Basic Data Structures

链表(Linked list)

链表的定义:

```
class ListNode:
    def __init__(self, val, next=None):
        self.val = val
        self.next = next
```

虚拟头结点:

链表的一大问题就是操作当前节点必须要找前一个节点才能操作。这就造成了,头结点的尴尬,因为头结点没有前一个 节点了。

每次对应头结点的情况都要单独处理,所以使用虚拟头结点的技巧,就可以解决这个问题。

```
dummy_head = ListNode(next=head) #添加一个虚拟节点
```

链表的增删改查

经典题

203.Remove Linked List Elements

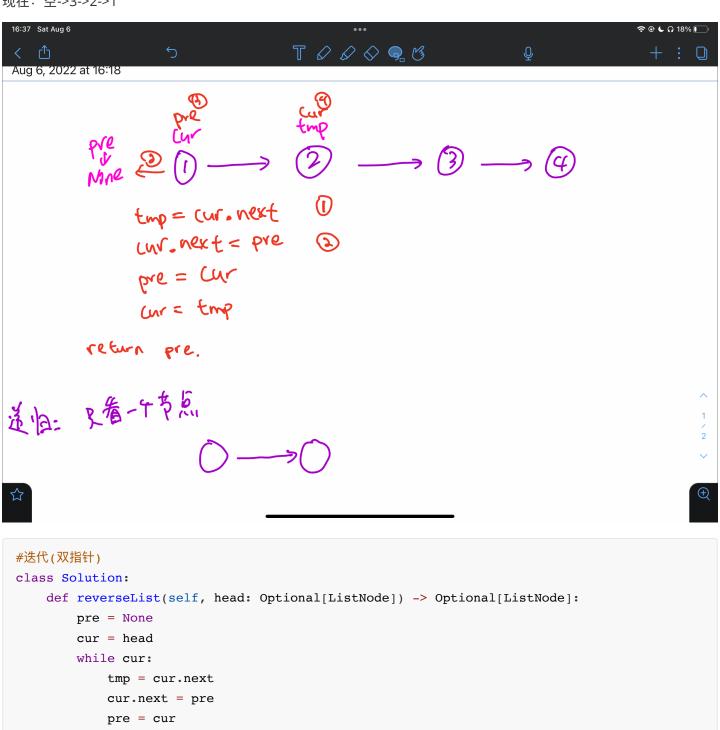
206.Reverse Linked List

cur = tmp

return pre

原: 1->2->3->空

现在: 空->3->2->1

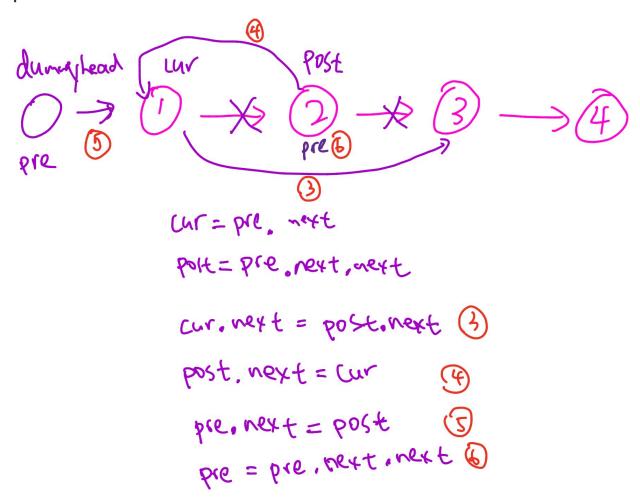


```
#递归
class Solution:
    def reverseList(self, head: ListNode) -> ListNode:
        def reverse(pre,cur):
            if not cur:
                return pre

        tmp = cur.next
        cur.next = pre

        return reverse(cur,tmp) #重点在这里
```

24. Swap Nodes in Pairs



```
class Solution:
    def swapPairs(self, head: Optional[ListNode]) -> Optional[ListNode]:
    #pre用来控制跳一格 巧妙
    dummyhead = ListNode(next=head)
    pre = dummyhead
    while pre.next and pre.next.next:
        cur = pre.next
        post = pre.next.next
        cur.next = post.next
        post.next = cur
        pre.next = post
        pre = pre.next.next
        return dummyhead.next
```

19.Remove Nth Node From End of List

快慢指针。需要注意的是n的边界问题,是大于等于还是大于。

```
class Solution:

def removeNthFromEnd(self, head: Optional[ListNode], n: int) -> Optional[ListNode]:

#快慢指针 快指针先走n步 当快指针到达末尾时 慢指针的next就是要删除的
dummyhead = ListNode(next=head)
slow = dummyhead
fast = dummyhead
while n>0:
    fast = fast.next
    n -= 1
while fast.next:
    slow = slow.next
    fast = fast.next
slow.next = slow.next.next
return dummyhead.next
```

160.Intersection of Two Linked Lists

```
class Solution:
    def getIntersectionNode(self, headA: ListNode, headB: ListNode) -> Optional[ListNode]:
    # 难度在于两个链的长度不相等
    # 遍历2次,第一次到达结尾后去到另一条链遍历,这样就能弥补长度差
    p = headA
    q = headB
    while p!=q:
        if not p:
        p = headB
```

```
else:
    p = p.next
if not q:
    q = headA
else:
    q = q.next

while p!=q:
    p = p.next if p else headB
    q = q.next if q else headA

'''
return p
```

141.Linked List Cycle

```
class Solution:

def hasCycle(self, head: Optional[ListNode]) -> bool:
    #快慢指针
    #快慢指针    #快指针一次2步,慢指针一次一步,如果能相遇则有还
    dummyhead = ListNode(next=head)
    slow= dummyhead
    fast = dummyhead
    while fast and fast.next:
        fast = fast.next.next
        slow = slow.next
        if slow == fast:
            return True
    return False
```

142.Linked List Cycle II(求入口)

```
class Solution:
    def detectCycle(self, head: Optional[ListNode]) -> Optional[ListNode]:
        slow = head
        fast = head

#找到相交的地方
while fast and fast.next:
        slow = slow.next
        fast = fast.next.next
        if slow == fast:
            break

#无环 返回None
```

```
if not fast or not fast.next:
    return None

#快指针留下 慢指针重新开始 最终他们会在环的开头相交

#为什么会这样? 看图
slow = head
while slow != fast:
    slow = slow.next
    fast = fast.next
return slow
```

练习

21. Merge Two Sorted Lists

```
class Solution:
    def mergeTwoLists(self, list1: Optional[ListNode], list2: Optional[ListNode]) ->
Optional[ListNode]:
        #不在原链上修改 很复杂
        #新开一个头结点 接在后面即可
       dummyhead = ListNode(next = None)
       p = dummyhead
       headA = list1
       headB = list2
       while headA and headB:
           if headA.val>headB.val:
               p.next = headB
               headB = headB.next
           else:
               p.next = headA
               headA = headA.next
           p = p.next
        if headA:
           p.next = headA
        if headB:
           p.next = headB
       return dummyhead.next
```

86. Partition List

```
class Solution:
   def partition(self, head: Optional[ListNode], x: int) -> Optional[ListNode]:
       # 新开一个头结点small 比x小的插入
       # 新开一个头结点big 比x大的插入
       # big插到small后面即可
       small = ListNode(next=None)
       smallhead = small
       big = ListNode(next=None)
       bighead = big
       #这样写会断链,但是无所谓,我们只关心head后面的结点,他们会以此存放到small或者big
       while head:
           if head.val < x:</pre>
               smallhead.next = head
               smallhead = smallhead.next
           else:
               bighead.next = head
               bighead = bighead.next
           head = head.next
       bighead.next = None
       smallhead.next = big.next
       return small.next
```

876.Middle of the Linked List

```
class Solution:
    def middleNode(self, head: Optional[ListNode]) -> Optional[ListNode]:
    #快慢指针 快指针两倍速走 慢指针指向的就是中间
    fast = head
    slow = head
    #fast and fast.next处理单双数问题
    while fast and fast.next:
        slow = slow.next
        fast = fast.next.next
    return slow
```

同时考察206, 21, 876. 非常好的题目

```
class Solution:
   def reorderList(self, head: Optional[ListNode]) -> None:
       Do not return anything, modify head in-place instead.
       #思路
       1.1.1
       1、快慢指针找到中间节点
       2、反转后半段链表
       3、合并前半段链表和后半段链表
       def find mid(head):
           slow = head
           fast = head
           while fast.next and fast.next.next:
               slow = slow.next
               fast = fast.next.next
           return slow
       def reverse link(head):
           pre = None
           cur = head
           while cur:
               tmp = cur.next
              cur.next = pre
               pre = cur
              cur = tmp
           return pre
       #因为链表长度只相差1或者0,所以可以这样搞。但是对于长度差距大于1的链表需要用21的方法
       def merge_list(11,12):
           while 11 and 12:
               tmp1 = 11.next
               tmp2 = 12.next
               11.next = 12
               11 = tmp1
               12.next = 11
               12 = tmp2
       "##四中四
```

```
if not head:
    return None

mid = find_mid(head)

l1 = head

l2 = mid.next

mid.next=None

l2 = reverse_link(l2)

merge_list(l1,l2)
```

328. Odd Even Linked List

```
class Solution:
   def oddEvenList(self, head: Optional[ListNode]) -> Optional[ListNode]:
       结点1作为奇数链的头 结点2作为偶数链的头
       从第3个点开始遍历,依次轮流附在奇、偶链的后面
       遍历完后,奇数链的尾连向偶链的头,偶链的尾为空,返回奇数链的头
       if not head:
          return None
       odd = head
       even = head.next
       even_head = head.next
       while even and even.next:
           odd.next = even.next
          odd = odd.next
          even.next = odd.next
          even = even.next
       odd.next = even head
       return head
```

92. Reverse Linked List II

```
class Solution:
    def reverseBetween(self, head: Optional[ListNode], left: int, right: int) ->
Optional[ListNode]:
    #难点在于边界
    # 题目给的left right是真实的下标+1
    dummyhead = ListNode(next=head)
    pre = dummyhead

for __in renge(1 left):
```

```
pre = pre.next

cur = pre.next

for _ in range(left,right):
    tmp = cur.next
    cur.next = tmp.next
    tmp.next = pre.next
    pre.next = tmp
```

237.Delete Node in a Linked List

题目限制 不能访问head

把下一node的值复制到当前node, 然后跳过下一节点

```
class Solution:
    def deleteNode(self, node):
        """
        :type node: ListNode
        :rtype: void Do not return anything, modify node in-place instead.
        """
        node.val = node.next.val
        node.next = node.next.next
```

234. Palindrome Linked List

注意:当长度是奇数的时候,mid应该是中间节点再右移一位

```
# Definition for singly-linked list.

# class ListNode:

# def __init__(self, val=0, next=None):

# self.val = val

# self.next = next

class Solution:

def isPalindrome(self, head: Optional[ListNode]) -> bool:

# 找到中间

# 将后面的翻转看是否等于前面

slow = head

fast = head

while fast.next and fast.next.next:
```

```
DIOM - DIOM . HEVE
    fast = fast.next.next
if fast != None:
    slow = slow.next
def reverse(head):
    if not head:
        return None
    pre = None
    cur = head
    while cur:
        tmp = cur.next
        cur.next = pre
        pre = cur
        cur = tmp
    return pre
left = head
right = reverse(slow)
while right:
    if right.val != left.val:
        return False
    right = right.next
    left = left.next
return True
```

148. Sort List

考点: 合并两个有序列表, 找中间节点, 归并排序

```
# Definition for singly-linked list.
# class ListNode:
# def __init__(self, val=0, next=None):
# self.val = val
# self.next = next
class Solution:
def sortList(self, head: Optional[ListNode]) -> Optional[ListNode]:
# 链表的归并排序
# 考点: 合并两个有序列表, 找中间节点, 归并排序
if not head or not head.next:
    return head
    mid = self.find_mid(head)
    right_head = mid.next

mid_next = None
```

```
left_head = head
    left = self.sortList(left_head)
    right = self.sortList(right_head)
    new_list = self.merge(left,right)
    return new list
def find_mid(self,head):
    if not head or not head.next:
        return None
    slow = head
    fast = head
    while fast.next and fast.next.next:
        slow = slow.next
        fast = fast.next.next
    return slow
def merge(self,list1,list2):
    dummyhead = ListNode(next = None)
    p = dummyhead
    headA = list1
    headB = list2
    while headA and headB:
        if headA.val>headB.val:
            p.next = headB
            headB = headB.next
        else:
            p.next = headA
            headA = headA.next
        p = p.next
    if headA:
        p.next = headA
    if headB:
        p.next = headB
    return dummyhead.next
```

配合画图

```
class Solution:
    def rotateRight(self, head: Optional[ListNode], k: int) -> Optional[ListNode]:
       #不停地把最后一个节点放在头结点前面
       if not head:
           return None
       #统计链表有多少个结点
       total = 1
       p = head
       while p.next:
           p = p.next
           total += 1
       #真正的k
       real_k = k % total
       if real_k == 0:
           return head
       # p此时在链表末尾 将收尾相连
       p.next = head
       for _ in range(total-real_k):
           p = p.next
       newhead = p.next
       p.next = None
       return newhead
```

82. Remove Duplicates from Sorted List II

```
class Solution:
    def deleteDuplicates(self, head: Optional[ListNode]) -> Optional[ListNode]:
    #把有重复项的元素全部删掉
    #画图

dummyhead = ListNode(next=head)
    pre = dummyhead
    cur = dummyhead.next
    while cur:
    while cur.next and cur.val == cur.next.val:
    cur = cur.next
```

```
if pre.next == cur:
    pre = pre.next
else:
    pre.next = cur.next

cur = cur.next
return dummyhead.next
```

```
# 递归

class Solution(object):

def deleteDuplicates(self, head):
    if not head or not head.next:
        return head

if head.val != head.next.val:
        head.next = self.deleteDuplicates(head.next)

else:
    move = head.next
    while move and head.val == move.val:
        move = move.next
    return self.deleteDuplicates(move)

return head
```

25. Reverse Nodes in k-Group

递归:对于有序的才可以

```
#翻转a, b之间的元素
    newhead = self.reverse(a,b)
    a.next = self.reverseKGroup(b,k)
    return newhead
#左闭右开的翻转
def reverse(self,a,b):
    if not a:
        return None
    pre = None
   cur = a
   tmp = a
   while cur != b:
       tmp = cur.next
       cur.next = pre
        pre = cur
        cur = tmp
    return pre
```

数组(array)

二分查找(Binary search algorithm)

704. Binary Search(模版)

```
# 普通二分查找
class Solution:
    def search(self, nums: List[int], target: int) -> int:
        left, right = 0, len(nums) - 1

    while left <= right:
        #'left + right' may cause the Integer Overflow, meaning that left+right >

2147483647

    middle = left + (right-left) // 2
    if nums[middle] < target:
        left = middle + 1
        elif nums[middle] > target:
            right = middle - 1

    elif nums[middle]==target:
```

```
return middle
        return -1
# 寻找左边界的二分查找
        def search_left(left,right):
            while left <= right:</pre>
                mid = left + (right-left)//2
                if nums[mid] < target:</pre>
                    left = mid + 1
                elif nums[mid] > target:
                    right = mid - 1
                elif nums[mid] == target: #区别
                    right = mid - 1
            if left>=len(nums) or nums[left] != target: #若越界或匹配不到 返回-1
            return left
# 寻找右边界的二分查找
        def search right(left,right):
            while left <= right:
                mid = left + (right-left)//2
                if nums[mid] < target:</pre>
                    left = mid + 1
                elif nums[mid] > target:
                    right = mid - 1
                elif nums[mid] == target:#区别
                   left = mid + 1
            if right < 0 or nums[right]!=target: #若越界或匹配不到 返回-1
                return -1
            return right
```

注意: 只有数组是单调函数才能用二分查找!

69. Sqrt(x)

```
def mySqrt(self, x: int) -> int:
    #二分查找 寻找左边界版

1, r, ans = 0, x, -1
    while 1 <= r:
        mid = (1 + r) // 2
        if mid * mid <= x:
            ans = 1
        1 = mid + 1
        else:
        r = mid - 1
    return ans
```

双指针(快慢指针 & 左右指针)

26. Remove Duplicates from Sorted Array

```
class Solution:
    def removeDuplicates(self, nums: List[int]) -> int:
    #快慢指针, 快指针比慢指针快1步
    slow = 0
    fast = 1
    while fast<len(nums):
        if nums[fast] != nums[slow]:
            nums[slow+1] = nums[fast]
            slow += 1
        fast += 1
    return slow+1 #返回数组长度
```

写法二

```
class Solution:
    def removeDuplicates(self, nums: List[int]) -> int:
        j = 2
        for i in range(2, len(nums)):
            if nums[i] != nums[j - 2]:
                nums[j] = nums[i]
                j += 1
        return j
```

27. Remove Element

由于题目要求删除数组中等于 val 的元素,因此输出数组的长度一定小于等于输入数组的长度,我们可以把输出的数组直接写在输入数组上。可以使用双指针:右指针 right 指向当前将要处理的元素,左指针 left 指向下一个将要赋值的位置。

- 如果右指针指向的元素不等于 *val*,它一定是输出数组的一个元素,我们就将右指针指向的元素复制到左指针位置,然后将左右指针同时右移;
- 如果右指针指向的元素等于 val, 它不能在输出数组里, 此时左指针不动, 右指针右移一位。

整个过程保持不变的性质是:区间 [0, left) 中的元素都不等于 val。当左右指针遍历完输入数组以后,left 的值就是输出数组的长度。

这样的算法在最坏情况下(输入数组中没有元素等于 val),左右指针各遍历了数组一次。

head

公众号: labuladong

283. Move Zeroes

```
class Solution(object):
    def moveZeroes(self, nums):
        """
    :type nums: List[int]
    :rtype: None Do not return anything, modify nums in-place instead.
```

```
if not nums:
    return 0
# 两个指针i和j
j = 0
for i in range(len(nums)):
    # 当前元素!=0, 就把其交换到左边, 等于0就不动, 相当于交换到右边
    if nums[i]!=0:
        nums[j],nums[i] = nums[j],nums[j]
        j += 1
```

75. Sort Colors

思想同283, 只不过这题要使用3指针

```
class Solution:
    def sortColors(self, nums: List[int]) -> None:
        Do not return anything, modify nums in-place instead.
        0.000
       #三指针
       if len(nums) < 2:</pre>
           return
        i = 0 #用来交换的
        j = len(nums)-1 #用来交换的
       k = 0 #用来遍历的
       #元素是0 换到左边 元素是2 换到右边
       while k<=j:
           if nums[k]==0:
               nums[k],nums[i] = nums[i],nums[k]
               i += 1
               k += 1
           elif nums[k]==1:
               k += 1
           else:
               nums[k],nums[j] = nums[j],nums[k]
               j -= 1
```

```
class Solution:
    def sortedSquares(self, nums: List[int]) -> List[int]:
       #思路一 先平方后排序 直接做
       #思路二 见代码 双指针
       i = 0
       j = len(nums)-1
       p = len(nums)-1
       res = [0] * len(nums)
       while i <= j:
           if abs(nums[i]) > abs(nums[j]):
               res[p] = nums[i] * nums[i]
               i += 1
           else:
               res[p] = nums[j] * nums[j]
               j -= 1
           p = 1
       return res
```

88. Merge Sorted Array

```
class Solution:
    def merge(self, nums1: List[int], m: int, nums2: List[int], n: int) -> None:
       Do not return anything, modify nums1 in-place instead.
        0.000
       # 三指针
       # 从前面往后会导致覆盖, 所以选择从后往前
       i = m-1
        j = n-1
       p = len(nums1)-1
       while i \ge 0 and j \ge 0:
           if nums1[i] > nums2[j]:
               nums1[p] = nums1[i]
               i -= 1
           else:
               nums1[p] = nums2[j]
               j -= 1
           p = 1
```

```
# 因为元素本身存放在nums1, 只需考虑nums2是否遍历完while j >= 0:
    nums1[p] = nums2[j]
    j -= 1
    p -= 1
```

11. Container With Most Water

贪心问题:为什么要移动较低的那一边?因为移动会底边会减小,只有当高增大可以抵消边的缩小才能面积变大。所以 移动较少的那边,尝试找到一个可以高到抵消到移动减小边的高度。

```
class Solution:

def maxArea(self, height: List[int]) -> int:

# 面积等于 (j-i) * min(i,j)

res = -1

left = 0

right = len(height)-1

while left < right:

area = (right-left) * min(height[right], height[left])

res = max(res, area)

#移动较低一边,尝试找到比当前更高的

if height[left] < height[right]:

left += 1

else:

right -= 1

return res
```

42. Trapping Rain Water

```
class Solution:
    def trap(self, height: List[int]) -> int:
        # 对于i 能装的水为 min(左边最高,右边最高)-height[i]
        # 双指针 边遍历边更新

i = 0
        j = len(height) - 1
        max_left = 0
        max_right = 0
        res = 0
        while i < j:
        max_left = max(height[i],max_left)

        max_right = max(height[j],max_right)
```

```
#res = min(max_left,max_right) - height[i]
# 更新小的 思想同11
#将上式分成两部
if max_left < max_right:
    res += max_left - height[i]
    i += 1
else:
    res += max_right - height[j]
    j -= 1
return res</pre>
```

滑动窗口(Sliding window)

常用于字符串问题

所谓滑动窗口,就是不断的调节子序列的起始位置和终止位置,从而得出我们要想的结果。

滑动窗口算法的思路是这样:

1、我们在字符串 s 中使用双指针中的左右指针技巧,初始化 left = right = 0 ,把索引**左闭右开**区间 [left, right) 称为一个「窗口」。

PS: 理论上你可以设计两端都开或者两端都闭的区间,但设计为左闭右开区间是最方便处理的。因为这样初始化 left = right = 0 时区间 [0, 0) 中没有元素,但只要让 right 向右移动(扩大)一位,区间 [0, 1) 就 包含一个元素 0 了。如果你设置为两端都开的区间,那么让 right 向右移动一位后开区间 (0, 1) 仍然没有元素;如果你设置为两端都闭的区间,那么初始区间 [0, 0] 就包含了一个元素。这两种情况都会给边界处理带来不必要的麻烦。

- 2、我们先不断地增加 right 指针扩大窗口 [left, right], 直到窗口中的字符串符合要求(包含了 T 中的所有字符)。
- 3、此时,我们停止增加 right ,转而不断增加 left 指针缩小窗口 [left, right) ,直到窗口中的字符串不再符合要求(不包含 T 中的所有字符了)。同时,每次增加 left ,我们都要更新一轮结果。
- 4、重复第2和第3步, 直到 right 到达字符串 s 的尽头。

这个思路其实也不难,**第2步相当于在寻找一个「可行解」,然后第3步在优化这个「可行解」,最终找到最优解**,也就是最短的覆盖子串。左右指针轮流前进,窗口大小增增减减,窗口不断向右滑动,这就是「滑动窗口」这个名字的来历。

所需要的变量:

哈希表 needs 和 window 相当于计数器,分别记录 I (要匹配的字符)中字符出现次数和「窗口」中的相应字符的出

现次数。

valid: 合法字符数量。当window['a']达到某个数值时, valid+=1

left, right: 左右指针, 用于全局遍历

start, length: 用于记录符合条件的区间【start: start+length】

现在开始套模板、只需要思考以下几个问题:

- 1、什么时候应该移动 right 扩大窗口?窗口加入字符时,应该更新哪些数据?
- 2、什么时候窗口应该暂停扩大, 开始移动 left 缩小窗口? 从窗口移出字符时, 应该更新哪些数据?
- 3、我们要的结果应该在扩大窗口时还是缩小窗口时进行更新?

如果一个字符进入窗口,应该增加 window 计数器;如果一个字符将移出窗口的时候,应该减少 window 计数器;当 valid 满足 need 时应该收缩窗口;应该在收缩窗口的时候更新最终结果。

76. Minimum Window Substring(模版)

```
class Solution:
   def minWindow(self, s: str, t: str) -> str:
       if len(s) < len(t):
          return ''
       need = defaultdict(int)
       for char in t:
          need[char] += 1
       window = defaultdict(int) #统计窗口内每个字符串的个数
       left,right = 0,0
       valid = 0 #合法字符数量 如need[a] = 3,当window[a] = 3时, valid +=1
       start, length = 0, inf
       while right < len(s):
          c = s[right] #c是移入窗口的字符
          right += 1
                       #扩大窗口
          #如果进去的是需要的字符,就进行更新(哈希表window和合法字符数量valid的更新)
          if need.get(c):
              window[c] += 1
              if window[c] == need[c]:
                  valid += 1
          # 判断左边窗口是否要收缩
          while valid == len(need): #收缩窗口的条件(全部字符都满足要求了)
              #更新start和length
              if right - left < length:</pre>
                  start = left
                  # 由于每次扩大窗口都对right+1了 所以此时right是目标区间的后一个元素下标
```

```
# 对应切片互闭石井 此处个用+1
length = right - left

#缩小窗口
d = s[left]
left += 1
#如果出去的是需要的字符,就进行更新(哈希表window和合法字符数量valid的更新)
if need.get(d):
    if window[d] == need[d]:
        valid -=1
    window[d] -= 1

return '' if length == inf else s[start:start+length]
```

3. Longest Substring Without Repeating Characters

```
class Solution:
    def lengthOfLongestSubstring(self, s: str) -> int:
       #哈希表记录滑动窗口内子串
       window = defaultdict(int)
       start, length = 0, 0
       left,right = 0, 0
       res = 0
       while right < len(s):</pre>
           #先扩大窗口再移入
           c = s[right]
           right += 1
           window[c] += 1
           #收缩
           while window[c] > 1:
               d = s[left]
               left += 1
               window[d] -= 1
           res = max(res,right-left)
        return res
```

```
class Solution:
   def findAnagrams(self, s: str, p: str) -> List[int]:
       need = defaultdict(int)
       window = defaultdict(int)
       res = []
       start = 0
       left,right = 0, 0
       valid = 0
       for ss in p:
           need[ss] += 1
       while right < len(s):
           c = s[right]
           right += 1
           if need.get(c):
               window[c] += 1
               if window[c] == need[c]: #字符在窗口出现的次数和需要的次数一样 valid才加
                   valid += 1
           while right - left >= len(p): #左边收缩条件
               if valid == len(need):
                   res.append(left)
               d = s[left]
               left += 1
               #注意和右边扩展有区别
               #先-valid再-window
               if need.get(d):
                   if window[d] == need[d]:
                       valid -= 1
                   window[d] = 1
       return res
```

\567. Permutation in String

基本同438

```
class Solution:
    def checkInclusion(self, s1: str, s2: str) -> bool:
        left, right = 0, 0
        need = defaultdict(int)
        for s in s1:
            need[s] += 1
        window = defaultdict(int)

        valid = 0
```

```
res = False
while right < len(s2):
    c = s2[right]
    right += 1
    if need.get(c):
        window[c] += 1
        if window[c] == need[c]:
            valid += 1
    while right - left >= len(s1):
        if valid == len(need):
            return True
        d = s2[left]
        left += 1
        if need.get(d):
            if window[d] == need[d]:
                valid -= 1
            window[d] = 1
return False
```

\239. Sliding Window Maximum

最直接的做法, 但是超时

```
class Solution:
    def maxSlidingWindow(self, nums: List[int], k: int) -> List[int]:
        left, right = 0, k-1
        res = []
        while right < len(nums):
            max_value = max(nums[left:right+1])
        res.append(max_value)
        right += 1
        left += 1
        return res</pre>
```

使用单调队列:所有队列里的元素都是按递增(递减)的顺序队列,这个队列的头是最小(最大)的元素。

队列储存的是数组的下标!!!

```
class Solution:
    def maxSlidingWindow(self, nums: List[int], k: int) -> List[int]:
    # 初始化队列和结果,队列存储数组的下标
    queue = []
```

```
for i in range(len(nums)):
    # 如果当前队列最左侧存储的下标等于 i-k 的值, 代表目前队列已满。
    # 但是新元素需要进来,所以列表最左侧的下标出队列
    if queue and queue[0] == i - k:
        queue.pop(0)

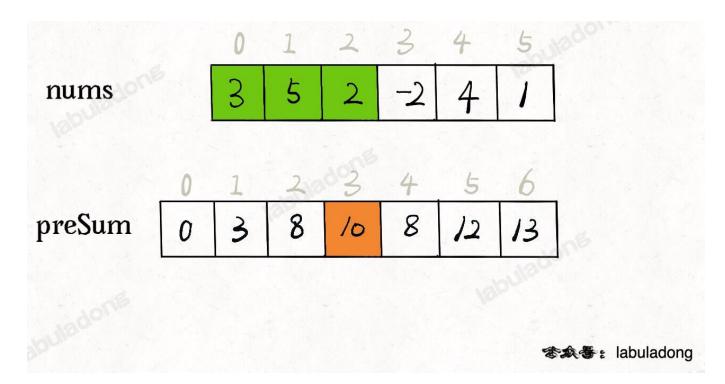
# 对于新进入的元素,如果队列前面的数比它小,那么前面的都出队列
    # 目的是维护第一个数永远是最大的(单调队列的定义)
while queue and nums[queue[-1]] < nums[i]:
        queue.pop()
    # 新元素入队列,进来的是下标
    queue.append(i)

# 当前的大值加入到结果数组中
    if i >= k-1:
        res.append(nums[queue[0]])

return res
```

前缀和数组(Prefix sum)

前缀和技巧适用于快速、频繁地计算一个索引区间内的元素之和。



303. Range Sum Query - Immutable

```
class NumArray:

def __init__(self, nums: List[int]):
    # self.preSum里先放一个0, 即列表里总共放n+1个元素
    self.preSum = [0]
    for num in nums:
        self.preSum.append(self.preSum[-1] + num)

def sumRange(self, left: int, right: int) -> int:
    # 查询闭区间的累加和
    # 在self.preSum里, index=right + 1时, 是加了index=right这个元素的值
    # 因此闭区间[left, right]的累加和等于下面
    return self.preSum[right + 1] - self.preSum[left]
```

排序

快速排序

nums: 5 2 9 1 3 6

选择基准值 step 0:最左端元素选为pivot 2 pivot=5 nums: pivot left right step 1: right从右往左滑动, 1 3 6 pivot=5 找到一个小于pivot的元素, nums: 并将其放置于`空`出来的left处 left right step 2: left从左往右滑动, 3 9 6 2 1 pivot=5 找到一个大于pivot的元素, nums: 并将其放置于`空`出来的right处 left right 分割 step n: 重复step 1和step 2, 9 6 nums: 3 2 pivot=5 partition 直至left=right left --right : final step: left=right, 6 2 nums: pivot=5 将pivot置于left处 left -right 5 6 nums: 临时结果: 对pivot两侧的子数组 递归排序 2 6 递归排序, 子数组 分别重复上述步骤 子数组1 子数组2 2 3 5 6 9 最终结果: nums: #quick_sort class Solution:

def sortArray(self nums. List(intl) => List(intl.

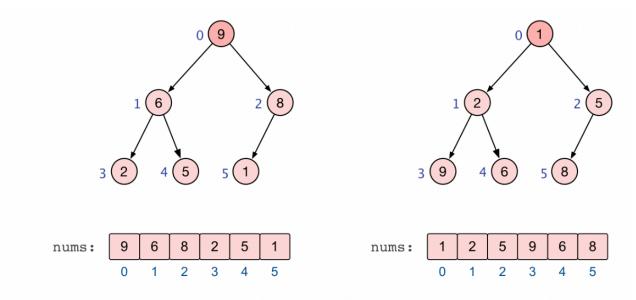
```
ATTEM (DOTE) HAMD. DIDCLING), . DIDCLING.
       def partition(arr, low, high):
           pivot = arr[low] #固定pivot
           left = low
           right = high
           while left < right:
               #从右边往左边找到第一个小于pivot的元素 放在left位置
               while left<right and arr[right] >= pivot: #注意这里是大于等于
                   right -= 1
               arr[left] = arr[right]
               #从左边往右边找到第一个大于pivot的元素 放在right位置
               while left<right and arr[left] <= pivot: #注意这里是小于等于
                   left += 1
               arr[right] = arr[left]
           #最终将pivot元素放在中间 即left==right处
           arr[left] = pivot
           return left
       def random select pivot(arr,low,high):
           pivot_idx = random.randint(low,high)
           arr[low], arr[pivot_idx] = arr[pivot_idx], arr[low]
           return partition(arr,low,high)
       def quicksort(arr,low,high):
           if(low >= high): #注意这里是大于等于
               return
           #mid = partition(arr,low,high)
           mid = random_select_pivot(arr,low,high)
           quicksort(arr,low,mid-1)
           quicksort(arr,mid+1,high)
       quicksort(nums, 0, len(nums)-1)
       return nums
# quick selection 快速选择(用于topK split问题)
#quick select, find a pivot, move all nums that is smaller than it to it's left, all nums
> pivot to it's right. check how many on the left. if count == k, return, if count > k, it
must be on the left, otherwise on the right.
```

```
# merge sort
#思想比较简单
class Solution:
   def sortArray(self, nums: List[int]) -> List[int]:
        def merge sort(arr,low,high):
            if low >= high:
               return
           mid = low + (high-low)//2
            merge_sort(arr,low,mid) #右边到mid
            merge_sort(arr,mid+1,high)
            left, right = low, mid+1
            tmp = [] #记录结果
            while left <= mid and right <= high:
                if arr[left] <= arr[right]:</pre>
                    tmp.append(arr[left])
                   left += 1
                else:
                    tmp.append(arr[right])
                   right += 1
            #比较完了 还没放入tmp的直接放入
            while left <= mid:
                tmp.append(arr[left])
                left += 1
            while right <= high:
                tmp.append(arr[right])
                right += 1
            arr[low:high+1] = tmp #注意是浅复制
        merge_sort(nums,0,len(nums)-1)
        return nums
```

堆排序

● 基本定义:大根堆和小根堆

大根堆 小根堆



大根堆/大顶堆:每个节点的值均大于等于其左右孩子节点的值;小根堆/小顶堆:每个节点的值均小于等于其左右孩子节点的值。

● 如何进行堆排序?

\215. Kth Largest Element in an Array

Hash Table

- 用数组作为哈希表
- 用dict作为哈希表{key:value}

python中: collections.defaultdict 和 collections.Counter的使用

https://docs.python.org/3/library/collections.html#counter-objects

242. Valid Anagram

用两个dict,分别插入字符串,比较两个dict是否相等即可。

```
def isAnagram(self, s: str, t: str) -> bool:
    a = defaultdict(int)
    b = defaultdict(int)
    for char in s:
        a[char] += 1
    for char in t:
        b[char] += 1
    if a == b:
        return True
    return False
```

1002. Find Common Characters

难点:

1. Counter 的定义和用法

```
from collections import Counter
s = 'evergfsafcage'
h = Counter(s)
print(h)
#Counter({'e': 3, 'g': 2, 'f': 2, 'a': 2, 'v': 1, 'r': 1, 's': 1, 'c': 1})
#elements()
print(list(h.elements()))
#['e', 'e', 'e', 'v', 'r', 'g', 'g', 'f', 'f', 's', 'a', 'a', 'c']
```

2. python中集合取交集的取法: &=

```
class Solution:
    def commonChars(self, words: List[str]) -> List[str]:
        res = Counter(words[0])
        for i in words:
            res &= Counter(i)
        return list(res.elements())
```

349. Intersection of Two Arrays

```
class Solution:
   def intersection(self, nums1: List[int], nums2: List[int]) -> List[int]:
     return set(nums1) & set(nums2)
```

49. Group Anagrams

哈希表的使用

将每个字母出现的次数使用字符串表示,作为哈希表的键

```
class Solution:
    def groupAnagrams(self, strs: List[str]) -> List[List[str]]:
        mp = defaultdict(list)
        for s in strs: #遍历字符串
            count = [0] * 26
            for i in s: #遍历每个字符
                  count[ord(i)-ord('a')] += 1
                 mp[tuple(count)].append(s) #需要将 list 转换成 tuple 才能进行哈希
        return list(mp.values())
```

73. Set Matrix Zeroes

集合的使用, 比较简单粗暴

注意第二次循环使用的技巧 in

```
class Solution:
    def setZeroes(self, matrix: List[List[int]]) -> None:
        """

    Do not return anything, modify matrix in-place instead.
        """
```

```
#先找到所有o的下标 在删
zero_row = set()
zero_col = set()
m = len(matrix)
n = len(matrix[0])
for i in range(m):
    for j in range(n):
        if matrix[i][j]==0:
            zero_row.add(i)
            zero_col.add(j)

for i in range(m):
    for j in range(n):
        if i in zero_row or j in zero_col:
            matrix[i][j] = 0
```

N sum 问题

1. Two Sum

难点

- 1. 使用 for idx, val in enumerate(num):同时遍历编号和数值
- 2. 一边遍历一边加入字典

```
class Solution:
    def twoSum(self, nums: List[int], target: int) -> List[int]:
    #遍历列表同时查字典

#字典{数值: 编号} 因为要输出的是编号
    res = dict()
    for idx,val in enumerate(nums):
        tar = target - val
        if tar in res:
            return [res[tar],idx]
        else:
        res[val] = idx
```

本题还可以用双指针做,但是排序会造成序号的改变, 所以我们需要记住排序中记住序号

记住当前序号并排序

```
nums = sorted(enumerate(nums), key= lambda x: x[1])
# res: [(4, 1), (3, 2), (1, 3), (0, 4), (2, 6)]
```

```
class Solution:
    def twoSum(self, nums: List[int], target: int) -> List[int]:
        return self.solve(nums, target)

def solve(self, nums, target):
    nums = sorted(enumerate(nums), key= lambda x: x[1])
    l, r = 0, len(nums)-1

while 1 < r:
    twosum = nums[1][1] + nums[r][1]
    if twosum == target:
        return [nums[1][0], nums[r][0]]
    elif twosum < target:
        1 += 1
    elif twosum > target:
        r -= 1
```

15. 3Sum

关键在于如何去重:无论是对于nums[i],left,right,要是当前的元素和前一个元素相等,直接跳过!

```
class Solution:
   def threeSum(self, nums: List[int]) -> List[List[int]]:
       #双指针
       res = []
       nums.sort()
       for i in range(len(nums)):
           #left 和 right需要写在循环里面, 因为是不断变化的
           left = i+1
           right = len(nums)-1
           if nums[i]>0: #left和right都在nums[i]右边
           if i >= 1 and nums[i] == nums[i - 1]: # 当前元素和前一个元素相等,会出现重复,跳过
              continue
           while left<right:
               s = nums[i]+nums[left]+nums[right]
              if s < 0:
                  left += 1
               elif s > 0:
                  right -= 1
               else:
                  res.append([nums[i],nums[left],nums[right]])
                  #去重 无论是left和right,要是当前元素和前一个元素相等,都跳过
```

```
while left != right and nums[left] == nums[left + 1]: left += 1
    while left != right and nums[right] == nums[right - 1]: right -= 1
    left += 1
    right -= 1
return res
```

16. 3Sum Closest

```
class Solution:
    def threeSumClosest(self, nums: List[int], target: int) -> int:
        #不用去重
        res = inf
        nums.sort()
        for i in range(len(nums)):
            left = i+1
            right = len(nums)-1
            while left < right :</pre>
                 cur_sum = nums[i] + nums[left] + nums[right]
                 if abs(target - cur_sum) < abs(target - res):</pre>
                     res = cur_sum
                 if cur_sum < target:</pre>
                     left += 1
                 elif cur_sum > target:
                     right -= 1
                 else:
                     return res
        return res
```

18. 4Sum

```
#去重
        if j>i+1 and nums[j] == nums[j-1]:
            continue
        left = j+1
        right = n-1
        while left < right:
            s = nums[i] + nums[j] + nums[left] + nums[right]
            if s < target:</pre>
                left += 1
            elif s > target:
                 right -= 1
            else:
                 res.append([nums[i],nums[j],nums[left],nums[right]])
                 #去重
                 while(left < right) and nums[left]==nums[left+1]:</pre>
                 while(left < right) and nums[right] == nums[right-1]:</pre>
                     right -= 1
                 left += 1
                 right -= 1
return res
```

字符串(string)

python 常用字符串方法(纯字符串题目)

```
对于字符串s
#1. 翻转
s = s[::-1]
#2. 切片
#3. 先用list变成列表,再用''.join(s)恢复成字符串(如151题)
res = list(i for i in s.split(' ') if len(i)>0)
#4. 替换(用xx替换空格或类似的)
'xx'.join(s.split(' '))
#5. 先局部翻转再全部翻转 & 先整体翻转再局部翻转
```

344. Reverse String

```
class Solution:
    def reverseString(self, s: List[str]) -> None:
        """

        Do not return anything, modify s in-place instead.
        """

        #双指针
        left = 0
        right = len(s)-1
        while left < right:
            s[left], s[right] = s[right],s[left]
        left += 1
        right -= 1
```

最简单的方法如下:

需要注意的需要用s[:]浅拷贝。

浅拷贝是在另一块地址中创建一个新的变量或容器,但是容器内的元素的地址均是源对象的元素的地址的拷贝。也就是 说新的容器中指向了旧的元素(新瓶装旧酒)。

深拷贝是在另一块地址中创建一个新的变量或容器,同时容器内的元素的地址也是新开辟的,仅仅是值相同而已,是完全的副本。也就是说(新瓶装新酒)。

```
class Solution:
    def reverseString(self, s: List[str]) -> None:
        """
        Do not return anything, modify s in-place instead.
        """
        s[:] = s[::-1]
```

541. Reverse String II

使用上一题的函数进行翻转

要注意的是把字符串弄成列表操作,然后再把它变回字符串。

注意range的第三个参数是步长

```
class Solution:
   def reverseStr(self, s: str, k: int) -> str:
        def reverse_substring(text):
```

```
left, right = 0, len(text) - 1
while left < right:
    text[left], text[right] = text[right], text[left]
    left += 1
    right -= 1
    return text

res = list(s)

for cur in range(0, len(s), 2 * k):
    # 如果k是2, 则cur分别是0,2,4,6, 。。。
    res[cur: cur + k] = reverse_substring(res[cur: cur + k])

return ''.join(res)</pre>
```

151. Reverse Words in a String

```
class Solution:
    def reverseWords(self, s: str) -> str:
        #转成列表
        #翻转列表
        #恢复成字符串
        res = list(i for i in s.split(' ') if len(i)>0)
        res = res[::-1]
        return ' '.join(res)
```

栈和队列(stack and queue)

20. Valid Parentheses

非常经典的使用栈的问题!

思路:遍历字符串,当遇到左括号时,往栈中插入相应的右括号;如果遇到右括号,比较此时栈中最顶元素是否为同一个右括号,不是直接返回false,此外,当字符串还没被遍历完,栈已经空了,说明连续出现了两个一样的右括号,返回false。当字符串被遍历完时,如果栈也为空,则返回true,否则返回false。

```
elif item == '[':
          stack.append(']')
elif item == '{':
          stack.append('}')
elif not stack or item != stack[-1]:
          return False
else:
          stack.pop()
return True if not stack else False
```

1047. Remove All Adjacent Duplicates In String

```
class Solution:
    def removeDuplicates(self, s: str) -> str:
    #对对碰
    stack = []
    for item in s:
        if stack and item == stack[-1]:
            stack.pop()
        else:
            stack.append(item)

return ''.join(stack)
```

150. Evaluate Reverse Polish Notation

Matrix模拟

48. Rotate Image

```
class Solution:
    def rotate(self, matrix: List[List[int]]) -> None:
        """

    Do not return anything, modify matrix in-place instead.
        """

    # 有点智力题的意思

    # 先上下翻转 再沿着45度对角线翻转
    row = len(matrix)

    col = len(matrix[0])
```

```
# 上下翻转
i = 0
j = row-1
while i<j:
    matrix[i][:], matrix[j][:] = matrix[j][:], matrix[i][:]
    i += 1
    j -= 1

# 对角线
for i in range(row):
    for j in range(col):
        if i==j:
             break
        matrix[i][j], matrix[j][i] = matrix[j][i], matrix[i][j]

return matrix
```

54. Spiral Matrix

```
class Solution:
    def spiralOrder(self, matrix: List[List[int]]) -> List[int]:
        #设计上下左右边界 当初碰到边界时
        row = len(matrix)
        col = len(matrix[0])
        left bound = 0
        right_bound = col-1
        up\_bound = 0
       down bound = row-1
        nums = []
        while len(nums) < row*col:</pre>
            # 在最上边从左到右
            if up_bound <= down_bound:</pre>
                for j in range(left_bound,right_bound+1):
                    nums.append(matrix[up_bound][j])
                up_bound += 1
            # 在最右边从上到下
            if right_bound >= left_bound:
                for j in range(up_bound,down_bound+1):
                    nums.append(matrix[j][right_bound])
                right bound -= 1
```

```
# 在最下的从石到左
if down_bound >= up_bound:
    for j in range(right_bound,left_bound-1,-1):
        nums.append(matrix[down_bound][j])
    down_bound -= 1

#在最左边从下到上
if left_bound <= right_bound:
    for j in range(down_bound,up_bound-1,-1):
        nums.append(matrix[j][left_bound])
    left_bound += 1

return nums
```

59. Spiral Matrix II

54的代码稍微修改即可

```
class Solution:
    def generateMatrix(self, n: int) -> List[List[int]]:
        #和54题同理 这次是填入数字
       matrix = [[0 for _ in range(n)] for _ in range(n)]
        left bound = 0
        right\_bound = n-1
        up\_bound = 0
        down bound = n-1
        nums = 1
        while nums <= n*n:
            # 在最上边从左到右
            if up_bound <= down_bound:</pre>
                for j in range(left_bound,right_bound+1):
                    matrix[up_bound][j] = nums
                    nums += 1
                up bound += 1
            # 在最右边从上到下
            if right bound >= left bound:
                for j in range(up_bound,down_bound+1):
                    matrix[j][right_bound] = nums
                    nums += 1
```

```
right_bound -= 1

# 在最下面从右到左
if down_bound >= up_bound:
    for j in range(right_bound,left_bound-1,-1):
        matrix[down_bound][j] = nums
        nums += 1
    down_bound -= 1

#在最左边从下到上
if left_bound <= right_bound:
    for j in range(down_bound,up_bound-1,-1):
        matrix[j][left_bound] = nums
        nums += 1
    left_bound += 1
```

73. Set Matrix Zeroes

```
class Solution:
    def setZeroes(self, matrix: List[List[int]]) -> None:
        Do not return anything, modify matrix in-place instead.
       #先找到所有0的下标 在删
       zero_row = set()
       zero col = set()
       m = len(matrix)
        n = len(matrix[0])
        for i in range(m):
            for j in range(n):
                if matrix[i][j]==0:
                    zero_row.add(i)
                    zero_col.add(j)
        for i in range(m):
            for j in range(n):
                if i in zero_row or j in zero_col:
                    matrix[i][j] = 0
```

```
class Solution:
    def searchMatrix(self, matrix: List[List[int]], target: int) -> bool:
    # 有序 想到二分查找 可以转成1维搜
    # 也可以从右上角开始搜
    row = 0
    col = len(matrix[0])-1
    while row <= len(matrix)-1 and col >= 0:
        if matrix[row][col] < target:
            row += 1
        elif matrix[row][col] > target:
            col -= 1
        elif matrix[row][col] == target:
            return True
    return False
```

240. Search a 2D Matrix II

```
class Solution:

def searchMatrix(self, matrix: List[List[int]], target: int) -> bool:
# 每一行递增
# 每一列递增

# 还是从右上角开始遍历 大了往左移动 小了往右移动
row = 0
col = len(matrix[0])-1
while row < len(matrix) and col >= 0:
    if matrix[row][col] > target:
        col -= 1
    elif matrix[row][col] < target:
        row += 1
    elif matrix[row][col] == target:
        return True
return False
```

378. Kth Smallest Element in a Sorted Matrix

思想: 和74,240差不多, 也是从左下或者右上开始走

走法演示如下,依然取 mid = 8:

_					A	
	1	3	5	7	9	11
	2	4	6	8 —	10	12
	3	5	7	9	11	13
	4	6	8 -	10	12	14
	5	7	9	11	13	15
	6 -	8 -	10	12	14	16

可以这样描述走法:

- 初始位置在 matrix[n-1][0] (即左下角);
- 设当前位置为 matrix[i][j]。若 $matrix[i][j] \le mid$,则将当前所在列的不大于 mid 的数的数量 (即 i+1)累加到答案中,并向右移动,否则向上移动;
- 不断移动直到走出格子为止。

我们发现这样的走法时间复杂度为 O(n),即我们可以线性计算对于任意一个 mid,矩阵中有多少数不大于它。这满足了二分查找的性质。

不妨假设答案为 x, 那么可以知道 $l \le x \le r$, 这样就确定了二分查找的上下界。

每次对于「猜测」的答案 mid, 计算矩阵中有多少数不大于 mid:

- 如果数量不少于 k, 那么说明最终答案 x 不大于 mid;
- 如果数量少于 k,那么说明最终答案 x 大于 mid。

这样我们就可以计算出最终的结果 x 了。

class Solution:
 def kthSmallest(self, matrix: List[List[int]], k: int) -> int:
 n = len(matrix)

def check(mid):
 # nums >= k 目标值<=k
nums <= k 日标值>k

```
# Hums ~ K , 口小川上/K
    i, j = n - 1, 0
    num = 0
    while i \ge 0 and j < n:
        if matrix[i][j] <= mid:</pre>
            num += i + 1
            j += 1
        else:
            i -= 1
    return num >= k
left, right = matrix[0][0], matrix[-1][-1]
while left <= right:</pre>
   mid = (left + right) // 2
    if check(mid):
       right = mid - 1
    else:
        left = mid + 1
return left
```