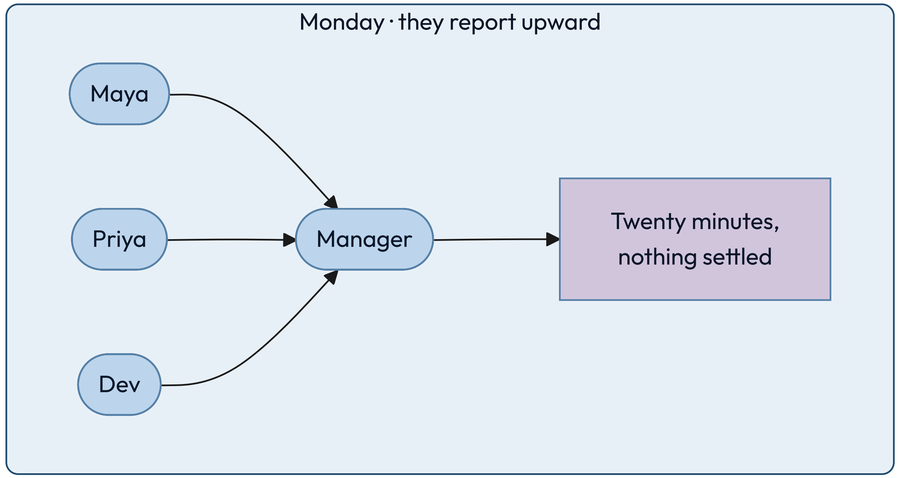
Shower, train, laptop, registry. Anyone can write that list.

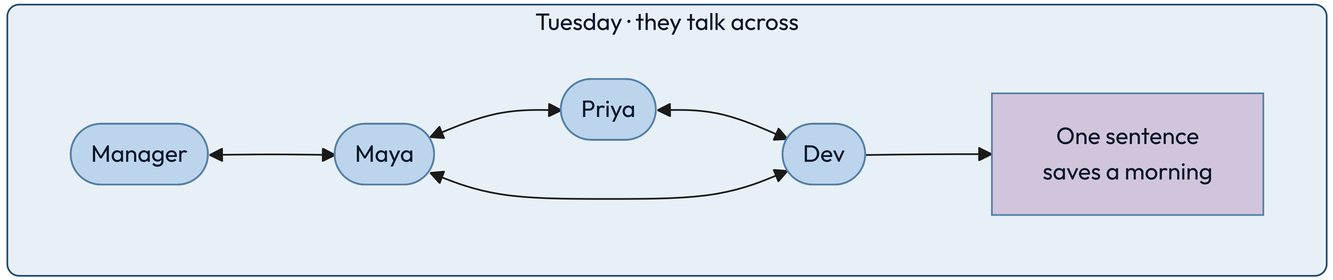
THE CONNECTIONS ARE THE SYSTEM

Change what depends on what, and the behavior changes with the same parts.

09:15 · THE SAME FOUR PEOPLE

Rewire a system without changing a single part, and it behaves differently.





WORD TWO · PURPOSE

A purpose becomes visible at the moment you miss it.

Maya wrote it on a sticky note at nine thirty: get 482 merged before four. At four o'clock it had not merged, and the gap between those two facts is the clearest thing in her whole day.

IN MAYA'S DAY

She wrote hers down at 09:30, which is why she could tell at 16:00 that she had missed it.

WRITE IT DOWN OR YOU WILL DRIFT

An unwritten purpose gets quietly replaced by whatever arrived most recently.

A SYSTEM'S PURPOSE IS WHAT IT DOES

Not what its owners say it does. Watch the outputs, not the mission statement.

WORD THREE · BOUNDARY

The boundary is the line between what you change and what you ask for.

At ten fifteen another team asked Maya to fix a flaky test in their repository. She said no. That refusal is the boundary, and she could feel exactly where it was because saying no was uncomfortable.

IN MAYA'S DAY

She could decline the favor. She could not decline the release window.

INSIDE IS WHAT YOU CHANGE

Your code, your schema, your deploys, the alerts that wake you.

OUTSIDE IS WHAT YOU NEGOTIATE

The reviewer's time zone, the registry, the build, another team's repository.

WORD FOUR · ENVIRONMENT

The environment is everything that arrives without asking.

A signal failure held her train for nine minutes. A page pulled her into an outage in a service she does not own. Neither was load, neither was a bug, and neither was hers.

IN MAYA'S DAY

She had no say over either. She did have a say over what happened to 482 while she was gone, and had not decided it in advance.

IT IS NOT THE SAME AS LOAD

Regulation, a partner's outage and a colleague's time off are environment too.

YOU DESIGN FOR IT, NOT AGAINST IT

You cannot stop the page. You can decide what happens to 482 when it comes.

WORD FIVE · PROCESS

A process is a repeatable transformation, and it runs at a rate.

Push, wait twelve minutes for the build, read a comment, fix, push again. Maya ran that loop twice. Every process has inputs, a transformation, outputs, a rate, and something it leaves behind.

IN MAYA'S DAY

Twenty-four minutes of her day were spent waiting for a machine, and she chose none of them.

THE RATE IS PART OF THE DEFINITION

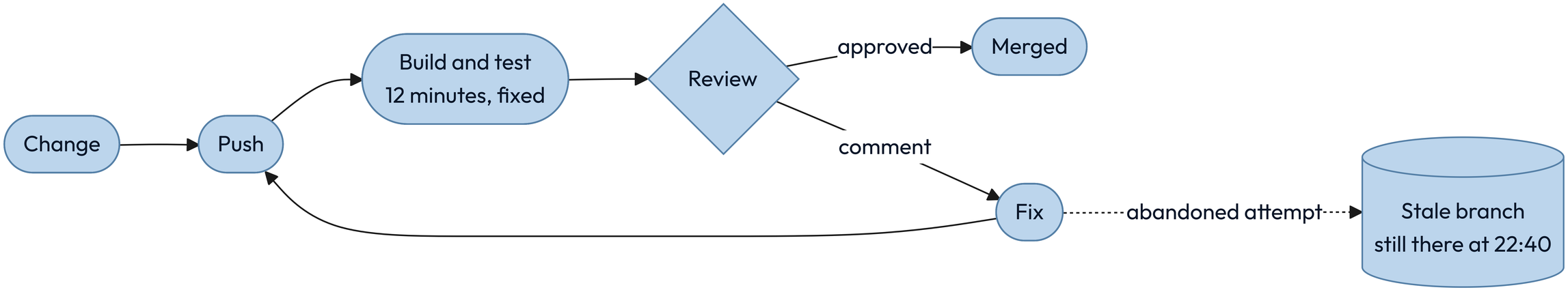
"Run the tests" is not a process until you say twelve minutes, and how often.

LEFTOVER STATE IS WHERE BUGS LIVE

Her abandoned branch is the same shape as a half-applied database migration.

11:00 · THE LOOP, DRAWN

Every process leaves something behind, and that is the part nobody draws.



WORD SIX · MODEL

A model leaves things out on purpose. What beats you is the thing you left out by accident.

Maya's plan for Tuesday was a model. It held the build time, the review, and the four o'clock window. It did not hold the fact that her reviewer sleeps six time zones away, and that one gap decided the day.

IN MAYA'S DAY

Her plan was right about everything in it and silent about the thing that beat her.

EVERY DIAGRAM IN THIS DECK IS A MODEL

Each one throws away something true so you can see one thing clearly.

SAY WHAT YOU LEFT OUT

An omission you can name is a simplification. One you cannot name is a bug.

WORD SEVEN · CONSTRAINT

A constraint is a limit that removes options rather than adding caveats.

Four limits shaped Maya's day, and each came from somewhere different. A twelve-minute build. A team of three. A reviewer asleep six time zones away. A rule about production data.

IN MAYA'S DAY

Physical, economic, human and legal limits, all four of them landing before ten in the morning.

REAL CONSTRAINTS CARRY NUMBERS

"The build is slow" is a complaint. "Twelve minutes" is something you can design against.

SOME YOU CHOSE, AND CAN REVISIT

A team of three is a decision. The speed of the build is arithmetic.

WORD EIGHT · INVARIANT

An invariant is a claim that stays true on the day everything else fails.

Main was deployable every minute of Maya's Tuesday. It was true while the registry was down, while she was paged, and at four o'clock when 482 did not land.

IN MAYA'S DAY

She missed her goal and broke no invariant, and those are two different kinds of bad day.

WRITE THEM AS SENTENCES

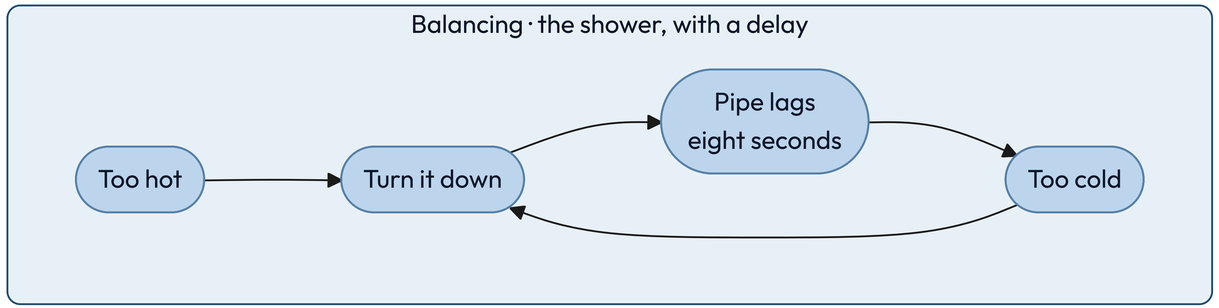
"Main is always deployable." "A balance is never negative." Then name the alert.

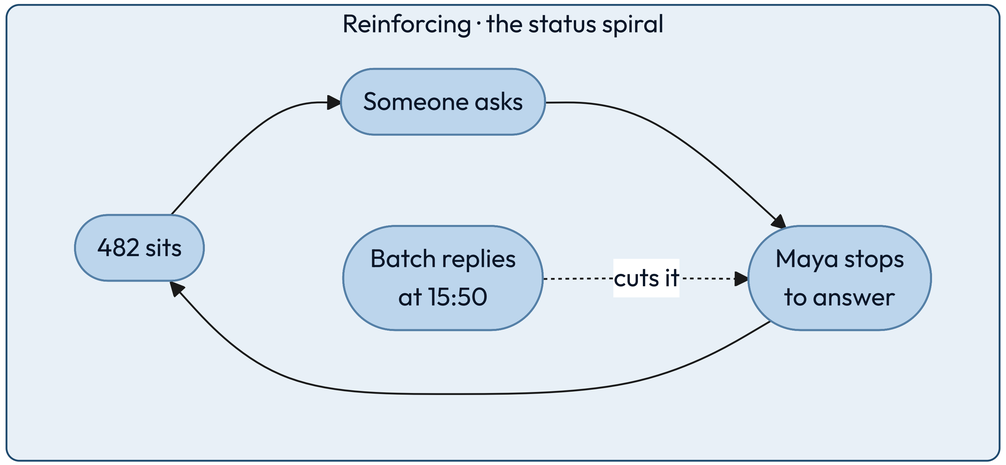
THEY CONSTRAIN OTHER PEOPLE

A real invariant changes what your teammates are allowed to merge.

06:55 AND 15:30 · TWO LOOPS

Two loops look identical until you find the sign.





✦ *A balancing loop pushes back toward a target, and the eight-second delay in the pipe is exactly why she overshoots. A reinforcing loop pushes harder in the direction it is already going.*

Infrastructure is what you only notice on the day it stops.

Wifi, the VPN, the package registry, the build fleet, the identity provider. Maya used every one of them before nine o'clock and thought about none of them until 09:05, when the registry went down and three teams stopped.

THE TEST

If it vanishing stops several unrelated things at once, it is infrastructure.

BORING IS THE REQUIREMENT

It earns its place by being predictable, never by being interesting.

IT IS SOMEONE ELSE'S SYSTEM

Your infrastructure is another team's product, with its own boundary and invariants.

17:00 · THE LAST WORD

Emergence is a pattern the parts fall into, that none of them contains.

Nothing merges on Thursdays. Maya checked six weeks. It is a rhythm, and nobody built a rhythm.

Four things feed each other rather than add up. Work that misses the window waits at the front of tomorrow. A review that comes back with comments sends the same change to the back of the queue. The queue is served one hour a day. And nobody ships into a weekend, so Thursday is the last window of the week — the day the backlog is deepest and the day it has to clear.

No policy names Thursday. You find it by watching the queue for six weeks.

Five words describe the thing itself.

You did not learn these from a definition. You watched each one happen first, which is the order that sticks.

01	**System**	Parts, connected, with a purpose.	The whole morning.
02	**Purpose**	What it does, not what it says.	The gap at four o'clock.
03	**Boundary**	What you change, not ask for.	The favor she declined.
04	**Environment**	What arrives uninvited.	The page.
05	**Process**	A transformation with a rate.	Push, build, review.

Three more are the tools you think with.

You never handle the system itself. You handle a drawing of it, its limits and its promises — standing on infrastructure you notice only when it stops.

01	**Model**	A simplification you chose.	Her plan for the day.
02	**Constraint**	A limit that removes options.	Twelve minutes.
03	**Invariant**	What must never go false.	Main is deployable.

THE TRANSLATION

Every move in Maya's day has a name in software.

IN HER TUESDAY	IN A SYSTEM	WHAT IT DECIDES
Five pull requests, one reviewer	The bottleneck	Where any improvement has to land
Declining the fourth status ask	Admission control	Whether load sheds or the system collapses
A twelve-minute build	A fixed cost per attempt	How many attempts a day can hold
The abandoned branch	Leftover state	What a retry finds when it arrives
Nothing merges on Thursdays	Emergence	What no single owner can fix

YOUR TURN

Take the turnstile at a station and name its parts before you turn the page.

You already know how one works. Write down five things: its purpose, its boundary, one constraint, one invariant, and the infrastructure it stands on. Two minutes, on paper, and cover the next slide until you have them.

Do not hunt for a clever answer. The point is that you now have words for a thing you have walked past a thousand times without ever describing.

ONE ANSWER

Here is a turnstile, in the words you now have.

01	**Purpose**	Let paying people through, stop everyone else. Counting riders is a side effect.
02	**Boundary**	The gate, its reader, the local rules. Not the fare service it asks.
03	**Constraint**	One person at a time, about a second each. That sizes the hall.
04	**Invariant**	Never open without a valid fare. Never close on a person.
05	**Infrastructure**	Power, the link to the fare service, the floor.

02

PART TWO

Every design starts with someone who wants something, and something in the way.

THE FRAME

Name the person and the force, and everything downstream has an answer.

Juniors skip both, almost every time. They write "users" instead of one person with one goal, and they write "scale" instead of a force with a number on it. Neither produces a single design decision.

The protagonist is who the system is for. The antagonist is what makes serving them hard. Everything downstream — the purpose, the boundary, the invariants, the first move — falls out of that pair.

THE DRILL

Two sentences, ninety seconds, and four things you no longer have to guess.

```
Protagonist:  <name> wants to <do X> so that <Y>.
Antagonist:   But <W> — and W is <a number>.

  Purpose       one sentence: what counts as this working
  Boundary      what is mine when W happens, and what I only ask for
  Invariant     what must stay true or <name> stops trusting this
  First move    what W does to the cheapest design that could work
```

✦ *Run it on a turnstile, then a cash machine, then the thing on your screen right now. It gets fast.*

THE DRILL, ON TUESDAY

Maya is the protagonist of her own day, and the interruptions are the antagonist.

MAYA WANTS 482 MERGED BEFORE FOUR O'CLOCK

So that a bug affecting real users stops affecting them, and so she is not carrying it into next week.

BUT SHE IS INTERRUPTED SIX TIMES

A page, a favor, four status requests. Not one of them unreasonable, and each costs the reload, not just the minute.

THE TRAP

Most systems have several protagonists, and you must say which one loses.

Instagram has at least four: the reader opening the app, the ordinary poster, the celebrity with five hundred million followers, and the engineer carrying the pager. They want incompatible things.

Naming one as the protagonist decides who waits. Say who you are not designing for, out loud, and most of the arguments later in the design turn out to be about that.

THE JOIN

The antagonist chooses which kind of solution you are allowed to build.

THE ANTAGONIST IS…	YOU ARE BUILDING	BECAUSE
Nobody knows if they exist	An MVP	The risk is demand, not load
Growth — the protagonist multiplies	A scaled system	The limit is real and namable
Cost on a path you have profiled	An optimized system	The measurement came first
Physics — you are near a real bound	An optimal system	Only here is a proof worth it
A person who wants in	Security work, at any rung	Required at every rung

INSTAGRAM · THE CASTING

Our protagonist reads on a phone and gives us about a second.

She opens the app a dozen times a day, on a cellular network, usually while doing something else. She wants photographs from the people she chose, recent, ranked, and hers. The poster is a supporting character: he tolerates a spinner, and every time the design has a choice, work goes onto his side.

THE READER'S DEMAND

A page on her phone in under a second, which is 200 ms of server time plus the network.

THE POSTER CAN WAIT

Uploading, transcoding and delivery may all take their time. Nobody is watching.

CASTING DECIDES THE DESIGN

It is why the feed is built when someone posts, not when someone reads.

The antagonist is not scale. It is the shape of the follow graph.

Scale is a quantity, and quantities have a price you can pay. The thing you cannot buy your way out of is a distribution. Follower counts are heavy-tailed: almost everyone has a few hundred, and a handful have hundreds of millions. No single account looks like the average, and the far end sits six orders of magnitude from the middle.

THE NUMBER THAT MATTERS

The median account has about 150 followers. The largest has five hundred million.

ONE ALGORITHM CANNOT SERVE BOTH

That gap is why the design ends up with two paths instead of one.

HOLD ON TO THIS

Every hard decision in Part five is this one fact again.

YOUR TURN

Cast the last app you opened, in the same two sentences.

Pick something ordinary: a maps app, a chat client, the thing your team ships. Fill in both lines, exactly as the drill has them:

```
Protagonist:  <name> wants to <do X> so that <Y>.
Antagonist:   But <W> — and W is <a number>.
```

The second line is the hard one. If you cannot put a number on W, you have just found the first thing worth going and measuring.

ONE ANSWER

A maps app casts the same way, in two sentences.

THE PROTAGONIST

A driver already moving wants the next turn early enough to take it, so that the road and the screen never compete. Everyone else — the person searching, the person saving a place, the person reading reviews — is slower and can wait.

THE ANTAGONIST

The signal drops in the tunnel, and the tunnel is ninety seconds long. That number is the design: ninety seconds of route has to be on the phone before the phone goes quiet.

03

PART THREE

"Design Instagram" is five different questions.

BEFORE THE BOXES

Before you design anything, say which kind of answer is wanted.

The same words — design a photo sharing app — have five legitimate answers that share almost no architecture. A build that takes a week and a build that takes two years are both correct, for different questions.

Gall's law says it plainly: a complex system that works is invariably found to have evolved from a simple system that worked. You climb this ladder. You do not parachute onto it.

Each rung costs more and commits harder than the one below it.

You move up when the rung you are standing on stops paying, and you move up one rung at a time.

01	**MVP**	Buys information fast.	Does anyone want this?
02	**Scaled**	Holds up as load grows.	Will it survive success?
03	**Optimized**	Cuts cost on a known path.	Where is the money going?
04	**Optimal**	Provably best, for one model.	What is the real limit?
05	**Specialized**	An advantage nobody can copy.	What can only we do?

RUNG ONE · MVP

An MVP is an experiment wearing a product's clothes.

Its job is to produce a decision, not to last. Every hour spent making it durable is an hour spent on a system you may correctly delete next month.

THE TELL

The riskiest thing is whether anyone wants it, and being wrong is cheap.

BUY SIMPLICITY, NOT CAPACITY

One database, one service, one region, boring technology, nothing clever anywhere.

NAME THE EXIT IN ADVANCE

Write down the number that means "stop, this now has to be built properly."

RUNG TWO · SCALED

A scaled system survives ten times the load without a rewrite.

You have stopped buying information and started buying headroom. The question moves from "does it work" to "what breaks first, and how do I move that limit."

THE TELL

Growth is real, and the current design has a ceiling you can point at.

A BIGGER BOX FIRST, THEN THE FOUR MOVES

Vertical scaling is not one of the four; it is what buys time before you need them, and it is reversible. Then reduce, duplicate, defer — and spread last, because partitioning is the one you cannot undo cheaply.

COST PER UNIT STARTS COUNTING

Cost per request stops being noise and becomes the second constraint on every choice.

RUNG THREE · OPTIMIZED

Optimizing means moving one measured number on one hot path.

Optimization without a profile is decoration. You need the measurement first, the target second, and a willingness to accept the complexity you are about to add.

THE TELL

A profile shows which tenth of the work is most of the cost.

PREMATURE IS HALF THE QUOTE

Knuth also wrote that we should not pass up the critical three percent. Both halves.

COMPLEXITY IS THE INVOICE

Each optimization narrows the assumptions the system is allowed to break.

RUNG FOUR · OPTIMAL

Optimal means provably best against an objective you wrote down.

Rarer than it sounds. It needs a stated objective, a stated model, and a proof or a bound — and it is optimal only for the assumptions you fixed in place.

THE TELL

The objective is a function and the constraints are inequalities, on paper.

THE PROOF IS AGAINST A MODEL

Change the assumptions and the optimal answer changes underneath you.

IT AGES BADLY

An optimal design pinned to last year's hardware is a legacy system with a certificate.

A specialized system trades generality for something nobody can copy.

Custom silicon, a purpose-built storage engine, a scheduler that knows your physics. You give up flexibility, portability and hiring pool for a capability the market cannot sell you.

THE SIGNAL

The advantage is durable, measurable, and central to why customers choose you.

THE COST NEVER ENDS

You now maintain what everyone else gets free from a vendor, forever.

ALMOST NOBODY IS HERE

Most teams reaching for this rung needed the optimized one and got excited.

THE RULE

Climb one rung at a time, and only on evidence.

MOVE UP WHEN THE CURRENT RUNG FAILS ON A MEASUREMENT

A number, a profile, a named risk. Not a feeling that things are getting big.

MOVE UP EXACTLY ONE RUNG

Jumping from MVP to optimal buys rigor for assumptions nobody has tested yet.

MOVE BACK DOWN WHEN THE EVIDENCE CHANGES

A rewrite that simplifies is a legitimate move.

A design is finished when taking one more thing out would break it.

Antoine de Saint-Exupéry wrote that perfection arrives not when there is nothing left to add, but when there is nothing left to take away. Dieter Rams said it in three words: less, but better. Both are tests you can run on a whiteboard at four in the afternoon.

RUN IT ON EVERY BOX

Delete it on paper and follow what happens. If nothing downstream changes, it was never holding anything up.

RUN IT ON EVERY NUMBER

A figure nobody can act on is decoration. Turn it into a threshold or take it out.

LESS, BUT BETTER, IS NOT LESS

Rams designed for decades of use. Removing something the system needs is not restraint, it is a defect.

YOUR TURN

Three requests arrive this week. Name the rung each one is asking for.

One: a founder wants to know whether anybody will pay for something that does not exist yet. Two: a service that works fine is about to take ten times the traffic, and nobody knows what gives first. Three: the bill for a single endpoint is now larger than the team that owns it.

Write a rung for each and one sentence on what it costs. Then turn the page.

ONE ANSWER

The question sets the rung. The size of the company does not.

01	**Will anybody pay?**	MVP. You are buying information and paying for it with everything you will throw away.
02	**Ten times the traffic**	Scaled. You are buying headroom and paying in machines and the coordination they need.
03	**The bill is too large**	Optimized. You are buying back cost and paying in flexibility. Profile first, or you are decorating.

04

PART FOUR

Six kits: fifteen things you can reach for, and the invariants behind all of them.

HOW TO READ A KIT

Reach for it when, walk away when, and the constraint you inherit.

One shape for every entry you choose between — a store, a runtime, a delivery tier, a quota — so you can hold two options side by side without re-reading a manual. Scale and reliability are mostly practices rather than choices, so those two kits run on diagrams, patterns and invariants instead. Every kit opens with a diagram and closes with its invariants.

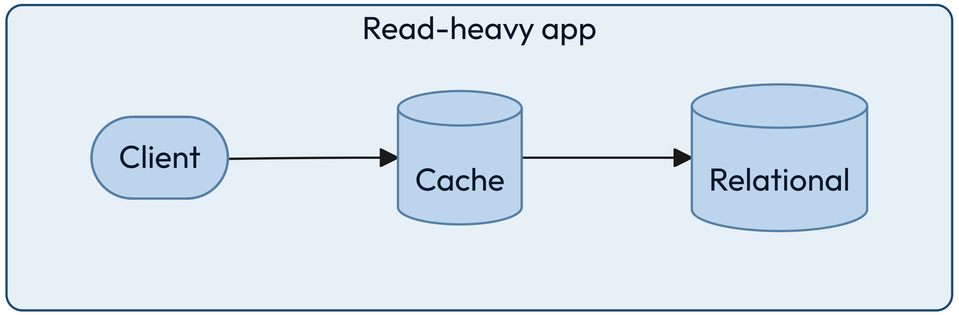
The entries name concepts, not products. Products turn over every few years. The constraint that an append-only store puts on your reads does not.

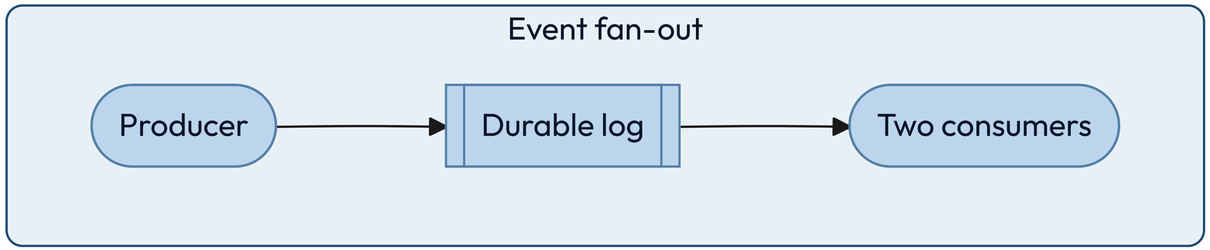
DATA

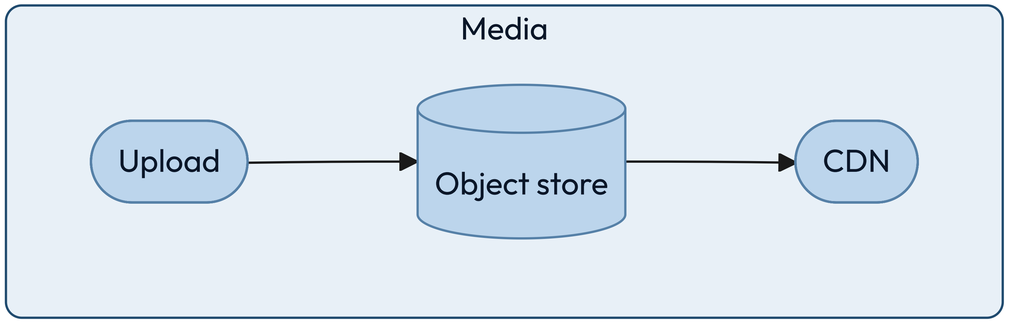
Every storage choice is a bet about how you will read it later.

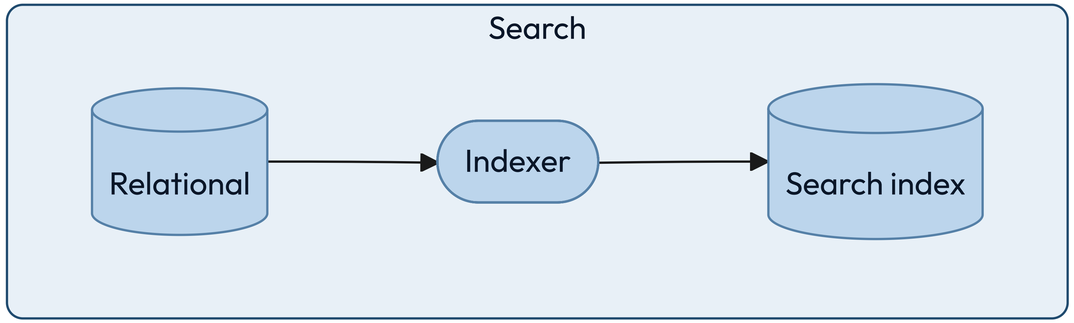
DATA KIT · THE PATTERNS

Almost everything you will ever store fits one of a few shapes.









The product name is the last thing you decide about a store.

Answer these before anyone says a brand. Most database arguments are really a disagreement about question two.

01	**Shape**	How the data is structured.	Rows, documents, or edges?
02	**Access**	How it is read and written.	Lookups, ranges, or scans?
03	**Consistency**	What a reader may see.	Must it be the latest?
04	**Scale**	Size, throughput, growth.	Does it fit one machine?
05	**Cost**	Money, operations, skill.	Who runs it at 3am?

DATA KIT · THE THREE PASSES

Those five questions do not all fire at once. The choice runs in three passes.

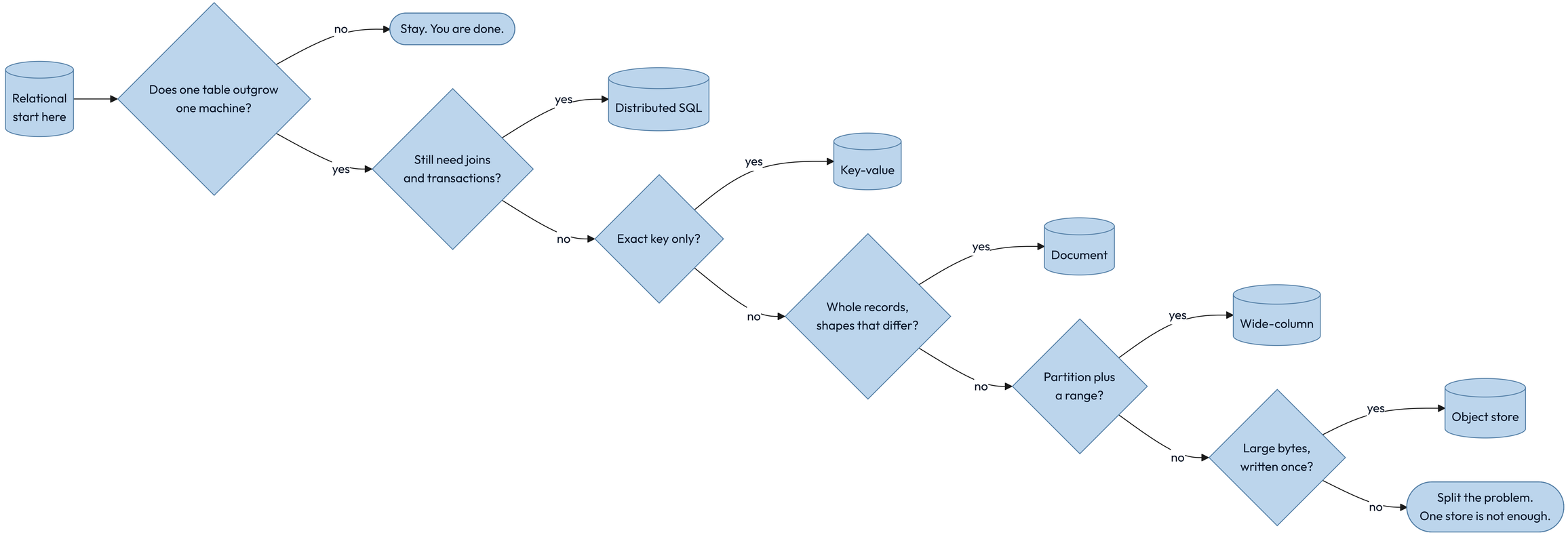
Pass one asks about shape and access, and it picks the store you start from. The next slide draws that pass as a tree, and for most systems it is the whole decision.

Pass two asks whether anything you need is a capability no shape provides: similarity, ranked text, proximity, a push the instant a row changes. That last one can delete a tier — a store that streams to the client is the socket service you were about to write.

Pass three runs only when two candidates both fit. It settles the tie on how each behaves at 3am, not on what it stores. Open at pass three and you will argue about pricing before you have agreed what the data is.

DATA KIT · PASS ONE

Start at relational. Every branch below is a reason to leave it.



DATA KIT · RELATIONAL

A relational store is the default until you can name why it is not.

Reach for it when

Entities relate, writes touch several at once, and the questions will keep changing.

Walk away when

One table outgrows one machine's writes, or the schema genuinely differs per row.

The constraint you inherit

Joins and transactions need a coordinator. Self-sharding loses both; a distributed engine charges a round trip.

DATA KIT · DISTRIBUTED SQL

An engine that partitions itself is the first answer now, not the last resort.

Reach for it when

One table outgrew one machine and you still want joins, transactions and a schema.

Walk away when

Every query is a single exact key, or the budget cannot carry cross-partition coordination.

The constraint you inherit

A transaction across partitions pays a round trip, so related rows still belong together.

DATA KIT · KEY-VALUE

A key-value store is a hash map with an operations team.

Reach for it when

You always know the exact key, and handing the value back is the store's only job.

Walk away when

You need to ask any question at all about what is inside the value.

The constraint you inherit

There is no second way in. Every new query means a new key you write and maintain yourself.

DATA KIT · DOCUMENT

Choosing a document store means choosing to keep one object whole.

Reach for it when

A read fetches one self-contained object, and the shape varies between records.

Walk away when

The same fact lives in many documents and has to stay consistent across them.

The constraint you inherit

Denormalized copies. Every update to a shared fact becomes a fan-out — one write you now owe to many places.

DATA KIT · WIDE-COLUMN

A wide-column store buys enormous write throughput and freezes your queries.

Reach for it when

Writes are relentless, and every read is a partition key plus a sorted range.

Walk away when

Query patterns are still moving, or you need a transaction across partitions.

The constraint you inherit

The primary key is the schema. Changing how you query means rewriting the data.

DATA KIT · OBJECT STORE

An object store is the cheapest home for a write-once file.

Reach for it when

Items are large, written once, fetched by key, and must survive for a decade.

Walk away when

You need to modify part of an object, or to list and filter by what is inside it.

The constraint you inherit

Listing is slow and expensive. The index of what you stored belongs somewhere else.

DATA KIT · PASS TWO

A capability no shape provides means adding a store beside the source, not swapping the source out.

CAPABILITY	THE QUESTION IT ANSWERS	WHAT YOU ADD
Similarity	What is closest in meaning to this?	A vector index
Ranked text	Which document best matches these words?	A search index
Proximity	What is within two kilometers?	A geospatial index
Live push	What changed, the moment it changed?	A store that streams
Retention	What did this metric do all month?	A time-series store
Traversal	Who is reachable from here?	A graph store

DATA KIT · GRAPH

A graph store is for questions that traverse, not questions that filter.

Reach for it when

The query walks relationships of unknown depth: reachability, paths, recommendations.

Walk away when

You have relationships but only ever join two hops. A relational store does that faster.

The constraint you inherit

Traversals resist partitioning, so scaling out is genuinely harder here than anywhere else.

HALFWAY

Two stores you filed as derived are holding truth.

An object store and a time-series store are usually where a fact first lands — nothing regenerates a photograph or last Tuesday's CPU samples. They are sources of truth wearing the clothes of a derived tier.

The genuinely derived stores are the search index, the vector index, the cache and the warehouse. Anything derived must be rebuildable, must be allowed to lag, and must never be the only copy.

DATA KIT · CACHE

A cache converts a correctness problem into a timing problem.

Reach for it when

Reads repeat, the source is expensive, and slightly stale is genuinely acceptable.

Walk away when

Staleness is unsafe, or the working set is larger than the cache will ever be.

The constraint you inherit

Invalidation. You now own a second copy whose wrongness is measured in seconds.

DATA KIT · DURABLE LOG

A durable log turns "do it now" into "do it reliably, soon."

Reach for it when

Producers outpace consumers, or several systems need to see the same events.

Walk away when

The caller needs the result inside the same request.

The constraint you inherit

At-least-once delivery. Every consumer must be idempotent or you will charge somebody twice.

DATA KIT · THE SCAN

Six stores, side by side, on the four things that usually decide it.

STORE	ACCESS	CONSISTENCY	SCALES BY	WEAK AT
Relational	Key, range, join	Strong on the leader	Replicas, then partitioning	Cross-shard writes
Distributed SQL	Key, range, join	Strong across partitions	Horizontal	Cross-partition transactions
Key-value	Exact key	Engine-specific	Horizontal	Rich queries
Document	Key, secondary index	Engine-specific	Horizontal	Facts split across documents
Wide-column	Partition plus range	Tunable	Horizontal	New query patterns
Graph	Traversal	Engine-specific	Horizontal, with effort	Cutting the graph across machines

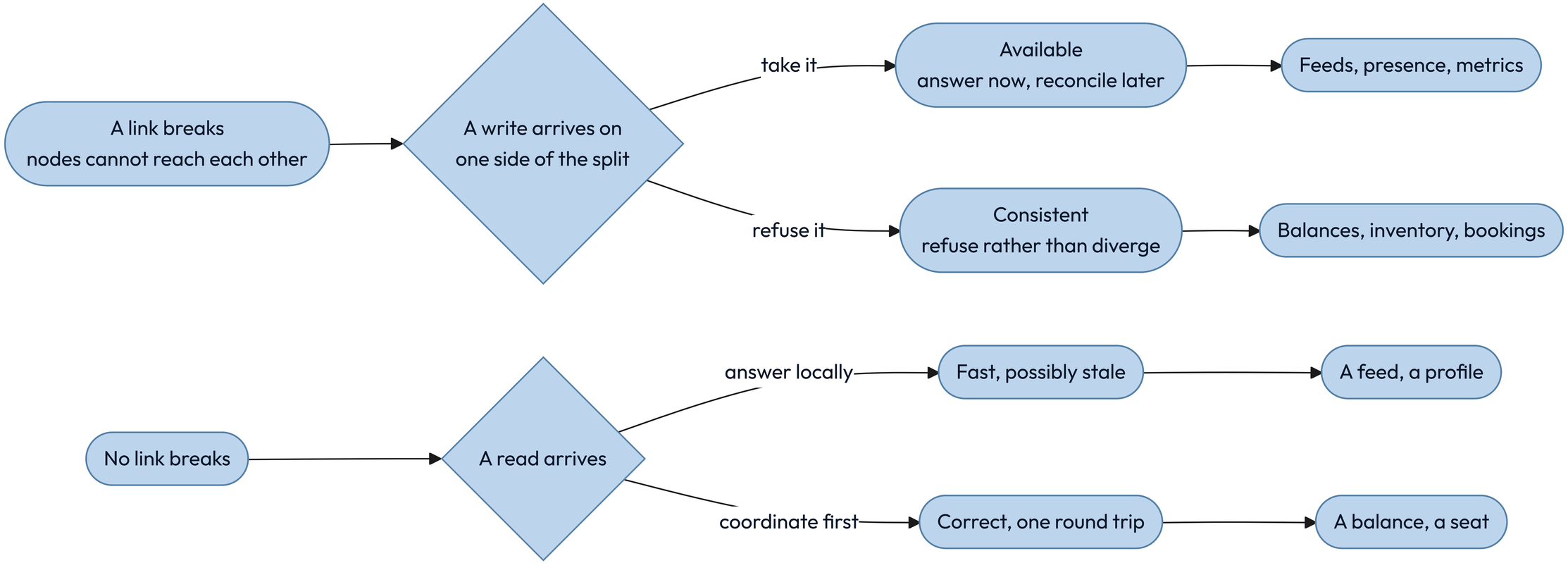
DATA KIT · PASS THREE

When two stores both fit, what settles it is how each one behaves on a bad night.

ASK	WHY IT SETTLES THE TIE
Who carries the pager?	A managed service, or your own team at 3am
What does a lost node cost?	Some engines lose capacity, some lose writes in flight
How does the bill grow?	Per gigabyte, per request, or per machine-hour
How does data get out?	A dump, a change stream, or nothing at all
How long is a restore?	Backups are easy. The restore is the number
Who already runs one here?	A store your team knows beats a better one nobody has run

DATA KIT · UNDER A PARTITION

CAP prices a partition. Consistency charges you the rest of the time.



DATA KIT · CONSISTENCY

Consistency is not one thing, and most arguments are about which kind you meant.

01	**Linearizable**	Every read sees the latest write. Costs a coordination round trip, every time.
02	**Read your writes**	You always see your own edits. Anyone who has just typed something expects at least this.
03	**Eventual**	Replicas agree in the end. Cheapest, and correct for feeds, counters and presence.

DATA KIT · TWO PHILOSOPHIES

ACID and BASE answer different questions.

Ask whether coordinating now costs less than reconciling later.

I
ACID

Coordinate first, so the data is never observably wrong. You pay in latency and in how far you can partition.

→

II
BASE

Diverge now, converge later. You pay in application code, because every reader must tolerate stale state.

An index is a second copy of your data, sorted for exactly one question.

Indexes make reads fast by making writes slower and storage larger. Each one is a promise to maintain a sorted structure on every single insert, update and delete, forever.

THE TELL

You can name the exact query each index exists to serve, out loud, right now.

THE BILL ARRIVES ON WRITES

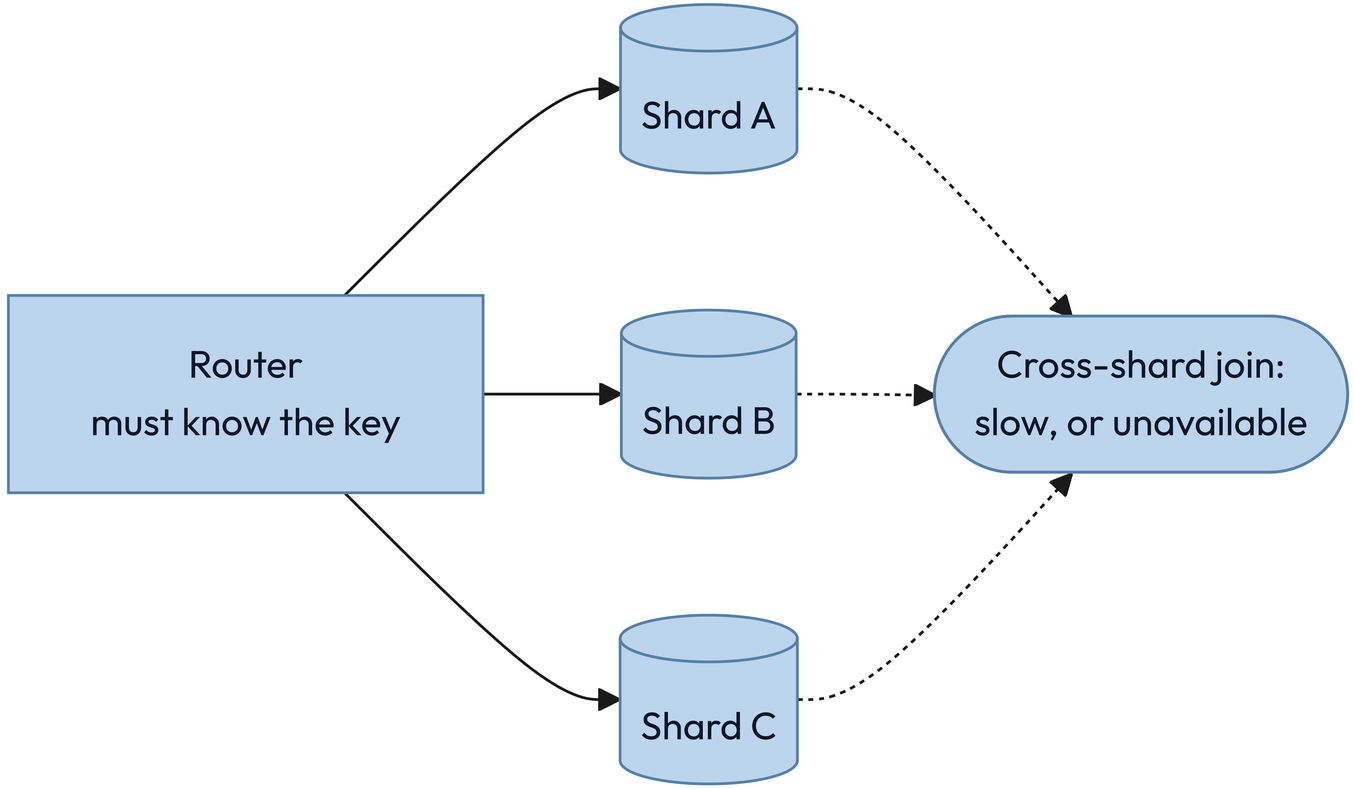
Five indexes mean five extra structures touched on every insert.

HOW MANY ROWS IT REMOVES, NOT HOW MANY VALUES IT HAS

An index earns its keep by how rare a match is. Three evenly spread values narrow nothing; three where one is rare narrow almost everything.

DATA KIT · PARTITIONING

Sharding buys throughput by giving up the questions that cross shards.



DATA KIT · THE SHARD KEY

The shard key is the decision you cannot undo cheaply.

By hash of the key

Spreads load evenly and destroys range queries. The safe default.

By range

Keeps ranges fast and invites a hot spot on the newest range.

By tenant or region

Matches both the access pattern and the law, until one tenant grows enormous.

KEY INSIGHT

Name the query that will now cross shards. If you cannot, you have not chosen a key — you have chosen a hash.

DATA KIT · ORDERING AND IDENTITY

Two machines never agree on the time, so never let a clock alone decide an order.

Reach for it when

Anything is sorted, paged or deduplicated. A sortable id carries time and writer in one value.

Walk away when

One writer owns the sequence and a unique id breaks the ties. Then a clock is enough.

The constraint you inherit

An id you sort by can never change, and every page cursor comes to depend on it.

DATA KIT · CHANGING A SCHEMA

Changing a column on a running system takes five steps, not one.

STEP 01

Expand

Add the field. Nothing writes it, nothing reads it, both versions still work.

STEP 02

Write both

New code fills it alongside the old one. Nothing reads it, so mistakes are cheap.

STEP 03

Backfill

Fill rows written before step two, in restartable batches.

STEP 04

Cutover

Move reads across. Both are still written, so going back costs nothing.

STEP 05

Contract

Stop writing the old field, then drop it.

DATA KIT · WHY THOSE FIVE STEPS

A deploy is never atomic, so every change must tolerate the version beside it.

For the minutes or hours a rollout takes, two versions of your code run against one store. Both have to work. That is the entire reason for the five steps: each one is safe while the step before it is still live, which is exactly what a rollout cannot promise you about any bigger jump.

The same rule governs an API other people call. Add fields, never repurpose them. Make anything new optional. Remove nothing until you can show that nobody calls it — and if you cannot show that, you have not earned the right to remove it.

DATA KIT · DELETING FOR REAL

A delete is a fan-out that reaches every copy you ever made.

The row itself

A tombstone, not a gap. Replicas have to learn the row is gone.

Every derived copy

Caches, feeds, search indexes, aggregates. Each holds its own copy and each needs telling.

The backups

You cannot rewrite a backup. Set a retention window and let the copy expire instead.

The bytes at the edge

A CDN serves what it cached. Purging is eventual, so signed URLs need short lives.

KEY INSIGHT

Design the delete when you design the write. Retrofitting one across six stores is a quarter of somebody's year.

DATA KIT · THE INVARIANTS

Four sentences hold, or the data design is not one you can defend.

01

One source of truth per fact

Every other copy is derived and says so.

02

Every derived copy rebuilds — and deletes

Unattended from the source, and gone from all of them on request.

03

Every queue consumer is idempotent

At-least-once is the only delivery you get.

04

Every write path states its consistency

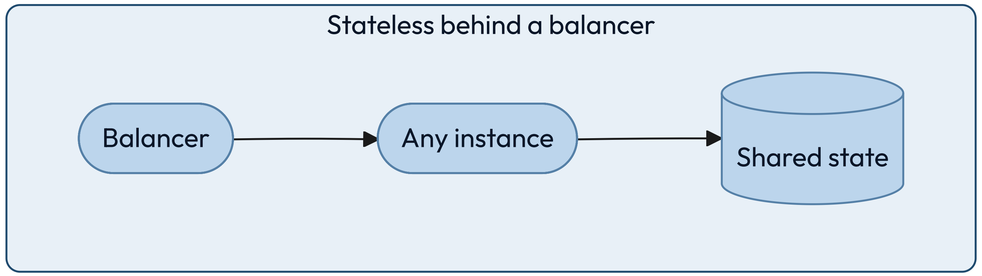
"Whatever the database does" is not a level.

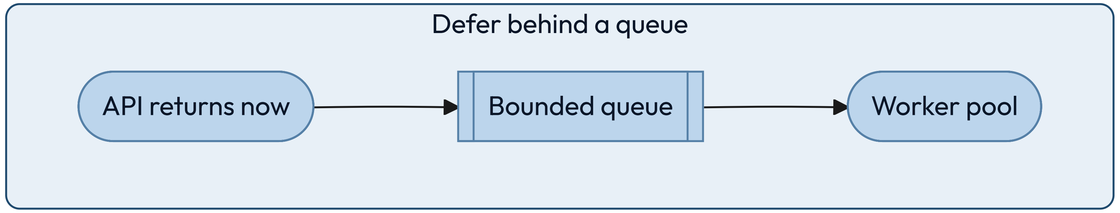
COMPUTE

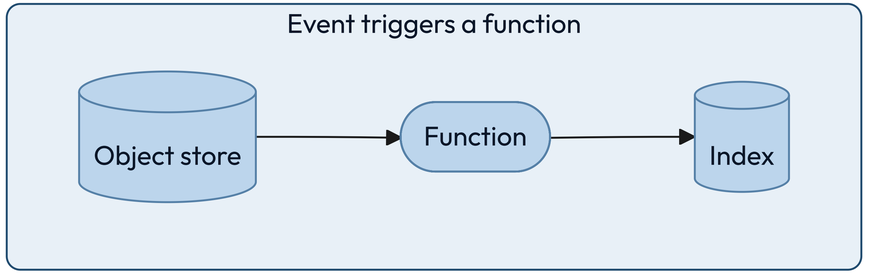
Choosing compute is choosing how much of the machine you still own.

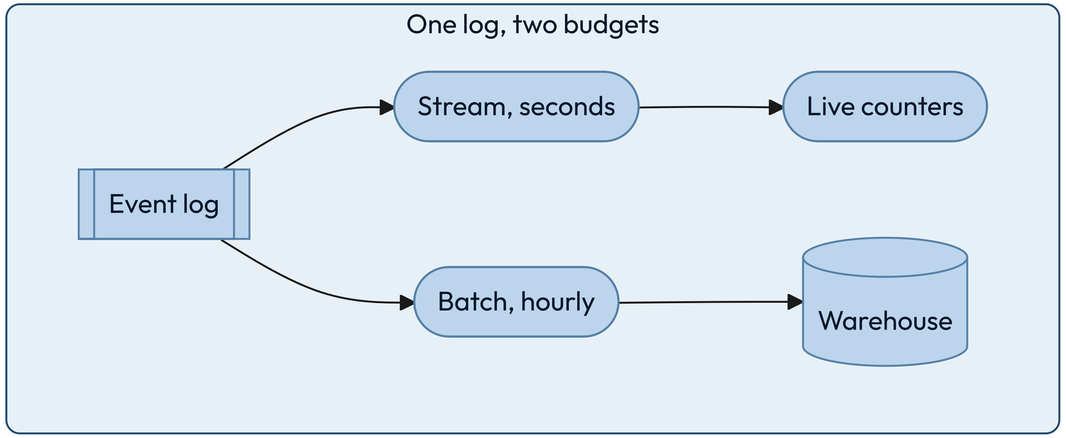
COMPUTE KIT · THE PATTERNS

Every compute choice hands more of the machine to somebody else.









COMPUTE KIT · VIRTUAL MACHINES

A virtual machine is what you reach for when the process outlives the request.

Reach for it when

The workload runs continuously, holds state in memory, or needs unusual kernel settings.

Walk away when

Traffic is spiky and idle machines start to dominate the bill.

The constraint you inherit

Patching, capacity planning, and the gap between staging and production are now yours.

COMPUTE KIT · CONTAINERS

A container makes the deployable unit identical everywhere it runs.

Reach for it when

You run many services, deploy often, and want one packaging story across all of them.

Walk away when

You run one small service. The scheduler that places containers on machines costs more than it saves at that size.

The constraint you inherit

The orchestrator is now infrastructure, with its own failure modes and its own pager.

COMPUTE KIT · FUNCTIONS

A function is compute you rent by the millisecond, so idle costs you nothing.

Reach for it when

Traffic is spiky or rare, the work is short, and per-request isolation is welcome.

Walk away when

Runs are long, or steady traffic makes per-request pricing more expensive than a machine.

The constraint you inherit

Cold starts, unless you pay to keep instances warm — which is paying for idle again.

COMPUTE KIT · THE PROPERTY

Statelessness is what makes a machine replaceable.

A STATEFUL INSTANCE

It holds something no other instance has: a session, a lock, a warm cache, an open connection. Scaling means moving that state, and a crash means losing it.

›

A STATELESS INSTANCE

Any instance serves any request because the state lives somewhere else. Scaling is arithmetic and failure is a routing change.

COMPUTE KIT · THE INVARIANTS

Anything you run should satisfy all four before it meets real traffic.

01

Any instance can be killed without a customer noticing

If that is false, you have state you have not named yet.

02

Deployments are reversible

A rollback is a routine operation, not an incident response.

03

Capacity is a number somebody owns

Not "it autoscales" — the ceiling, the cost, and who gets paged at it.

04

Startup does not depend on startup order

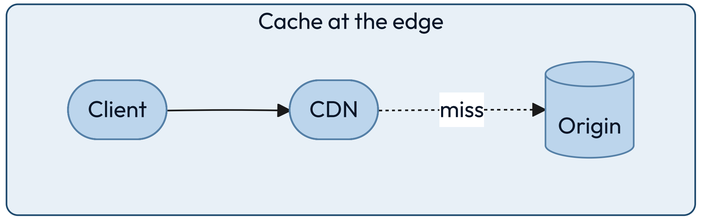
Services retry into each other instead of requiring a sequence.

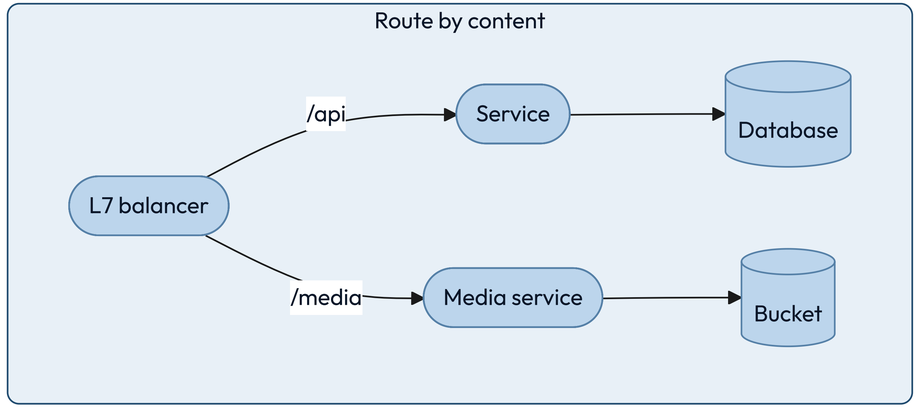
NETWORK

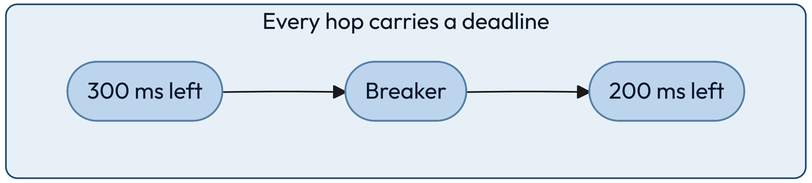
Distance is the one cost you cannot optimize away.

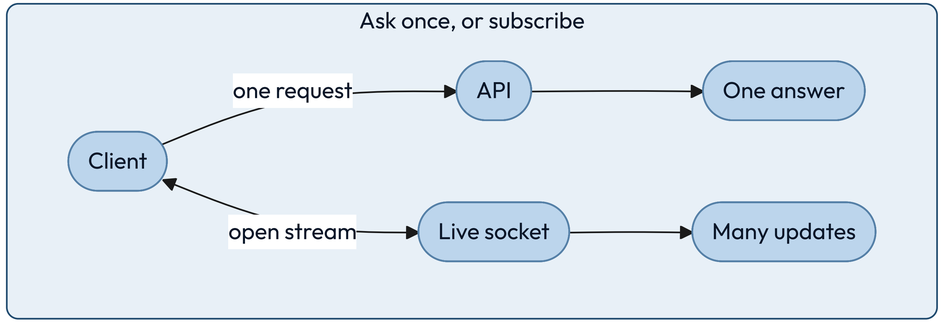
NETWORK KIT · THE PATTERNS

Distance charges a different price at every hop.



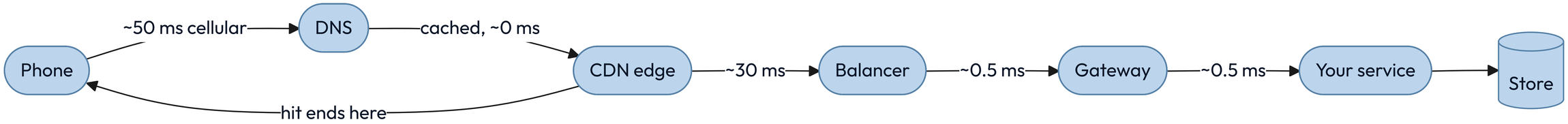






NETWORK KIT · THE REQUEST PATH

A tap crosses several hops before your code sees it, and each one spends budget.



NETWORK KIT · THE NUMBERS

Most latency arguments end the moment somebody says the actual numbers.

01	**Memory read**	100 ns
02	**Datacenter hop**	0.5 ms
03	**Cross-continent round trip**	150 ms

WHY THAT NUMBER IS FINAL

Light in fiber covers about 200 kilometers per millisecond.

Nothing changes that. The measured 150 milliseconds is well above the straight-line floor, because packets do not travel in straight lines and every hop queues.

A design that needs three sequential intercontinental round trips has spent half a second before it executes an instruction. Replication, caching and edge delivery all exist to buy that distance back, and none of them makes it free.

NETWORK KIT · THE CDN

A CDN moves bytes closer to readers, and ends the connection there too.

Reach for it when

The content is large, popular, and identical for many people.

Walk away when

Nothing is shared and nothing is far. A personal response still wins the handshake back.

The constraint you inherit

A second copy with its own staleness. Version immutable URLs; purge the rest, and time the purge.

NETWORK KIT · TIMEOUTS

A request with no timeout is a resource leak waiting for a bad afternoon.

Every waiting request holds a connection, a thread and some memory. Under a slow dependency, unbounded waits turn one struggling service into a queue of stuck callers — which is the shape of the page that pulled Maya in at twenty to eleven.

WHAT GOOD LOOKS LIKE

Every outbound call has a deadline, and the deadline shrinks as it propagates.

RETRIES NEED A BUDGET

Retrying without a cap turns a brief failure into a flood you built yourself.

BACKOFF MUST BE RANDOM

Synchronized retries arrive together and rebuild the spike you just survived.

NETWORK KIT · THE INVARIANTS

A call that leaves your process owes you four things.

01

Every remote call has a deadline

Inherited from the caller, and always shorter than the caller's own.

02

Every retry has a budget and jitter

Bounded attempts, randomized delays, and a breaker when the target is down.

03

Every write is idempotent or keyed

The network will deliver your request twice. Decide now what that means.

04

Distance appears in the design

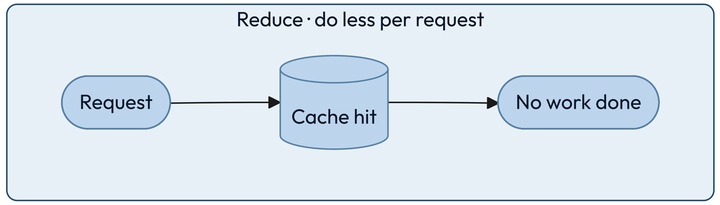
Round trips between regions are counted on purpose, not discovered in production.

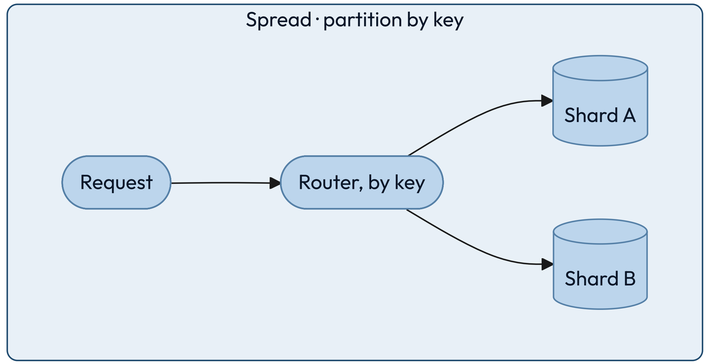
SCALE

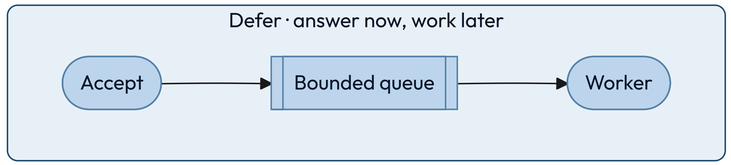
Every scaling change is one of four moves.

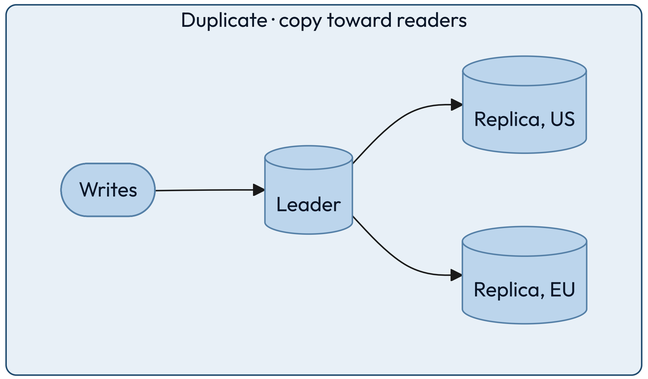
SCALE KIT · THE PATTERNS

The four moves, and what each one costs you.









SCALE KIT · THE ONE FORMULA

Little's law ties concurrency, throughput and latency together.

$$L = \lambda W$$

WHERE

- L — requests in flight at once
- λ — arrivals per second
- W — time each one spends inside

USING IT

The formula sizes your thread pool before anyone guesses.

At 2,000 requests per second averaging 50 milliseconds each, 100 requests are in flight at any moment. That is your minimum concurrency, and no tuning makes it smaller while the other two numbers hold.

Maya's Tuesday is the same arithmetic at human scale. Two builds a day, twelve minutes each, leave the pipeline idle for almost the whole day, so nothing she pushes ever queues — which is how you know the twelve minutes was never her bottleneck. The reviewer was. Find your own numbers — arrivals on the load balancer, duration in the handler.

WHEN IT GOES WRONG

Tripled latency breaks a bounded pool and an unbounded one in opposite directions.

Leave the pool unbounded and the arithmetic runs forward: latency triples, so the requests in flight triple with it, and a machine sized for the good day runs out of memory.

Bound it at 100 and concurrency cannot rise, so throughput falls instead. `100/0.15s` is about 667 a second against the 2,000 still arriving, and the queue in front grows without limit until something sheds it.

Neither is a tuning problem. The first is why pools have ceilings, and the second is why every queue behind one needs a ceiling too.

SCALE KIT · TAIL LATENCY

The average request is a fiction, and your users live in the tail.

A page that makes 100 parallel calls waits for the slowest one. Give each call a one-percent chance of being slow. Then 63 percent of pages hit at least one slow call. That is *1 - 0.99^100*, and you can redo it on a napkin.

THE CHECK

A percentile tells you what one person felt. An average tells you what the fleet must carry. Report the 99th per dependency, and keep the average for capacity only.

FAN-OUT AMPLIFIES IT

More parallel calls turn a rare slow response into a common slow page.

HEDGING BUYS IT BACK, ON A BUDGET

Send a duplicate after the 95th percentile and take whichever answers first. A hedge is a second request, so cap it — a few percent of traffic. Unbudgeted, it is the reinforcing loop again, arriving exactly when you are already slow.

SCALE KIT · CACHING

Each caching pattern owns a different failure.

Cache-aside

The app fills the cache on a miss. Simple, and it stampedes on a cold key unless one reader fills it while the rest wait on that one fill.

Read-through

The cache fetches for you. Cleaner code, and now the cache is on the critical path.

Write-through

Write both together. Always fresh, and every write pays the cache's latency.

Write-behind

Write the cache, flush later. Fastest writes, and a crash loses them.

KEY INSIGHT

Choose by which failure you can survive, not by which pattern reads best in code.

SCALE KIT · IDEMPOTENCY

Idempotency is what makes a retry safe, and retries are not optional.

Networks duplicate, clients retry, queues redeliver. The only question is whether the second delivery is harmless or charges somebody twice.

THE RULE

Every mutating endpoint takes a client-supplied key and deduplicates on it.

THE KEY COMES FROM THE CALLER

Generated before the first attempt, reused on every retry of that same intent.

STORE THE RESULT, NOT THE FACT

A repeat returns the original answer, not an error saying it already happened.

SCALE KIT · THE INVARIANTS

Nothing here is scalable until the last of the four is true.

01

The bottleneck is named and measured

Not suspected. A number, a graph, and the resource it belongs to.

02

Every queue is bounded

An unbounded queue turns a throughput problem into a memory outage.

03

Admission control sheds load before the system collapses

A balancing loop from Part one, and the move Maya made at 15:50 when she stopped answering.

04

Adding a machine is routine

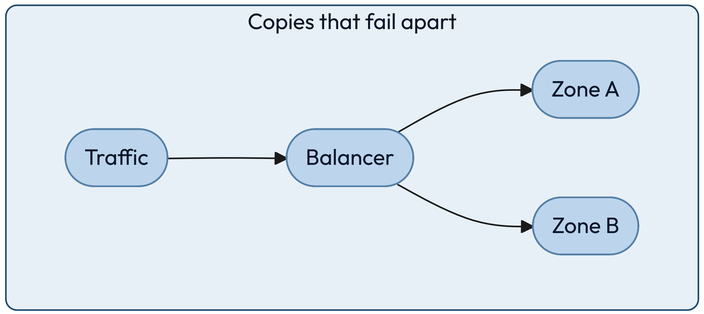
No manual steps, no rebalancing outage, no cold-cache stampede.

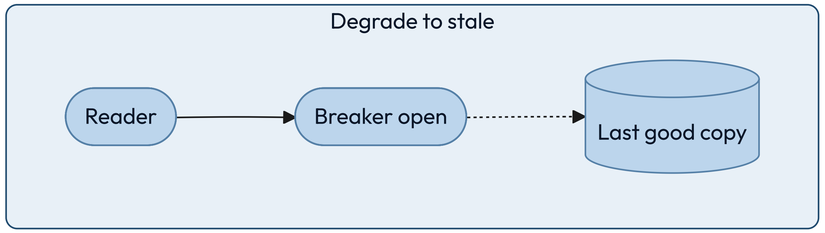
RELIABILITY

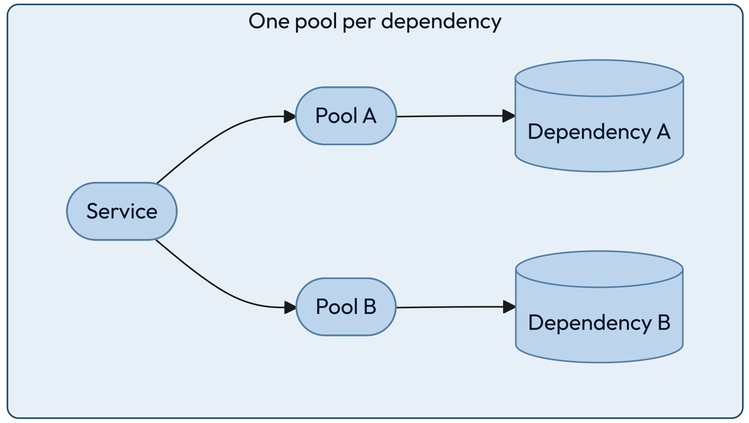
Failure is the environment, not the exception.

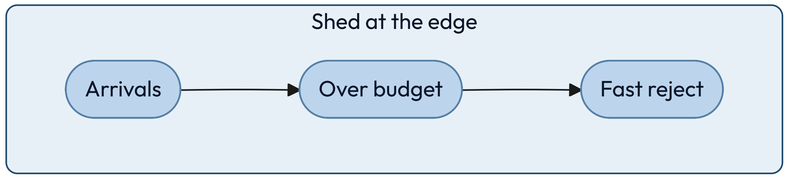
RELIABILITY KIT · THE PATTERNS

Each pattern keeps one failure from becoming every failure.



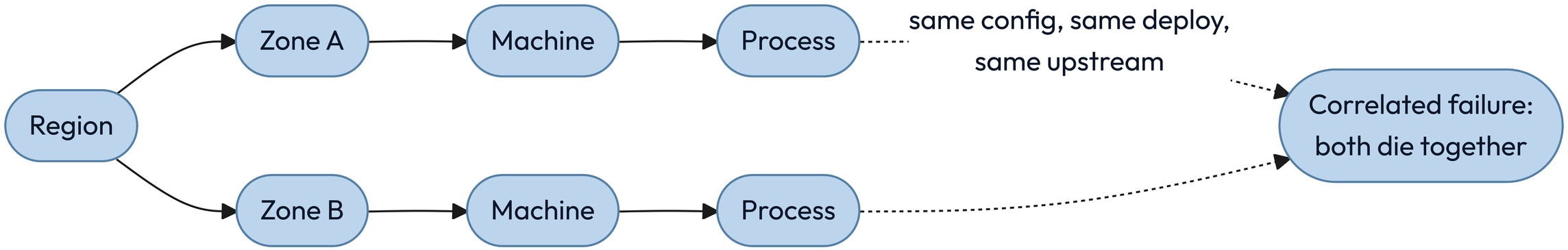






RELIABILITY KIT · FAILURE DOMAINS

Redundancy only helps when the copies can fail apart.



RELIABILITY KIT · CONTAINMENT

A slow dependency takes the whole building unless something bounds the wait.

STEP 01

Timeout

Bound the wait, so a slow callee cannot hold your resources.

STEP 02

Retry with backoff

Recover from a blip, with a budget and jitter. Retries without them are a reinforcing loop.

STEP 03

Circuit breaker

Stop calling something that is clearly down, and let it recover.

STEP 04

Bulkhead

Give each dependency its own pool, so one queue cannot drain them all.

RELIABILITY KIT · THE OBJECTIVES

Four words people use interchangeably decide whether you may ship this week.

01	**SLI**	The measurement: the share of requests served under 300 milliseconds.
02	**SLO**	Your internal target for it, such as 99.9 percent over 28 days.
03	**SLA**	The external promise, with money attached, and always looser than the SLO.
04	**Budget**	What the objective lets you spend. Exhausted, feature work stops.

RELIABILITY KIT · THE NINES

An availability target is a budget that shrinks fast.

01	**99 percent**	3.65 days per year
02	**99.9 percent**	8.8 hours per year
03	**99.99 percent**	53 minutes per year
04	**99.999 percent**	5.3 minutes per year

READING THAT TABLE

Past three nines the humans stop being fast enough.

Fifty-three minutes a year is about one incident with a page, a login, a look at a dashboard and a decision. So four nines does not mean a faster on-call engineer. It means automatic failover and automatic rollback, because a person is no longer in the loop.

Each nine roughly multiplies the cost. Pick the number the business actually needs, and write down what you are choosing not to buy.

RELIABILITY KIT · OBSERVABILITY

Each signal answers a different question, and none of them replaces another.

Metrics

Cheap, aggregate, always on. They tell you that something is wrong.

Logs

Detailed, expensive, one event at a time. They tell you what happened in one case.

Traces

One request across every service. They tell you where the time actually went.

KEY INSIGHT

If you cannot answer "which dependency is slow" inside a minute, you have metrics, not observability.

RELIABILITY KIT · DEGRADATION

A well-designed system gets worse in an order somebody chose.

Under pressure something has to give. Either you decided in advance which features degrade, or the system decides at random and drops the one that takes money.

THE RULE

Every feature has a stated tier, and the lowest tier fails first by design.

READ PATHS OUTLIVE WRITE PATHS

Serving something slightly stale beats serving an error page, for almost every product.

THE FALLBACK IS EXERCISED

An untested degraded mode is untested code running in your worst hour.

RELIABILITY KIT · THE INVARIANTS

Production earns trust one of these at a time.

01

Every dependency has a defined failure behavior

Written down: degrade, queue, or fail fast. Never "we will see."

02

Redundant copies fail independently

Different zones, different deploys, different upstreams. Otherwise it is one copy.

03

Recovery is practiced

Restores, failovers and rollbacks happen on a schedule, not for the first time at 3am.

04

The system tells you before a user does

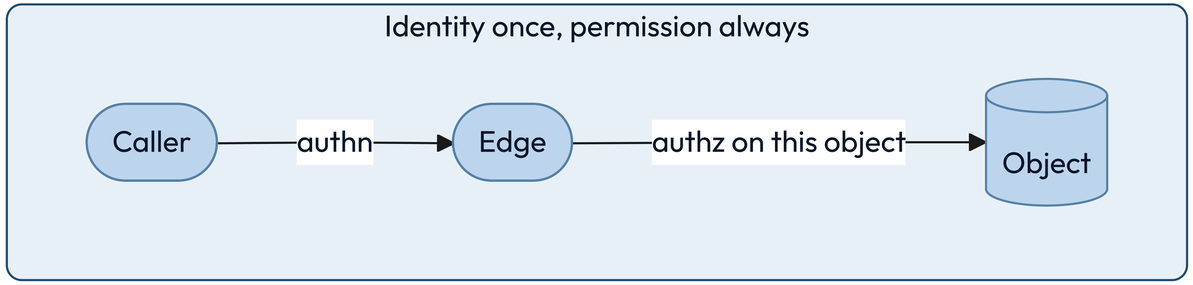
Alerts fire on the indicator, never on the complaint.

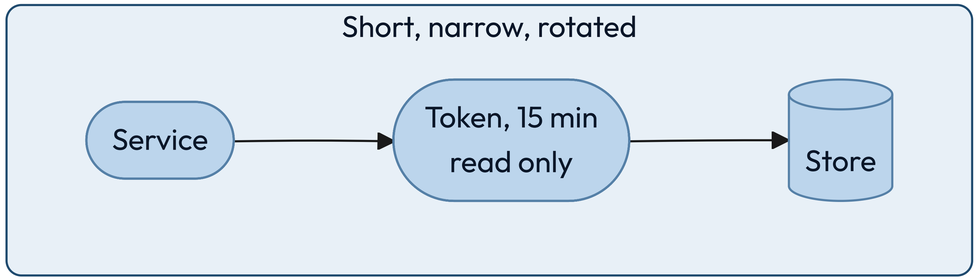
SECURITY

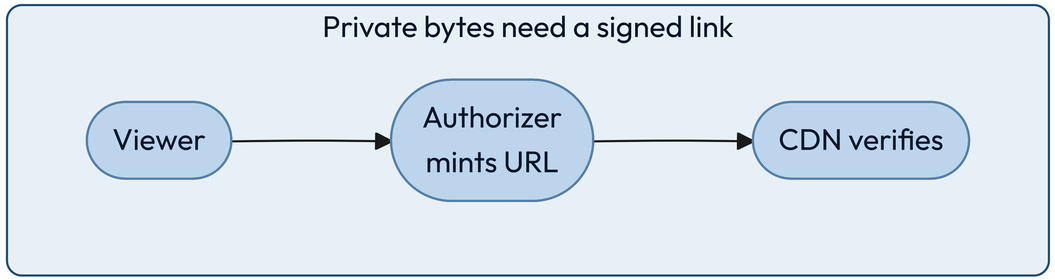
Assume the boundary is already crossed.

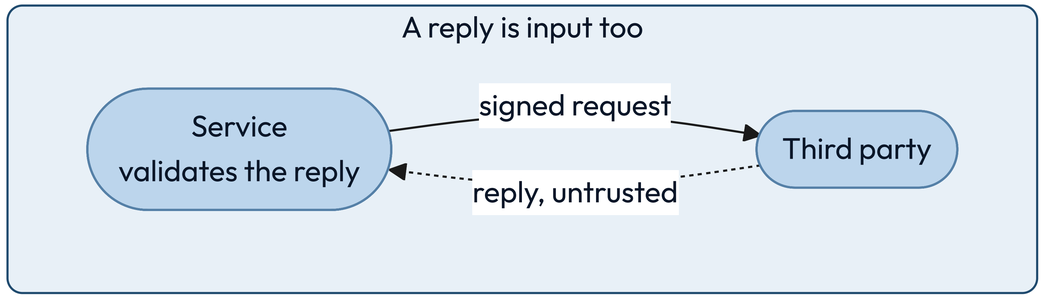
SECURITY KIT · THE PATTERNS

All four defenses assume somebody is already inside.



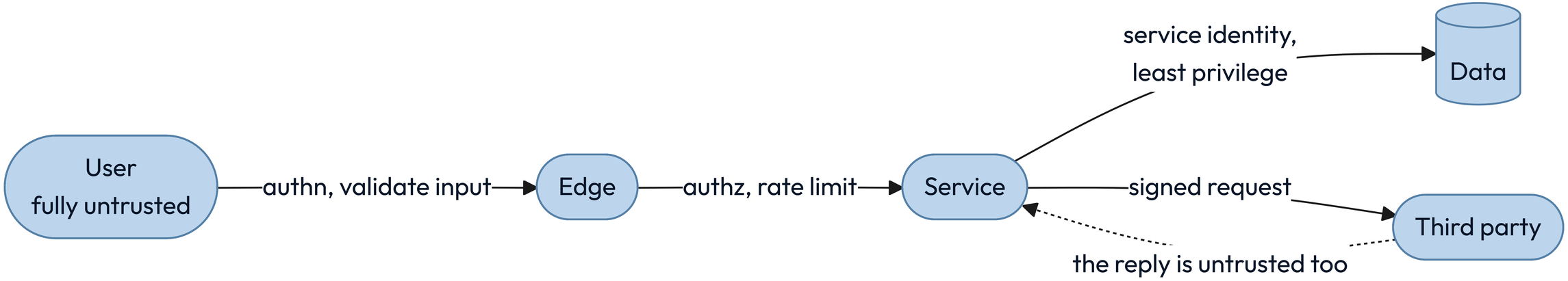






SECURITY KIT · TRUST BOUNDARIES

Every arrow that crosses a trust boundary needs a check on the far side.



SECURITY KIT · THE PAIR

Authentication asks who you are. Authorization asks what you may touch.

AUTHENTICATION

Establishes identity once, at the edge, and hands down a claim anyone downstream can verify. Getting it wrong lets the wrong person in.

AUTHORIZATION

Decides, on every single request, whether this identity may touch this specific object. Getting it wrong lets the right person read somebody else's data.

SECURITY KIT · LEAST PRIVILEGE

Least privilege is measured by what one stolen credential can reach.

Assume any single credential eventually leaks. The design question is not whether that happens; it is how far the person holding it can travel, and for how long.

THE CHECK

You can say in one sentence what each service account could reach if stolen.

SCOPE IT AND EXPIRE IT

Narrow permissions, short lifetimes and automatic rotation beat a careful human.

SEPARATE READ FROM WRITE

Most services need one or the other. Almost none needs to delete.

SECURITY KIT · ENCRYPTION

Encryption has three states, and a fourth thing that outranks all of them.

01	**Transit**	TLS everywhere, including between your own services. Networks are untrusted.
02	**Rest**	Disk and backup encryption. Stops a stolen device, not a stolen credential.
03	**Use**	Decrypted in memory to be processed. The hardest state, and the one attacked.
04	**Keys**	A managed store or an HSM, rotated on a schedule, never in the repository.

SECURITY KIT · THREAT MODELING

Six questions find the holes before somebody else does.

- Can someone pretend to be another identity? — spoofing
- Can data be changed in transit or at rest? — tampering
- Can an actor deny having done something? — repudiation
- Can private data reach the wrong reader? — disclosure
- Can one caller exhaust a shared resource? — denial
- Can a caller gain rights nobody granted? — elevation

SECURITY KIT · YOUR DEPENDENCIES

Most of the code you ship, you did not write and have not read.

Maya's Tuesday stopped when a package registry went down. That same registry is also the shortest path into her build: whoever can publish a version she installs can run their code on her machine, and then inside her deploy.

Pin exact versions and commit the lock file, so the build you tested is the build that ships. Keep an inventory of what you actually depend on, including everything your dependencies pulled in behind you. Take updates on a schedule you chose rather than the moment they appear, and read the diff of anything that touches a credential, a network call, or a build step.

SECURITY KIT · QUOTAS

An authorized caller can still take the whole system down.

Reach for it when

Anything is shared: a database, a queue, a paid third party, an image decoder.

Walk away when

The caller is a trusted internal batch you would rather see fail loudly than throttled.

The constraint you inherit

A limit needs an identity to count against. Use the account, the key and the address.

SECURITY KIT · THE INVARIANTS

Sign your name to a design only when you can say all four.

01

Every request is authorized against the specific object

Not the endpoint. Object-level checks are where the breaches happen.

02

Secrets never enter the repository or a log line

They live in a managed store, are injected at runtime, and they rotate.

03

All input is untrusted, including from your own services

A compromised internal caller is the ordinary case. Design for it.

04

Every privileged action is attributable

An immutable record of who, what, when, and from where.

05

Every dependency is pinned, inventoried and rate-limited

You did not write most of your code, and an authorized caller can still exhaust you.

YOUR TURN

Pick a store for each of these, and name the query that will hurt.

One: a table of orders, ten thousand a day, where support answers questions nobody has asked yet. Two: a session token, looked up on every request and never scanned. Three: eight years of sensor readings, written once, read by device and by day.

Write the store and the one query that will cross a partition. Then turn the page.

ONE ANSWER

One of these never leaves relational. The other two were never relational to begin with.

01	**Ten thousand orders a day**	Relational, for years. Ten thousand a day is nothing, and the unasked questions are the requirement. Nothing crosses a partition, because nothing is partitioned.
02	**A session token**	Key-value. Exact key, no scan, expiry built in. The query that hurts is "which sessions belong to this user".
03	**Eight years of readings**	Wide-column, partitioned by device and ranged by time. The query that hurts is any question about one day across every device.

05

PART FIVE

Now we design the thing Maya opened while the train sat.

INSTAGRAM · THE WORKSHEET

You fill in all of these before the first box, except the eighth. That one you discover.

Protagonist	A reader on a phone. Twelve opens a day, on cellular.
Antagonist	The shape of the follow graph — a tail at 5×10^8 edges.
Purpose	A fresh, personal, ranked page in under a second.
Boundary	In: feed, posts, edges, media. Out: the phone, the network, the law.
Environment	Evening peak ~2.5x the mean, partitions, duplicate deliveries, people who want in.
Constraints	200 ms p99 · 60 opens per write · media durable forever
Invariants	A post reaches eligible followers · media never lost · a like counts once
Bottleneck	Suspected: filling in each post. Confirmed at ~80M reads/s once the read path exists.
Solution type	Scaled. Not an MVP, not optimal.

INSTAGRAM · WHAT IT MUST DO

The whole product rests on four operations.

01

Post a photo

Upload media, attach a caption, publish it to the people who follow you.

02

Follow an account

Build the directed graph that decides whose posts you are eligible to see.

03

Read a feed

A ranked, paginated page of recent posts from the accounts you follow.

04

React

Like and comment, with counts visible on every post.

INSTAGRAM · WHAT IT MUST GUARANTEE

What you promise decides the architecture more than what you build.

01

The feed serves in under 200 milliseconds

Server-side, 99th percentile, measured from the reader's own region.

02

Reads dominate writes by about sixty to one

Counted in opens. Push fan-out nearly erases that underneath: 330K feed writes a second against 420K feed reads.

03

Posts, media and the graph are durable forever

Lose any of the three and nothing rebuilds them. Only derived data rebuilds.

04

Eventual consistency is fine on the feed

A late post and a lagging count are fine. A blank count is not.

INSTAGRAM · THE ESTIMATE, WORKED

Two assumptions and some division give you every number that matters.

Assume 500 million daily users opening the feed twelve times a day, and 100 million photos posted a day. The rest is division. Twelve opens each, times five hundred million people, over the seconds in a day: about `70K` feed opens a second. A hundred million posts over the same day: `1.2K` writes. At two megabytes each: `200 TB` a day. The ratio lands near `60:1`.

One of those numbers does less work than it looks. Change the 500 million, holding opens and posts per person fixed, and the ratio does not move at all.

INSTAGRAM · WHAT EACH NUMBER DECIDES

Four numbers decide something here. The 500 million users decide nothing.

01	**Writes · 1.2 K/s**	One machine could take the posts. Each one becomes hundreds of feed writes, which is the hard part.
02	**Feed opens · 70 K/s**	The number everyone sizes from, and an open is not a request. The real one comes later.
03	**Media · 200 TB a day**	An object store, not a database. Settled right here.
04	**Followers · up to 500 M**	The only number with no ceiling. Two read paths come from this.

YOUR TURN

Redo it for a smaller product before you turn the page.

Fifty million daily users, opening the app four times a day, posting two million photos. Work out reads per second, writes per second, and the ratio between them.

Do it now, on paper, in under a minute. Cover the next slide until you have one. The point is not the answer; it is that you can produce one at all, in the meeting where it is being decided.

THE ANSWER

The ratio moved, and the reason it moved is the lesson.

`50M × 4 = 200M reads/day ÷ 86,400 ≈ 2.3K reads/s` · `2M ÷ 86,400 ≈ 23 writes/s` · `ratio ≈ 100:1`

A thirtieth of the read traffic, and the ratio went from 60:1 to 100:1. It moved because you changed how often each person posts, not just how many people there are. The ratio is opens per user divided by posts per user: it survives being wrong about population, never about behavior.

INSTAGRAM · WHAT THE AVERAGE HIDES

A daily average is not a capacity number, and two multipliers separate them.

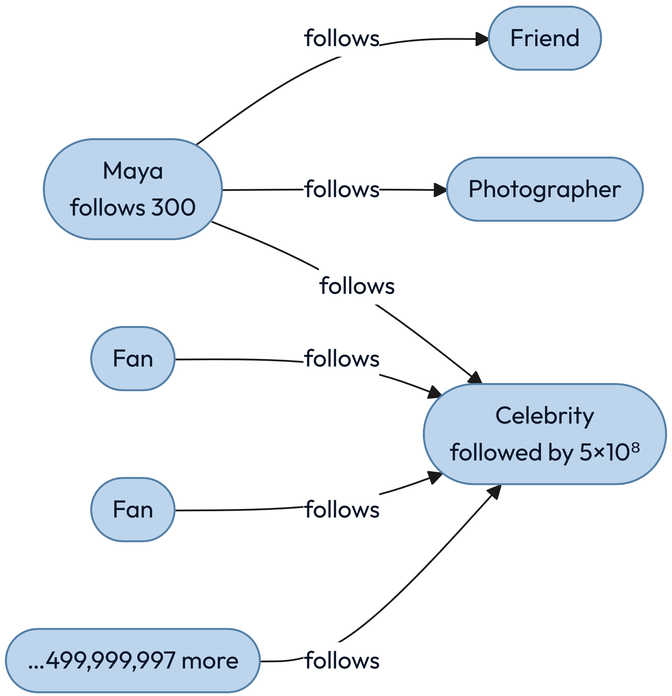
Seventy thousand a second counts opens, averaged across a whole day. A fleet serves neither opens nor averages.

An open is a session: the reader scrolls, and every screenful is another fetch. Call it six, which already puts the mean read rate at 420K a second. Traffic is not flat either — the evening peak runs a few times the daily mean. Call it two and a half.

`70K × 6 × 2.5 ≈ 1M page-fetches a second.` Size against that, and write both multipliers next to it. Leaving them out is how a fleet gets built an order of magnitude too small, and this pair is the pair most often left out.

INSTAGRAM · THE GRAPH

The follow graph is directed, and the two directions are not the same size.



✦ *The number of accounts you follow is capped at 7,500. The number who follow you is capped by nothing. Every hard problem in this design sits on that second number.*

Your following list is capped. Your followers list is capped by nothing.

Instagram caps your following list at seven and a half thousand accounts. Nobody limits how many people may follow you. An account on the far end of that second number — hundreds of millions of followers — is called a **super node**, and every hard problem here comes from one.

FOLLOWING, PER PERSON

At most 7,500. A read of that list is one partition and a sorted range.

FOLLOWERS, PER PERSON

About 150 for most people. Five hundred million at the top. A million times apart.

EVERYTHING AFTER THIS

Follows from that asymmetry. Nothing else in the design spans that range.

INSTAGRAM · THE STORAGE

The edge is stored twice, because it answers two different questions.

```
following                                  followers
  PK  (source_id)                            PK  (target_id, bucket)
  CK  (target_id)                            CK  (created_at DESC, source_id)
  val bucket                                 the follower list, newest first
  "does A follow B?" is a point read         bucket = hash(source) % B(target)
  capped at 7,500 rows                       B = clamp(followers / 10_000, 1, 512)
```

✦ *The forward edge is the source of truth, and it stores the bucket and timestamp its reverse row was written with — together those address the row exactly. That is what lets an unfollow find the reverse edge years later: `B` has grown since, so recomputing `hash(source) % B` now lands somewhere else and the edge would never be deleted. A background job rebuilds the reverse table from the forward edges and repairs drift. `B` is a high-water mark stored on the account, not the live follower count: counts fall when people unfollow, and a smaller `B` would strand every edge in the buckets above it. It ratchets up and never down, and a fan-out worker reads it fresh rather than from the cached count a profile page shows, because a stale, smaller `B` scans too few buckets and silently skips the newest followers. Early edges were written when `B` was smaller, so the low-numbered buckets hold several times the average until `B` reaches its cap.*

INSTAGRAM · WHAT EACH READ COSTS

Four of these reads are cheap. The fifth is the whole problem.

THE READ	THE PATH	WHAT IT COSTS
Does A follow B?	`following (A)` then `B`	One point read
Who does A follow?	`following (A)` range	7,500 rows, one partition
Who follows B, below the bucket line?	`followers (B, 0)` range	~10^4 rows, one partition
Follower count	`user_edge_stats`	One point read, cached
Who follows B, celebrity B?	`followers (B, 0..511)`	**5×10^8 rows, 512 partitions**

LET US GET THIS WRONG

Design it the obvious way: when you post, write into every follower's feed.

The reader then does almost nothing. One lookup, one page, no merging, no ranking across sources. It is fast, it is simple, and for an account with two hundred followers it is unambiguously correct: two hundred small writes, and every reader gets a page in a millisecond.

Hold that design in your head. Now a single account with five hundred million followers posts one photograph. Before you turn the page, predict what happens.

ONE POST, ONE CELEBRITY

25 GB

Five hundred million feed inserts at roughly fifty bytes each, from a single API call.

500 seconds to drain, alone

At a million inserts a second with nothing else arriving. Ordinary posts keep coming at 330K, so the real clear is `5×10⁸ / (1M - 330K)`, about thirteen minutes.

330K inserts a second

The platform's ordinary load: 1.2K posts times a mean fan-out near 280. Totals need the mean. The median is far lower and would halve your estimate.

25 minutes of everyone else's work

Not a wait — a quantity. One post equals what the whole platform normally produces in twenty-five minutes.

ONE SUPER NODE, FOUR FIRES

Writing 25 GB from one call is only the first thing that breaks.

The fan-out queue

25 GB of inserts from one call, in front of everybody else's posts.

The hot partition

Unbucketed, those edges are one partition key on one replica set.

The cache

25 GB of churn evicts the working set of readers who did nothing wrong.

The thundering herd

One cache key, invalidated on publish, missed by a million readers at once.

KEY INSIGHT

A skewed key produces a system where one machine is on fire and the rest are idle.

WHY THE HYBRID IS FORCED

Neither strategy is wrong. Each one is wrong somewhere.

Pushing costs one write per follower, and the follower count is unbounded. Pulling costs one read per followee, and the followee count is capped at seven and a half thousand. So push is cheap exactly where the unbounded side is small, and pull is cheap exactly where the bounded side is what you walk.

There is a second reason, and it is the better one: a celebrity's recent-posts list is written once and read five hundred million times. It is the most cacheable object in the system.

DECISION

Push below the threshold, pull above it, and merge at read.

The two strategies fail at opposite ends of the same graph, so the design uses each one where the other breaks.

ONE STRATEGY FOR EVERYONE

- Simple to explain, guaranteed to fail at one end. Push drowns on celebrities; pull costs hundreds of reads per page for everyone else.

SPLIT AT THE THRESHOLD

- Ordinary posts land in a page that is already built; the celebrities she follows are fetched on demand and merged in, so no writer ever fans out to millions. The read costs one lookup plus the celebrity pulls.

RECOMMENDATION

The threshold sits where two cost curves cross, and the graph draws it for you.

INSTAGRAM · THE THRESHOLD

Fifty thousand followers is the line, and the honest predicate is a rate.

Fifty thousand puts a fraction of a percent of accounts on the pull path, yet someone following three hundred typically has ten to thirty above it — the accounts a person picks are not a random sample.

Thirty pulls at a one-percent slow call puts the page in trouble, so **cap the pulls at twenty**, newest first. The next slide checks whether that works.

The better predicate is fan-out work per day. Fifty thousand followers posting forty times a day costs more than two hundred thousand posting weekly, and the product is what you pay for.

INSTAGRAM · CHECKING THE FIX

Capping the pulls at twenty does not meet the promise it was chosen to meet.

At thirty pulls and a one-percent slow call, *1 - 0.99^30* puts twenty-six percent of pages on a slow dependency. Cap at twenty and it is eighteen percent — still eighteen times a 99th-percentile budget. The cap sounded like the answer and fails its own arithmetic, which is exactly why you run the check on your own fixes.

THE CACHE IS WHAT BOUNDS IT

Those pulls hit a cache, so their slow rate is far below one percent — measure it rather than assuming, because at 0.2 percent twenty pulls are still four percent of pages.

THE DEADLINE IS THE BACKSTOP

Fire every pull, render at eighty milliseconds with whatever arrived, and let the stragglers land on the next fetch.

A CAP STILL COSTS SOMEONE

A reader following seven and a half thousand accounts can have hundreds above the line. For them twenty is not trimming the quietest — it is most of their pull path.

The celebrity is a protagonist of the product and the antagonist of your design.

The product exists partly for her. She is the reason a hundred million people opened the app this morning, and the reason the write path cannot be one code path. Both things are true at once, and holding both is what designing feels like.

FOR THE PRODUCT

She is the most valuable account on the platform, and she must post instantly.

FOR THE DESIGN

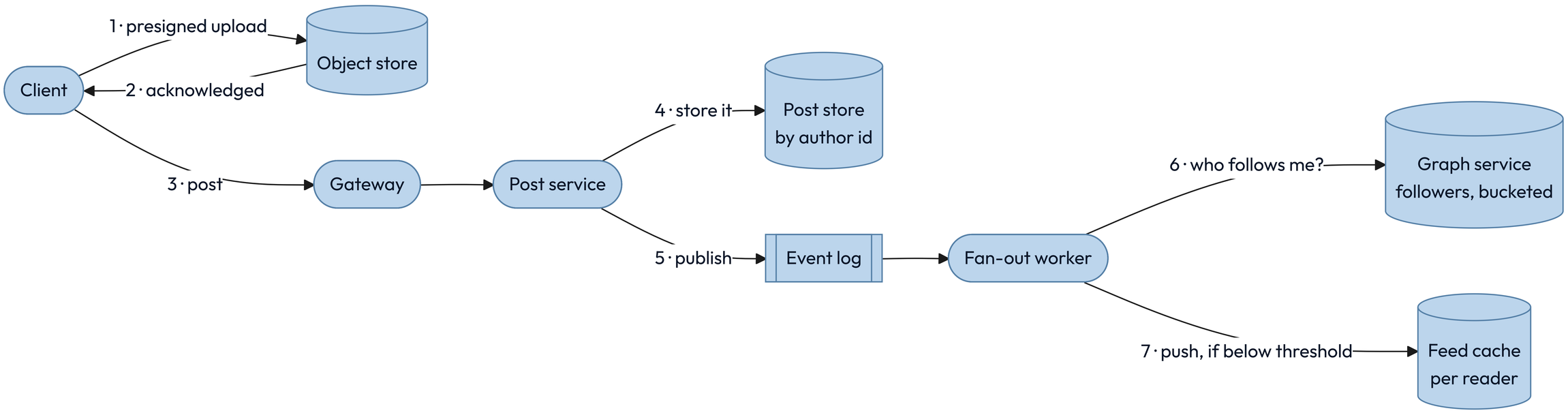
She is a five-hundred-million-edge node that breaks every uniform assumption.

WHAT YOU DO ABOUT IT

You do not resolve the tension. You build a second path and name why it exists.

INSTAGRAM · THE WRITE PATH

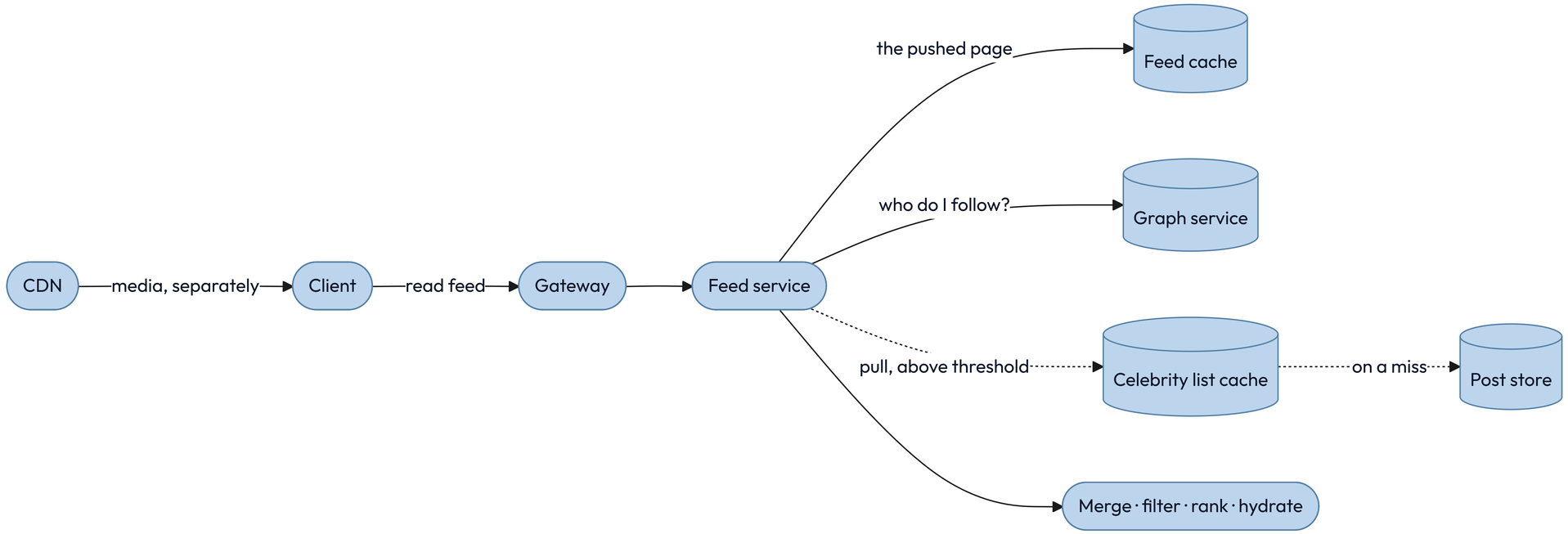
The graph service is the box everyone forgets to draw.



✦ *Step 3 comes after step 2 on purpose. Publish before the bytes are acknowledged and the post becomes readable while its media does not exist, so every reader who reaches it sees a broken tile. Transcoding may still be running — a variant that is not ready falls back to the original — but the original must be there before anybody is told the post exists.*

INSTAGRAM · THE READ PATH

The read path is one lookup, plus whatever the celebrities added.



✦ *The pull path goes through a cache, not straight to the store. That box is the whole reason the hybrid is affordable: one celebrity's recent-posts list is written once and read five hundred million times, so it is the most cacheable object in the system. Draw it, or the design you hand someone points every one of those reads at the post store.*

INSTAGRAM · ASSEMBLING A PAGE

Hydration costs more than the other three stages together.

STEP 01

Fetch and merge

Read the precomputed page, then pull recent posts from the celebrities she follows.

STEP 02

Filter

Drop blocked accounts, deleted posts and private accounts she does not follow.

STEP 03

Rank

Score the candidates. Paginate on a stable key, never on the score.

STEP 04

Hydrate

Fetch each post, its media URLs, its counts, and whether she already liked it.

THE NUMBER WE MISSED

Every page you serve is eighty lookups behind it, so the read number is eighty million a second.

A twenty-item page needs four lookups an item — the post row, the media variants, the counts, and whether this reader liked it. Eighty a page, against the million page-fetches a second the peak estimate gave you, is eighty million.

Batch them all. Do not fold the count into the feed entry — that entry is written at post time, when the count is zero. Fold the like check the other way instead: one batched read across the twenty candidates.

Batching is not only about throughput. Eighty parallel calls put fifty-five percent of pages on a slow dependency — that is `1 - 0.99^80`, against a 99th-percentile promise.

INSTAGRAM · THE LIKELY BUG

Your own post must appear instantly, or people think the upload failed.

The feed is eventually consistent, which is correct for everyone else's posts and completely wrong for your own. A person who posts and does not see it reads that as data loss, not as staleness, and posts again.

WHERE IT COMES FROM

Read-your-writes, the consistency level from Part four, applied to one reader's own posts.

WRITE YOUR OWN FEED SYNCHRONOUSLY

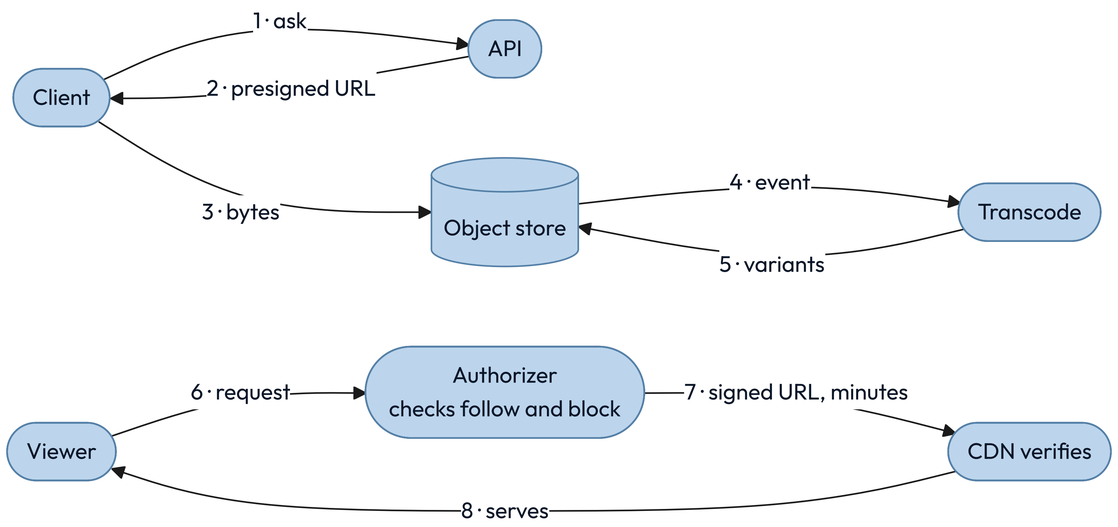
Inside the POST request, before it returns. One extra write, on one key.

AND LET THE CLIENT HELP

It inserts the post it just created optimistically, and reconciles on the next fetch.

INSTAGRAM · THE MEDIA PATH

Media never touches your API servers, and private media never touches an open URL.



✦ *Step 2 is the one place an untrusted client writes straight into your storage. So the grant names the exact key it may write, and carries four limits besides: one content type, a size ceiling, a short expiry, and a rate per account. Let the client choose the key and it writes over somebody else's media. Without them, anyone can fill your storage at your expense, and step 4 hands bytes they chose to an image decoder. All input is untrusted, including input you asked for.*

INSTAGRAM · WHAT VIDEO CHANGES

Every number in this design assumed a two-megabyte photo.

A phone uploads video in parts, because one request that dies at ninety percent would otherwise start over. So the presigned grant covers a multi-part upload rather than a single PUT.

Transcoding stops being a few still sizes and becomes a ladder of bitrates cut into short segments, so a player can step down when the signal weakens. The CDN then serves thousands of small objects per video instead of one.

The fallback goes too. A photo whose variant is not ready can be served at full size; a source video cannot, so the post stays unreadable until one rendition finishes. And the 200 TB a day becomes petabytes, which is where "durable forever" starts to cost real money.

INSTAGRAM · THE PATH NOBODY DESIGNS

A million likes on one post is the follower list again, on a different axis.

Every like is a write against the same post id — one key, one partition, one leader — arriving tens of thousands a second. That is the hot spot the data kit warned about, and the queue will sometimes deliver the same like twice.

THE ROW IS THE TRUTH

Store the pair once, keyed by reader and post. The count is an aggregate you rebuild from it, which is what survives a redelivery or an unlike.

SHARD THE HOT COUNTER ONLY

Increment a random one of a hundred sibling keys, but only when the row insert actually created a row. A blind increment double-counts every redelivery, forever.

COMMENTS ARE THE SECOND UNBOUNDED LIST

Bounded per reader, unbounded per post. Page them; never load them with the post.

INSTAGRAM · THE CHANGES NOBODY DRAWS

Every design gets drawn at rest. These four happen while it is running.

A new follow

Their earlier posts fanned out before you followed them. Backfill a few, or state the gap.

An unfollow

Pushed entries do not remove themselves. Filter at read against the current edge.

The same post twice

A crossed threshold leaves it in the pushed page and in the pull. Dedupe at the merge.

A cold celebrity key

A million readers miss it at once. One reader refills it while the rest wait on that fill.

KEY INSIGHT

A feed is a cache of a relationship. Change the relationship, leave the cache alone, and the reader is shown something untrue.

INSTAGRAM · WHERE IT BREAKS

Three of these are certainties and one is only likely. Each already has its answer.

Celebrity post

The fan-out queue floods. Answer: the pull path above the threshold.

Feed cache eviction

Cold readers cost a rebuild. Answer: rebuild lazily from posts and edges.

Redelivery

The queue delivers twice. Answer: the feed entry is keyed on reader and post.

Zone loss

A whole zone disappears. Answer: media replicated, every derived store rebuildable.

INSTAGRAM · THE INVARIANTS

Six sentences have to hold, or the design is not finished.

- ○ A post is visible to every eligible follower who reaches it — except past the pull cap, where the same quiet accounts lose every page. — rotate the cap, or say it out loud
- ✓ Media is never lost once an upload is acknowledged. — durable
- ✓ A like counts exactly once per reader and post. — idempotent
- ✓ A feed page never repeats or skips an item. — stable cursor + dedupe at merge
- ○ A blocked account is filtered at read — but a signed media URL outlives the check. — keep the TTL short
- ✓ Every derived store rebuilds from posts and edges. — rebuildable

06

PART SIX

Now run the whole method again, on something you could ship this month.

PARKING · THE ASK

A lot owner wants drivers to scan a sticker on the bay and pay to park.

No app to download. No account to create. A driver walks up, points a phone at a sticker, pays, and walks away.

Instagram was one design at one rung, and it was already enormous when we met it. This one starts at nothing and climbs, which is what your first year actually looks like.

PARKING · THE CAST

The driver has one hand free and about thirty seconds of patience.

THE PROTAGONIST

A driver who has already parked, standing in the rain with a phone in one hand. They want to pay and walk away. They will not install anything and they will not make an account.

THE ANTAGONIST

The garage is underground and the signal is poor. A driver whose page hesitates taps Pay again, so the same park can be charged twice.

PARKING · THE WORKSHEET, FILLED

Nine fields get answered before a single box goes on the board.

```
Protagonist   A driver at the bay. One hand, thirty seconds, no app.
Antagonist    Poor signal underground, and a second tap on Pay.
Purpose       The car is paid for before the driver walks away.
Boundary      In: bays, sessions, payments, enforcement. Out: the card network.
Environment   Rain, cold hands, a low battery, a sticker somebody peeled.
Constraints   physical: no signal   economic: card fees
              human: thirty seconds  legal: refunds on request
Invariants    One park charges once · an expiry never extends by accident
Bottleneck    Suspected: none yet. Confirmed once wardens start asking.
Solution type MVP. Nobody knows yet whether drivers scan the sticker.
```

PARKING · RUNG ONE, THE MVP

Rung one is one lot, a printed sticker per bay, and a page that takes a card.

What you build

A sticker on every bay carrying a link with the lot and bay in it. The provider's own form takes the card, so the number never touches your server.

What it buys

The only answer you need this month: do drivers scan the sticker, and do they finish paying.

What it charges

A warden walks the lot typing bay numbers into a phone. That holds at one lot and gives out at ten.

PARKING · RUNG ONE, THE DATABASE

One table holds all of it: no partitioning, no cache, no queue.

```
sessions
  id            uuid
  idem_key      unique. one park, one key, however many taps
  status        pending, then paid
  lot_id, bay   which sticker was scanned
  plate         typed by the driver
  started_at    set when the payment clears
  expires_at    started_at plus the minutes bought
  amount_cents  what you charged
  payment_ref   the provider's id, filled on the way out
```

Two hundred lots of forty bays, turning over four times a day, is 32,000 rows. One machine holds that for years.

PARKING · WHAT BREAKS FIRST

Traffic breaks nothing here. Two other things break anyway.

Thirty-two thousand rows a day is a rounding error, so scale is not your problem and will not be for a long time. Say that out loud, because it stops a team building for a load that never arrives.

The first break is a double charge. The page hesitates on a weak signal, the driver taps Pay again, and one park costs them twice. That reaches a human the same day.

The second is the signal itself. The card form sits on the far side of a network that keeps disappearing.

PARKING · RUNG ONE, CHARGING ONCE

Two taps on Pay have to produce one charge, and a retry must not add another.

The second tap is a different request that means the same thing. So the page mints one key for this park and keeps it across reloads, and that key decides — not your code.

INSERT FIRST, IN A PENDING STATE

The key sits in a unique column. If the insert conflicts, read that row and hand back what it already says.

CHARGE WITH THE SAME KEY

Payment providers take an idempotency key too. After a crash, a retry returns the first charge instead of making a second.

NEVER READ, THEN DECIDE, THEN WRITE

Two requests both read "not charged yet" and both charge. Let the unique index refuse the second instead.

PARKING · RUNG ONE, THE WEAK SIGNAL

The phone can drop off after the card is charged, and it often does.

The card network answers your payment provider, not the driver's phone. So the provider calls you back on a webhook, and that call is what marks the session paid. A lost signal then costs the driver a spinner rather than a park.

That endpoint is on the public internet. Check the signature your provider sends before you believe a field in it, or you have built a free parking machine. And it will arrive more than once: key it on the payment reference and let the repeat land on nothing.

YOUR TURN

Wardens arrive. Say what they ask, how often, and what answers it.

Two hundred lots have gone live. A warden walks a lot of forty bays and needs to know which cars are paid for right now.

Write down the one question they ask the system, roughly how often it is asked, and the one index that answers it. Then turn the page.

ONE ANSWER

The warden asks one small question, and it stays small.

01	**The question**	Is bay 12 in lot 40 paid at this moment. One row, never a list.
02	**How often**	Two bays a minute per warden, two hundred lots: about 400 a minute. Under ten a second.
03	**What answers it**	An index on lot, bay and expiry. The question is a point read and it stays one.

PARKING · RUNG TWO, SCALED

Two hundred lots, and the manual parts give out before the machine does.

What actually changed

Not the traffic. The manual work. One warden typing bay numbers held at one lot and gives out at two hundred.

Give the warden a list, not a keyboard

Their phone asks the indexed question once for the whole lot and shows the bays that are not paid.

Move slow work off the path a driver waits on

Owner reports go to a read replica. Receipts and nightly payouts go behind a bounded queue.

PARKING · RUNG THREE, OPTIMIZED

13%

What the card fee takes from a three-dollar park, at thirty cents plus 2.9 percent.

Servers cost pennies

Thirty-two thousand rows a day runs on the smallest machine sold. Tuning it saves nothing worth having.

The fee is the bill

Thirty-nine cents of every three dollars, and most of it is a flat charge per transaction.

So batch the charges

Hold a regular's four parks and charge once: 65 cents of fees, not a dollar fifty-five.

PARKING · WHAT WE REFUSED

A junior would build these three first, and each one is still waiting on something.

REFUSED	WHY	WHAT WOULD EARN IT
A mobile app	A driver in the rain will not install one	Regulars who park daily, once they exist
Accounts and login	A screen between the sticker and the money	Batched charges, which need a known driver
A live map of free bays	Needs a sensor in every bay	Somebody willing to pay for the sensors

PARKING · THE LADDER, CLIMBED

One product climbed three rungs, and nothing was skipped on the way up.

01	**MVP**	One lot, one table, a card form. Bought the answer to "will anyone scan this".
02	**Scaled**	Two hundred lots. A replica for reports, a queue for receipts. Nothing sharded.
03	**Optimized**	Batched charges, because the profile said the fee was the bill and the servers never were.

PARKING · WHERE THE KITS LANDED

Five kit entries carried this design, and you met all five before you saw the problem.

Relational, because nothing here outgrows one machine and the questions keep changing. An index, on the one question a warden asks. Idempotency twice over: on the driver's second tap, and on a webhook your provider will send again. A read replica, to keep reports off the path a driver waits on. A bounded queue, for the work nobody is waiting for.

Not one of those is a product name, and not one of them was a guess.

07

PART SEVEN

Every kit entry landed somewhere in the feed design, including the one we refused to use.

INSTAGRAM · STORE AND RUN

Every store-and-run entry landed somewhere specific in the feed design.

KIT ENTRY	WHERE IT LANDED	WHY THAT ONE
Wide-column	The two edge tables	One key, then a sorted list of who is on the other end
Object store	Photo bytes and variants	Large, immutable, fetched by key, kept for years
Key-value	The feed cache	The key is known exactly and nothing is asked of the value
Durable log	Post events into fan-out	Producers outpace consumers, deliberately
Stateless services	Feed and post service	Any instance serves any reader; state lives elsewhere
Bounded queue	Fan-out and transcode	Unbounded, one celebrity means an outage

INSTAGRAM · SCALE AND DEFENSE

The network, scale, reliability and security kits landed here too.

KIT ENTRY	WHERE IT LANDED	WHY THAT ONE
Reduce	One cached list per celebrity	Written once, read five hundred million times
Spread	Posts by author, edges by bucket	The bucket bounds the biggest list we still fan out
Defer	Fan-out on write, below the threshold	Moves the cost out of the reader's 200 milliseconds
CDN with versioned URLs	Every photo variant	Same bytes for many readers, and purging is eventual
Bulkhead	Fan-out workers kept off the read path	A burst of ordinary fan-out must not starve the pool serving feeds
Object-level authorization	Every hydration	The endpoint check passes; the object check stops the breach

The most useful entry here is the one we refused.

Most juniors asked to design Instagram reach for a graph database, because the words "social graph" are right there. The data kit already answered it, sixty slides before this design began.

WHAT THE CARD SAID

Walk away when you have relationships but only ever join two hops.

WHAT THE FEED ACTUALLY DOES

Walks one graph hop, me to the people I follow, then a keyed lookup that is not a traversal.

WHY THE REFUSAL MATTERS MOST

A kit that only says yes is a catalog. One that says no is a tool.

Take one box out on paper, and follow what happens to the rest.

Part three set the test. Saying where each piece landed proves nothing about whether it is needed. Deleting pieces on paper is what tells you the design is finished.

THE CELEBRITY LIST CACHE

Remove it and every reader of every celebrity post reads the store directly. It stays.

THE EVENT LOG BEFORE FAN-OUT

Remove it and the write waits for the whole fan-out before returning. It stays.

THE FOLLOWER COUNT CACHE

Remove it and every display of a profile counts rows. Keep it for display, and read `B` fresh.

THE ARTIFACT

This is the whole method, and it fits on one page.

```
Protagonist   ____________________  wants to _____________  so that _______
Antagonist    But ______________  —  and it is this big: ___________________

Purpose       ____________________________________________________________
Boundary      In: _________________________  Out: _________________________
Environment   ____________________________________________________________
Constraints   physical ______  economic ______  human ______  legal ______
Invariants    1 ___________________  2 ___________________  3 ____________
Bottleneck    ____________________________  measured at ___________________
Solution type MVP  ·  scaled  ·  optimized  ·  optimal  ·  specialized
```

THE REVIEW

Run these six on any design, starting with your own.

- ◯ Who is the protagonist, and who are we not designing for? — casting
- ◯ What is the antagonist, and what number describes it? — evidence
- ◯ Which rung of the ladder is this, and does everyone agree? — frame
- ◯ What breaks first at ten times the load? — scale
- ◯ What happens when each dependency is slow rather than down? — failure
- ◯ What can one stolen credential reach? — security

WHAT TO DO ON MONDAY

The last of these four is the one that teaches you.

STEP 01

Pick the unglamorous system

Not a famous one. The service you were debugging on Thursday, or the pipeline nobody wants to own.

STEP 02

Cast it before you draw it

Protagonist and antagonist first, one sentence each. If you cannot name them, you do not understand it yet.

STEP 03

Fill the other seven fields

Let them argue with you. A field you cannot answer is the design question you have been avoiding.

STEP 04

Bring the page to your one-on-one

Ask the person across from you which field you got wrong. That conversation is the whole point of learning this.

HOW TO THINK ABOUT SYSTEMS

Name the person, name the force, and the design follows.

Maya never designed anything on Tuesday, and she used every word in Part one before she went to bed.

Glossary

A – T

TERM	DEFINITION
ACID	All of it happens or none does, the rules hold, and a commit survives a crash.
API	The contract one system offers another: the operations and what they promise.
BASE	ACID's loose counterpart: answer under failure, let replicas converge later.
CAP	Split by a network fault, answer with stale data or refuse to answer.
CDN	Caches near readers, so bytes travel a short distance instead of an ocean.
MVP	The smallest build that puts a real answer in front of a real user.
SLA	The external promise, with money attached, always looser than the objective.
SLI	The measurement: the share of requests served inside your target.
SLO	Your internal target for that measurement, over a stated window.
TLS	The encryption any connection crossing a network you do not own should have.