

Detail Computational Results for the Paper: Neighborhood Combination Search for Single-machine Scheduling with Sequence-dependent Setup Time

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Abstract In a local search algorithm, one of its most important features is the definition of its neighborhood which is crucial to the algorithm's performance. In this paper, we present an analysis of neighborhood combination search for solving the single-machine scheduling problem with sequence-dependent setup time with the objective of minimizing total weighted tardiness (SMSWT). We first propose a new neighborhood structure named Block Swap which can be considered as an extension of the previously widely used Block Move neighborhood, and a fast incremental evaluation technique to enhance its evaluation efficiency. Then, we present two kinds of neighborhood combinations of the Block Swap and Block Move neighborhoods, called neighborhood union and token-ring search, and incorporate them into two representative metaheuristic algorithms: iterated local search (ILS_{new}) and hybrid evolutionary algorithm (HEA_{new}) to investigate their performance. Extensive experiments show the competitiveness of the token-ring search combination mechanism of the two neighborhoods. Tested on the 120 public benchmark instances, our HEA_{new} algorithm has a highly competitive performance in solution quality and computational time compared with both the exact algorithms and recent metaheuristics. We have also tested the HEA_{new} algorithm with the selected neighborhood combination search to deal with the 64 public benchmark instances of the single-machine scheduling problem with sequence-dependent setup time. Our HEA_{new} algorithm is able to match the optimal or the best known results for all the 64 instances. In particular, the computational time for reaching the best well-known results for five challenging instances is significantly reduced.

This is a supplementary file to the paper where the detail computational results are provided. Tables 1-3 reports the comparison results of the best solutions obtained by HEA_{new}, ILS_{new}, and the reference algorithms on 40 instances p1-p40, p41-p80, p81-p120 of the SMSWT problem, respectively. Table 4 reports the comparison results of the best solutions obtained by HEA_{new} and the reference algorithms on the 64 instances of the the SMST problem.

Keywords single-machine scheduling, sequence-dependent setup time, neighborhood combination search, token-ring search, hybrid evolutionary algorithm

Table 4. Comparison of the Best Solutions Obtained by HEA_{new} and the Reference Algorithms on the 64 Instances of the the SMST Problem

Instance	<i>N</i>	OPT	GVNS	HEA	ILS- RVND _{SBP}	GA-VND	IPBVND	HEA _{new}
p401	15	90	90	90	90	90	90	90
p402	15	0	0	0	0	0	0	0
p403	15	3418	3418	3418	3418	3418	3418	3418
p404	15	1067	1067	1067	1067	1067	1067	1067
p405	15	0	0	0	0	0	0	0
p406	15	0	0	0	0	0	0	0
p407	15	1861	1861	1861	1861	1861	1861	1861
p408	15	5660	5660	5660	5660	5660	5660	5660
p501	25	261	261	261	261	261	261	261
p502	25	0	0	0	0	0	0	0
p503	25	3497	3497	3497	3497	3497	3497	3497
p504	25	0	0	0	0	0	0	0
p505	25	0	0	0	0	0	0	0
p506	25	0	0	0	0	0	0	0
p507	25	7225	7225	7225	7225	7225	7225	7225
p508	25	1915	1915	1915	1915	1915	1915	1915
p601	35	12	12	12	12	12	12	12
p602	35	0	0	0	0	0	0	0
p603	35	17587	17587	17587	17587	17587	17587	17587
p604	35	19092	19092	19092	19092	19092	19092	19092
p605	35	228	228	228	228	228	228	228
p606	35	0	0	0	0	0	0	0
p607	35	12969	12969	12969	12969	12969	12969	12969
p608	35	4732	4732	4732	4732	4732	4732	4732
p701	45	97	99	97	97	98	97	97
p702	45	0	0	0	0	0	0	0
p703	45	26506	26506	26506	26506	26506	26506	26506
p704	45	15206	15206	15206	15206	15206	15206	15206
p705	45	200	202	200	200	200	200	200
p706	45	0	0	0	0	0	0	0
p707	45	23789	23789	23789	23789	23789	23789	23789
p708	45	22807	22807	22807	22807	22807	22807	22807
p551	55	183	194	183	185	183	183	183
p552	55	0	0	0	0	0	0	0
p553	55	40498	40540	40498	40498	40498	40498	40498
p554	55	14653	14653	14653	14653	14653	14653	14653
p555	55	0	0	0	0	0	0	0
p556	55	0	0	0	0	0	0	0
p557	55	35813	35830	35813	35830	35813	35813	35813
p558	55	19871	19871	19871	19871	19871	19871	19871
p651	65	247	264	247	259	247	247	247
p652	65	0	0	0	0	0	0	0
p653	65	57500	57515	57500	57508	57500	57500	57500
p654	65	34301	34301	34301	34301	34301	34301	34301
p655	65	0	4	2	4	0	1	0
p656	65	0	0	0	0	0	0	0
p657	65	54895	54895	54895	54895	54895	54895	54895
p658	65	27114	27114	27114	27114	27114	27114	27114
p751	75	225	241	229	237	229	231	237
p752	75	0	0	0	0	0	0	0
p753	75	77544	77627	77544	77559	77544	77544	77544
p754	75	35200	35219	35200	35209	35203	35200	35200
p755	75	0	0	0	0	0	0	0
p756	75	0	0	0	0	0	0	0
p757	75	59635	59716	59635	59644	59635	59635	59635
p758	75	38339	38339	38339	38339	38339	38339	38339
p851	85	360	402	381	381	369	375	381
p852	85	0	0	0	0	0	0	0
p853	85	97497	97595	97497	97497	97522	97497	97497
p854	85	79042	79271	79086	79090	79054	79086	79042
p855	85	258	280	270	274	256	266	274
p856	85	0	0	0	0	0	0	0
p857	85	87011	87075	87011	87064	87011	87011	87011
p858	85	74739	74755	74739	74739	74739	74739	74739
Sum			46	59	51	57	59	61

bold: optimal solution, *italic:* best known solution.