

**CMPS 101 - Programming assignment 3 - W13**  
**Solving the Jumping Pegs Problem**  
**via Backtracking**  
**while using Hashing for Memoization**  
 Handed out Tu 2-19-13    Due Th 2-28-13 in locker

February 22, 2013

1. Design a program for finding a solution to the Jumping Pegs Problem via Backtracking
2. Use Hashing to avoid solving the same board a second time.

Rules of the problem

- | denotes a peg
- X denotes a double peg
- A peg | must jump over 2 pegs (over| | or over X) and land on a peg
- $\begin{array}{ccccccc} | & | & | & | & \text{---}> & | & | & X \\ \wedge & & & & & \wedge & & \\ | & X & | & & \text{---}> & X & X \\ \wedge & & & & & \wedge & \end{array}$
- A peg can jump either left or right
- Double pegs can't move and spaces are ignored
- Cant land on an empty spot or double peg.

Input: Start configuration of single and double pegs

Question: Is there a sequence of jumps s.t. all pegs are doubled?

—If not then output “no solution”.

—If yes then output a sequence of boards or moves shows how the solution was achieved.

Begin by:

1. Solving small problems by hand.  
Try to solve the following:

```

| | X | | | | |      ---> no solution
| | | | | |          ---> no solution
| | | | | | | | | |  ---> X X X X X

```

2. Write a recursive backtracking routine for solving the problem

Hashing:

1. Represent the “board” as a sequence of bits: 0 encodes | and 1 encodes X.  
When empty spots appear then you need to left shift.
2. Interpret bit sequences as keys represented in binary that you can hash on.
3. Whenever a board is known to have no solution, then hash it.
4. Before recursing on a new board, check whether it is not in your hash table.  
If it is, then you can skip this recursion.
5. You only need to implement **Insert** and **Search** but not **Delete**.
6. Use open addressing. Keep track of your load factor  $\alpha$ . If  $\alpha > .2$ , then rehash the whole table content into one **4** times as larger. For the initial hash table size we suggest to use:  $m = 2^{12} = 4096$ . You are welcome to vary the **.2**, **4**, and  $2^{12}$  constants in sensible ways.
7. Start with the division method and linear probing.
8. Once this works, then use double hashing:  $h(k, i) = (h_1(k) + h_2(k)) \bmod m$ .  
 $\text{frac} := kA \bmod 1, \text{mfrac} := m \text{frac}, h_1(k) := \lfloor \text{mfrac} \rfloor, h_2(k) := \lfloor m(\text{mfrac} \bmod 1) \rfloor$ .
9. Finally implement the enlarging your hash table procedure.
10. Input format: A sequence of bits.