

Upgrading Aragon Voting

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EthCC - March 6th, 2019



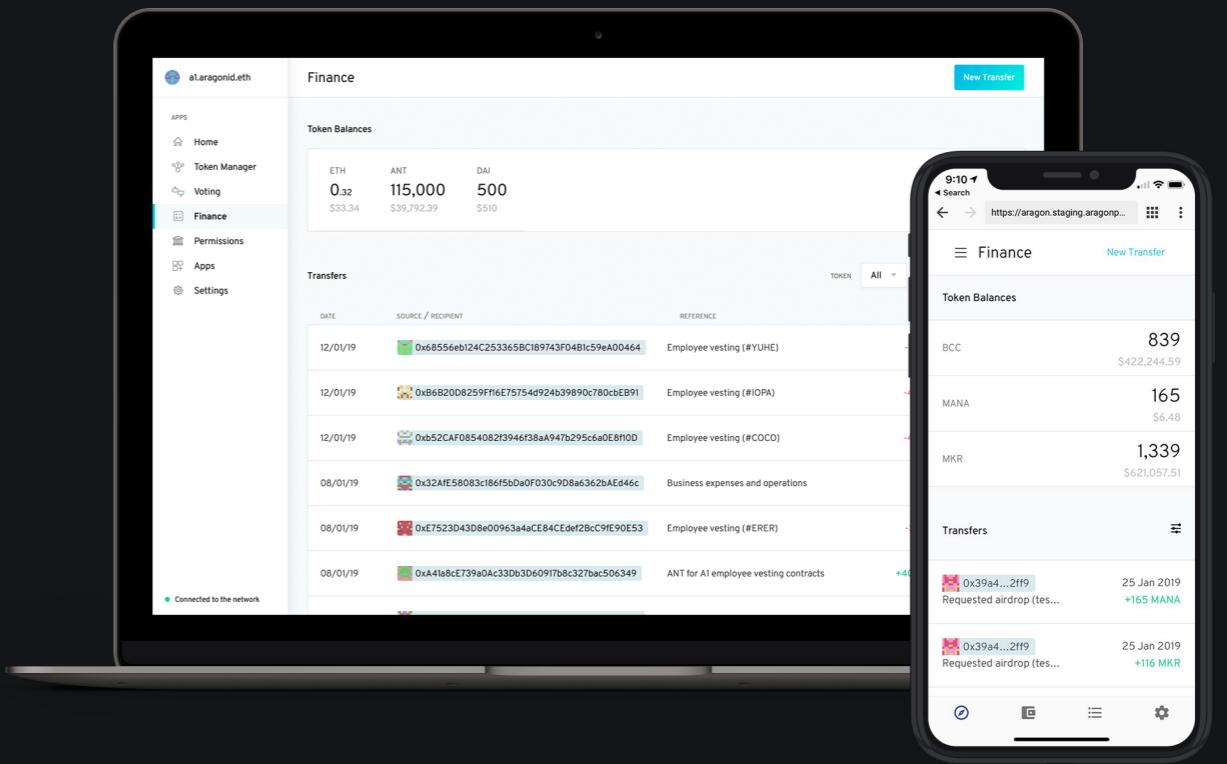
Aragon v0.5 beta
launch in Rinkeby

NEXT
WEEK

Ethereum Community Conference
March 8 - 10, 2018 | Paris, France

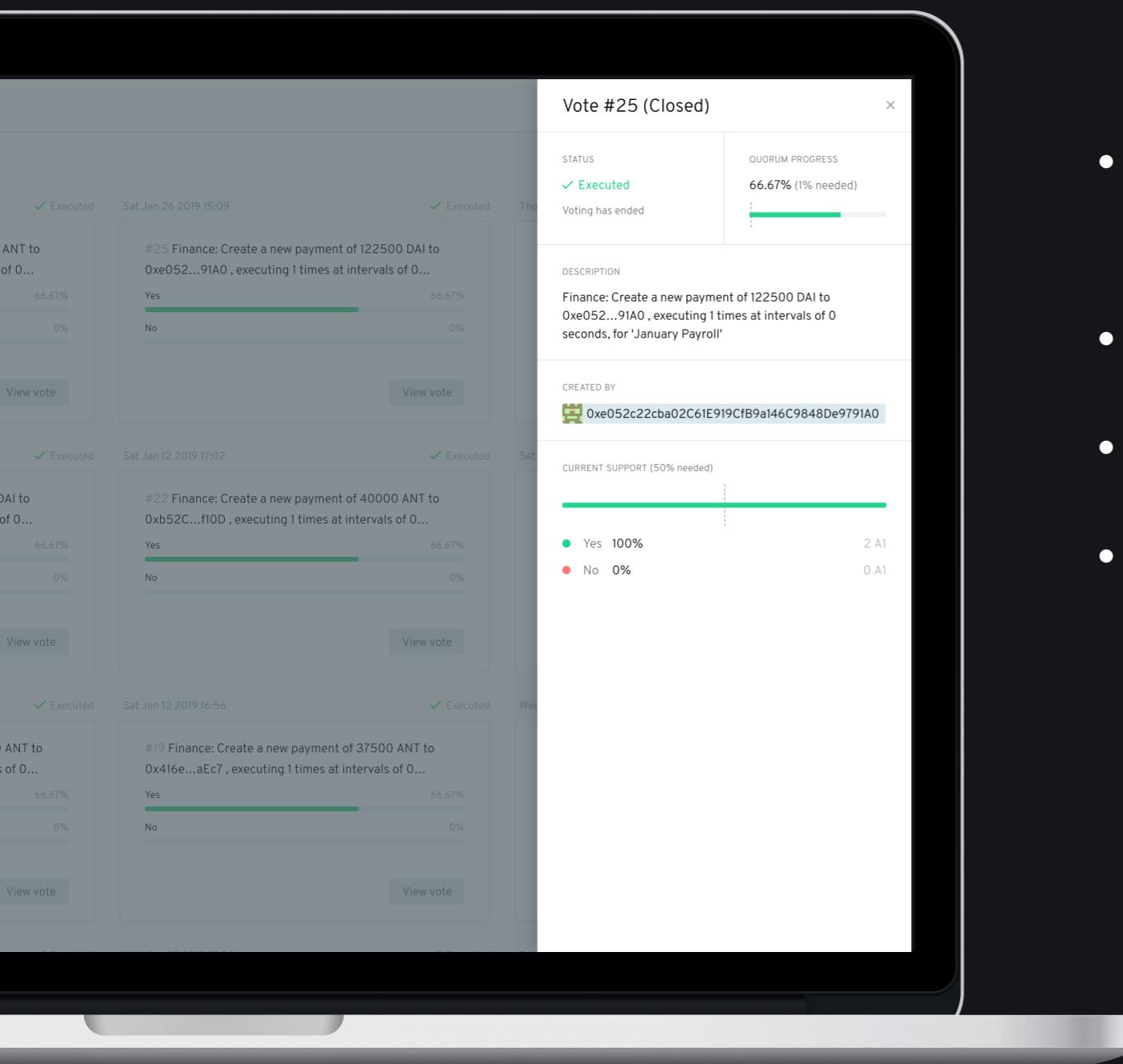
Aragon

- DApp to create and manage decentralized organizations
- Built on the Ethereum VM
- Available on the web, desktop and mobile
- 353 DAOs running on the Ethereum mainnet



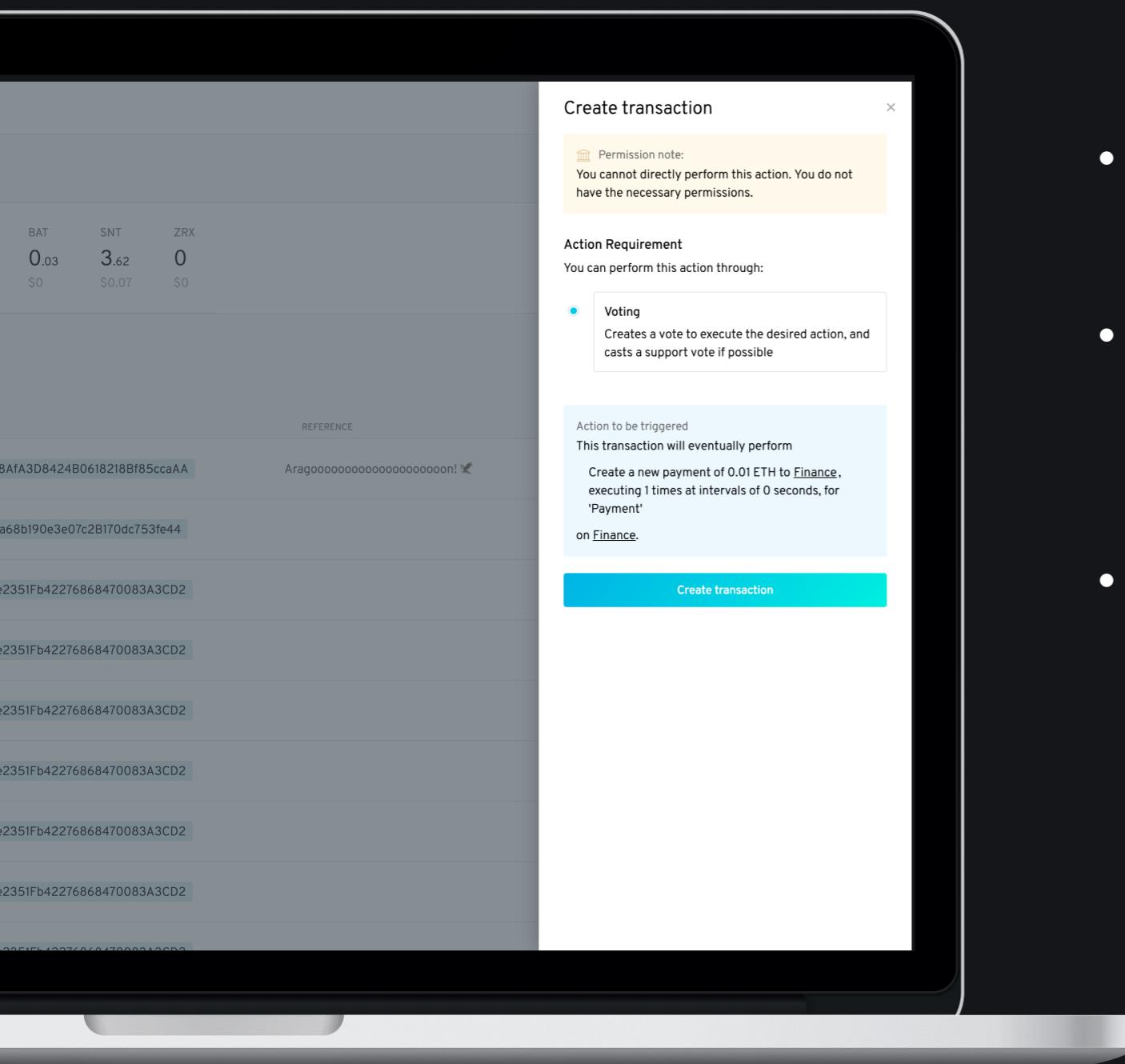
app.aragon.org

Aragon Voting v1



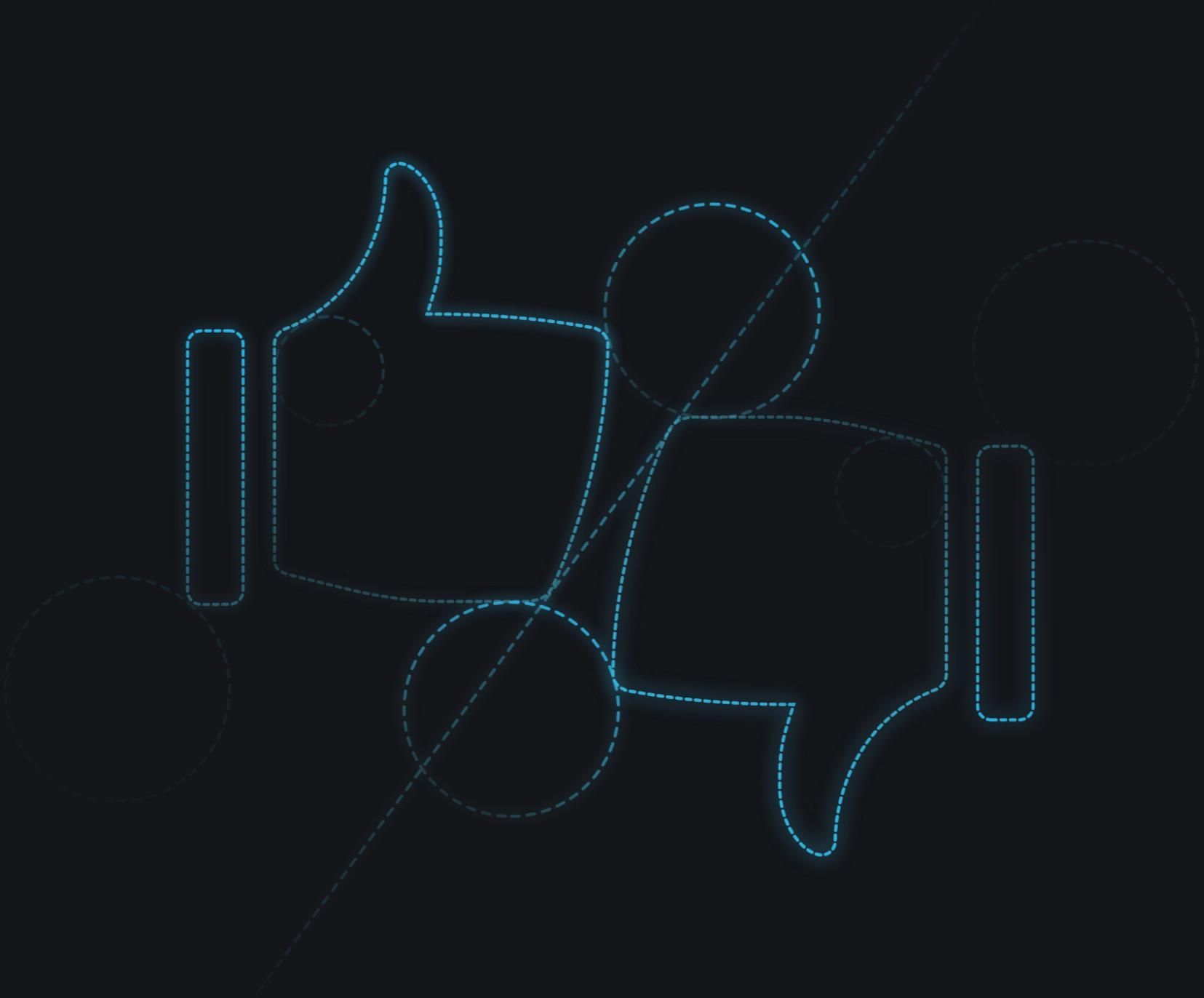
- Uses MiniMeToken as the only supported governance token
- Simple majority voting
- +\$1m secured with Voting
- Voting is the default root authority for Aragon DAOs

Forwarding



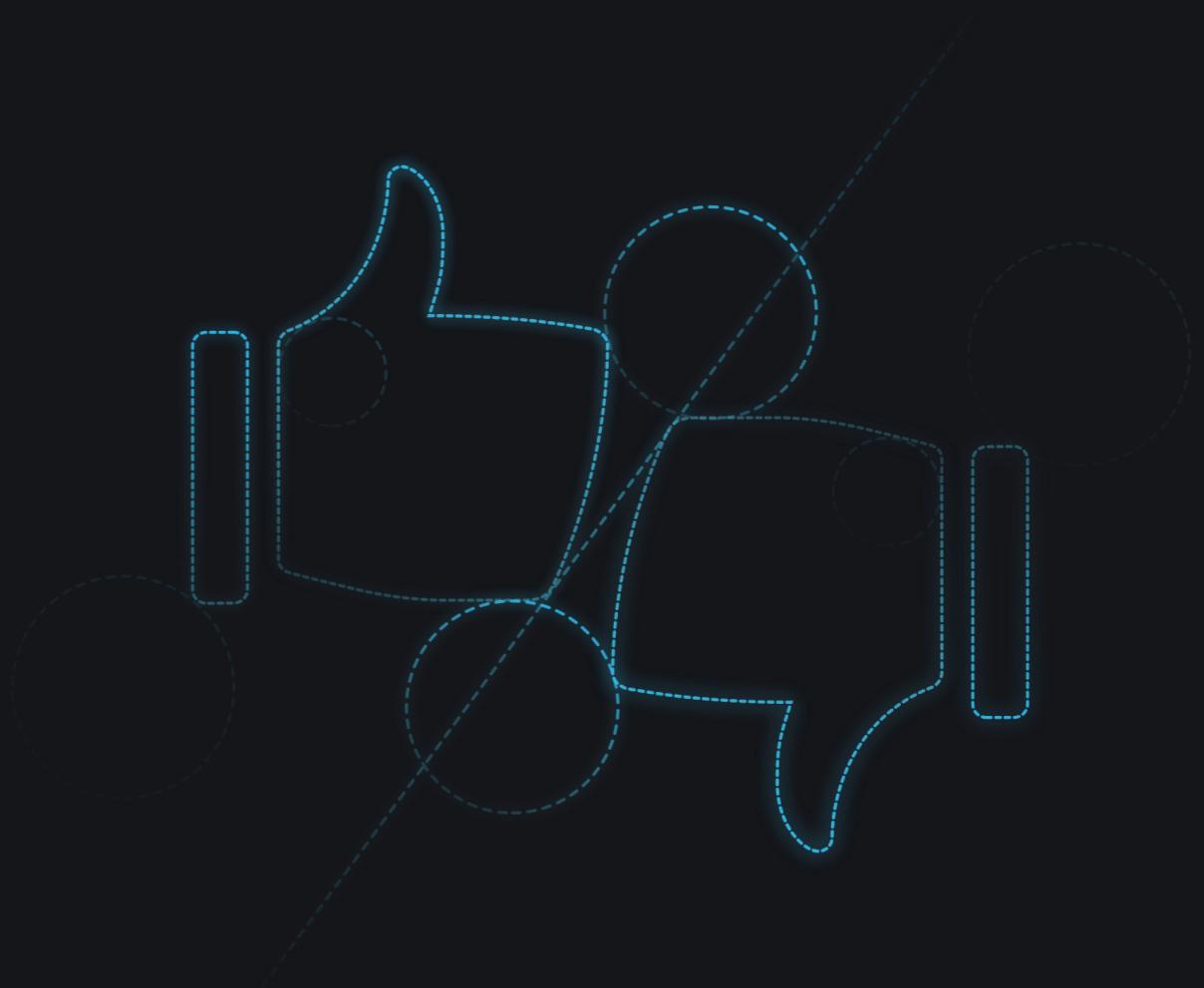
- aragonOS primitive for encoding complex permissions
- Users express their intent and Aragon figures out how to make it happen
- Forwarders can execute aragonOS' EVMScripts when some arbitrary logic checks out

Aragon Voting v2



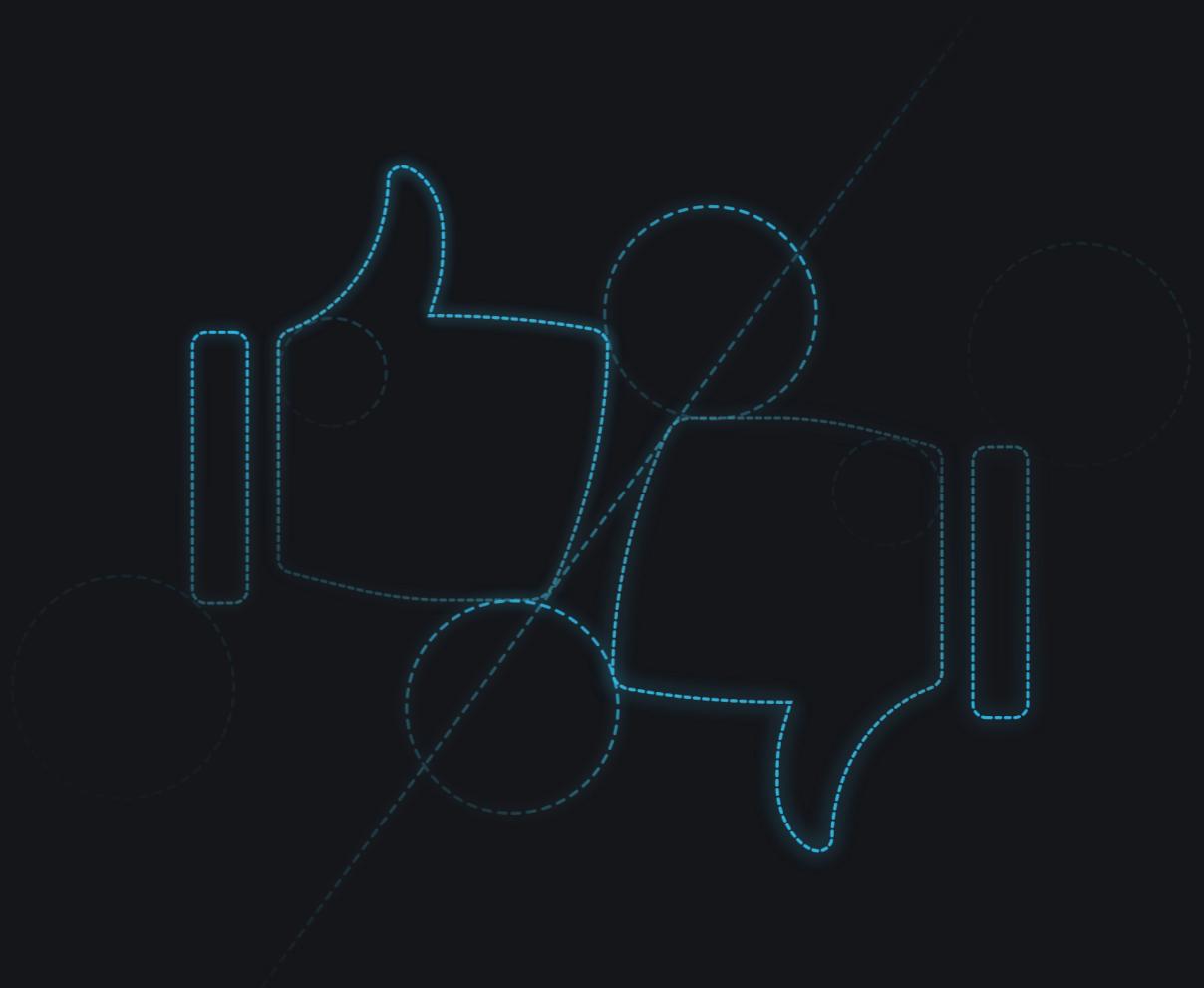
Aragon Voting v2

1. Remove MiniMe requirement



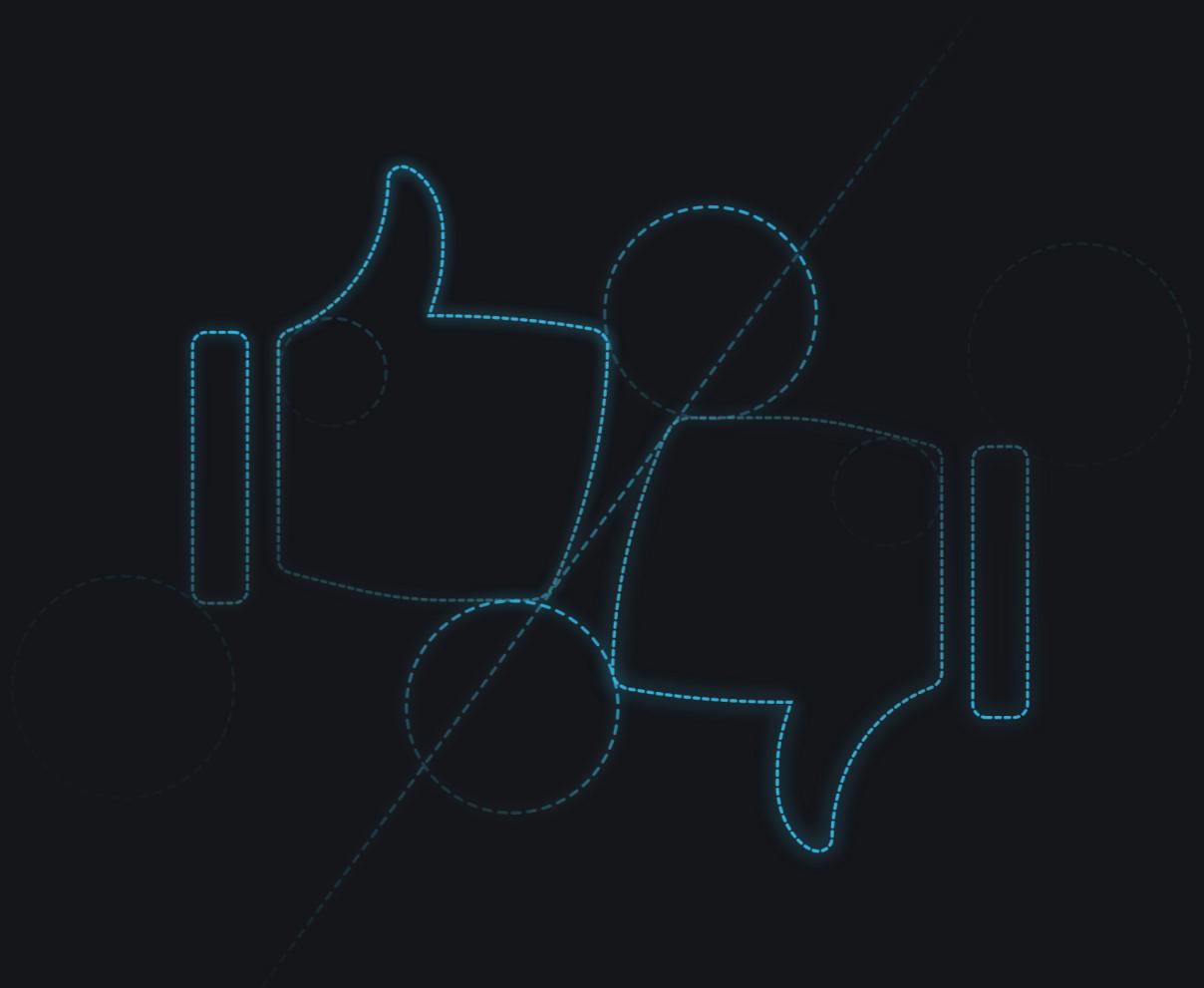
Aragon Voting v2

1. Remove MiniMe requirement
2. Really cheap voting



Aragon Voting v2

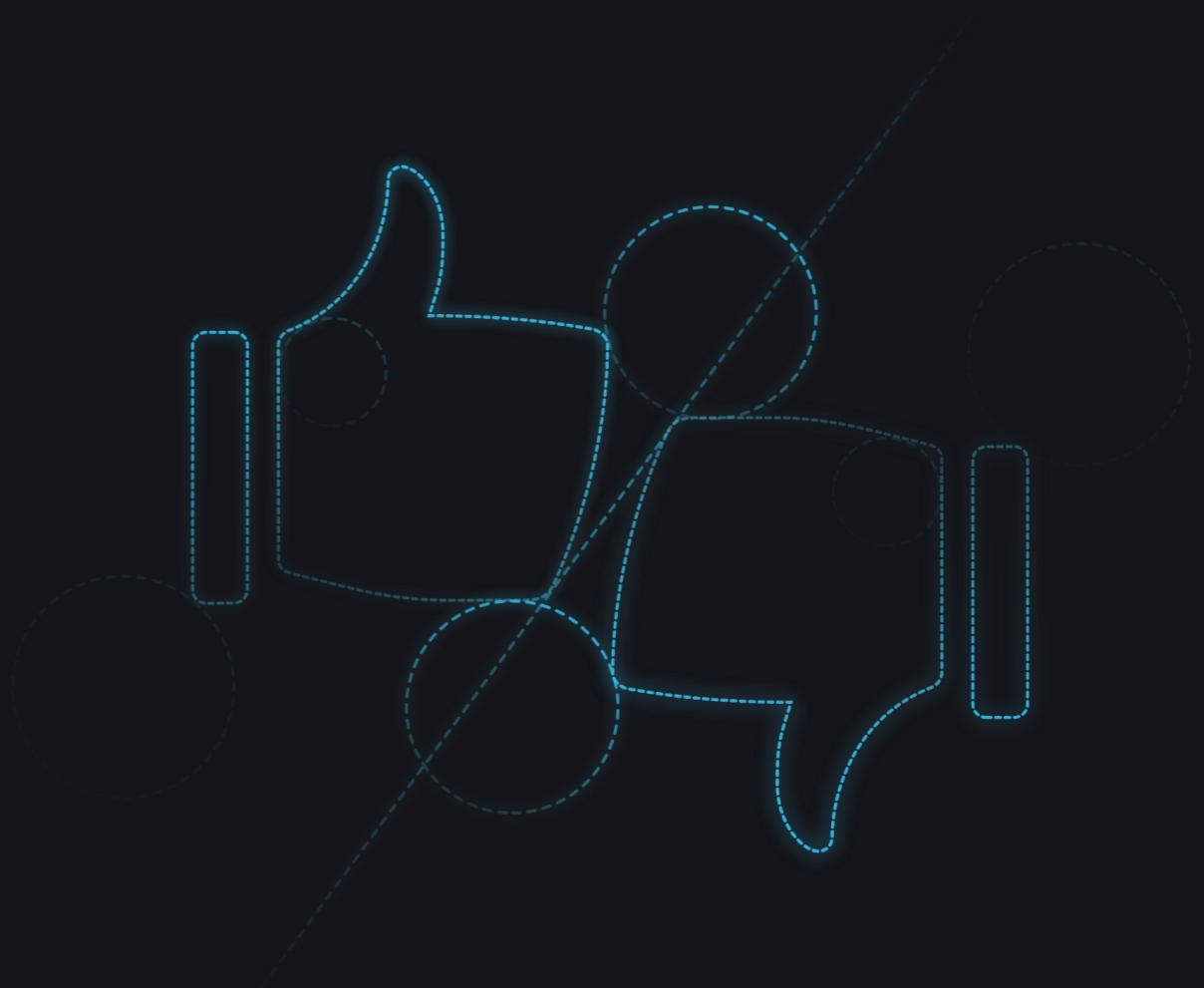
1. Support for all ERC20 tokens*



* tokens that use a standard data structure for storage

Aragon Voting v2

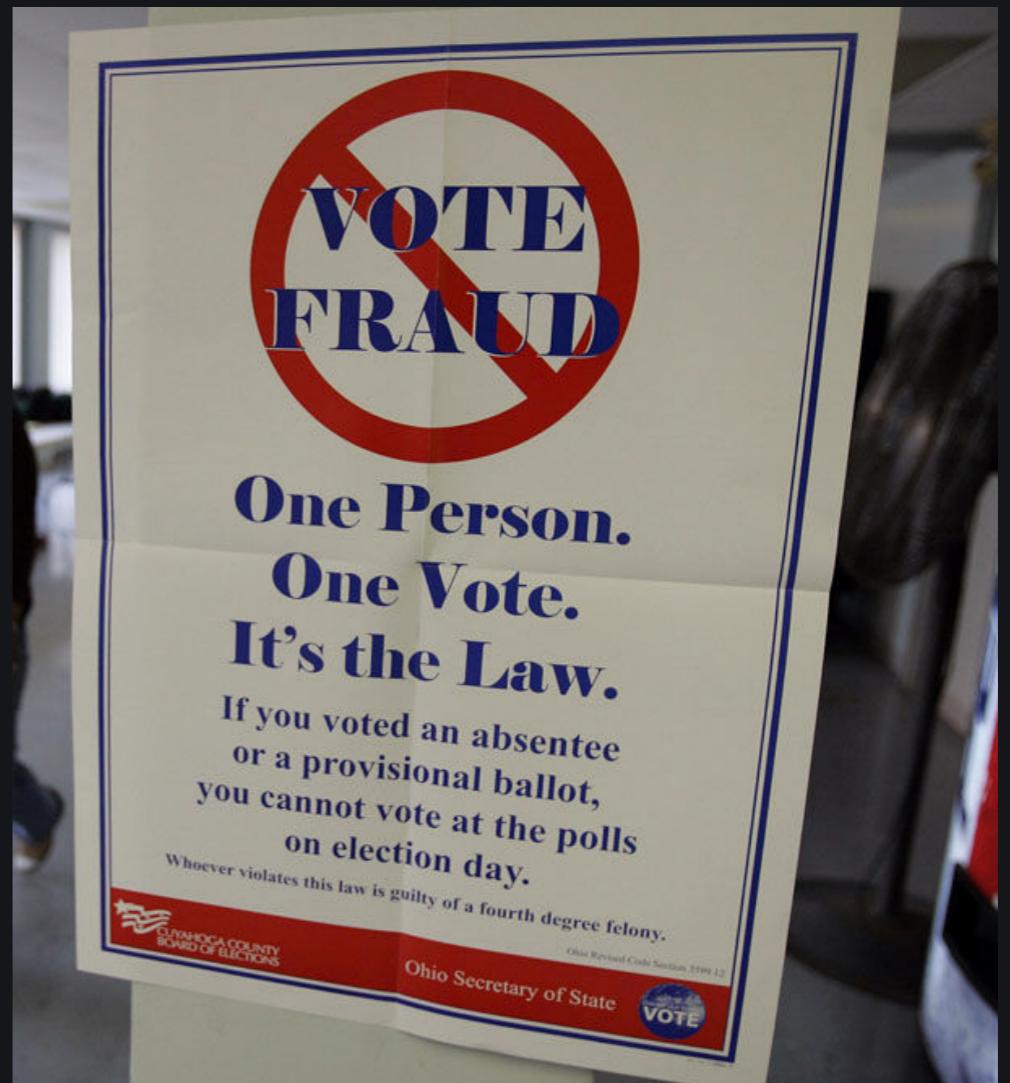
1. Support for all ERC20 tokens*
2. Off-chain vote aggregation



* tokens that use a standard data structure for storage

The double voting problem

- ERC20 tokens are fungible
- Can't detect if one particular token already voted, unless:
 - Tokens are locked during the vote location (staking)
 - Token balances can be snapshotted (MiniMeToken)



EVM Storage Proofs

- Prove what the value a given storage slot in a contract was at a past block height
- Derive storage values from a trusted block hash
- Uses `eth_getProof` (EIP1186)
- For Voting, prove the token balance of a holder at a snapshot block

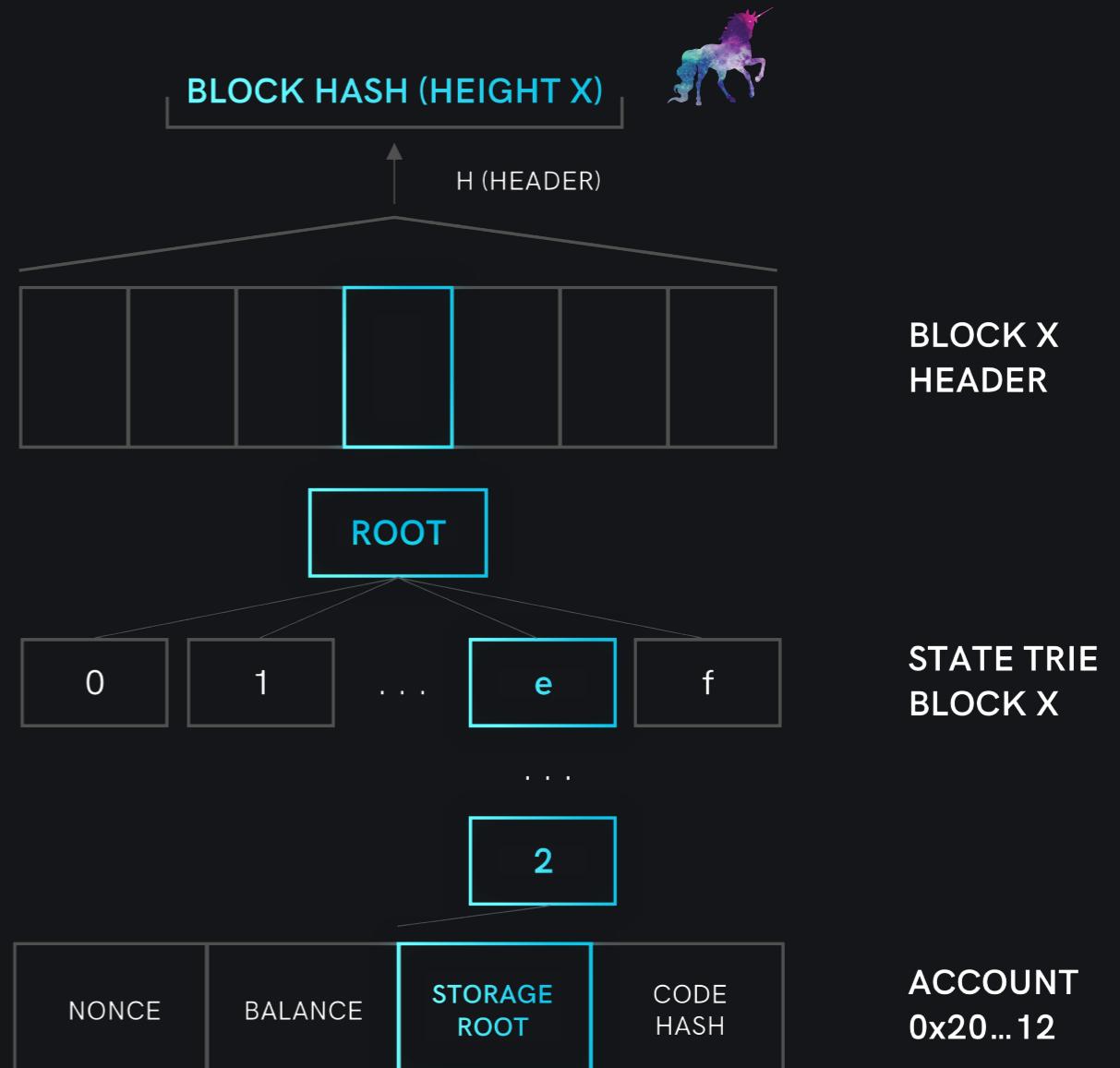


The Ethereum Storage Time Machine

Account proof

Block hash → storage root

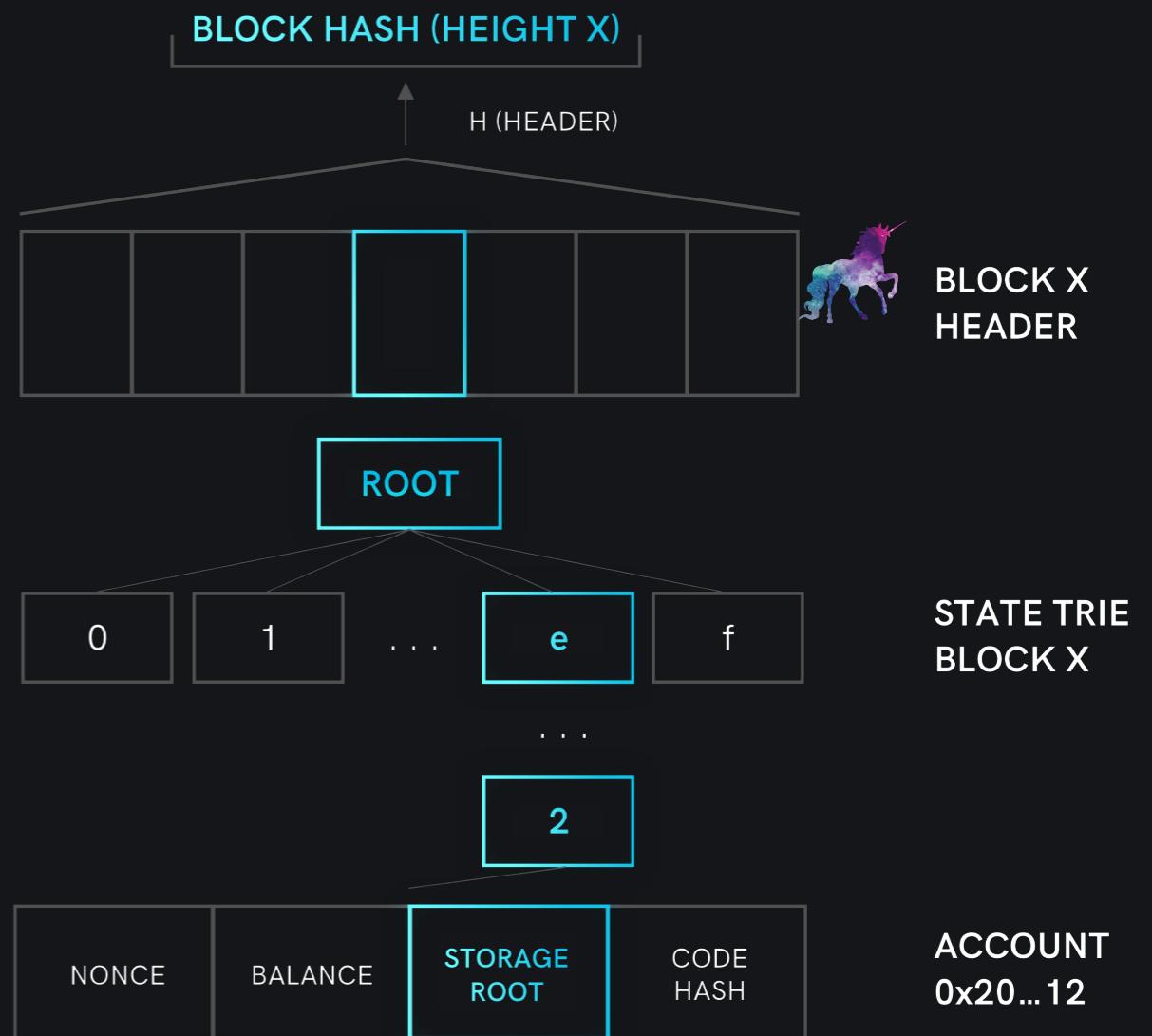
- Use **BLOCKHASH** to get block hash for block X



Account proof

Block hash → storage root

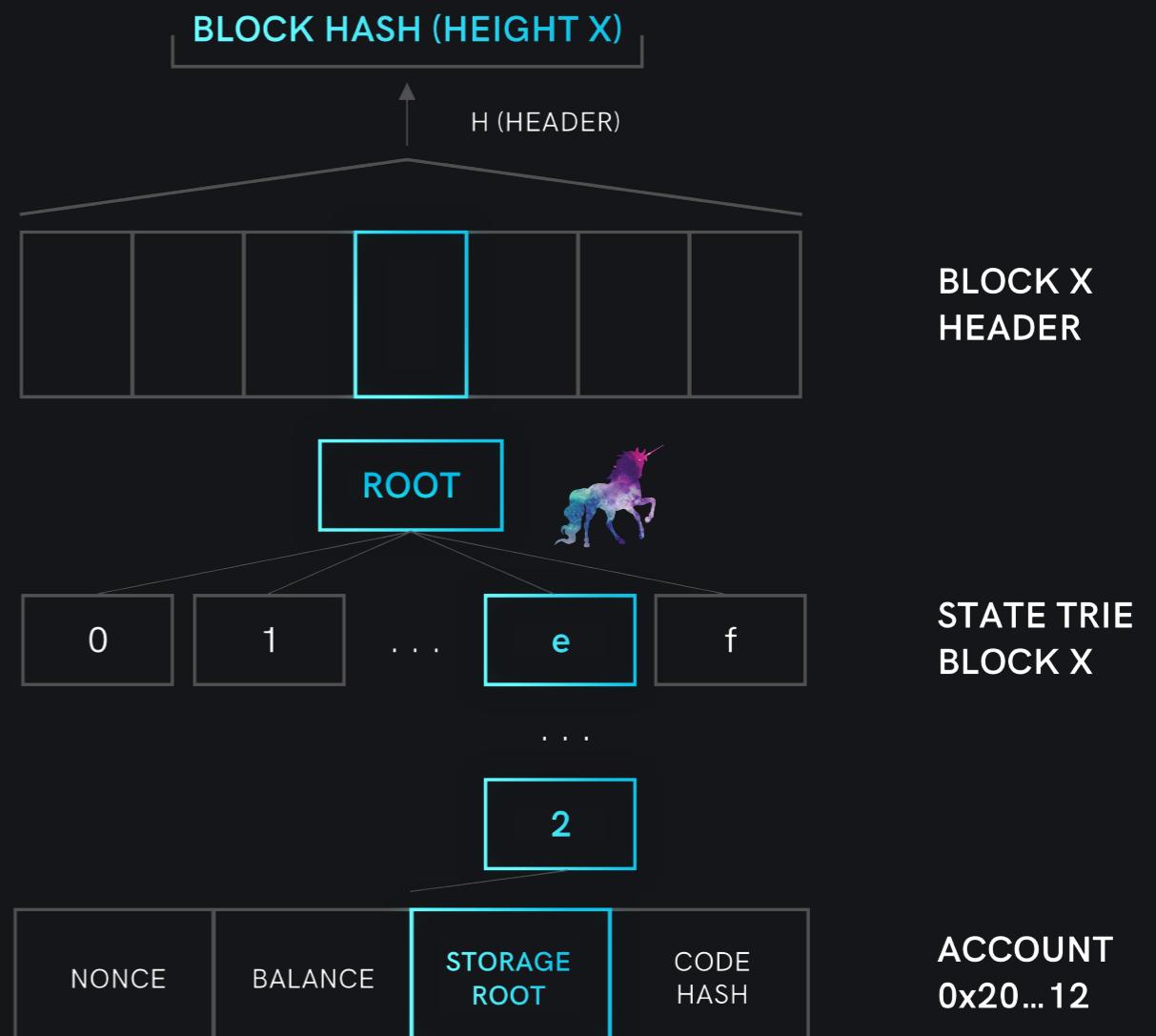
- Use **BLOCKHASH** to get block hash for block X
- Submit the block header and verify its hash matches



Account proof

Block hash → storage root

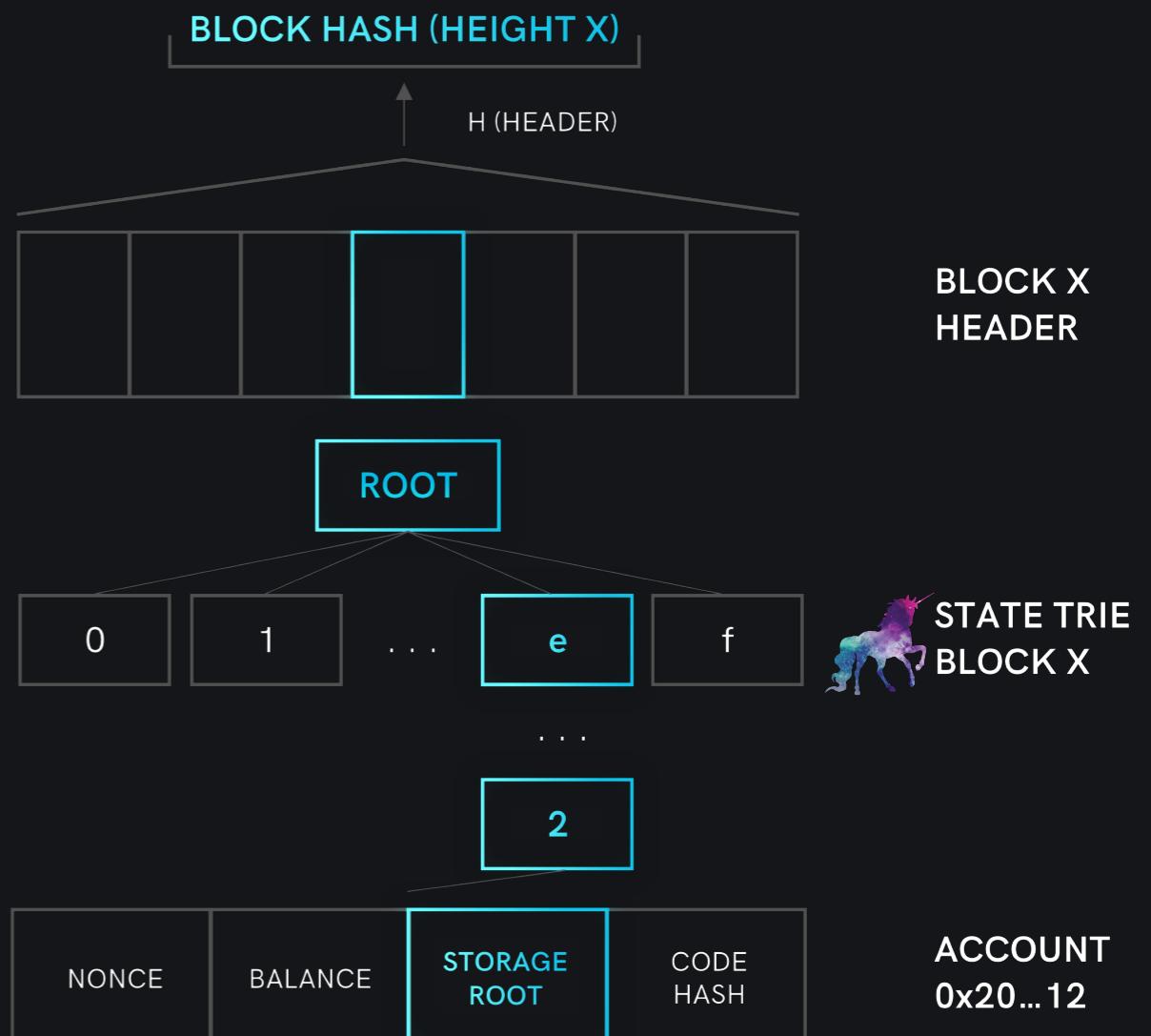
- Use **BLOCKHASH** to get block hash for block X
- Submit the block header and verify its hash matches
- Extract the state root from the block header



Account proof

Block hash → storage root

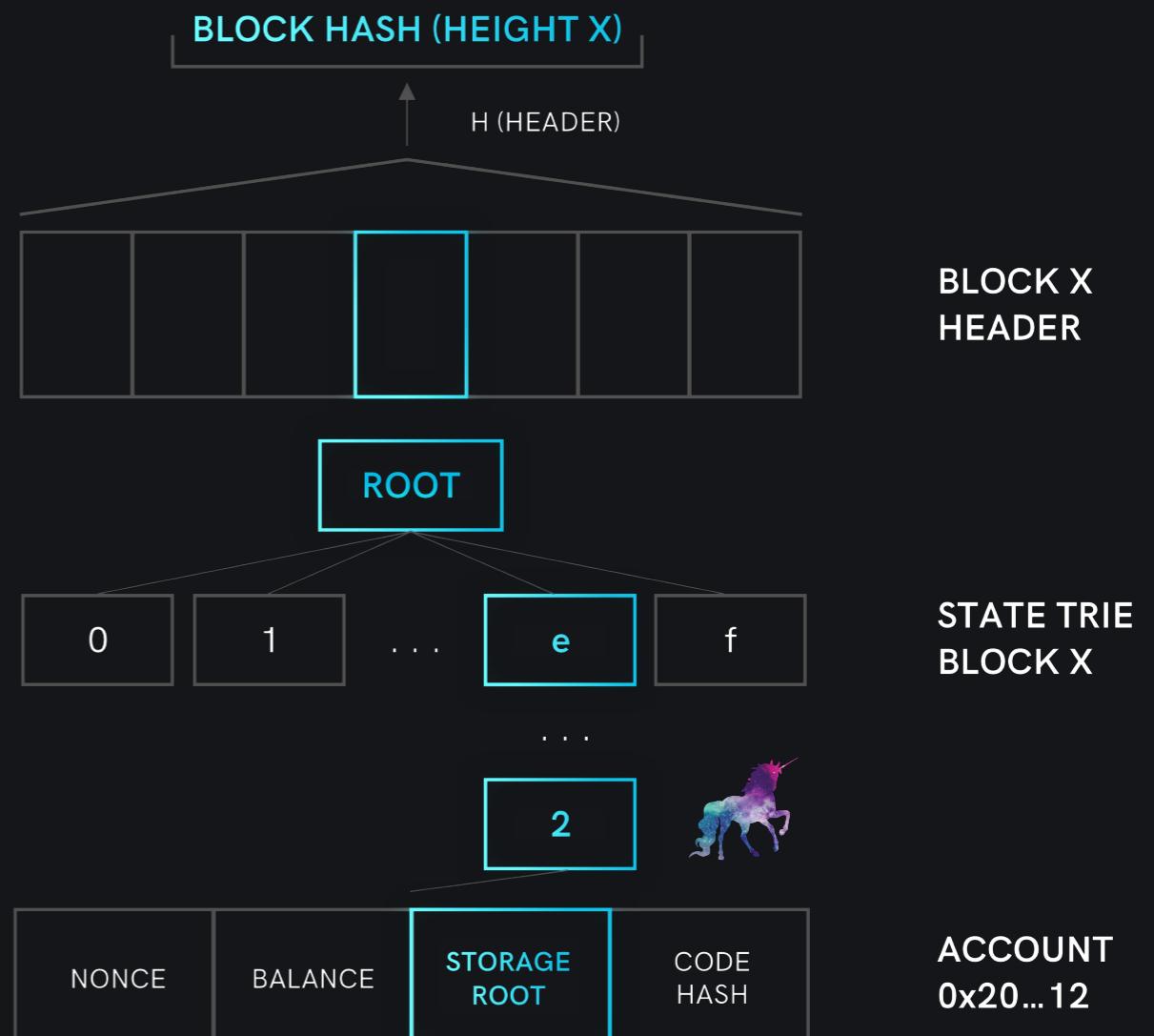
- Use **BLOCKHASH** to get block hash for block X
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- Extract the state root from the block header
- Submit merkle proof of the account state in the state trie



Account proof

Block hash → storage root

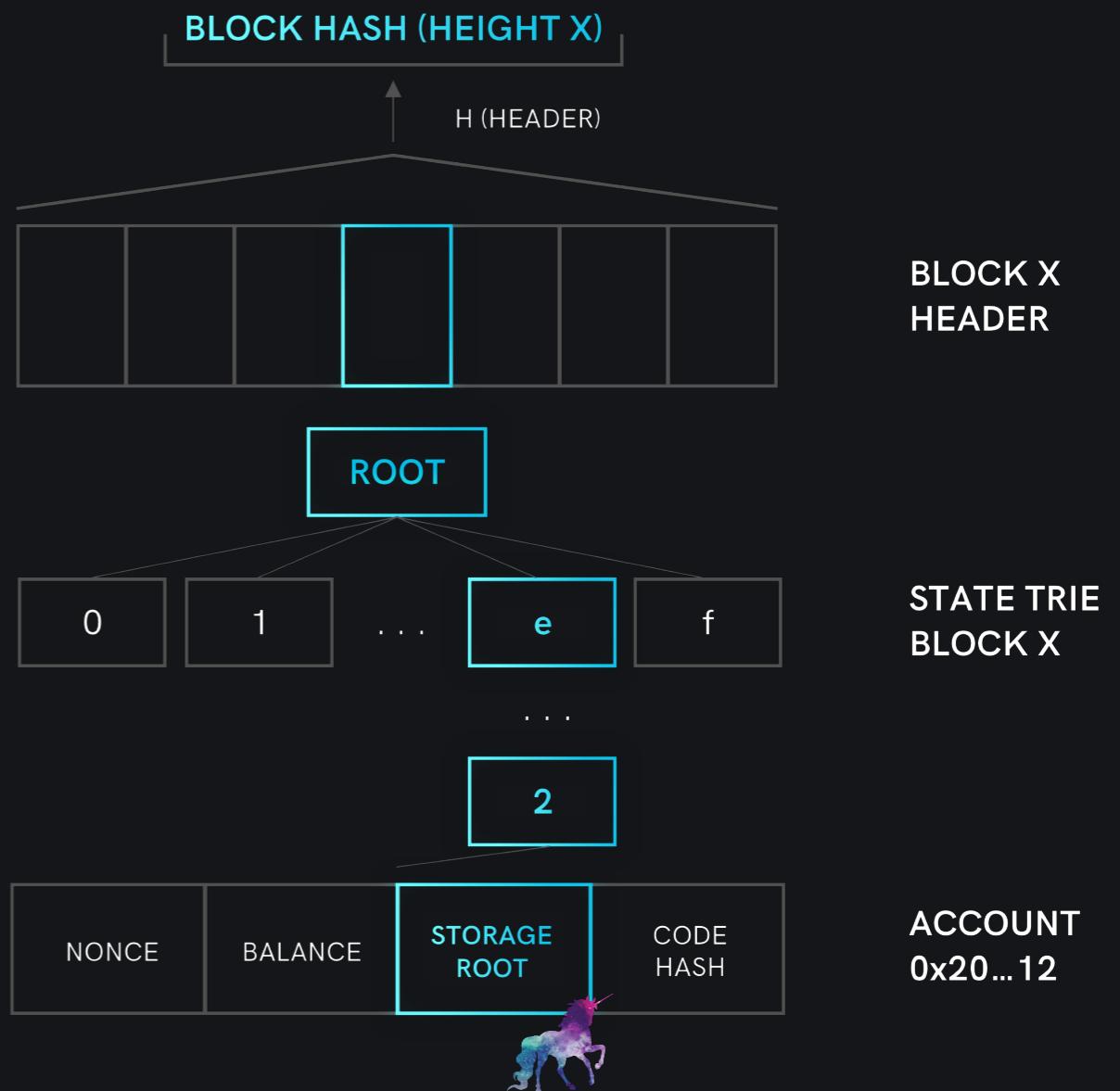
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- Verify merkle proof against the state root



Account proof

Block hash → storage root

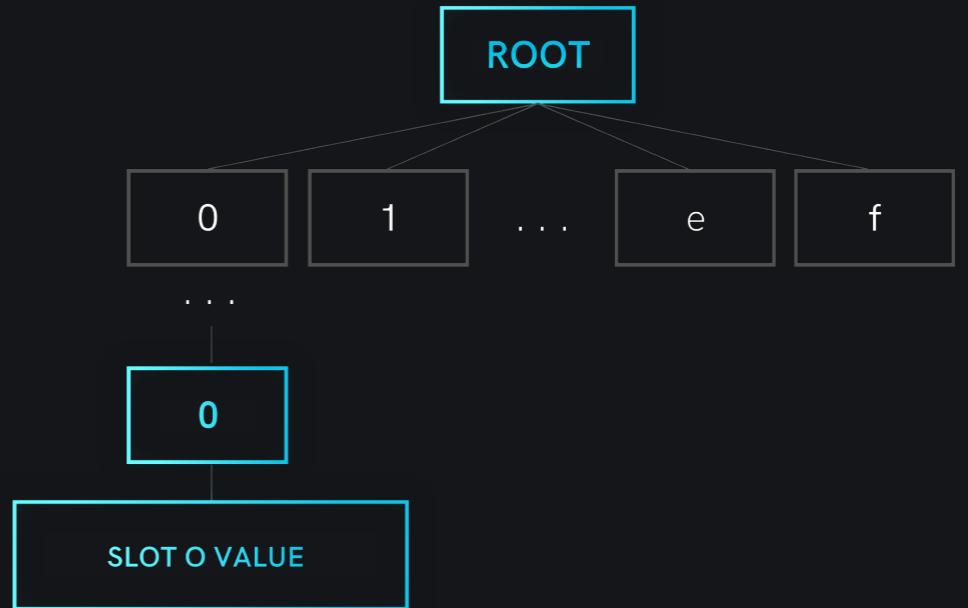
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- Submit the block header and verify its hash matches
- Extract the state root from the block header
- Submit merkle proof of the account state in the state trie
- Verify merkle proof against the state root
- Extract storage root from the account state



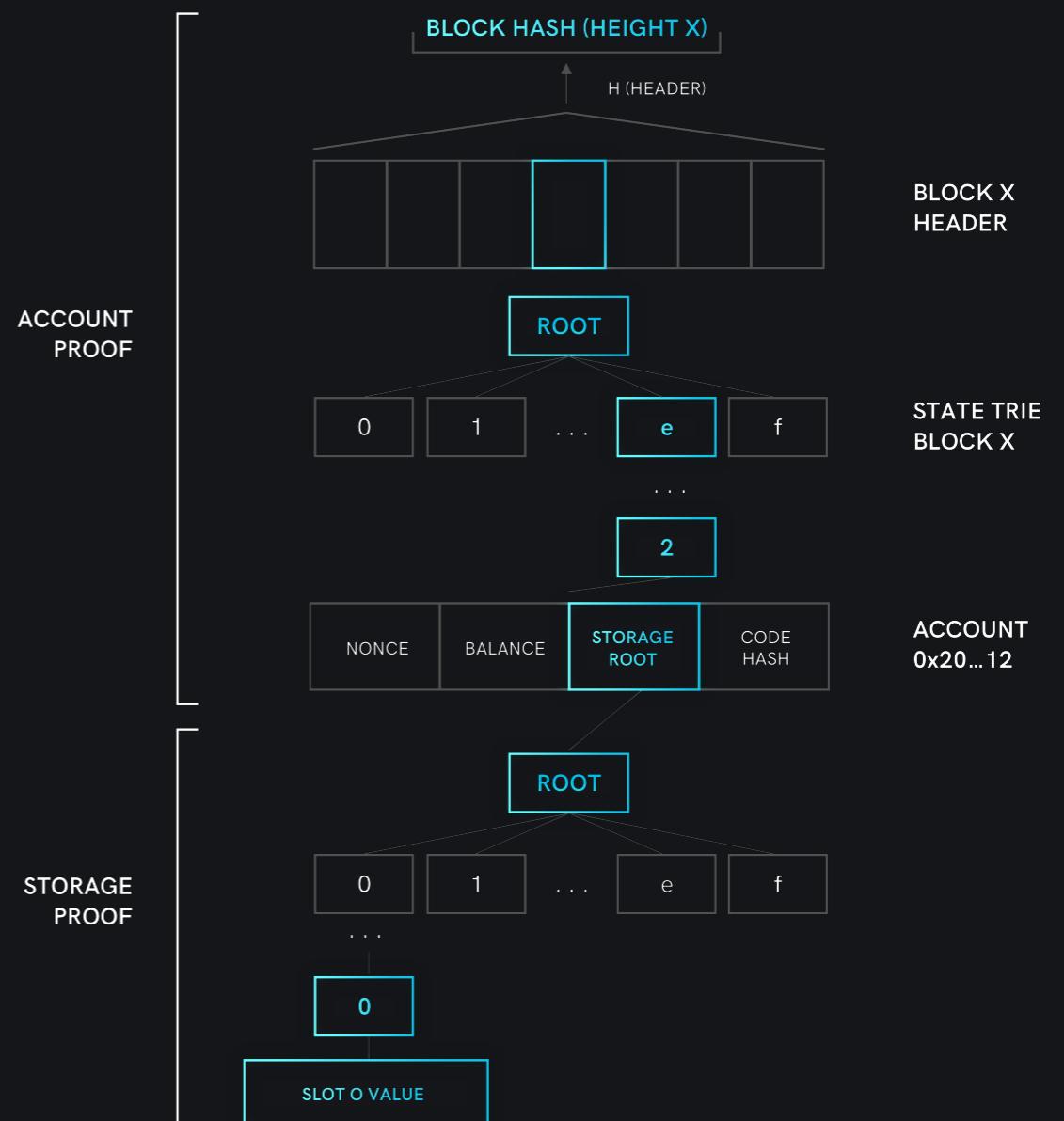
Storage proof

Storage root → slot value

- With an authenticated storage root
- Submit merkle proof of the storage slot value in the storage



EVM Storage Proofs



EVM Storage Proofs for ERC20s

Upsides

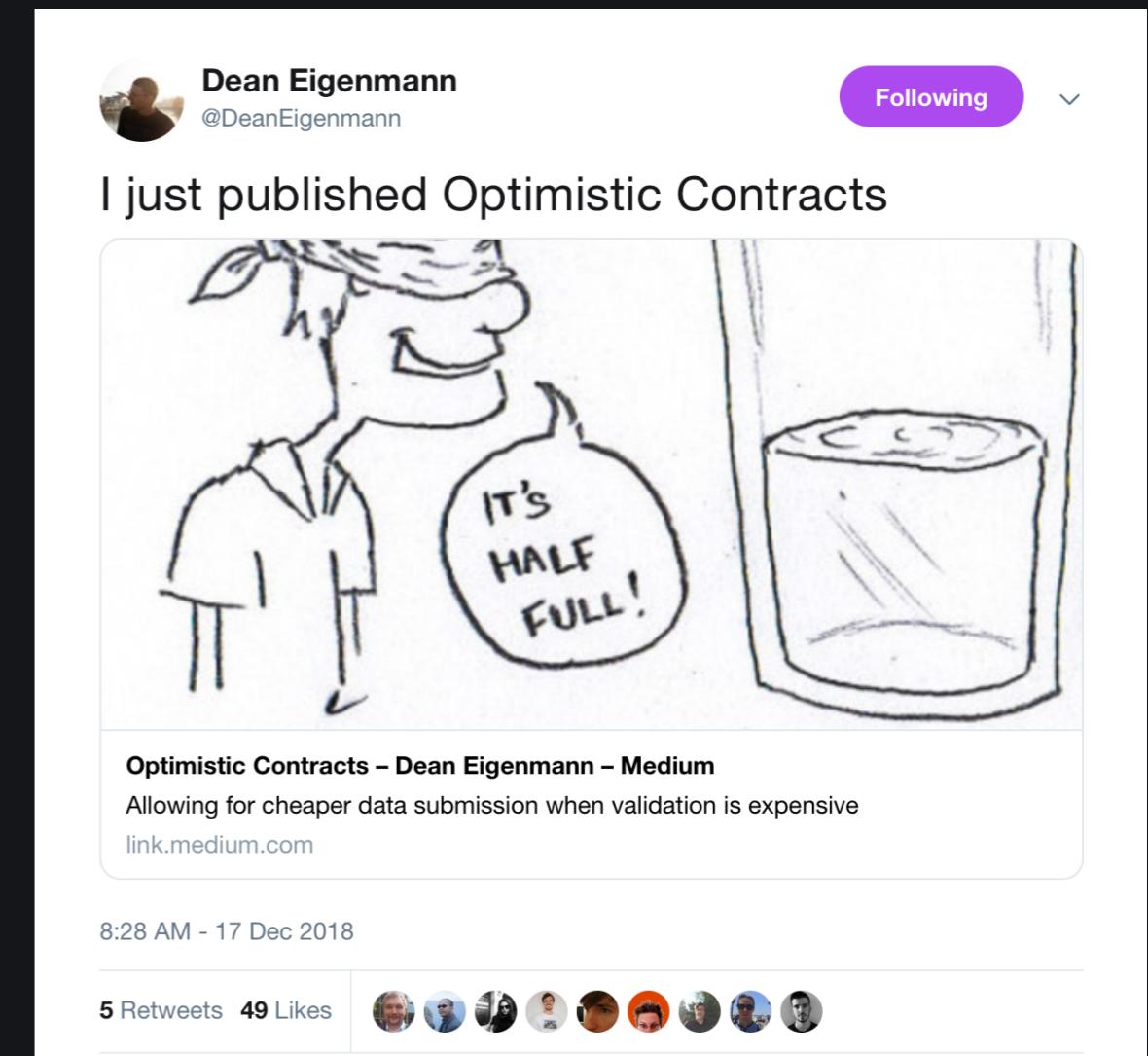
- Use all existing ERC20 tokens for Voting
- Snapshotting doesn't impose a cost for all transfers

Downsides

- Verifying a full proof uses ~500,000 gas
- **BLOCKHASH** can still only get the last 256 blocks
- Breaks EVM abstraction by checking storage directly, use carefully

Optimistic contracts

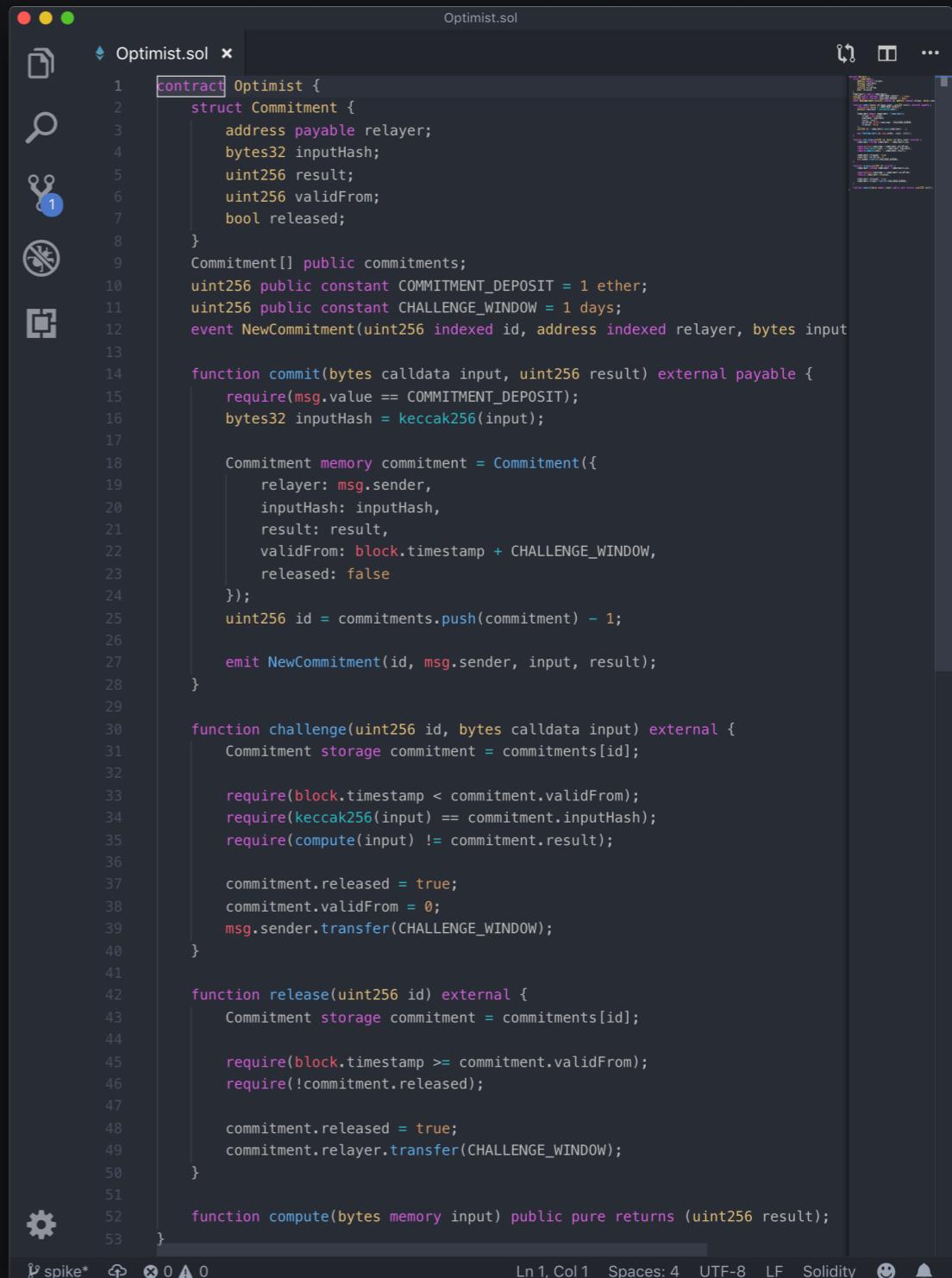
- Offloading computation heavy operations from the network
- Ensure computation correctness with bonded relayers instead of repeating it in every node
- Naïve implementation can scale up to 80x*: use 100k gas for computations that would take the block gas limit (8M)



* for an optimistic execution with no input

Optimistic in 50 LOC

- Commit to the result of executing a pure function with some input
- Challenges can be done during a certain time window, which executes the computation and compares to the committed result
- Challenger is awarded the deposit of the relayer on a valid challenge
- After the challenge window has passed, the result is taken as valid



```
contract Optimist {
    struct Commitment {
        address payable relayer;
        bytes32 inputHash;
        uint256 result;
        uint256 validFrom;
        bool released;
    }
    Commitment[] public commitments;
    uint256 public constant COMMITMENT_DEPOSIT = 1 ether;
    uint256 public constant CHALLENGE_WINDOW = 1 days;
    event NewCommitment(uint256 indexed id, address indexed relayer, bytes input);

    function commit(bytes calldata input, uint256 result) external payable {
        require(msg.value == COMMITMENT_DEPOSIT);
        bytes32 inputHash = keccak256(input);

        Commitment memory commitment = Commitment({
            relayer: msg.sender,
            inputHash: inputHash,
            result: result,
            validFrom: block.timestamp + CHALLENGE_WINDOW,
            released: false
        });
        uint256 id = commitments.push(commitment) - 1;

        emit NewCommitment(id, msg.sender, input, result);
    }

    function challenge(uint256 id, bytes calldata input) external {
        Commitment storage commitment = commitments[id];

        require(block.timestamp < commitment.validFrom);
        require(keccak256(input) == commitment.inputHash);
        require(compute(input) != commitment.result);

        commitment.released = true;
        commitment.validFrom = 0;
        msg.sender.transfer(CHALLENGE_WINDOW);
    }

    function release(uint256 id) external {
        Commitment storage commitment = commitments[id];

        require(block.timestamp >= commitment.validFrom);
        require(!commitment.released);

        commitment.released = true;
        commitment.relayer.transfer(CHALLENGE_WINDOW);
    }

    function compute(bytes memory input) public pure returns (uint256 result);
}
```

Voting Relay Protocols (VRP)

- Family of layer 2 protocols that aggregate votes off-chain and submit partial tally **optimistically**
- Update the tally once for many votes without verifying storage proofs
- Individual votes are multiple orders of magnitude cheaper to cast
- Vast design space with different trade-offs
- Voters just sign their vote message and send it to the relayer



Simple VRP (SVRP)

- VRP, except that **only one relayer** is allowed to relay votes at a time
- The relayer slot can be centrally appointed or auctioned
- The relayer stakes tokens as collateral that can be slashed if a fraud proof is submitted
- The relayer may ask for a fee payment to include a vote or sponsor voting

SVRP and centralization

- If the relayer doesn't return a signed inclusion commitment, the voter can submit the vote directly
- If the relayer commits to relaying a vote and they don't, they are slashed
- Data required to prove fraud is available through EVM logs
- If someone is monitoring, the relayer will always be slashed if submitting fraudulent votes

The SVRP bottleneck

- Naïve optimistic relies on the data required to submit fraud proofs being **available with EVM logs**
- In SVRP, the input data is the array of vote messages being relayed
- Cost per input byte = **88 gas**
- 1 vote message = **68 bytes**
- Hard cap of 938 votes relayed per batch (80% of the block gas limit)
- Compared to Voting v1, SVRP scales voting by **10.72x**

VRP messages

All messages must be RLP-encoded before sending them over any transport protocol or submitting them on-chain.

Vote message

Components:

- Voting app address identifier (4 bytes: `bytes4(hash(address))`)
- Proposal id (3 bytes: 14,706,125 proposals per Voting app)
- Casted vote (1 byte: 256 options)
- Voter balance (16 bytes, claimed user balance in the voting app's token at the snapshot block of vote id)
- Voter ECDSA signature for message hash (64 bytes, with the trailing bit trick)

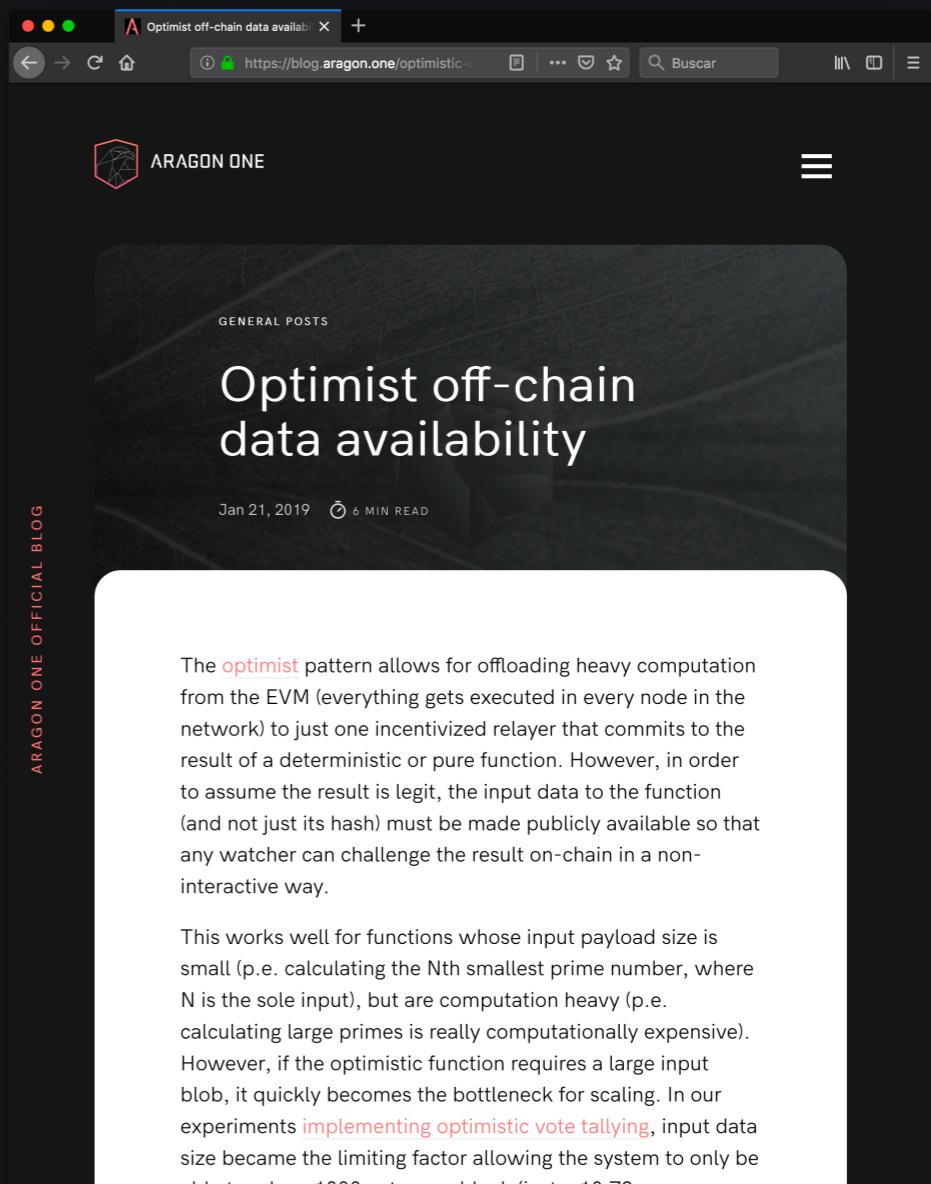
Total batch size: 88 bytes

The data availability problem

- Just submitting the hash on-chain and proving that the hashed data is available is impossible from a smart contract's worldview
- If the data is actually not made available, fraud couldn't be stopped, as the data is required for challenging
- To solve this issue, relayers could be required to submit on demand the full data blob to the chain so a challenge can happen, which can be used for griefing relayers

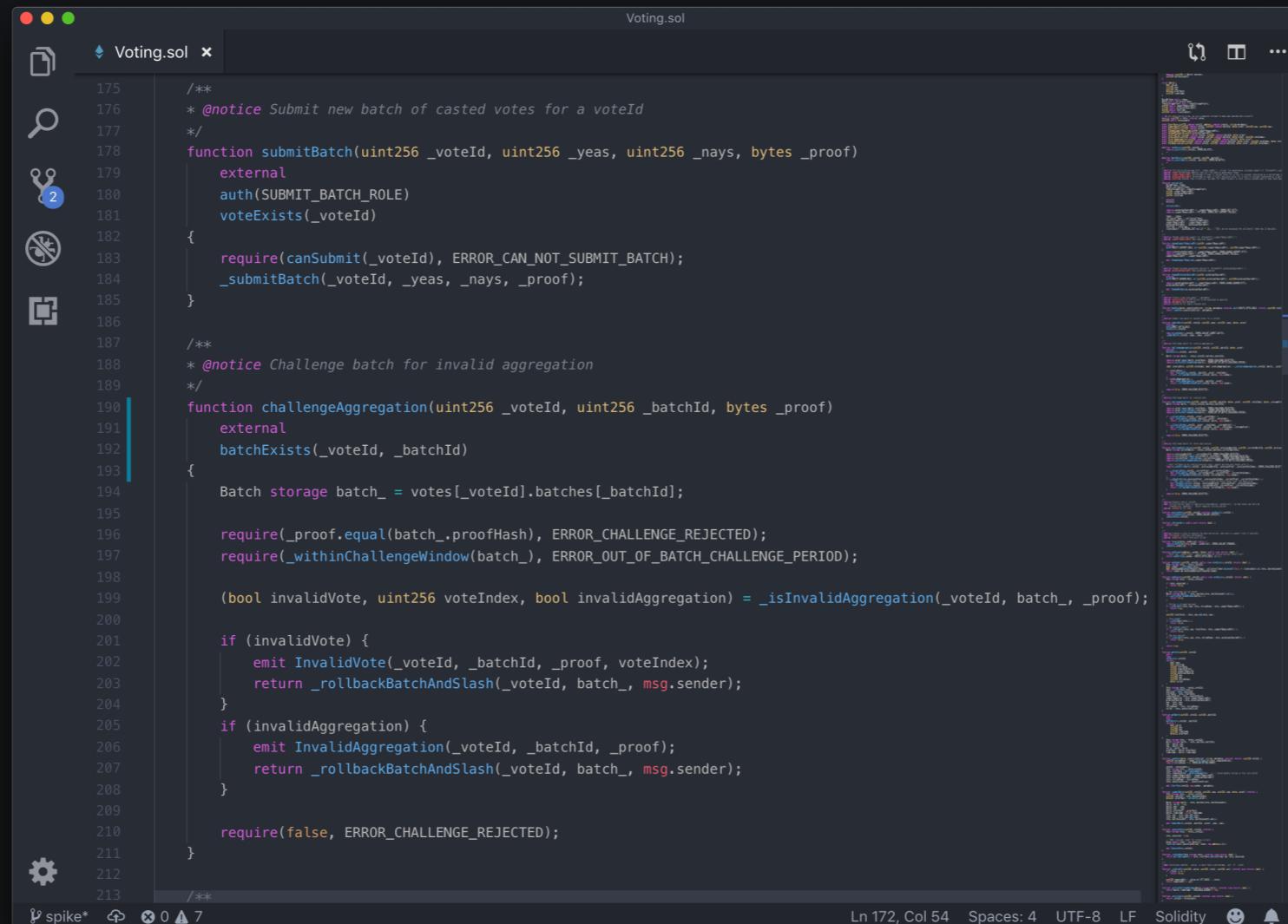
Optimistic data availability

- Used for **Availability VRP** (AVRP)
- A set of **N validators** is selected, but only one of them can relay at a time
- Validators receive relay requests and send them to other validators to receive a signed commitment that they have the data and result
- Later, a validator is randomly appointed as the relayer and they can relay it if they have signatures of **2/3 validators**
- Only one honest validator is needed for the integrity of the system (unless 2/3 validators are malicious and withhold data)



[blog.aragon.one](https://blog.aragon.one/optimistic-off-chain-data-availability/)

SVRP is now open source



A screenshot of a code editor displaying the `Voting.sol` Solidity smart contract. The code is written in Solidity and defines two main functions: `submitBatch` and `challengeAggregation`. The `submitBatch` function takes parameters for a vote ID, yeas, nays, and proof, and requires authentication and a valid vote ID. The `challengeAggregation` function takes a vote ID, batch ID, and proof, and checks if the proof is valid and within a challenge window. If either is invalid, it emits an event and rolls back the batch. The code also includes error handling for invalid aggregation.

```
175  /**
176   * @notice Submit new batch of casted votes for a voteId
177   */
178  function submitBatch(uint256 _voteId, uint256 _yeas, uint256 _nays, bytes _proof)
179      external
180      auth(SUBMIT_BATCH_ROLE)
181      voteExists(_voteId)
182  {
183      require(canSubmit(_voteId), ERROR_CAN_NOT_SUBMIT_BATCH);
184      _submitBatch(_voteId, _yeas, _nays, _proof);
185  }
186
187 /**
188 * @notice Challenge batch for invalid aggregation
189 */
190 function challengeAggregation(uint256 _voteId, uint256 _batchId, bytes _proof)
191     external
192     batchExists(_voteId, _batchId)
193  {
194     Batch storage batch_ = votes[_voteId].batches[_batchId];
195
196     require(_proof.equal(batch_.proofHash), ERROR_CHALLENGE_REJECTED);
197     require(_withinChallengeWindow(batch_), ERROR_OUT_OF_BATCH_CHALLENGE_PERIOD);
198
199     (bool invalidVote, uint256 voteIndex, bool invalidAggregation) = _isValidAggregation(_voteId, batch_, _proof);
200
201     if (invalidVote) {
202         emit InvalidVote(_voteId, _batchId, _proof, voteIndex);
203         return _rollbackBatchAndSlash(_voteId, batch_, msg.sender);
204     }
205     if (invalidAggregation) {
206         emit InvalidAggregation(_voteId, _batchId, _proof);
207         return _rollbackBatchAndSlash(_voteId, batch_, msg.sender);
208     }
209
210     require(false, ERROR_CHALLENGE_REJECTED);
211
212 /**
213 */

spike*  ↗ 0 ▲ 7  Ln 172, Col 54  Spaces: 4  UTF-8  LF  Solidity  ☺  📡
```

github.com/aragon/srvp

Voting v2: TLDR

- EVM Storage Proofs
 - Use any ERC20 to vote (no staking required)
 - Bring any existing token to Aragon for governance
 - github.com/aragon/evm-storage-proofs
- Voting Relay Protocols
 - Aggregate votes off-chain and relay partial tallies
 - Voters don't need ETH, and costs are reduced by 10-1000x
 - github.com/aragon/svrp
- Research discussions: forum.aragon.org
- Come work with us: wiki.aragon.org/jobs



Thanks

aragon.org

@AragonProject @AragonOneTeam @izqui9