

Evaluating NISQ Devices with Quadratic Nonresidues

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An unsolved problem since Gauss



Quadratic nonresidue problem (QNR):

Given a prime $p \equiv 1 \pmod 8$, find a y such that $x^2 \equiv y \pmod p$ has no solution.

Question: Is QNR in P?

Gauss [1] proved the first nontrivial upper bound for the least quadratic nonresidue showing that $y < 2\sqrt{p} + 1$. Current best analytic tools prove that $y < C \cdot p^\alpha$ for a non-zero α [2, p. 33].

QNR is in EQP_C

This new result shows that a quantum computer can solve QNR in quantum P.

Given $p \equiv 1 \pmod 8$, choose the smallest n such that $p < 2^n = N$. Let $\theta = \arccos\left(1 - \frac{2^n}{p-1}\right)$,

and $f(x) = \left[\left(\frac{x}{p}\right) = -1\right]$ and $0 \leq x < p$.

$[O(n)]$ Apply $H^{\otimes n}$ to $|0\rangle^{\otimes n}$ (Hadamard transform).

$$\frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} |x\rangle$$

$[O(n \log^2 n)]$ Compute Jacobi symbol indicator [3, 4].

$$\frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} |x\rangle \left[\left[\left(\frac{x}{p} \right) = -1 \right] \right]$$

$[O(n)]$ Compute $[x < p]$ indicator [5].

$$\frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} |x\rangle \left[\left[\left(\frac{x}{p} \right) = -1 \right] \right] | [x < p] \rangle$$

$[O(1)]$ Rotate odd QNRs less than p by -2θ .

$$\frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} e^{-i2\theta f(x)x_0} |x\rangle \left[\left[\left(\frac{x}{p} \right) = -1 \right] \right] | [x < p] \rangle$$

$[O(1)]$ Rotate all QNRs less than p by θ .

$$\frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} e^{i\theta f(x)(1-2x_0)} |x\rangle \left[\left[\left(\frac{x}{p} \right) = -1 \right] \right] | [x < p] \rangle$$

$[O(n \log^2 n)]$ Uncompute indicator functions.

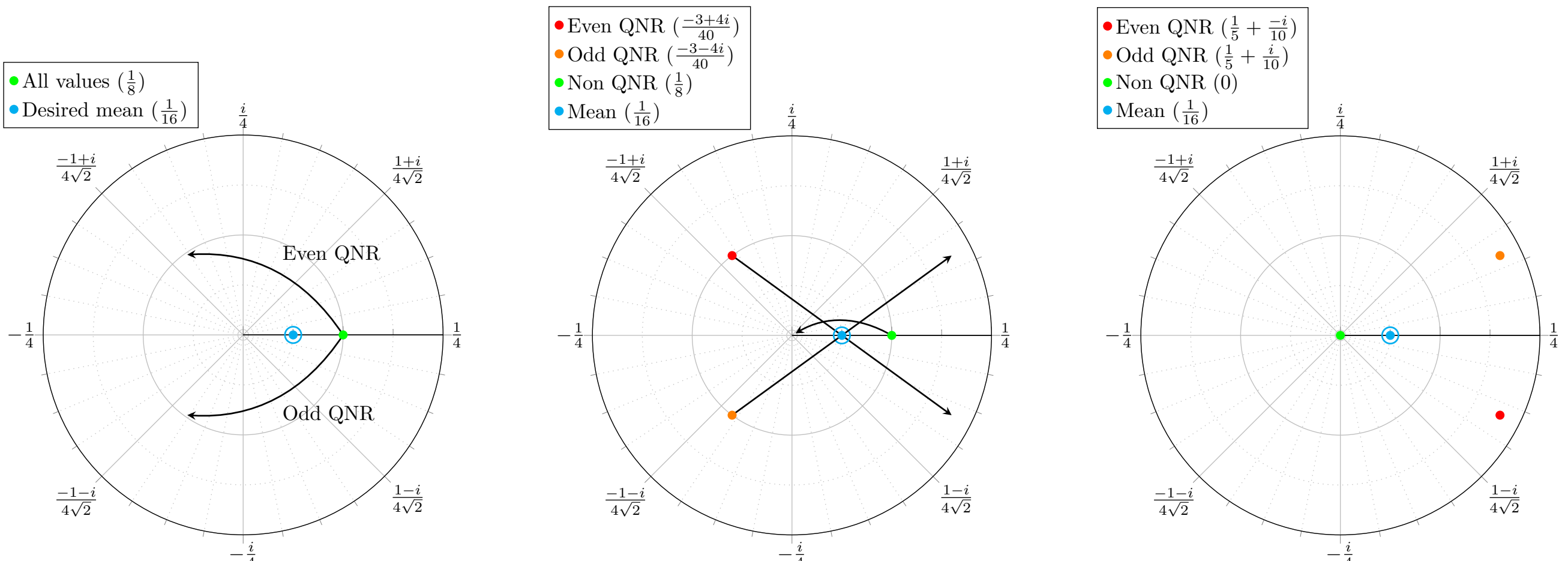
$$\frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} e^{i\theta f(x)(1-2x_0)} |x\rangle$$

$[O(n)]$ Use a Grover step to invert about the mean $\alpha = \frac{1}{2\sqrt{N}}$.

$$\frac{1}{\sqrt{N}} \sum_{x=0}^{N-1} \left(1 - e^{i\theta f(x)(1-2x_0)} \right) |x\rangle$$

$[O(n)]$ Observe a quadratic nonresidue modulo p .

Phase inversion in the QNR algorithm



Amplitude values for Quadratic Nonresidues for $p = 41$

Creating a NISQ test from the QNR algorithm

Using a single Jacobi symbol calculation, a quantum computer can find a QNR 100% of the time, whereas a classical computer can only succeed 75% of the time.

Even if we want to argue for a different classical bound, without a mathematical breakthrough, the provable success rate of any algorithm in P will always be less than 100%.

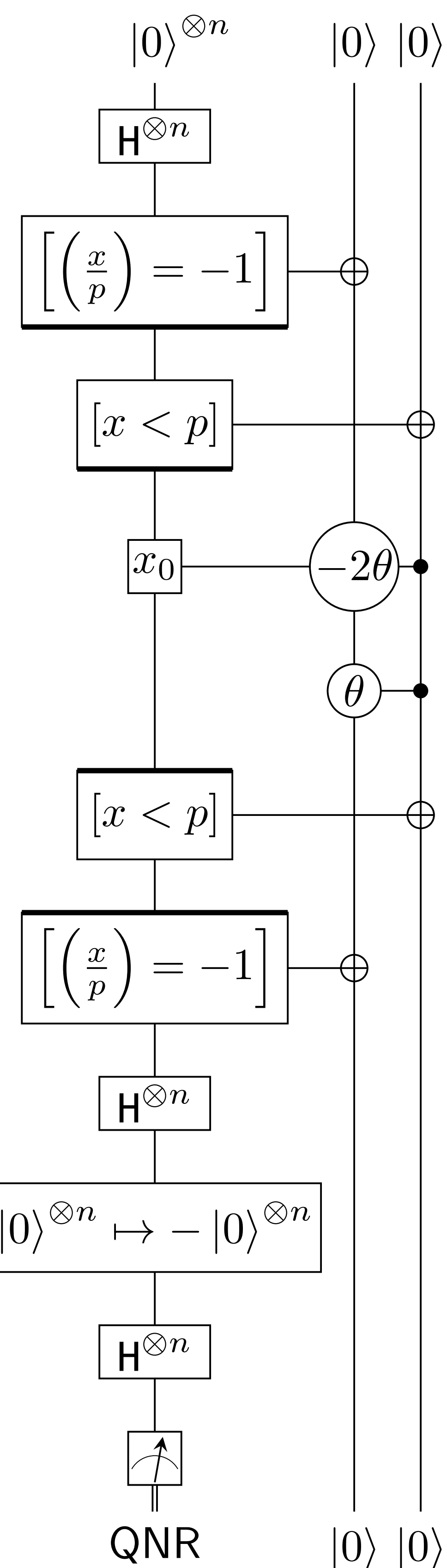
The QNR algorithm evaluates two NISQ properties:

- The rate of success
- The uniformity of the observations

QNR Test Advantages:

- Math inspired and implementation agnostic
- Infinite tests with $O(n \log^2 n)$ runtime
- Smallest tests are usable now
- Scores algorithmic success instead of a physical property

General QNR algorithm



Hadamard transform

Jacobi symbol indicator

Indicator for $[x < p]$

Rotate odd QNR less than p by -2θ

Rotate all QNR less than p by θ

Uncompute $[x < p]$

Uncompute Jacobi symbol

Hadamard transform

Grover gate

Hadamard transform

Measure random QNR

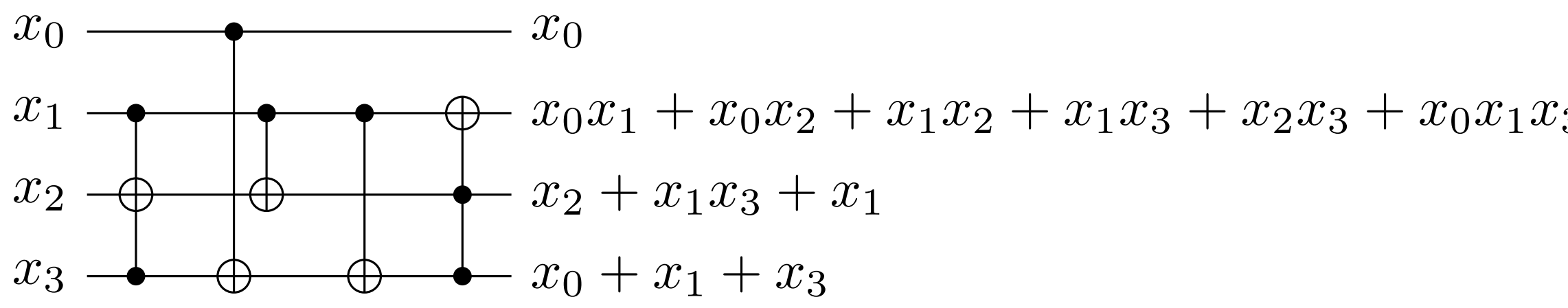
Note than when $p = 2^m + 1$, the inequality indicator $[x < p]$ can be left out since there are exactly 2^m nonzero residues. Additionally, since exactly half of the values less than 2^m will be nonresidues, a permutation of $0, 1, \dots, 2^m - 1$ will exist where one of the output wires is the parity bit, and another is the indicator function for the Jacobi symbol. This only happens for Fermat primes where $p = 2^{2^m} + 1$.

Design a QNR circuit for $p = 17$

Since $p = 17$ is a Fermat prime, the Jacobi symbol indicator is balanced.

$$\left[\left(\frac{x}{17} \right) = -1 \right] = x_0x_1 + x_0x_2 + x_1x_2 + x_1x_3 + x_2x_3 + x_0x_1x_3$$

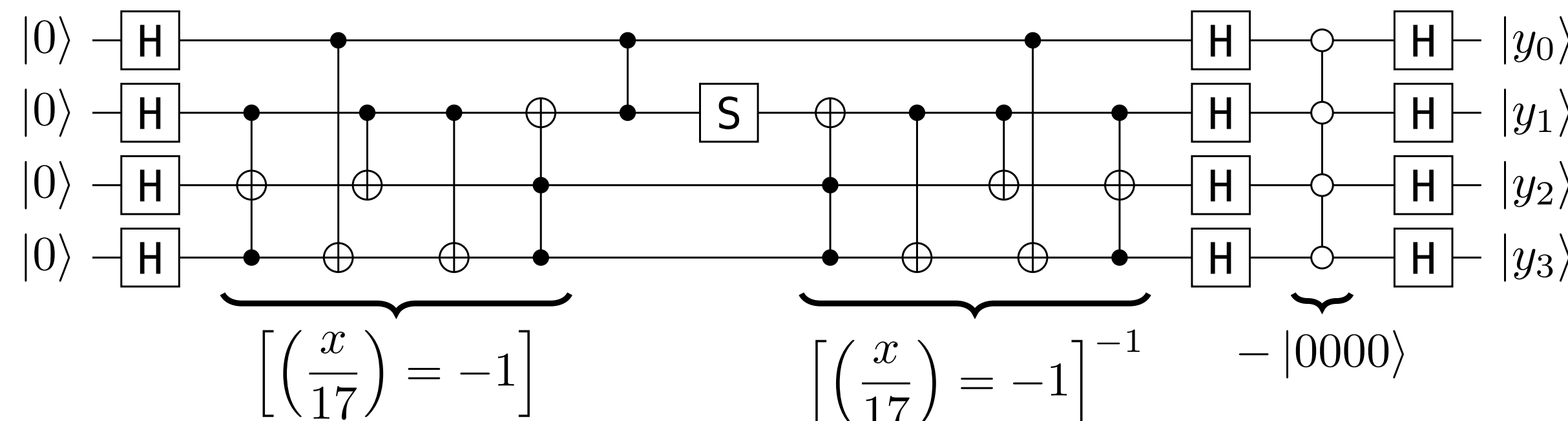
Jacobi symbol indicator permutation for $p = 17$



Ideally, the Jacobi symbol circuit should be derived from an algorithm, and not by computer search. Given current qubit fidelity levels, a full algorithm, even the $p = 17$ test is too hard for current NISQ devices to escape the noise floor.

As NISQ devices improve, a full algorithmic test for $p = 17$ will be meaningful. The next tests of interest will likely be $p = 41$ and $p = 257$.

Basic QNR circuit for $p = 17$



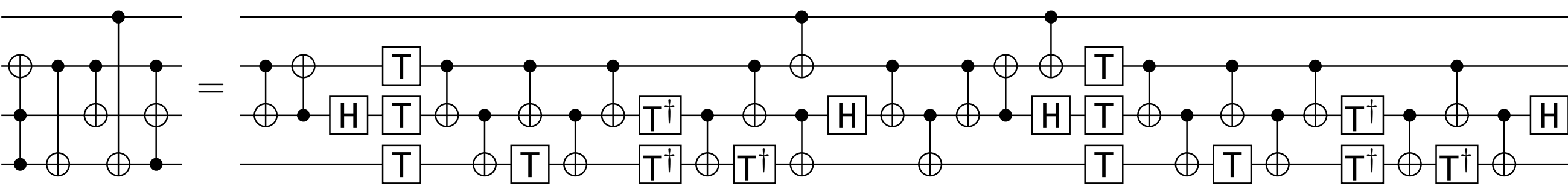
Running the QNR17 circuit on current NISQ devices

Two circuit properties are required to run on all of the tested NISQ devices:

- Only single qubit gates and CNOTs.
- All CNOTs must be linear nearest neighbor.

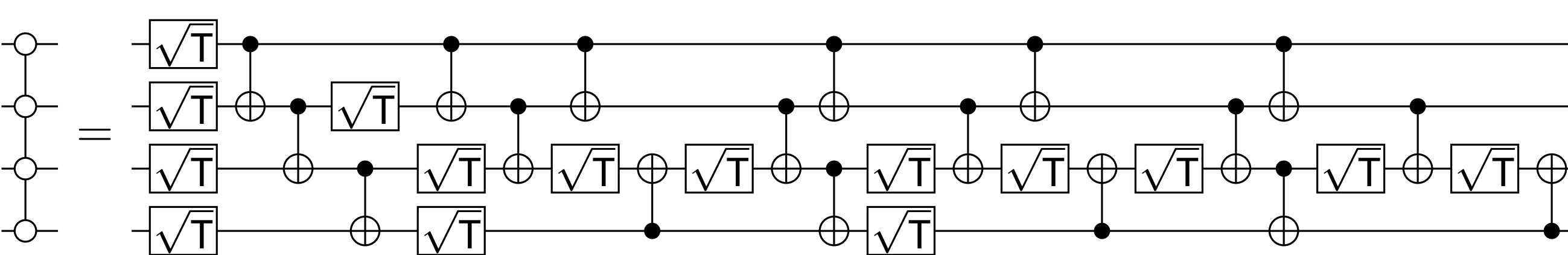
Inverse indicator permutation decomposition

Since a CNOT or CCNOT targeting a $|+\rangle$ does nothing, the entire forward pass of the indicator permutation can be removed. The backwards pass was constructed using ideas from [6].



Grover gate decomposition

A new method was discovered to flip the $|0000\rangle$ state without using an extra qubit.



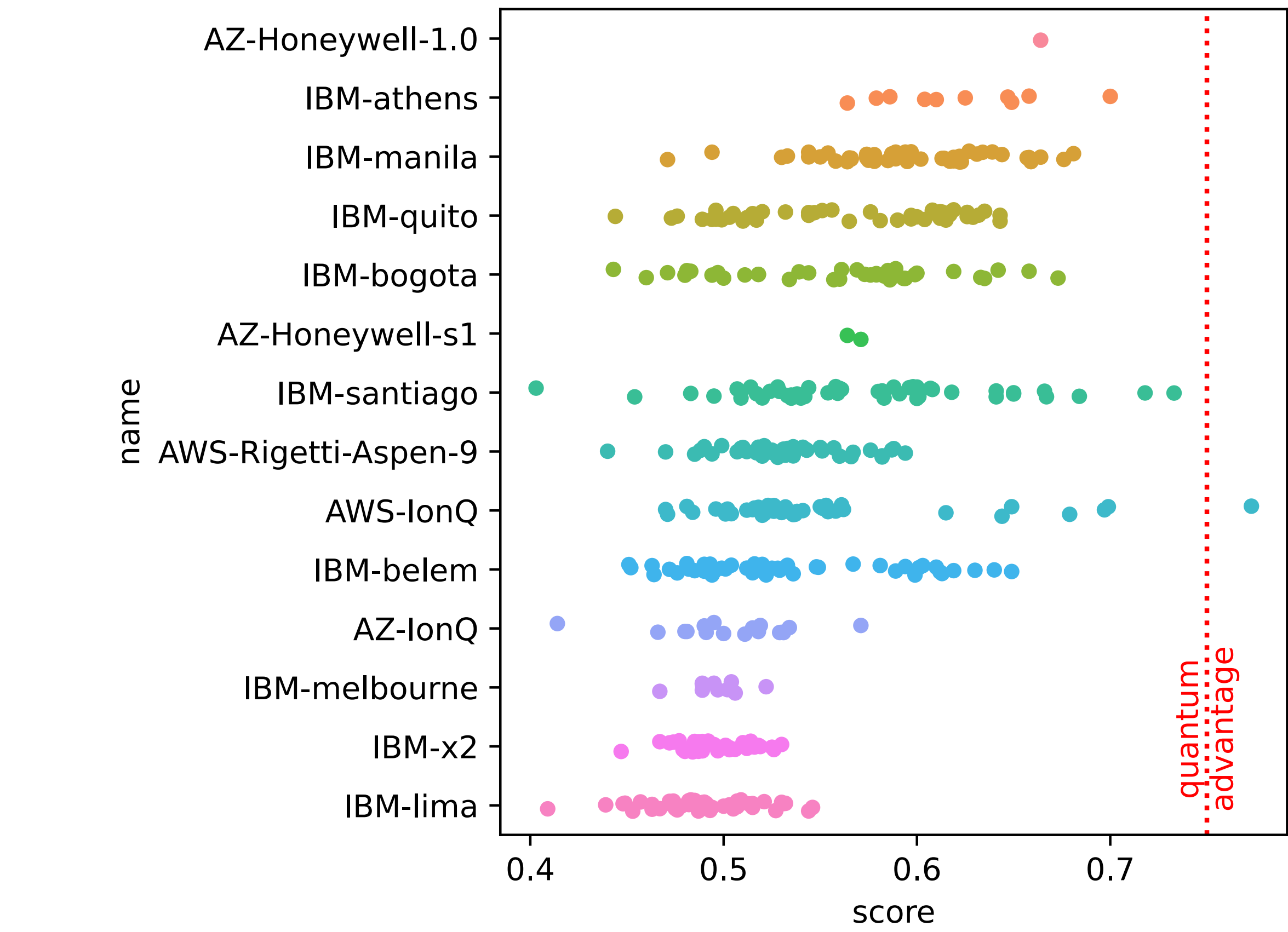
The vertical circuit on the far right shows the final transpiled circuit and a hyperlink to a Quirk circuit running it.

References

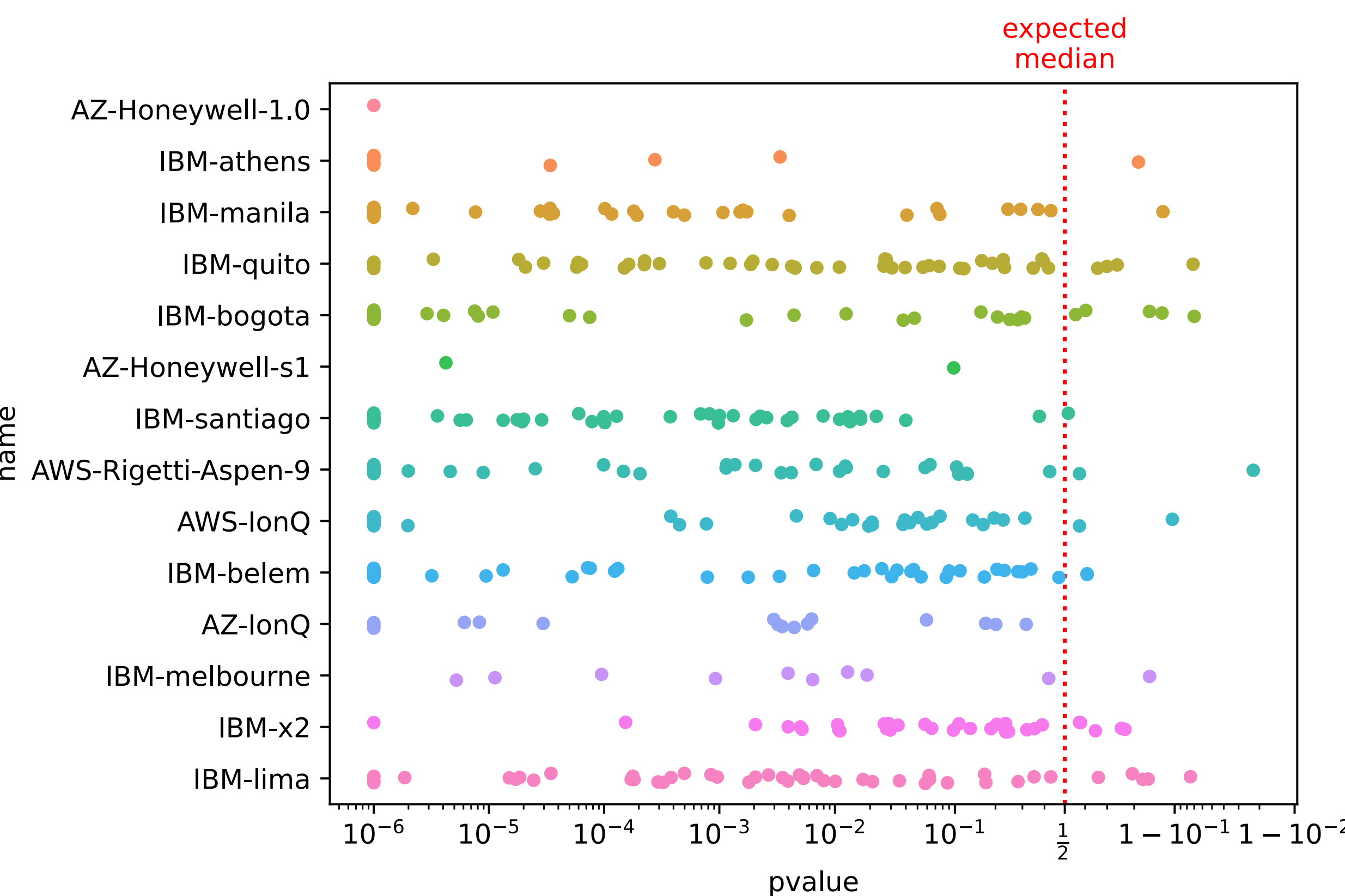
[1] Carl Friedrich Gauss, *Disquisitiones arithmeticae*, 1801.
[2] H. Cohen, *A course in computational algebraic number theory*, Springer-Verlag, Berlin, 1993.
[3] D. Harvey, J. Van Der Hoeven, *Integer multiplication in time $O(n \log n)$* , Annals of Mathematics, Princeton University, Department of Mathematics, 2020.
[4] R. P. Brent and P. Zimmerman, *An $O(M(n) \log n)$ algorithm for the Jacobi symbol*, CoRR, 2010.
[5] D. Oliveira, R. Ramos, *Quantum bit string comparator: Circuits and applications*, Quantum Computers and Computing, Vol.7, 2007.
[6] C. Gidney, *Minimum number of CNOTs for Toffoli with non-adjacent controls (answer:3964)*, https://quantumcomputing.stackexchange.com, 2018.

Test results for $p = 17$ (Jun-Aug 2021)

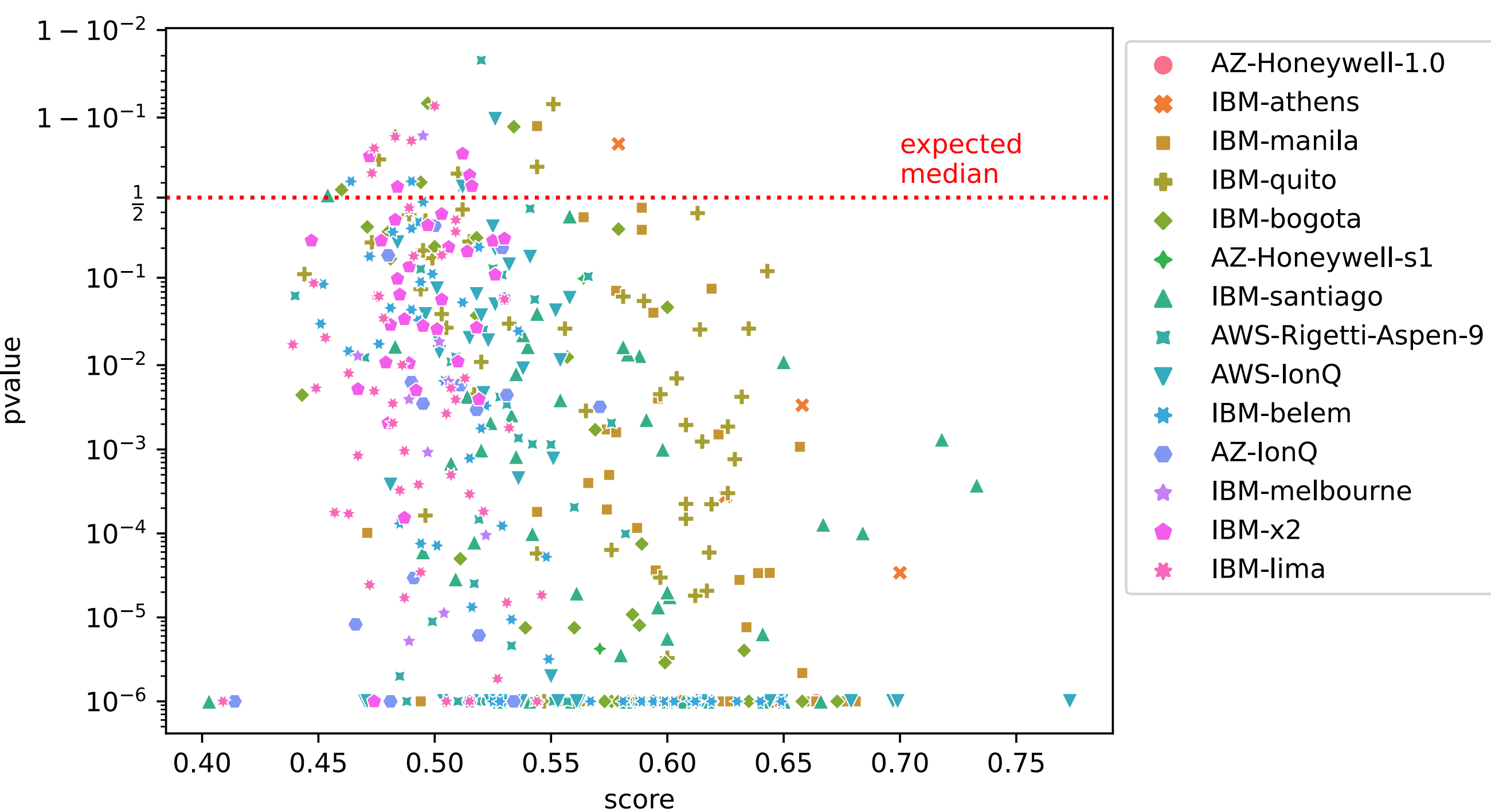
Success rate of 1000 shot runs



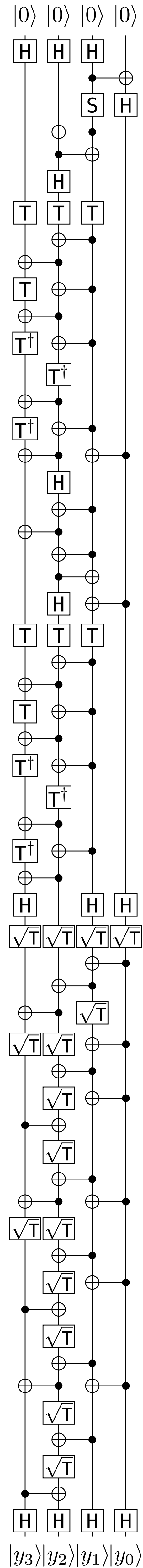
p -value of uniformity test for 1000 shot runs



Success rate vs. p -value



QNR17 test



Quirk circuit
tinyurl.com/
3zpvpdft

