# Reference Manual

PROGRAMMING LANGUAGE



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#### 1. INTRODUCTION

The design goal of Orange is to create a new language that is as simple as peeling an orange. Creating this new language will enable users to get an idea of how to implement the language methodically in a simple manner.

The first step in building Orange is to define the grammar of the language for which we use Flex as lexer/scanner and Bison as a parser.

Orange is a high-level language which is compiled, interpreted and also which generates intermediate code. This intermediate code is read line by line and executed in runtime environment.

# 2. COMPILING AND RUNNING ORANGE PROGRAMS

Orange source file should have .or as extension Intermediate file has .pulp as extension

## 3. OPERATORS OF ORANGE

OPERATORS	DESCRIPTION
"+"	Addition
"_"	Subtraction
٠٠**	Multiplication
··/>>	Division
"AND"	AND
"OR"	OR
"NOT"	NOT
"->"	Assignment operator
"~"	delimiter
"->->"	Comparison operator
"!->"	Not Equal to
"->>"	Greater than equal to
"-><"	Leasser than equal to
">"	Greater than
"<"	Lesser than
"YES"	True
"NO"	False
"LP"	Left parentheses
"RP"	Right parentheses
"{"	Block begin
"}"	Block end

### **Keywords**

Int	Integer
Float	Float
Boo	Boolean
Read	Input
write	output
nl	New line
is	if
isnot	else
stack	Stack
push	Push
pop	Pop
loop	while loop

#### 4. LEXICAL ANALYSIS

Lexical analysis is the first phase of compilation.

During this phase, source code received character-by-character is transformed into a sequence of "tokens."

For example, for the following expression in our language: write  $(3 + x *_2) \sim$ 

```
the resulting stream of tokens might be:
(keyword "write")
(delim "(")
(int 3)
(operator "+")
(identifier "x")
```

(operator "\*")

(int 2)

(delim ")")

(delim "~")

#### 5. SIMPLE ORANGE PROGRAMS

#### **I/O Operation**

Our language takes an input from the user using the "read" statement and displays output using the "write" statement as shown below.

```
int a~
read a~
write a~
```

#### **Output**

```
Enter value : 2

2

Expected intermediate code for the above is: decl a put a get a dsp end
```

#### **Conditional statements**

```
a=4~
is LP a > o RP
{
    write a
}
```

#### Output:

4

Expected Intermediate code for the above is:

```
get 4
put a
get a
get o
grt
bne 8
get a
dsp
end
```

#### **Looping construct**

Our language uses "loop" to define while loop as shown below.

```
read a~
loop LP a > 1 RP~

{~
    write a~
    a = a - 1;
}~
```

## **Output:**

Enter value: 2

2

Expected intermediate code for the above is:

amc

put a

get a

get 1

grt

get a

dsp

sub

get a

get 1

put a

bne 8

end

#### 6. GRAMMAR:

Parsing is done in accordance with BNF rules and using the following grammar in our language:

```
<PROGRAM> \rightarrow <BLOCK>
   \langle BLOCK \rangle \rightarrow \{\langle read \rangle '\sim' | \langle write \rangle '\sim' | \langle while \rangle '\sim' | \langle statement \rangle '\sim' | \langle write \rangle '\sim
   | <var> '~'} | <assignment>
   <assignment> → <identifier> '-->' <exp>
   <while> \rightarrow { 'loop' 'LP' <exp> 'RP' '{' <BLOCK> '}' }
   <var> → {'int'|'float'|'boo'} <identifier>
   \langle if \rangle \rightarrow \{ \text{ 'is' 'LP' } \langle exp \rangle \text{ 'RP' } \langle BLOCK \rangle \text{ ['isnot' } \langle BLOCK \rangle ] \}
   <read> \rightarrow 'read' <exp> '\sim'
   <write> → 'write' <exp> '~'
   \exp \rightarrow \{\exp_1 + \exp_2 \mid \exp_1 - \exp_2 \mid \exp_1 * \exp_2 \mid \exp_1 / \exp_2 \mid "("\exp_1")" \mid \exp_2 \mid \exp_1 + \exp_2 \mid "("exp_1")" \mid exp_2 \mid exp_1 + exp_2 \mid exp
   identifier | int | float | boo} '~'
   <identifier> \rightarrow [a-zA-Z]
 int \rightarrow \{D\}\{D\}^*
float \rightarrow \{D\} + "." + ?\{D\}^*
boo → 'YES' | 'NO'
D \rightarrow [o-9]
```