#### **Application 1:**

This application takes a number and puts it through a random function to calculate a floating-point value. The magic function, magicFunc(n), is equal to magicFunc(n - 1) - magicFunc(n - 2)/4. The initial value is set in the data section on the code. This first program will use a smaller cache size. The initial value is set to 15.

# Code: # ECE365 Project Phase 3 Application 1 # Alex Santana .data num: .word 15 # index .text main: \$a0, num lw magicFunc jal li \$v0, 2 # print float mtc1 \$zero, \$f6 # clear cvt.s.w \$f6, \$f6 # convert add.s \$f12, \$f6, \$f30 # print number return syscall #END PROGRAM---li \$v0, 10 syscall

```
magicFunc:
```

beq \$a0, \$zero, mfReturn0

slti \$t0, \$a0, 2

bne \$t0, \$zero, mfReturn1

j mfCalc

### mfReturn0:

add \$t0, \$zero, \$zero

j mfReturn

#### mfReturn1:

addi \$t0, \$zero, 1

j mfReturn

#### mfReturn:

mtc1 \$t0, \$f30

cvt.s.w \$f30, \$f30

jr \$ra

## mfCalc:

mtc1 \$zero, \$f6 # clear float

cvt.s.w \$f6, \$f6 # convert

addi \$sp, \$sp, -72

sw \$ra, 0(\$sp)

sw \$a0, 4(\$sp)

swc1 \$f0, 8(\$sp)

swc1 \$f2, 40(\$sp)

addi \$a0, \$a0, -1

jal magicFunc

add.s \$f0, \$f6, \$f30

lw \$a0, 4(\$sp)

addi \$a0, \$a0, -2

jal magicFunc

add.s \$f2, \$f6, \$f30

addi \$t0, \$zero, 4

mtc1 \$t0, \$f4

cvt.s.w \$f4, \$f4 # convert

div.s \$f2, \$f2, \$f4

sub.s \$f0, \$f0, \$f2

add.s \$f30, \$f6, \$f0

lw \$ra, 0(\$sp)

lw \$a0, 4(\$sp)

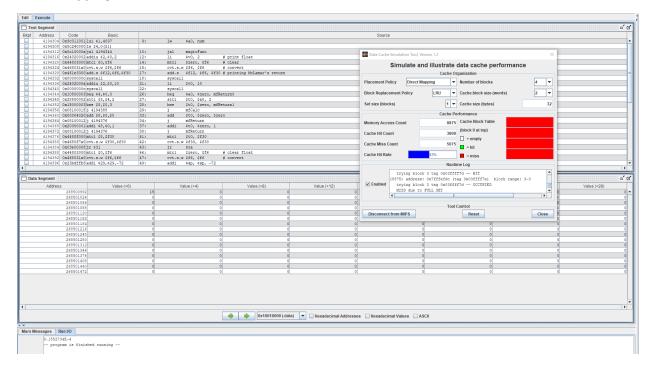
lwc1 \$f0, 8(\$sp)

lwc1 \$f2, 40(\$sp)

addi \$sp, \$sp, 72

jr \$ra

#### **Direct Mapping**



Number of blocks: 4

Cache block size: 2

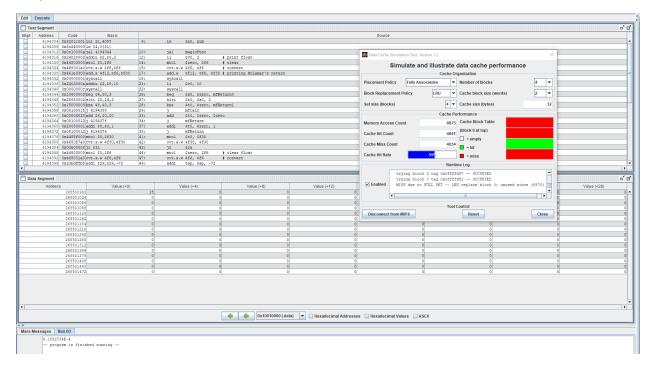
Memory access count: 8875

Cache hit count: 3800

Cache miss count: 5075

Cache hit ratio: 43%

#### **Fully Associative**



Number of blocks: 4

Cache block size: 2

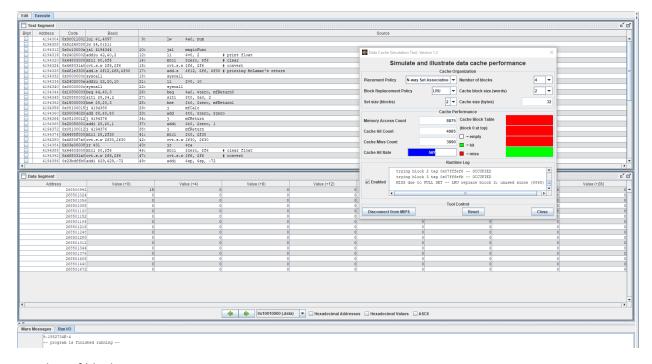
Memory access count: 8875

Cache hit count: 4841

Cache miss count: 4034

Cache hit ratio: 55%

#### **Set Associative**



Number of blocks: 4

Cache block size: 2

Memory access count: 8875

Cache hit count: 4985

Cache miss count: 3890

Cache hit ratio: 56%

#### Conclusion

This program help visualize how mapping affects performance. In this program the number 15 was put into a function. This allowed the program to return the same value after stepping through the same steps through all 3 mapping techniques. Because the function only did arithmetic, I set the blocks and size to a smaller amount so that it may use the cache more rigorously. This helped depict the different in hit ration among the set techniques and direct. Direct was obviously returned the lowest hit ratio. The difference between fully and set differed by only a percent as the instructions were not as complex. Even with these restrictions, the performance increase may still be seen amount the different techniques.

#### **Application 2:**

This application accesses an array and returns information such as steps size and rep count in order to better visualize cache operation and performance. Because of the random nature of the program, hit ratio was not constant but stayed in a range among different mappings.

## Code:

```
# ECE365 Project Phase 3 Application 2
```

# Alex Santana

.data

arr: .space 2048 # max array size in BYTES

.text

main:

```
li $a0, 256 # array size in BYTES
```

li \$a1, 8 # step size

li \$a2, 2 # rep count

li \$a3, 1

jal wordAccess

#END PROGRAM----

li \$v0, 10

syscall

#-----

## wordAccess:

la \$s0, arr # array pointer

addu \$s1, \$s0, \$a0 # array limit

sll \$t1, \$a1, 2 # inc step

wordLP:

move \$s6, \$a0

move \$s7, \$a1

move \$s5, \$v0

addiu \$a0, \$0, 100 # seed for random number generator

addiu \$a1, \$a1, 0

addiu \$v0, \$0, 42 # syscall 42 is random int range

syscall

sll \$a0, \$a0, 2 # offset

addu \$s4, \$s0, \$a0 # move pointer

sw \$0, 0(\$s4) # array[(index+offset)/4] = 0

move \$a0, \$s6

move \$a1, \$s7

move \$v0, \$s5

wordCheck:

addu \$s0, \$s0, \$t1 # increment ptr

blt \$s0, \$s1, wordLP

addi \$a2, \$a2, -1

bgtz \$a2, wordAccess

jr \$ra

byteAccess:

la \$s0, arr # array pointer

addu \$s1, \$s0, \$a0 # array limit

byteLP:

beq \$a3, \$0, byteZero

lbu \$t0, 0(\$s0) # inc index

addi \$t0, \$t0, 1

sb \$t0, 0(\$s0)

j byteCheck

byteZero:

sb \$0, 0(\$s0) # reset index

byteCheck:

addu \$s0, \$s0, \$a1 # inc pointer

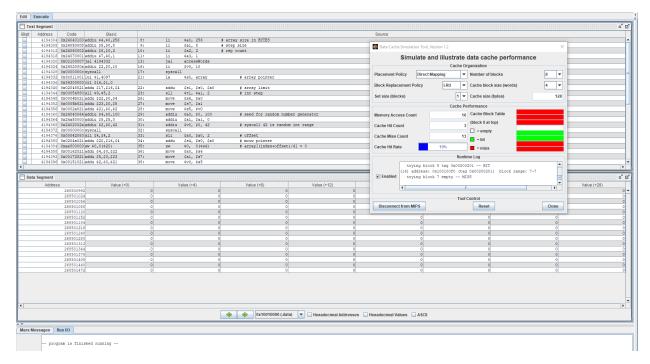
blt \$s0, \$s1, byteLP

addi \$a2, \$a2, -1

bgtz \$a2, byteAccess

jr \$ra

#### **Direct Mapping**



Number of blocks: 8

Cache block size: 4

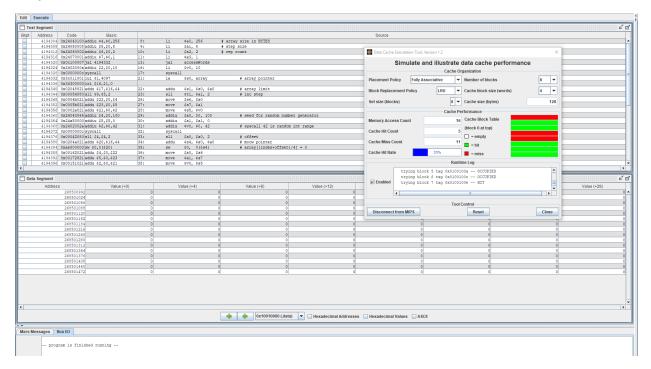
Memory access count: 16

Cache hit count: 3

Cache miss count: 13

Cache hit ratio: 19%

#### **Fully Associative**



Number of blocks: 8

Cache block size: 4

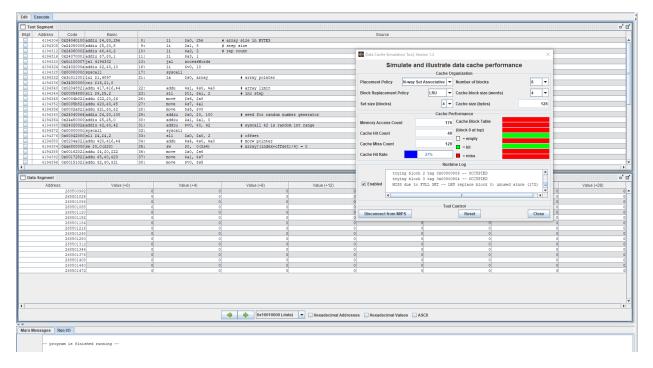
Memory access count: 16

Cache hit count: 5

Cache miss count: 11

Cache hit ratio: 31%

#### **Set Associative**



Number of blocks: 8

Cache block size: 4

Memory access count: 176

Cache hit count: 48

Cache miss count: 128

Cache hit ratio: 27%

#### Conclusion

I noticed the hit rate was non-deterministic. I believe this was due to the random nature of the program, but I was able to still conclude that improvement was attained when moving away from direct mapping. Because of the small number of blocks and size, direct mapping resulted in a hit ratio of under 20% as the rewriting of the black in the small case cause misses repeatedly. Upon changing mapping, =I noticed an increase in hit count and ratio as I was testing different number of sets. Fully associative and set associative mappings gave similar results but not as close as the previous program.